

# From application requests to Virtual IOPs: Provisioned key-value storage with Libra

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(\*now at Google)

# Shared Cloud



Google Cloud Platform

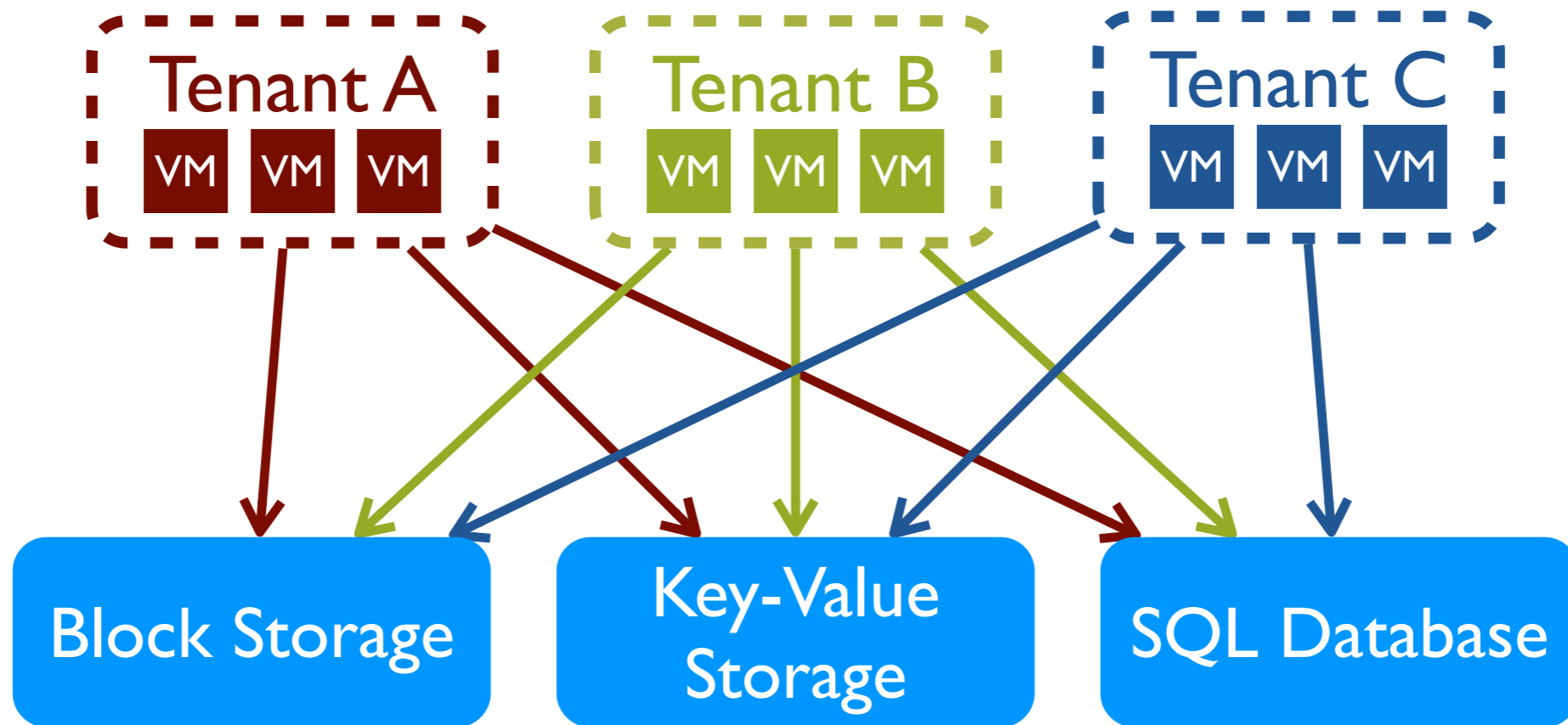


***rackspace***

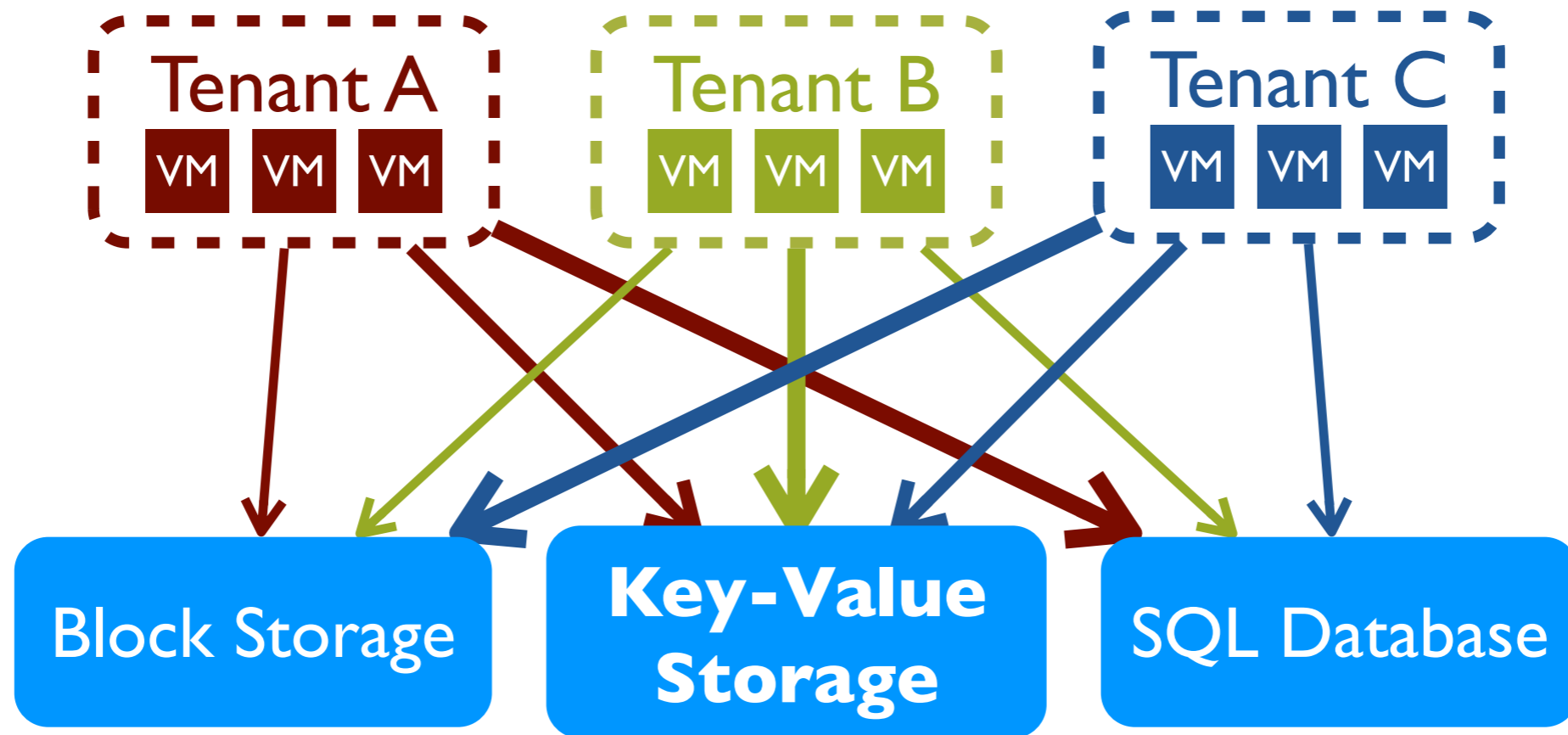


Windows Azure

# Shared Cloud Storage

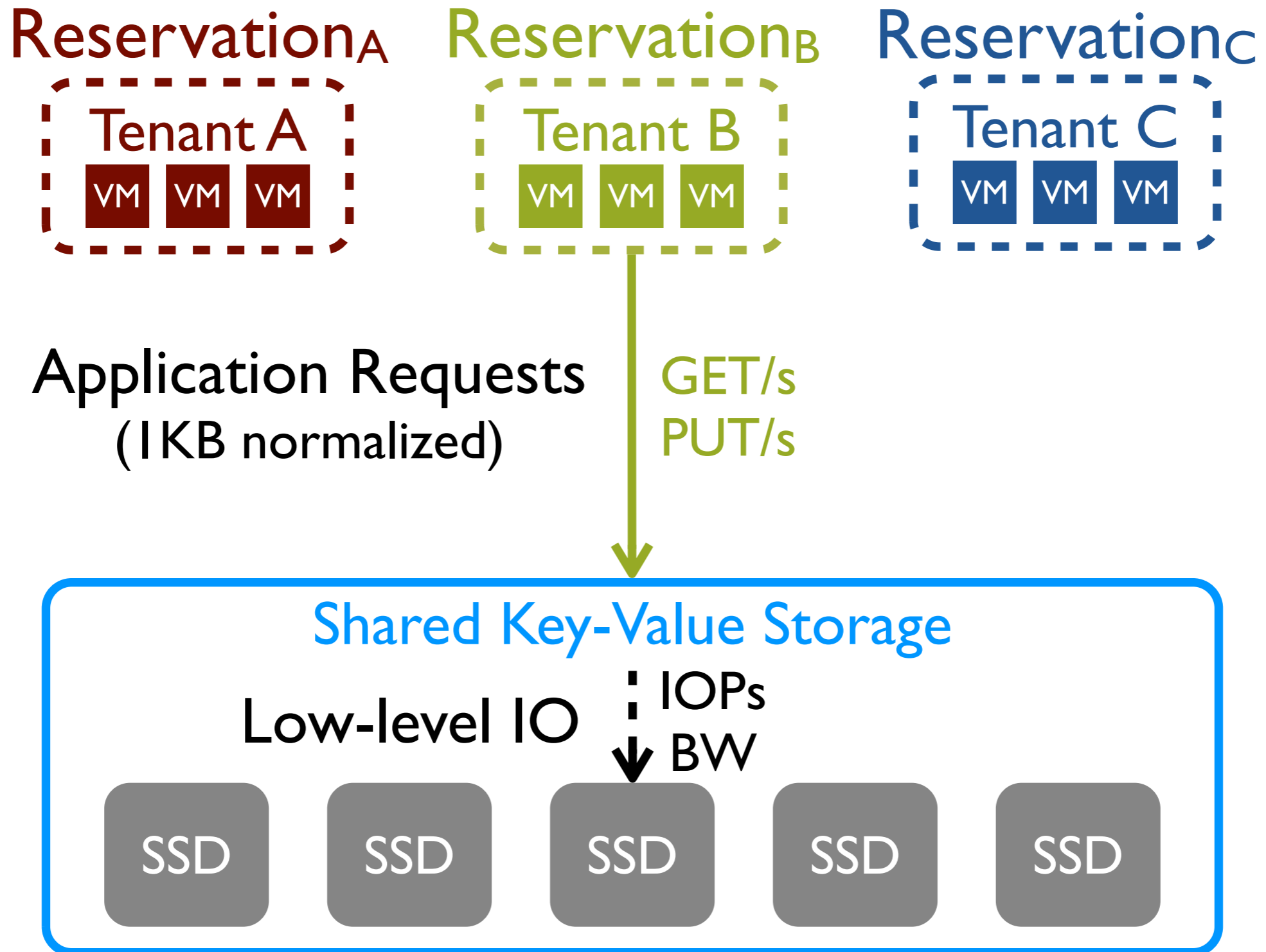


# Unpredictable Shared Cloud Storage



Disk IO-bound Tenants  
SSD-backed storage

# Provisioned Shared Key-Value Storage



# Libra Contributions

- **Libra IO Scheduler**

- Provisions *low-level IO allocations* for *app-request reservations* w/ high utilization.
