Compression; Error detection & correction

compression: squeeze out redundancy

- to use less memory or use less network bandwidth
- encode the same information in fewer bits
 - some bits carry no information
 - some bits can be computed or inferred from others
 - some bits don't matter to the recipient and can be dropped entirely

• error detection & correction: add redundancy

- to detect and fix up loss or damage
- add carefully defined, systematic redundancy
- with enough of the right redundancy, can detect damaged bits can correct errors

Compressing English text

- · letters do not occur equally often
- encode frequent letters with fewer bits, less frequent things with more bits (trades complexity against space)
 - e.g., Morse code, Huffman code, ...

run-length encoding

- encode runs of identical things with a count
- e.g., World Wide Web Consortium => WWWC => W3C
- words do not occur equally often
- encode whole words, not just letters
 - e.g., abbreviations for frequent words

Lempel-Ziv compression algorithm

- Abraham Lempel & Jacob Ziv, 1977
- "sliding window": when a sequence of input is the same as a previous sequence within a reasonable distance, replace the new sequence by (*length, offset*) of the previous one
 - each of the next *length* bytes is equal to the byte that is *offset* bytes earlier in the (uncompressed) stream
 - the (length, offset) will be shorter than the data it represents
- · compression adapts to properties of any particular input
 - algorithm itself is independent of nature of input
- lossless
 - compression followed by decompression reproduces the input exactly
- the basis of PKZip, Winzip
 - compresses Bible from 4.1 MB to 0.9 MB (typical for text)

JPEG (Joint Photographic Experts Group) picture compression

- lossy compression scheme, based on how human eyes work
- · digitize picture into pixels
- discard some color information (use fewer distinct colors)
 eye is less sensitive to color variation than brightness
- discard some fine detail
 - decompressed image is not quite as sharp as original
- discard some fine gradations of color and brightness
- use lossless compression techniques to compress resulting stream of numeric values
- can trade quality for space
 - compression usually 10:1 with little loss of quality
- used in web pages, digital cameras, ...
- works well for photos, less well for line art, drawings and solid blocks of color



MPEG (Moving Picture Experts Group) movie compression

- MPEG-2: lossy compression scheme, based on human perceptions
- uses JPEG for individual frames (spatial redundancy)
- adds compression of temporal redundancy
 - look at image in blocks
 - if a block hasn't changed, just transmit that fact, not the content
 - if a block has moved, transmit amount of motion
 - motion prediction (encode expected differences plus correction)
 - separate moving parts from static background

- ..

- used in DVD, high-definition TV, digital camcorders, video games
- rate is 3-15 Mbps depending on size, frame rate
 - 15 Mbps ~ 2 MB/sec or 120 MB/min ~ 100x worse than MP3
 - 3 Mbps ~ 25 MB/min; cf DVD 25 MB/min ~ 3000 MB for 2 hours
 - regular TV is ~ 15 Mbps, HDTV ~ 60-80 Mbps
 see www.bbc.co.uk/rd/pubs/papers/paper_14/paper_14.shtml

MP3 (MPEG Audio Layer-3) sound compression

- movies have sound as well as motion; this is the audio part
- 3 levels, with increasing compression, increasing complexity
- based on "perceptual noise shaping":
 - use characteristics of the human ear to compress better:
 - human ear can't hear some sounds (e.g., very high frequencies)
 - human ear hears some sounds better than others
 - louder sounds mask softer sounds
- break sound into different frequency bands
- encode each band separately
- encode 2 stereo channels as 1 plus difference
- gives about 10:1 compression over CD-quality audio
 - 1 MB/minute instead of 10 MB/minute
 - can trade quality against compression
- see http://www.oreilly.com/catalog/mp3/chapter/ch02.html

Other audio compression algorithms

- AAC (Advanced Audio Coding)
 - supposed to succeed MP3, used in Apple iTunes
- WMA (Windows Media Audio)
- Ogg Vorbis (open source)
- ...
 - maybe 10-20 times more compact than WAV format

• speech coding for cell phones, Internet telephony, etc.

- narrower frequency range (100 Hz 4 KHz)
- requires low delay
- uses a model of human vocal tract (specialized to voice) adapts to specific voice3
- much higher compression than for general audio

Summary of compression

- eliminate / reduce redundancy
 - more frequent things encoded with fewer bits
 - use a dictionary of encoded things, and refer to it (LZ)
 - encode repetitions with a count
- not everything can be compressed
 - something will be bigger

· lossless vs lossy compression

- lossy discards something that is not needed by recipient

tradeoffs

- encoding time and complexity vs decoding time and complexity
- encoding is usually slower and more complicated (done once)
- parameters in lossy compressions size, speed, quality

Error detection and correction

- systematic use of redundancy to defend against errors
- \cdot some common numbers have no redundancy
 - and thus can't detect when an error might have occurred
 - e.g., SSN -- any 9-digit number is potentially valid
- if some extra data is added or if some possible values are excluded, this can be used to detect and even correct errors
- common examples include
 - ATM & credit card numbers
 - ISBN for books
 - bar codes for products

ATM card checksum

- invented by Peter Luhn, IBM, 1954 (patented 1960)
- credit card / ATM card checksum:

starting at rightmost digit: multiply digit alternately by 1 or 2 if result is > 9 subtract 9 add the resulting digits sum should be divisible by 10

e.g., 12345678 is invalid 8 + (14-9) + 6 + (10-9) + 4 + 6 + 2 + 2 = 34 but 42345678 is valid 8 + (14-9) + 6 + (10-9) + 4 + 6 + 2 + 8 = 40

• defends against single digit errors and most transpositions

- these are the most common errors

ISBN checksum

```
checksum for 10-digit ISBN:
starting at leftmost digit:
multiply corresponding digit by 10, 9, 8, ... down to 1 inclusive
(a final X has value 10)
add the resulting numbers
result must be divisible by 11
e.g., 0-201-61586-X is valid
10*0 + 9*2 + 8*0 + 7*1 + 6*6 + 5*1 + 4*5 + 3*8 + 6*2 + 1*10 = 132 = 12*11
defends against transpositions and single digit errors
```

Parity & other binary codes

• parity bit: use one extra bit so total number of 1-bits is even

0110100 => 0110100<u>1</u> 0110101 => 01101010

- detects any single-bit error
- more elaborate codes can detect and even correct errors
- basic idea is to add extra bits systematically so that legal values are uniformly spread out, so any small error converts a legal value into an illegal one
 - some schemes correct random isolated errors
 - some schemes correct bursts of errors (used in CD-ROMs)
- no error correcting code can detect/correct all errors
 - a big enough error can convert one legal pattern into another one