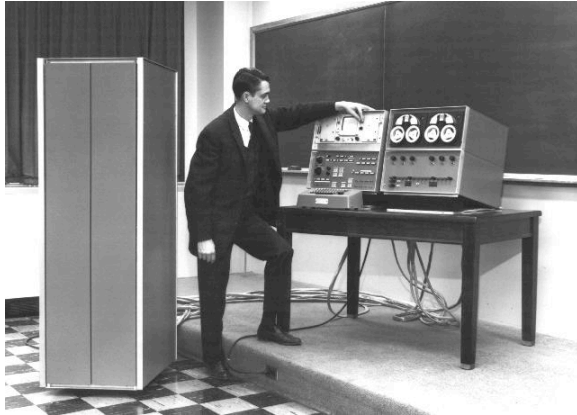


## TOY II



LINC

Introduction to Computer Science · Sedgewick and Wayne · Copyright © 2007 · <http://www.cs.Princeton.EDU/IntroCS>

Data representation. Binary and hex.

TOY.

- Box with switches and lights.
- 16-bit memory locations, 16-bit registers, 8-bit pc.
- $4,328 \text{ bits} = (255 \times 16) + (15 \times 16) + (8) = 541 \text{ bytes!}$
- von Neumann architecture.

TOY instruction set architecture. 16 instruction types.

TOY machine language programs. Variables, arithmetic, loops.



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## What We Do Today

Data representation. Negative numbers.

Input and output. Standard input, standard output.

Manipulate addresses. References (pointers) and arrays.

TOY simulator in Java and implications.



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## Data Representation

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## Digital World

Data is a sequence of bits. (interpreted in different ways)

- Integers, real numbers, characters, strings, ...
- Documents, pictures, sounds, movies, Java programs, ...

Ex. 01110101

- As binary integer:  $1 + 4 + 16 + 32 + 64 = 117$  (base ten).
- As character: 117<sup>th</sup> Unicode character = 'u'.
- As music: 117/256 position of speaker.
- As grayscale value: 45.7% black.



```
public class HelloWorld {
    public static void main(String[] args) {
        System.out.println("Hello, World!");
    }
}
```



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## Adding and Subtracting Binary Numbers

Decimal and binary addition.

$$\begin{array}{r}
 \text{1} \\
 013 \\
 + 092 \\
 \hline
 105
 \end{array}
 \qquad
 \begin{array}{r}
 \text{carries} \\
 \text{1 1} \\
 00001101 \\
 + 01011100 \\
 \hline
 01101001
 \end{array}$$

Subtraction. Add a negative integer.

e.g.,  $6 - 4 = 6 + (-4)$

Q. How to represent negative integers?

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## Representing Negative Integers

TOY words are 16 bits each.

- We could use 16 bits to represent 0 to  $2^{16} - 1$ .
- We want negative integers too.
- Reserving half the possible bit-patterns for negative seems fair.

Highly desirable property. If  $x$  is an integer, then the representation of  $-x$ , when added to  $x$ , is zero.

$$\begin{array}{r}
 x \\
 + (-x) \\
 \hline
 0
 \end{array}
 \qquad
 \begin{array}{r}
 00110100 \\
 + \text{? ? ? ? ? ? ?} \\
 \hline
 00000000
 \end{array}$$

$$\begin{array}{r}
 x \\
 + (-x) \\
 \hline
 0
 \end{array}
 \qquad
 \begin{array}{r}
 00110100 \\
 + 11001011 \\
 \hline
 11111111 \\
 + 1 \\
 \hline
 00000000
 \end{array}$$

-x: flip bits and add 1

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## Two's Complement Integers

To compute  $-x$  from  $x$ :

- Start with  $x$ .

$$\begin{array}{r}
 \text{leading bit determines sign} \\
 +4 \quad 000000000000001000
 \end{array}$$

- Flip bits.

$$\begin{array}{r}
 -5 \quad 111111111111110111
 \end{array}$$

- Add one.

$$\begin{array}{r}
 -4 \quad 111111111111110011
 \end{array}$$

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## Two's Complement Integers