- Supports arbitrary object distributions and workloads.

- **2 key mechanisms**

- Track per-tenant app-request resource profiles.
- Model IO resources with Virtual IOPs.

# Related Work

	Storage Type	App-requests	Work Conserving	Media
Maestro	Block	N	N	HDD
mClock	Block	N	Y	HDD
FlashFQ	Block	N	Y	SSD
DynamoDB	Key-Value	Y	N	SSD

# Provisioned Distributed Key-Value Storage

Global Reservation Problem  
[Pisces OSDI '12]

Reservation<sub>A</sub>



Reservation<sub>B</sub>



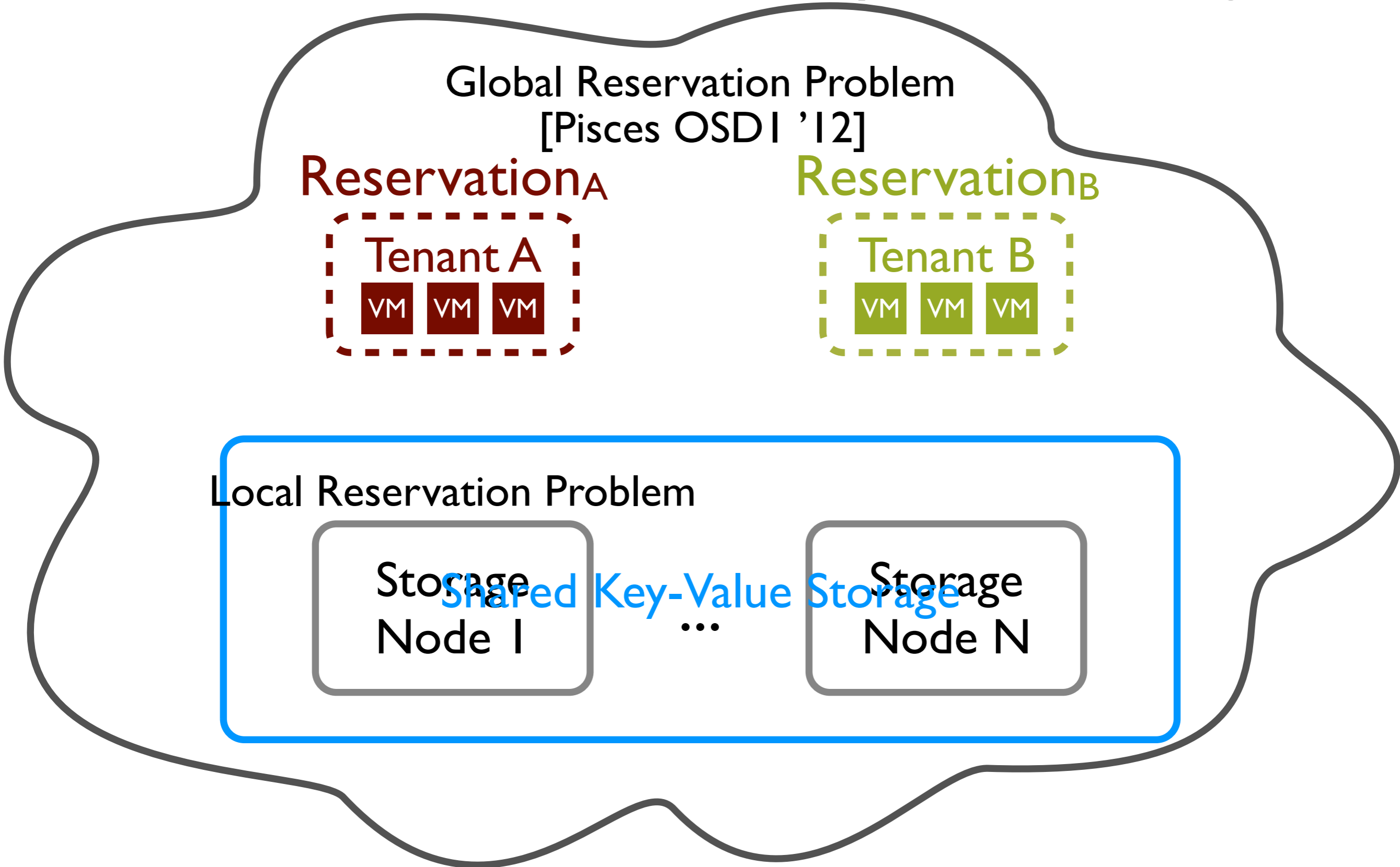
Local Reservation Problem

Storage  
Node I

Shared Key-Value Storage

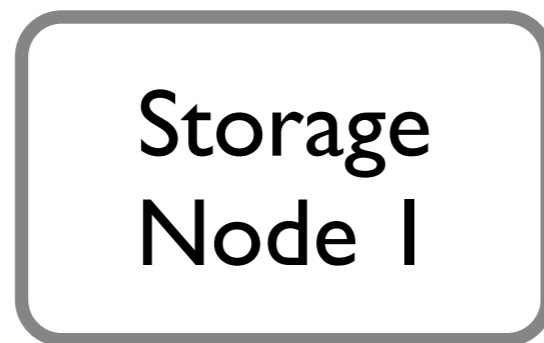
...

