



# Assembly Language: IA-32 Instructions

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## Goals of this Lecture

- **Help you learn how to:**
  - Manipulate data of various sizes
  - Leverage more sophisticated addressing modes
  - Use condition codes and jumps to change control flow
  - ... and thereby ...
  - Write more efficient assembly-language programs
  - Understand the relationship to data types and common programming constructs in high-level languages
- **Focus is on the assembly-language code**
  - Rather than the layout of memory for storing data
  - Precepts will cover that, assembler directives, etc.

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## Variable Sizes in High-Level Language



- C data types vary in size
  - Character: 1 byte
  - Short, int, and long: varies, depending on the computer
  - Float and double: varies, depending on the computer
  - Pointers: typically 4 bytes
- Programmer-created types
  - Struct: arbitrary size, depending on the fields
- Arrays
  - Multiple consecutive elements of some fixed size
  - Where each element could be a struct

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## Supporting Different Sizes in IA-32



- Three main data sizes
  - Byte (b): 1 byte
  - Word (w): 2 bytes
  - Long (l): 4 bytes
- Separate assembly-language instructions
  - E.g., addb, addw, and addl
- Separate ways to access (parts of) a register
  - E.g., %ah or %al, %ax, and %eax
- Larger sizes (e.g., struct)
  - Manipulated in smaller byte, word, or long units

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## Byte Order in Multi-Byte Entities



- Intel is a **little endian** architecture
  - Least significant byte of multi-byte entity is stored at lowest memory address
  - “Little end goes first”

The int 5 at address 1000:

1000	00000101
1001	00000000
1002	00000000
1003	00000000

- Some other systems use **big endian**
  - Most significant byte of multi-byte entity is stored at lowest memory address
  - “Big end goes first”

The int 5 at address 1000:

1000	00000000
1001	00000000
1002	00000000
1003	00000101

## Little Endian Example



```
int main(void) {
    int i=0x003377ff, j;
    unsigned char *p = (unsigned char *) &i;
    for (j=0; j<4; j++)
        printf("Byte %d: %x\n", j, p[j]);
}
```

Output on a  
little-endian  
machine

Byte 0: ff  
Byte 1: 77  
Byte 2: 33  
Byte 3: 0

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## IA-32 General Purpose Registers



31	15	8	7	0	16-bit	32-bit
	AH	AL			AX	EAX
	BH	BL			BX	EBX
	CH	CL			CX	ECX
	DH	DL			DX	EDX
		SI				ESI
		DI				EDI

General-purpose registers

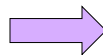
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## C Example: One-Byte Data



Global *char* variable *i* is in *%al*,  
the *lower byte* of the “A” register.

```
char i;  
...  
if (i > 5) {  
    i++;  
else  
    i--;  
}
```



```
    cmpb $5, %al  
    jle else  
    incb %al  
    jmp endif  
else:  
    decb %al  
endif:
```

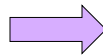
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## C Example: Four-Byte Data



Global *int* variable *i* is in *%eax*,  
the *full 32 bits* of the “A” register.

```
int i;  
...  
if (i > 5) {  
    i++;  
else  
    i--;  
}
```



```
    cmpl $5, %eax  
    jle else  
    incl %eax  
    jmp endif  
else:  
    decl %eax  
endif:
```

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## Loading and Storing Data



- Processors have many ways to access data
  - Known as “addressing modes”
  - Two simple ways seen in previous examples
- Immediate addressing
  - Example: `movl $0, %ecx`
  - Data (e.g., number “0”) embedded in the instruction
  - Initialize register ECX with zero
- Register addressing
  - Example: `movl %edx, %ecx`
  - Choice of register(s) embedded in the instruction
  - Copy value in register EDX into register ECX

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## Accessing Memory



- Variables are stored in memory
  - Global and static local variables in Data or BSS section
  - Dynamically allocated variables in the heap
  - Function parameters and local variables on the stack
- Need to be able to load from and store to memory
  - To manipulate the data directly in memory
  - Or copy the data between main memory and registers
- IA-32 has many different addressing modes
  - Corresponding to common programming constructs
  - E.g., accessing a global variable, dereferencing a pointer, accessing a field in a struct, or indexing an array

