

Princeton University

COS 217: Introduction to Programming Systems

SPARC “Leaf” Subroutine Calling Conventions

A subroutine can be a leaf subroutine only if it need not have its own stack frame.
That is, only if it calls no other subroutines and defines no local variables in memory.

A leaf subroutine may use only registers o0-o5 and g0-g1.

The calling subroutine need not know that the called subroutine is a leaf subroutine.

When subroutine f calls leaf subroutine g...

In f:

- (1) Store actual parameters 1 to 6 in o0-o5.
- (2) Store actual parameters 7, 8, ... in memory locations o6/sp + 92, o6/sp + 96, ...
- (3) Execute “call g”.

Store register pc in o7.

Note: Register o7 thus stores the address of the call instruction.
Jump to the instruction at label g.

- (4) But before executing the “jumped to” instruction, execute the delay instruction that follows the “call” instruction.

In g:

- (5) Use formal parameters in o0-o5 and o6/sp + 92, o6/sp + 96, ... to compute return value(s).
- (6) Store return values in registers o0-o5.
- (7) Execute “retl” (return from leaf).

Jump to o7 + 8

Jump to the instruction after the delay instruction after the call instruction.

- (8) But before executing the “jumped to” instruction, execute the delay instruction that follows the “retl” instruction, typically “nop”.

In f:

- (9) Retrieve g’s return values from o0-o5.