Storage  
Node N

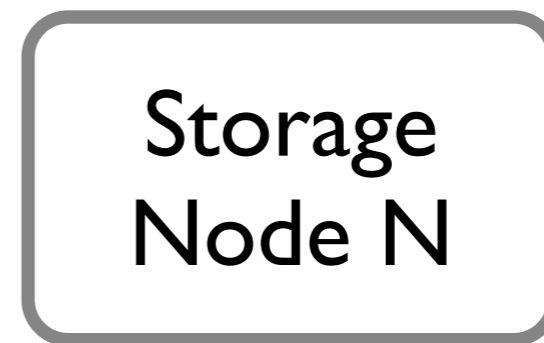




# Provisioned Distributed Key-Value Storage



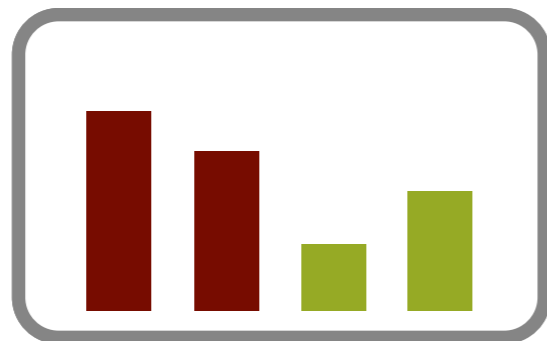
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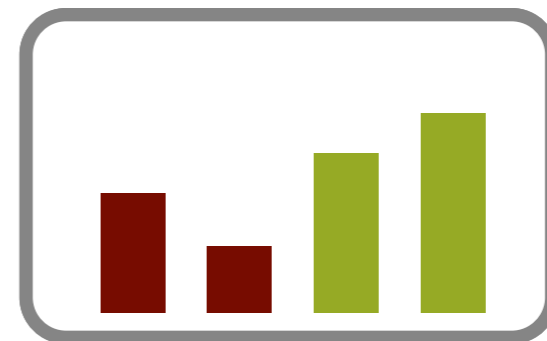
# Provisioned Distributed Key-Value Storage