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## Direct Addressing



- Load or store from a particular memory location
  - Memory address is embedded in the instruction
  - Instruction reads from or writes to that address
- IA-32 example: `movl 2000, %ecx`
  - Four-byte variable located at address 2000
  - Read four bytes starting at address 2000
  - Load the value into the ECX register
- Useful when the address is known in advance
  - Global variables in the Data or BSS sections
- Can use a label for (human) readability
  - E.g., “i” to allow “`movl i, %eax`”

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## Indirect Addressing



- Load or store from a previously-computed address
  - Register with the address is embedded in the instruction
  - Instruction reads from or writes to that address
- IA-32 example: `movl (%eax), %ecx`
  - EAX register stores a 32-bit address (e.g., 2000)
  - Read long-word variable stored at that address
  - Load the value into the ECX register
- Useful when address is not known in advance
  - Dynamically allocated data referenced by a pointer
  - The “(%eax)” essentially dereferences a pointer

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## Base Pointer Addressing



- Load or store with an offset from a base address
  - Register storing the base address
  - Fixed offset also embedded in the instruction
  - Instruction computes the address and does access
- IA-32 example: `movl 8(%eax), %ecx`
  - EAX register stores a 32-bit base address (e.g., 2000)
  - Offset of 8 is added to compute address (e.g., 2008)
  - Read long-word variable stored at that address
  - Load the value into the ECX register
- Useful when accessing part of a larger variable
  - Specific field within a “struct”
  - E.g., if “age” starts at the 8<sup>th</sup> byte of “student” record

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## Indexed Addressing



- Load or store with an offset and multiplier
  - Base address in a register
  - Fixed displacement embedded in the instruction
  - Offset computed by multiplying 2<sup>nd</sup> register with constant
  - Instruction computes the address and does access
- IA-32 example: `movl 2000(,%eax,4), %ecx`
  - No base register. Index register EAX (say, value of 10)
  - Multiplied by a multiplier of 1, 2, 4, or 8 (say, 4)
  - Added to a fixed displacement of 2000 (say, to get 2040)
- Useful to iterate through an array (e.g., `a[i]`)
  - Displacement is the start of the array (i.e., “a”); use register if need pointer dereferencing
  - Register is the index (i.e., “i”)
  - Multiplier is the size of the element (e.g., 4 for “int”)

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## Indexed Addressing Example



```
int a[20]; ← global variable
int i, sum=0;
for (i=0; i<20; i++)
    sum += a[i];
```

EAX: i  
EBX: sum  
ECX: temporary

```
movl $0, %eax
movl $0, %ebx
sumloop:
movl a(,%eax,4), %ecx
addl %ecx, %ebx
incl %eax
cmpl $19, %eax
jle sumloop
```



## Effective Address: More Generally



$$\text{Offset} = \left( \begin{array}{c} \text{eax} \\ \text{ebx} \\ \text{ecx} \\ \text{edx} \\ \text{esp} \\ \text{ebp} \\ \text{esi} \\ \text{edi} \end{array} \right) + \left( \begin{array}{c} \text{eax} \\ \text{ebx} \\ \text{ecx} \\ \text{edx} \\ \text{esp} \\ \text{ebp} \\ \text{esi} \\ \text{edi} \end{array} \right) * \left( \begin{array}{c} 1 \\ 2 \\ 4 \\ 8 \end{array} \right) + \left( \begin{array}{c} \text{None} \\ \text{8-bit} \\ \text{16-bit} \\ \text{32-bit} \end{array} \right)$$

Base                      Index                      scale                      displacement

- Displacement `movl foo, %ebx`
- Base `movl (%eax), %ebx`
- Base + displacement `movl foo(%eax), %ebx`  
`movl 1(%eax), %ebx`
- (Index \* scale) + displacement `movl foo(,%eax,4), %ebx`
- Base + (index \* scale) + displacement `movl foo(%edx,%eax,4), %ebx`