		15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
dec	hex	binary															
+32767	7FFF	0	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1
...																	
+4	0004	0	0	0	0	0	0	0	0	0	0	0	0	0	1	0	0
+3	0003	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	1
+2	0002	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	0
+1	0001	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1
+0	0000	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
-1	FFFF	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1
-2	FFFE	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	0
-3	FFFD	1	1	1	1	1	1	1	1	1	1	1	1	1	1	0	1
-4	FFFC	1	1	1	1	1	1	1	1	1	1	1	1	1	1	0	0
...																	
-32768	8000	1	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0

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## Properties of Two's Complement Integers

### Properties.

- Leading bit (bit 15) signifies sign.
- 0000000000000000 represents zero.
- Negative integer  $-x$  represented by  $2^{16} - x$ .
- Addition is easy.
- Checking for arithmetic overflow is easy.

**Not-so-nice property.** Can represent one more negative integer than positive integer.   
 $32,767 = 2^{15}-1$   $-32,768 = -2^{15}$

**Remark.** Java `int` data type is 32-bit two's complement integer.

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## Representing Other Primitive Data Types in TOY

**Bigger integers.** Use two 16-bit words per `int`.

### Real numbers.

- Use "floating point" (like scientific notation).
- Use four 16-bit words per `double`.

### Characters.

- Use ASCII code (8 bits / character).
- Pack two characters per 16-bit word.

**Note.** Real microprocessors add hardware support for `int` and `double`.

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## Standard Input and Output

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## Standard Output

### Standard output.

- Writing to memory location `FF` sends one word to TOY stdout.
- Ex. `9AFF` writes the integer in register `A` to stdout.

```
00: 0000 0
01: 0001 1

10: 8A00 RA ← mem[00]      a = 0
11: 8B01 RB ← mem[01]      b = 1
12: 9AFF write RA to stdout while(a > 0) {
13: 1AAB RA ← RA + RB       print a
14: 2BAB RB ← RA - RB       a = a + b
15: DA12 if (RA > 0) goto 12 b = a - b
16: 0000 halt
```

fibonacci.toy

standard  
output

```
0000
0001
0001
0002
0003
0005
0008
000D
0015
0022
0037
0059
0090
00E9
0179
0262
03DB
063D
0A18
1055
1A6D
2AC2
452F
6FF1
```

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## Standard Input

### Standard input.

- Loading from memory address `FF` loads one word from TOY stdin.
- Ex. `8AFF` reads an integer from stdin and store it in register `A`.

Ex: read in a sequence of integers and print their sum.

- In Java, stop reading when EOF.
- In TOY, stop reading when user enters `0000`.

```
while (!StdIn.isEmpty()) {
    a = StdIn.readInt();
    sum = sum + a;
}
StdOut.println(sum);
```

```
00: 0000 0
10: 8C00 RC ← mem[00]
11: 8AFF read RA from stdin
12: CA15 if (RA == 0) pc ← 15
13: 1CCA RC ← RC + RA
14: C011 pc ← 11
15: 9CFF write RC
16: 0000 halt
```

```
00AE
0046
0003
0000
00F7
```

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## Standard Input and Output: Implications

### Standard input and output enable you to:

- Get information out of machine.
- Put information from real world into machine.
- Process more information than fits in memory.
- Interact with the computer while it is running.

## Pointers

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## Load Address (a.k.a. Load Constant)

Load address. [opcode 7]

- Loads an 8-bit integer into a register.
- 7A30 means load the value 30 into register A.

Applications.

- Load a small **constant** into a register.
- Load a 8-bit **memory address** into a register.

register stores "pointer" to a memory cell

```
a = 0x30;
```

Java code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	1	1	1	0	1	0	0	0	1	1	0	0	0	0
7 <sub>16</sub>				A <sub>16</sub>				3 <sub>16</sub>				0 <sub>16</sub>			
opcode				dest d				addr							

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## Arrays in TOY

TOY main memory is a giant array.

- Can access memory cell 30 using load and store.
- 8C30 means load mem[30] into register C.
- Goal: access memory cell *i* where *i* is a variable.

...	...
30	0000
31	0001
32	0001
33	0002
34	0003
35	0005
36	0008
37	000D
...	...