Reservation<sub>A</sub>

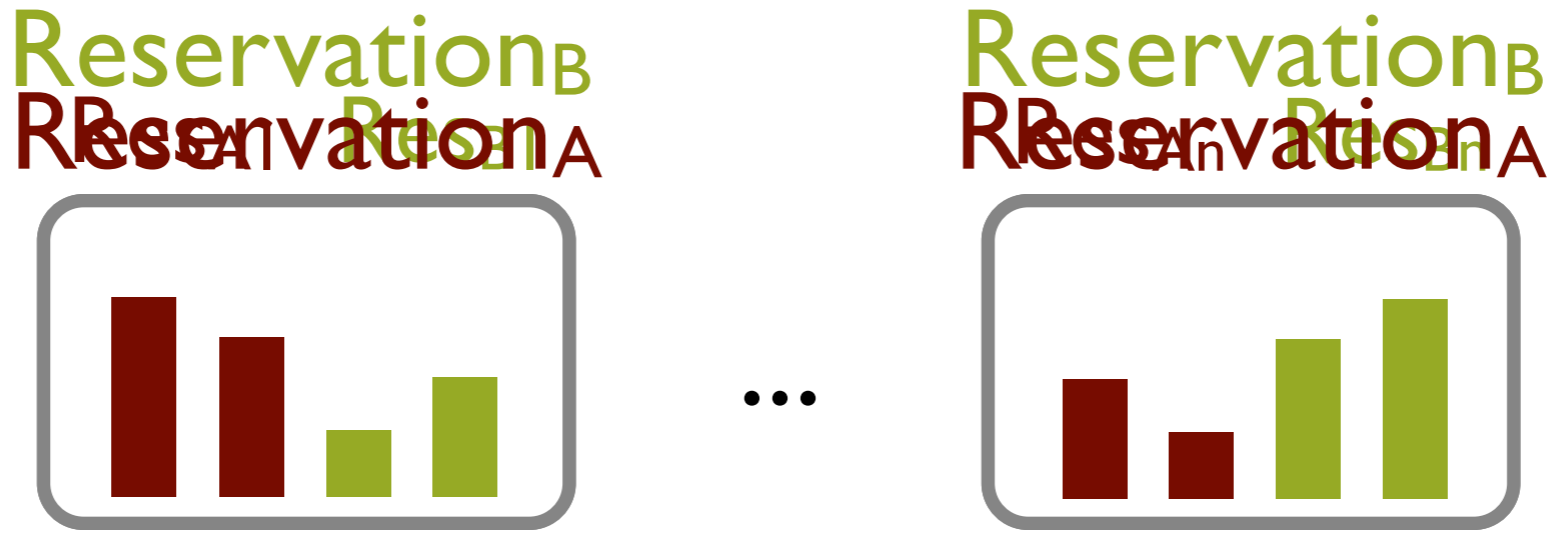
Reservation<sub>B</sub>



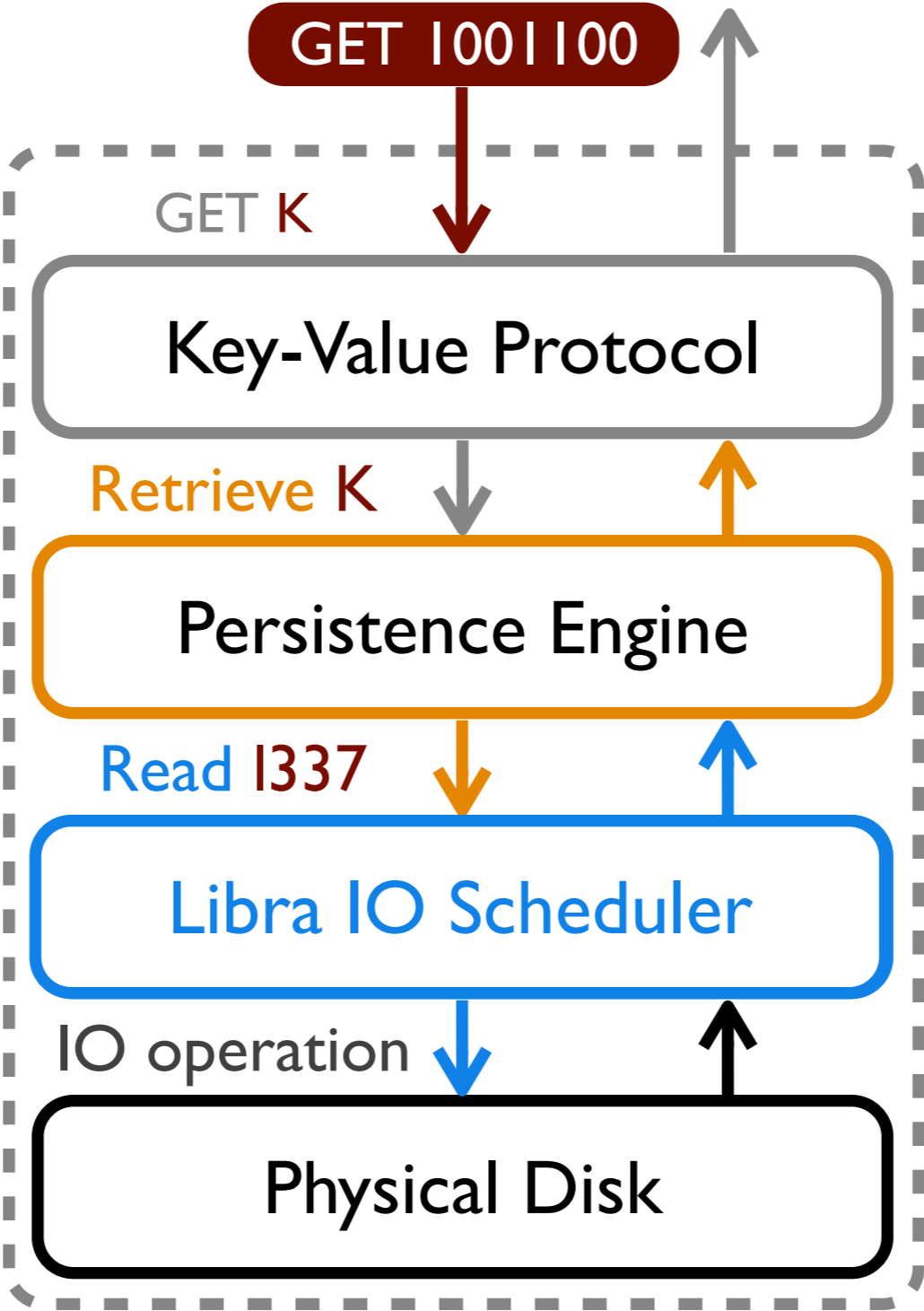
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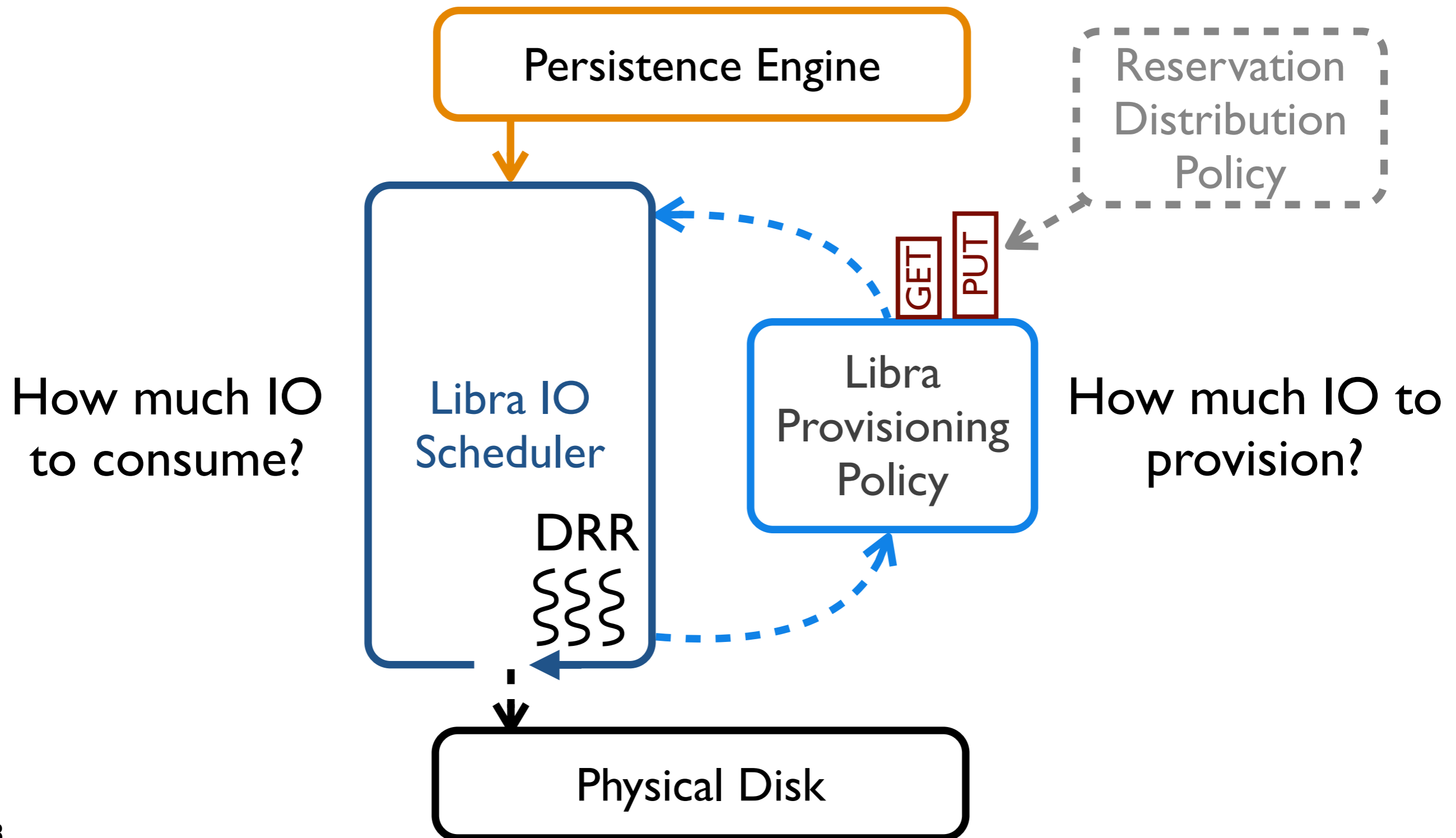
# Provisioned Distributed Key-Value Storage