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## Data Access Methods: Summary



- **Immediate addressing:** data stored in the instruction itself
  - `movl $10, %ecx`
- **Register addressing:** data stored in a register
  - `movl %eax, %ecx`
- **Direct addressing:** address stored in instruction
  - `movl foo, %ecx`
- **Indirect addressing:** address stored in a register
  - `movl (%eax), %ecx`
- **Base pointer addressing:** includes an offset as well
  - `movl 4(%eax), %ecx`
- **Indexed addressing:** instruction contains base address, and specifies an index register and a multiplier (1, 2, 4, or 8)
  - `movl 2000(%eax,1), %ecx`

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## Control Flow



- Common case
  - Execute code sequentially
  - One instruction after another
- Sometimes need to change control flow
  - If-then-else
  - Loops
  - Switch
- Two key ingredients
  - Testing a condition
  - Selecting what to run next based on result

```
    cmpl $5, %eax
    jle else
    incl %eax
    jmp endif
else:
    decl %eax
endif:
```

## Condition Codes



- 1-bit registers set by arithmetic & logic instructions
  - ZF: Zero Flag
  - SF: Sign Flag
  - CF: Carry Flag
  - OF: Overflow Flag
- Example: “addl Src, Dest” (“t = a + b”)
  - ZF: set if t == 0
  - SF: set if t < 0
  - CF: set if carry out from most significant bit
    - *Unsigned* overflow
  - OF: set if two’s complement overflow
    - (a>0 && b>0 && t<0)
    - || (a<0 && b<0 && t>=0)

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## Condition Codes (continued)



- Example: “`cmpl Src2,Src1`” (compare b,a)
  - Like computing a-b without setting destination
  - ZF: set if a == b
  - SF: set if (a-b) < 0
  - CF: set if carry out from most significant bit
    - Used for unsigned comparisons
  - OF: set if two's complement overflow
    - (a>0 && b<0 && (a-b)<0) || (a<0 && b>0 && (a-b)>0)
- Flags are *not* set by `lea`, `inc`, or `dec` instructions
  - Hint: this is useful for the extra-credit part of the assembly-language programming assignment! ☺

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## Example Five-Bit Comparisons



- Comparison: `cmp $6, $12`

01100		01100
- 00110	→	<u>+11010</u>
??		00110

  - Not zero: ZF=0 (diff is not 00000)
  - Positive: SF=0 (first bit is 0)
  - No carry: CF=0 (unsigned diff is correct)
  - No overflow: OF=0 (signed diff is correct)
- Comparison: `cmp $12, $6`

00110		00110
- 01100	→	<u>+10100</u>
??		11010

  - Not zero: ZF=0 (diff is not 00000)
  - Negative: SF=1 (first bit is 1)
  - Carry: CF=1 (unsigned diff is wrong)
  - No overflow: OF=0 (signed diff is correct)
- Comparison: `cmp $-6, $-12`

10100		10100
- 11010	→	<u>+00110</u>
??		11010

  - Not zero: ZF=0 (diff is not 00000)
  - Negative: SF=1 (first bit is 1)
  - Carry: CF=1 (unsigned diff of 20 and 28 is wrong)
  - No overflow: OF=0 (signed diff is correct)

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## Jumps after Comparison (cmpl)



- Equality
  - Equal: je (ZF)
  - Not equal: jne ( $\sim$ ZF)
- Below/above (e.g., unsigned arithmetic)
  - Below: jb (CF)
  - Above or equal: jae ( $\sim$ CF)
  - Below or equal: jbe (CF | ZF)
  - Above: ja ( $\sim$ (CF | ZF))
- Less/greater (e.g., signed arithmetic)
  - Less: jl (SF ^ OF)
  - Greater or equal: jge ( $\sim$ (SF ^ OF))
  - Less or equal: jle ((SF ^ OF) | ZF)
  - Greater: jg (!((SF ^ OF) | ZF))

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## Branch Instructions



- Conditional jump
  - $\{l,g,e,ne,\dots\}$  target      if (condition) {eip = target}