TOY memory

Load indirect. [opcode A]

a variable index

- AC06 means load mem[R6] into register C.

Store indirect. [opcode B]

a variable index

- BC06 means store contents of register C into mem[R6].

```
for (int i = 0; i < N; i++)
    a[i] = StdIn.readInt();
for (int i = 0; i < N; i++)
    StdOut.println(a[N-i-1]);
```

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## TOY Implementation of Reverse

TOY implementation of reverse.

- Read in a sequence of integers and store in memory 30, 31, 32, ...
- Stop reading if 0000.
- Print sequence in reverse order.

```

10: 7101  R1 ← 0001          constant 1
11: 7A30  RA ← 0030          a[]
12: 7B00  RB ← 0000          n

13: 8CFF  read RC              while(true) {
14: CC19  if (RC == 0) goto 19    c = StdIn.readInt();
15: 16AB  R6 ← RA + RB           if (c == 0) break;
16: BC06  mem[R6] ← RC           memory address of a[n]
17: 1BB1  RB ← RB + R1           a[n] = c;
18: C013  goto 13               n++;
                                }

```

read in the data

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## TOY Implementation of Reverse

TOY implementation of reverse.

- Read in a sequence of integers and store in memory 30, 31, 32, ...
- Stop reading if 0000.
- Print sequence in reverse order.

```

19: CB20  if (RB == 0) goto 20 while (n > 0) {
1A: 16AB  R6 ← RA + RB           address of a[n]
1B: 2661  R6 ← R6 - R1           address of a[n-1]
1C: AC06  RC ← mem[R6]           c = a[n-1];
1D: 9CFF  write RC               StdOut.println(c);
1E: 2BB1  RB ← RB - R1           n--;
1F: C019  goto 19                }
20: 0000  halt

```

print in reverse order

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## Unsafe Code at any Speed

Q. What happens if we make array start at 00 instead of 30?

```
10: 7101 R1 ← 0001      constant 1
11: 7A00 RA ← 0000      a[]
12: 7B00 RB ← 0000      n

13: 8CFF read RC          while(true) {
14: CC19 if (RC == 0) goto 19    c = StdIn.readInt();
15: 16AB R6 ← RA + RB          if (c == 0) break;
16: BC06 mem[R6] ← RC          address of a[n]
17: 1BB1 RB ← RB + R1          a[n] = c;
18: C013 goto 13              n++;
                               }
```

```
% more crazy8.txt
1 1 1 1 1 1 1 1
1 1 1 1 1 1 1 1
8888 8810
98FF C011
```

A. With enough data, becomes a **self-modifying program**

- can overflow buffer
- **and** run arbitrary code!

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## What Can Happen When We Lose Control (in C or C++)?

Buffer overrun.

- Array `buffer[]` has size 100.
- User might enter 200 characters.
- Might lose control of machine behavior.

```
#include <stdio.h>
int main(void) {
    char buffer[100];
    scanf("%s", buffer);
    printf("%s\n", buffer);
    return 0;
}
```

Consequences. Viruses and worms.

unsafe C program



22

## What Can Happen When We Lose Control (in C or C++)?

Buffer overrun.

- Array `buffer[]` has size 100.
- User might enter 200 characters.
- Might lose control of machine behavior.

```
#include <stdio.h>
int main(void) {
    char buffer[100];
    scanf("%s", buffer);
    printf("%s\n", buffer);
    return 0;
}
```

unsafe C program

Consequences. Viruses and worms.

## Buffer Overrun Example: JPEG of Death

Microsoft Windows JPEG bug. [September, 2004]

- Step 1. User views malicious JPEG in IE or Outlook.
- Step 2. Machine is Owned.
- Data becomes code by exploiting buffer overrun in GDI+ library.



Fix. Update old library with patched one.

but many applications install independent copies of GDI library

Java enforces security.

- Type safety.
- Array bounds checking.
- Not foolproof.



shine 50W bulb at DRAM  
[Appel-Govindavajhala '03]

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Moral.

- Not easy to write error-free software.
- Embrace Java security features.
- Don't try to maintain several copies of the same file.
- Keep your OS patched.