# Provisioned Local Key-Value Storage



# Libra Design



# Provisioning App-request Reservations is Hard

IO Amplification

Track the app-request resource profiles

IO Interference

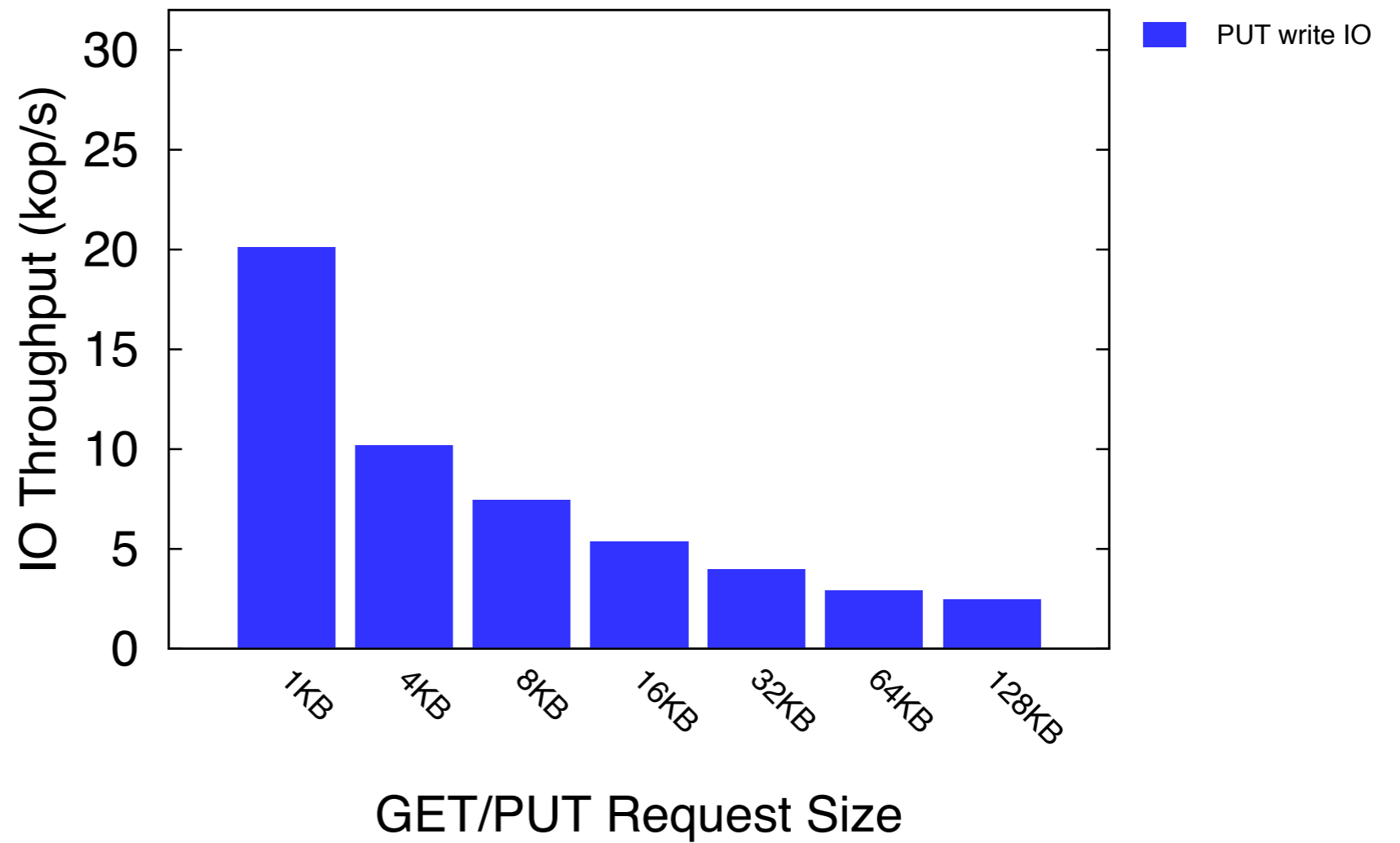
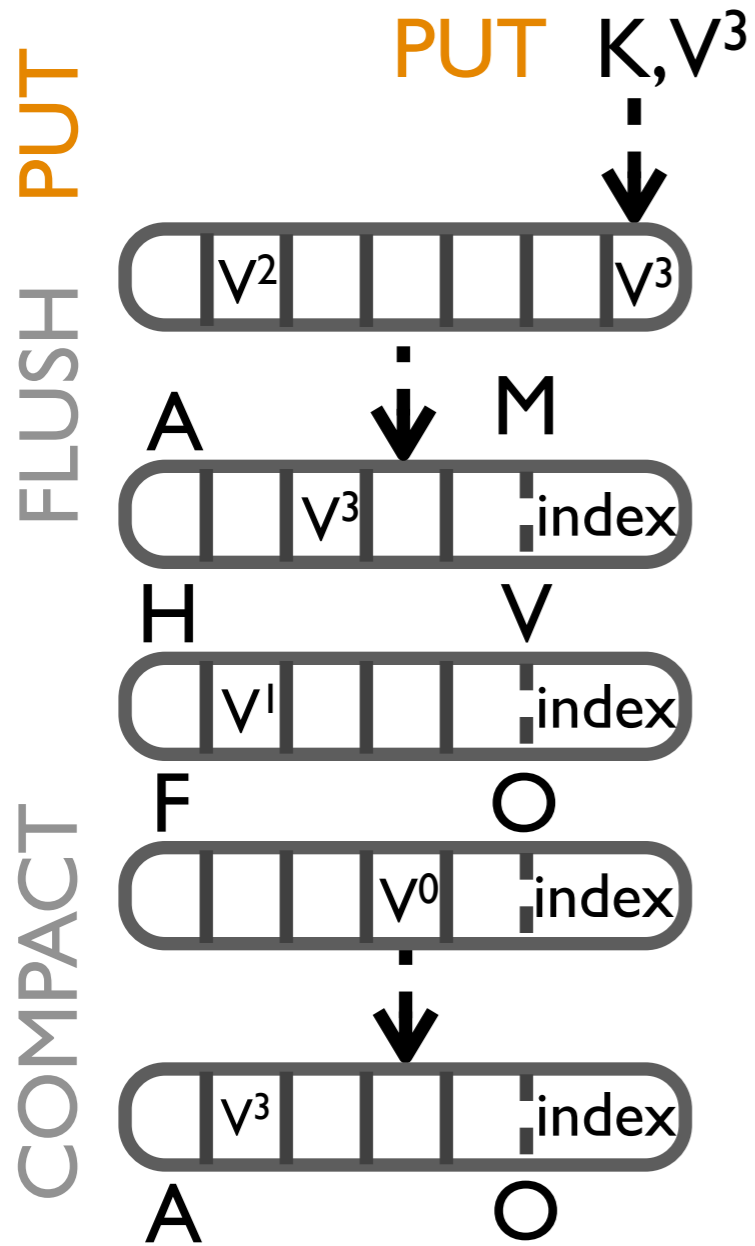
Understand that provisioning IO

Non-linear IO  
Performance

Non-IO without virtual KCPs

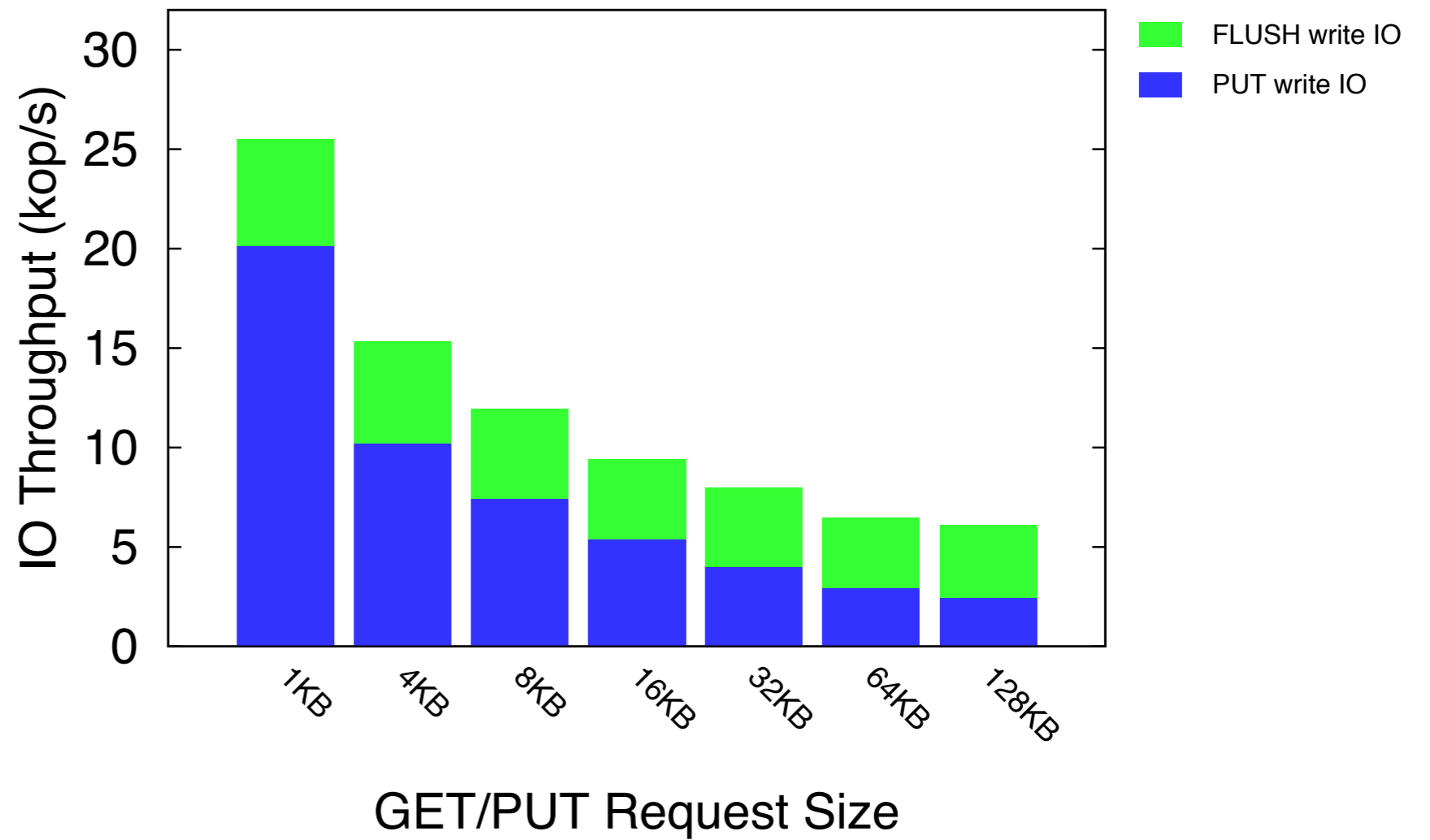
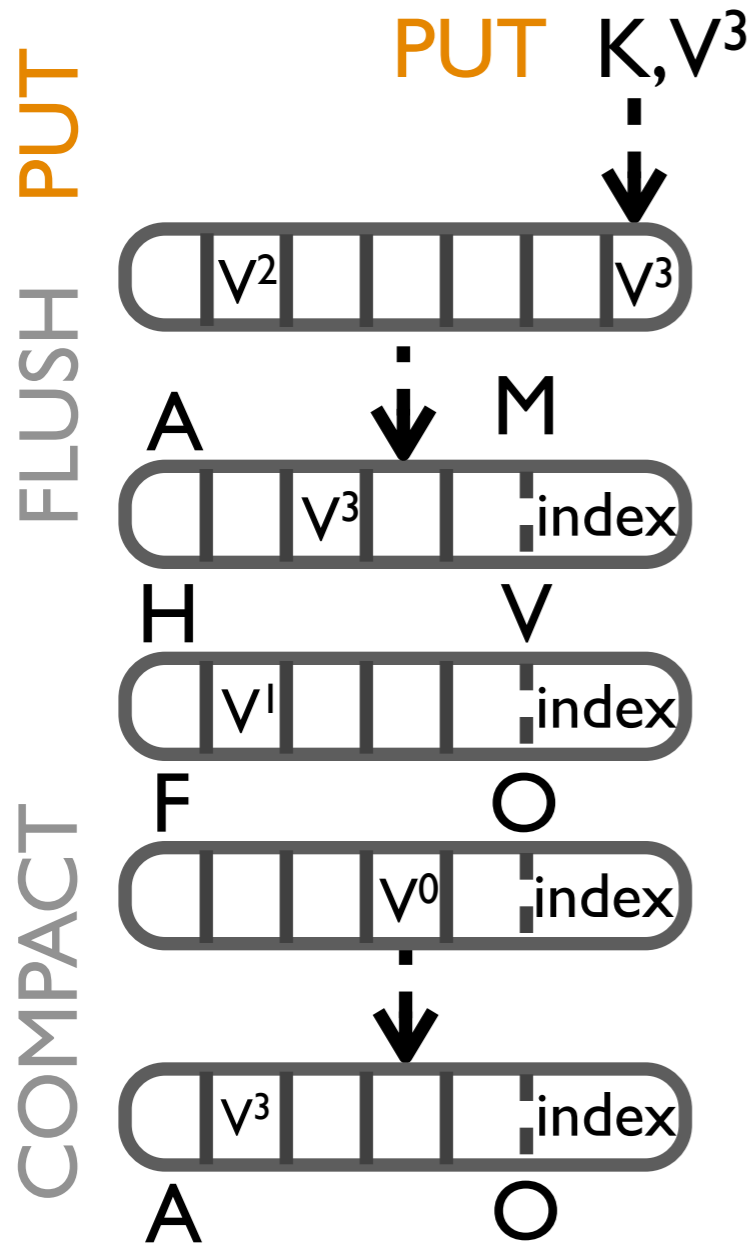
# Workload-dependent IO Amplification

