

# BICUBICAL DIRECTED TYPE THEORY

MATTHEW WEAVER

ABSTRACT. Many structures in mathematics and computer science include a notion of a directed transition (or morphism), and when working with such structures one intends to always preserve this directed structure as an invariant. Homotopy type theory provides a setting in which all constructions preserve morphism structure, but only supports invertible morphisms. Directed type theory is an extension of homotopy type theory that also accounts for directed morphisms. Bicubical directed type theory is a constructive model of type theory in which every type comes packaged with two special types: one of (undirected) paths in the type and another of morphisms (i.e. directed paths). This work is an extension of the cartesian cubical model of type theory [Angiuli et al., 2017] and the bisimplicial model of directed type theory [Riehl and Shulman, 2017]. For a type  $A$  and terms  $x, y : A$ , we write  $\mathbf{Path}_A x y$  to denote the type of paths from  $x$  to  $y$  and  $\mathbf{Hom}_A x y$  denotes the type of morphisms from  $x$  to  $y$ . The model is “bicubical” as the data for paths and morphisms in a type can each be represented by a cubical data-structure. We add this basic infrastructure to extensional type theory (i.e. dependent type theory with equality and uniqueness of identity proofs), and then use the internal logic of the type theory to describe and reason about the path and morphism structure of types—a technique first described in Orton and Pitts [2016].

In particular, we care about the various ways in which a dependent type can be fibrant: The general idea is that, given a type  $A$  and family of types  $B : A \rightarrow \mathbf{Type}$ , we want to know when path and/or morphism structure in  $A$  can be lifted to some structure in the type family  $B$ . When such structure is defined in  $B$  for every path in  $A$ , we call  $B$  a fibration. As an example, an important instance of this is called a coercion structure: If  $B : A \rightarrow \mathbf{Type}$  has a coercion structure, then any path  $p : \mathbf{Path}_A x y$  uniquely determines a function  $\mathbf{coe}_B p : B x \rightarrow B y$ .

Given each notion of fibration we define, we can internally construct a corresponding universe of types paired with the desired fibration structure using the construction described in Licata et al. [2018]. Take for example  $\mathbf{U}_{\mathbf{coe}}$ : Given  $A : \mathbf{Type}$ , any dependent type  $B : A \rightarrow \mathbf{U}_{\mathbf{coe}}$  is always equipped with a coercion structure  $\mathbf{coe}_B$ .

Lastly, we consider the different univalence structures each universe exhibits. Univalence structures provide an internal interpretation for paths and/or morphisms in the universe. For example, univalence for paths in a universe  $\mathbf{U}$  generally refers to the construction of an equivalence between the type of paths  $\mathbf{Path}_{\mathbf{U}} A B$  and the type of equivalences  $A \simeq B$  (i.e. pairs of invertible functions between  $A$  and  $B$ ). As our primary contribution, we construct a universe  $\mathbf{U}_{\mathbf{cov}}$  that is suitably fibrant with respect to both paths and morphisms, we prove it is closed under all the type-formers in the type theory, we build univalence for paths, and additionally we construct directed univalence: There is an equivalence between the function space  $A \rightarrow B$  and the morphism space  $\mathbf{Hom}_{\mathbf{U}_{\mathbf{cov}}} A B$  for every  $A$  and  $B$  in  $\mathbf{U}_{\mathbf{cov}}$ . As a caveat, the full proof of directed univalence is currently not entirely constructive; specifically, the proof that for every morphism  $p : \mathbf{Hom}_{\mathbf{U}_{\mathbf{cov}}} A B$  there is a path between  $p$  and its roundtrip through the equivalence uses a nonconstructive axiom.

Having presented this new model of type theory, we suggest how directed univalence can be useful for functional programming and when formalizing (meta)theory of computational structures. We have formalized all of our results using the Agda proof-assistant.

## REFERENCES

- C. Angiuli, G. Brunerie, T. Coquand, K.-B. Hou (Favonia), R. Harper, and D. R. Licata. Cartesian cubical type theory. Available from <https://github.com/dlicata335/cart-cube/blob/master/cart-cube.pdf>, 2017.
- D. R. Licata, I. Orton, A. M. Pitts, and B. Spitters. Internal universes in models of homotopy type theory. In *International Conference on Formal Structures for Computation and Deduction*, 2018.
- I. Orton and A. M. Pitts. Axioms for modelling cubical type theory in a topos. In *Computer Science Logic*, 2016.
- E. Riehl and M. Shulman. A type theory for synthetic  $\infty$ -categories. arXiv:1705.07442, 2017.