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## Dumping

Q. Work all day to develop operating system. How to save it?

A. Write short program `dump.toy` and run it to dump contents of memory onto tape.

```
00: 7001  R1 ← 0001
01: 7210  R2 ← 0010          i = 10
02: 73FF  R3 ← 00FF

03: AA02  RA ← mem[R2]      do {
04: 9AFF  write RA          a = mem[i]
05: 1221  R2 ← R2 + R1      print a
06: 2432  R4 ← R3 - R2      i++
07: D403  if (R4 > 0) goto 03 } while (i < 255)
08: 0000  halt
```

`dump.toy`

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## Booting

Q. How do you get it back?

A. Write short program `boot.toy` and run it to read contents of memory from tape.

```
00: 7001  R1 ← 0001
01: 7210  R2 ← 0010          i = 10
02: 73FF  R3 ← 00FF

03: 8AFF  read RA           do {
04: BA02  mem[R2] ← RA      read a
05: 1221  R2 ← R2 + R1      mem[i] = a
06: 2432  R4 ← R3 - R2      i++
07: D403  if (R4 > 0) goto 03 } while (i < 255)
08: 0000  halt
```

`boot.toy`

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## TOY Simulator

## TOY Simulator

Goal. Write a program to "simulate" the behavior of the TOY machine.

- TOY simulator in Java.
- TOY simulator in TOY!

```
public class TOY
{
    public static void main(String[] args)
    {
        int pc = 0x10; // program counter
        int[] R = new int[16]; // registers
        int[] mem = new int[256]; // main memory

        // READ .toy FILE into mem[10..]

        while (true)
        {
            int inst = mem[pc++]; // fetch and increment
            // DECODE
            // EXECUTE
        }
    }
}
```

```
% more add-stdin.toy
8C00 ← TOY program to load at 10
8AFF
CA15
1CCA
C011
9CFF
0000

% java TOY add-stdin.toy
00AE ← standard input
0046
0003
0000
00F7 ← standard output
```

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## TOY Simulator: Decode

Ex. Extract destination register of  $1_{CAB}$  by shifting and masking.

0	0	0	1	1	1	0	0	1	0	1	0	1	0	1	1	<code>inst</code>
$1_{16}$				$C_{16}$				$A_{16}$				$B_{16}$				
0	0	0	0	0	0	0	0	0	0	0	1	1	1	0	0	<code>inst &gt;&gt; 8</code>
$0_{16}$				$0_{16}$				1				$C_{16}$				
0	0	0	0	0	0	0	0	0	0	0	0	1	1	1	1	15
$0_{16}$				$0_{16}$				$0_{16}$				$F_{16}$				
0	0	0	0	0	0	0	0	0	0	0	0	1	1	0	0	<code>(inst &gt;&gt; 8) &amp; 15</code>
$0_{16}$				$0_{16}$				0				$C_{16}$				

```
int inst = mem[pc++];           // fetch and increment
int op  = (inst >> 12) & 15;    // opcode  (bits 12-15)
int d   = (inst >> 8) & 15;     // dest d  (bits 08-11)
int s   = (inst >> 4) & 15;     // source s (bits 04-07)
int t   = (inst >> 0) & 15;     // source t (bits 00-03)
int addr = (inst >> 0) & 255;   // addr   (bits 00-07)
```

## TOY Simulator: Execute

```
if (op == 0) break;           // halt

switch (op)
{
    case 1: R[d] = R[s] + R[t]; break;
    case 2: R[d] = R[s] - R[t]; break;
    case 3: R[d] = R[s] & R[t]; break;
    case 4: R[d] = R[s] ^ R[t]; break;
    case 5: R[d] = R[s] << R[t]; break;
    case 6: R[d] = R[s] >> R[t]; break;
    case 7: R[d] = addr; break;
    case 8: R[d] = mem[addr]; break;
    case 9: mem[addr] = R[d]; break;
    case 10: R[d] = mem[R[t]]; break;
    case 11: mem[R[t]] = R[d]; break;
    case 12: if (R[d] == 0) pc = addr; break;
    case 13: if (R[d] > 0) pc = addr; break;
    case 14: pc = R[d]; pc = addr; break;
    case 15: R[d] = pc; pc = addr; break;
}
```

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## TOY Simulator: Omitted Details

Omitted details.