LevelDB (LSM-Tree)



# Workload-dependent IO Amplification

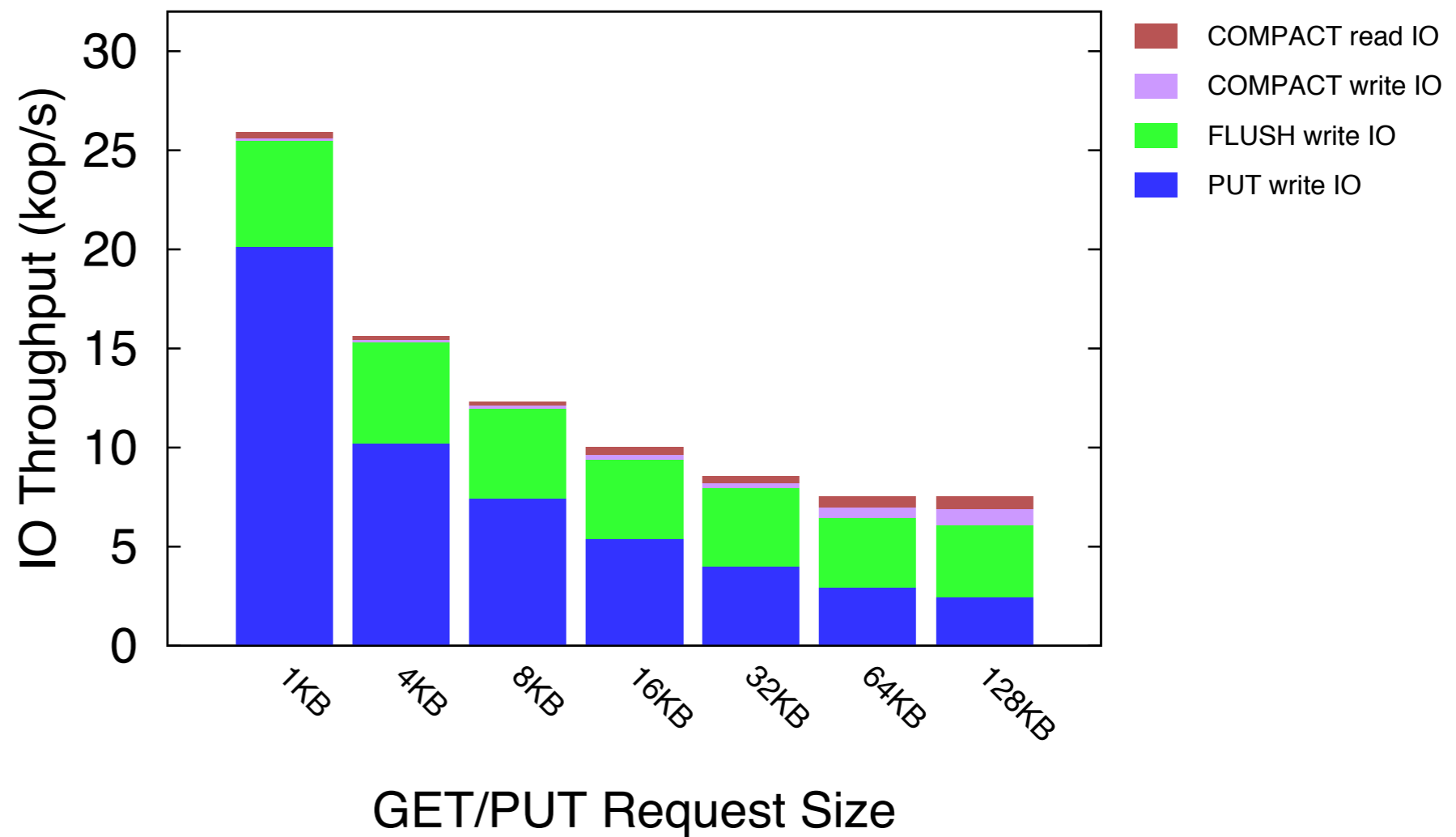
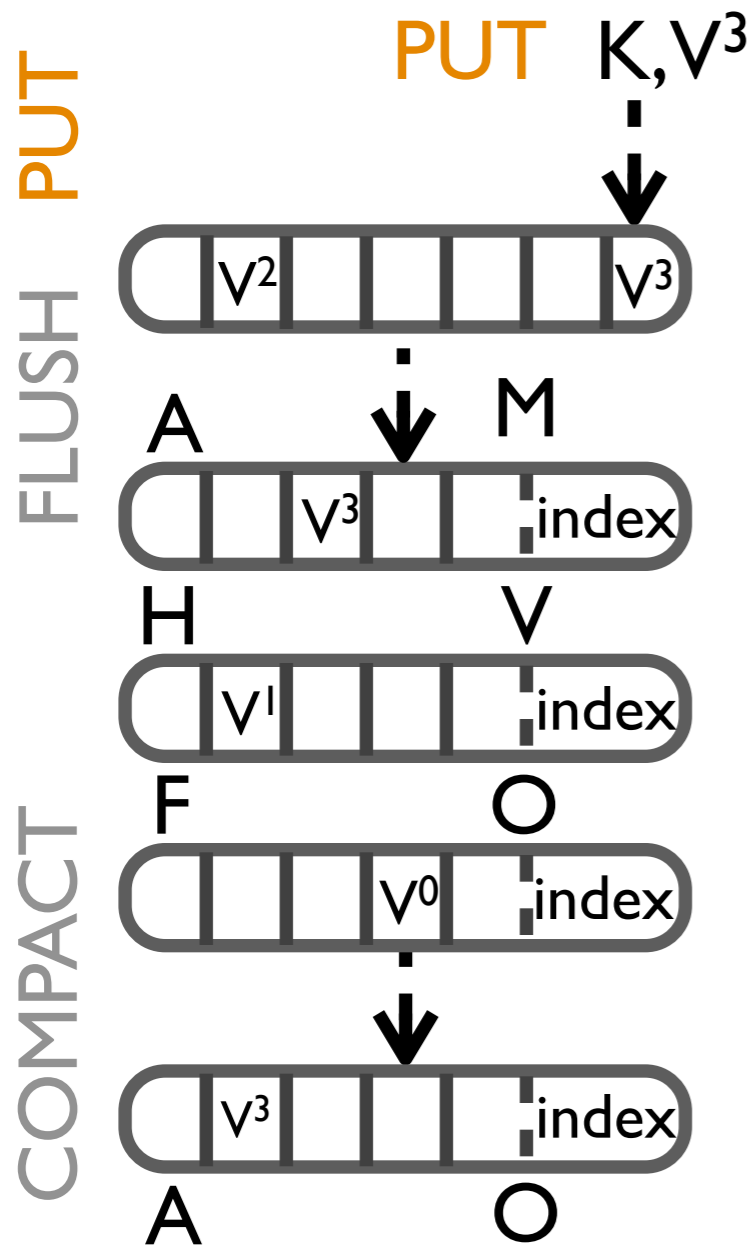
LevelDB (LSM-Tree)