Comparison	Signed	Unsigned	
=	e	e	"equal"
≠	ne	ne	"not equal"
>	g	a	"greater,above"
≥	ge	ae	"...-or-equal"
<	l	b	"less,below"
≤	le	be	"...-or-equal"
overflow/carry	o	c	
no ovf/carry	no	nc	

- Unconditional jump
  - jmp target
  - jmp \*register

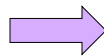
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## Jumping



- Simple model of a “goto” statement
  - Go to a particular place in the code
  - Based on whether a condition is true or false
  - Can represent if-the-else, switch, loops, etc.
- Pseudocode example: If-Then-Else

```
if (Test) {  
    then-body;  
} else {  
    else-body;
```



```
if (!Test) jump to Else;  
then-body;  
jump to Done;  
Else:  
    else-body;  
Done:
```

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## Jumping (continued)



- Pseudocode example: Do-While loop

```
do {  
    Body;  
} while (Test);
```



```
loop:  
    Body;  
    if (Test) then jump to loop;
```

- Pseudocode example: While loop

```
while (Test)  
    Body;
```



```
    jump to middle;  
loop:  
    Body;  
middle:  
    if (Test) then jump to loop;
```

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## Jumping (continued)



- Pseudocode example: For loop

```
for (Init; Test; Update)
    Body
```



```
Init;
if (!Test) jump to done;
loop:
Body;
Update;
if (Test) jump to loop;
done:
```

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## Arithmetic Instructions



- Simple instructions
  - `add{b,w,l} source, dest`      `dest = source + dest`
  - `sub{b,w,l} source, dest`      `dest = dest - source`
  - `inc{b,w,l} dest`              `dest = dest + 1`
  - `dec{b,w,l} dest`              `dest = dest - 1`
  - `neg{b,w,l} dest`              `dest = ~dest + 1`
  - `cmp{b,w,l} source1, source2`      `source2 - source1`
- Multiply
  - `mul (unsigned) or imul (signed)`  
`mul %ebx                    # edx, eax = eax * ebx`
- Divide
  - `div (unsigned) or idiv (signed)`  
`idiv %ebx                   # edx = edx, eax / ebx`
- Many more in Intel manual (volume 2)
  - `adc, sbb, decimal arithmetic instructions`

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## Bitwise Logic Instructions



- Simple instructions

<code>and{b,w,l} source, dest</code>	<code>dest = source &amp; dest</code>
<code>or{b,w,l} source, dest</code>	<code>dest = source   dest</code>
<code>xor{b,w,l} source, dest</code>	<code>dest = source ^ dest</code>
<code>not{b,w,l} dest</code>	<code>dest = ~dest</code>
<code>sal{b,w,l} source, dest (arithmetic)</code>	<code>dest = dest &lt;&lt; source</code>
<code>sar{b,w,l} source, dest (arithmetic)</code>	<code>dest = dest &gt;&gt; source</code>

- Many more in Intel Manual (volume 2)

- Logic shift
- Rotation shift
- Bit scan
- Bit test
- Byte set on conditions

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## Data Transfer Instructions

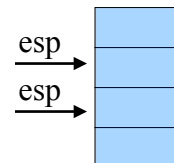


- `mov{b,w,l} source, dest`

- General move instruction

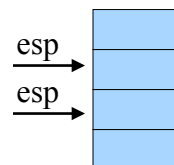
- `push{w,l} source`

```
pushl %ebx    # equivalent instructions
              subl $4, %esp
              movl %ebx, (%esp)
```



- `pop{w,l} dest`

```
popl %ebx    # equivalent instructions
              movl (%esp), %ebx
              addl $4, %esp
```



- Many more in Intel manual (volume 2)

- Type conversion, conditional move, exchange, compare and exchange, I/O port, string move, etc.

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## Conclusions



- **Accessing data**
  - Byte, word, and long-word data types
  - Wide variety of addressing modes
- **Control flow**
  - Common C control-flow constructs
  - Condition codes and jump instructions
- **Manipulating data**
  - Arithmetic and logic operations
- **Next time**
  - Calling functions, using the stack

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