- Register 0 is always 0.
  - reset  $R[0]=0$  after each fetch-execute step
- Standard input and output.
  - if `addr` is `FF` and opcode is load (indirect) then read in data
  - if `addr` is `FF` and opcode is store (indirect) then write out data
- TOY registers are 16-bit integers; program counter is 8-bit.
  - Java `int` is 32-bit; Java `short` is 16-bit
  - use casts and bit-whacking

Complete implementation. See `TOY.java` on booksite.

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## Simulation

Profound ideas stemming from simulation.

- Backwards compatibility
- Virtual machines
- Layers of abstraction

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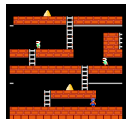


## Backwards Compatibility

Building a new computer? Need a plan for old software.

Two possible approaches

- Rewrite software (costly, error-prone, boring, and time-consuming).
- **Simulate old computer on new computer.**



Lode Runner



Apple IIe



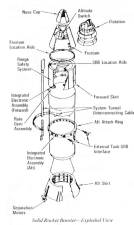
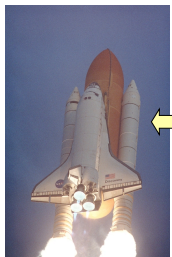
Mac OS X Apple IIe emulator widget running Lode Runner

Ancient programs still running on modern computers.

- Payroll
- Power plants
- Air traffic control
- Ticketron.
- Games.

## Effects of Backwards Compatibility: example 1

Q. Why is Space Shuttle SRB long and narrow?



A. Fits on standard US rail gauge.



...

A. Fits "back ends" of two war horses!



## Backwards Compatibility

Q. Why is standard US rail gauge 4 feet, 8.5 inches?



A. Same spacing as wheel ruts on old English roads.



Q. Why is wheel rut spacing 4 feet, 8.5 inches?

A. For Roman war chariots.



Q. Why is war chariot rut spacing 4 feet, 8.5 inches?

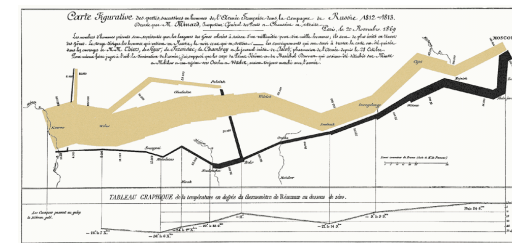
A. Fits "back ends" of two war horses!



## Effects of Backwards Compatibility: Example 2

Napoleon's march on Russia.

- Progress slower than expected.
- Eastern European ruts didn't match Roman gauge.
- Stuck in the field during Russian winter instead of Moscow.
- Lost war.



Lessons.

- Maintaining backwards compatibility can lead to inelegance and inefficiency.
- Maintaining backwards compatibility is Not Always A Good Thing.
- May need fresh ideas to conquer civilized world.

## Virtual machines

Building a new airplane? Simulate it to test it.

- Issue 1: Simulation may not reflect reality.
- Issue 2: May not be able to afford simulation.

Building a new computer? Simulate it to test it.

- Advantage 1: Simulation is reality (it defines the new machine).
- Advantage 2: Can develop software without having machine.
- Advantage 3: Can simulate machines you wouldn't build.

Example 1: An operating system can implement a **Virtual Memory** that is much larger than a real memory by simulating programs and going to disk or the web to reference "memory"

Example 2: An operating system can implement multiple **Virtual Machines** on the same real machine by keeping track of multiple PCs and rotating control to the different machines

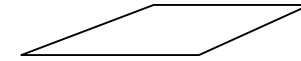
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## Layers of Abstraction

Is TOY real?

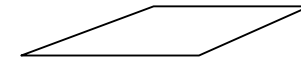


programmer



Java language specification

Is Java real?



Instruction set architecture



machine

Approaching a new problem?

- build an (abstract) language for expressing solutions
- design an (abstract) machine to execute the language
- food for thought: **Why build the machine?**

Examples: MATLAB, BLAST, AMP ....

simulate it instead

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