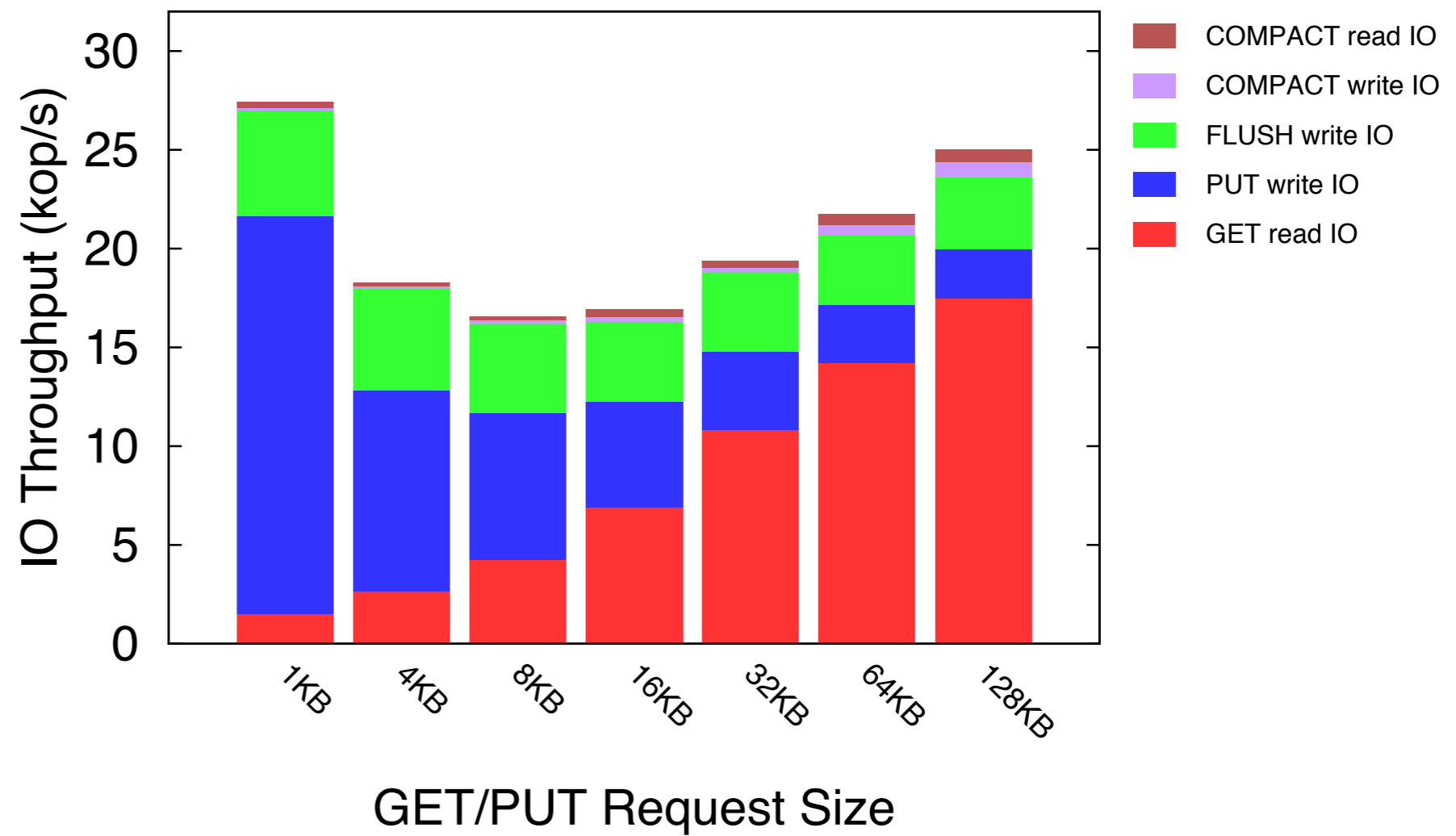
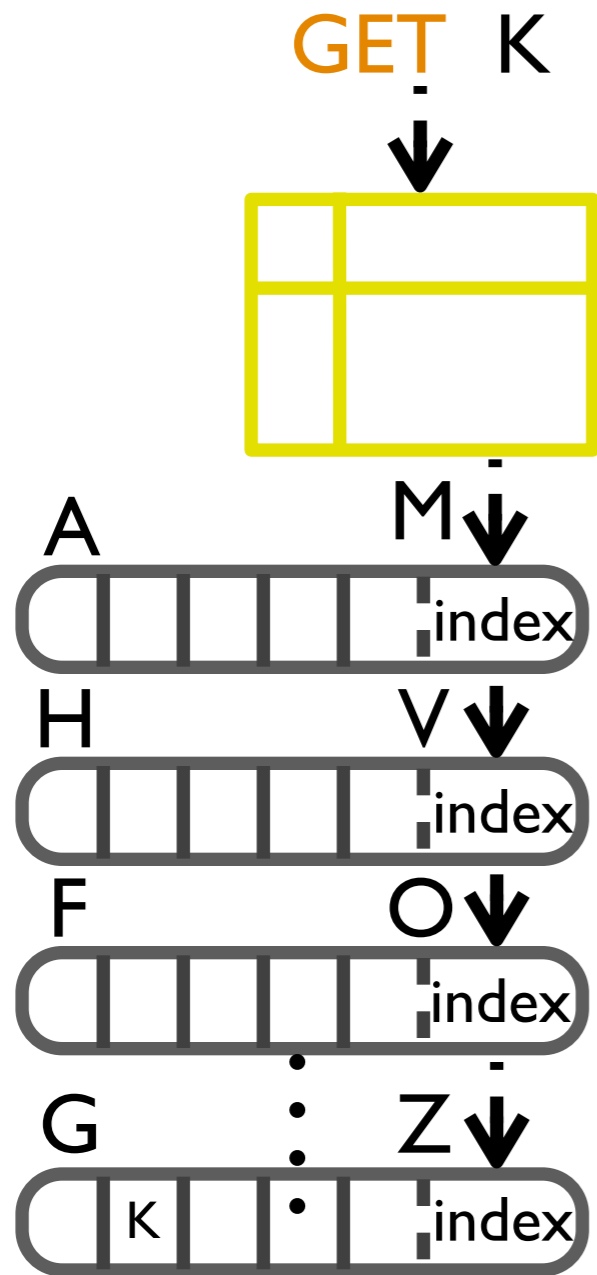


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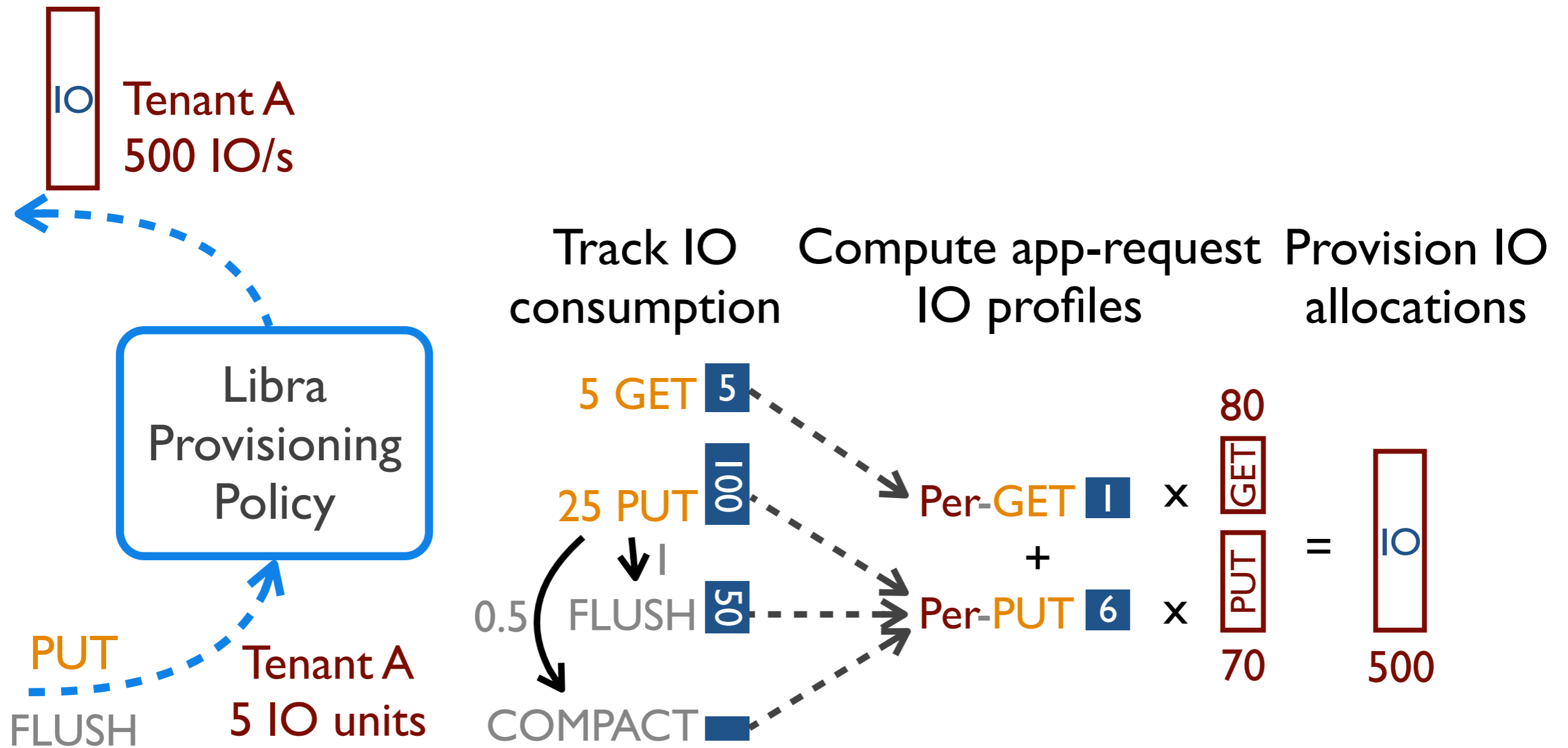
LevelDB (LSM-Tree)



# Workload-dependent IO Amplification



# Libra Tracks App-request IO Consumption to Determine IO Allocations

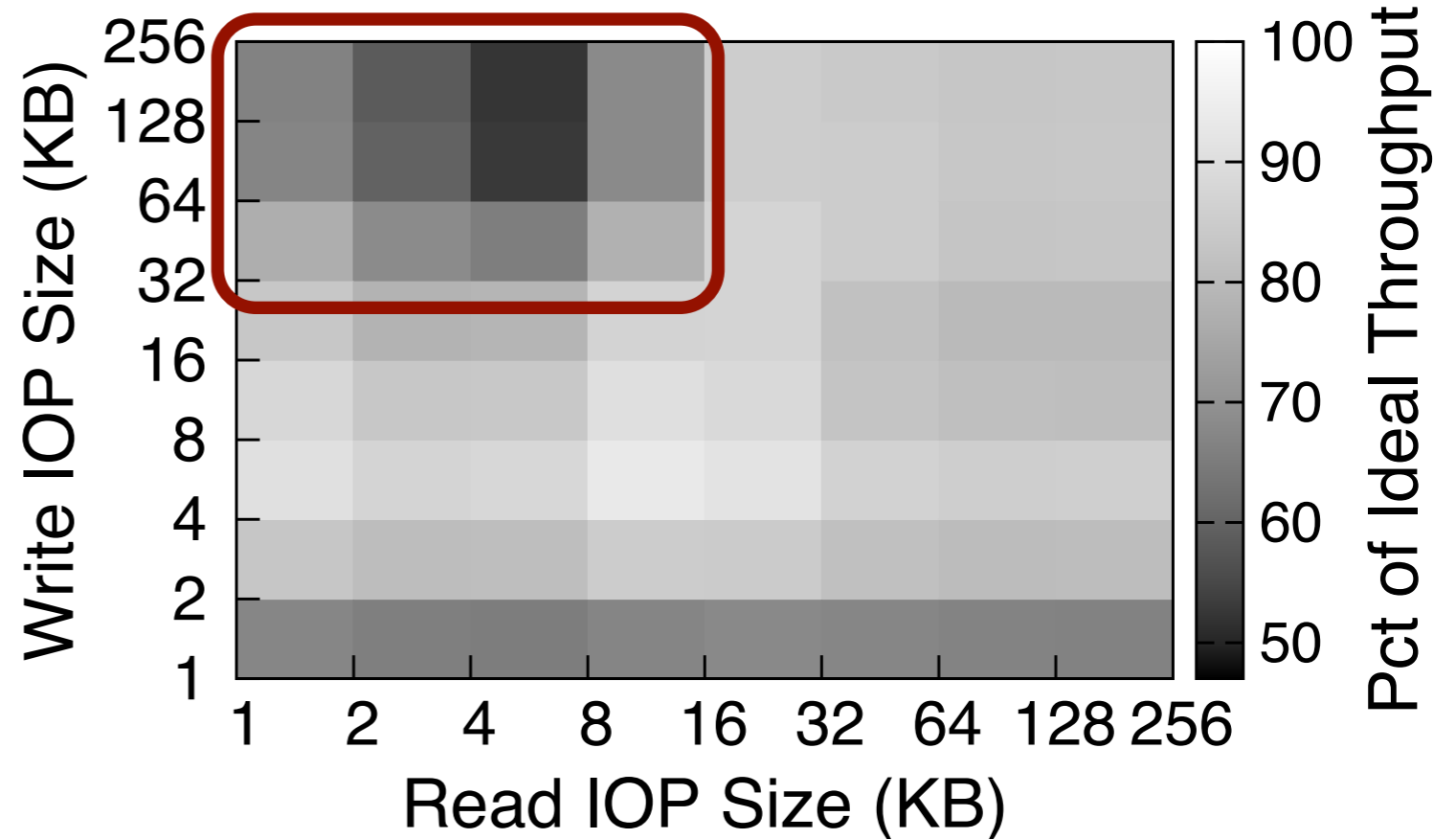


# Unpredictable IO Interference



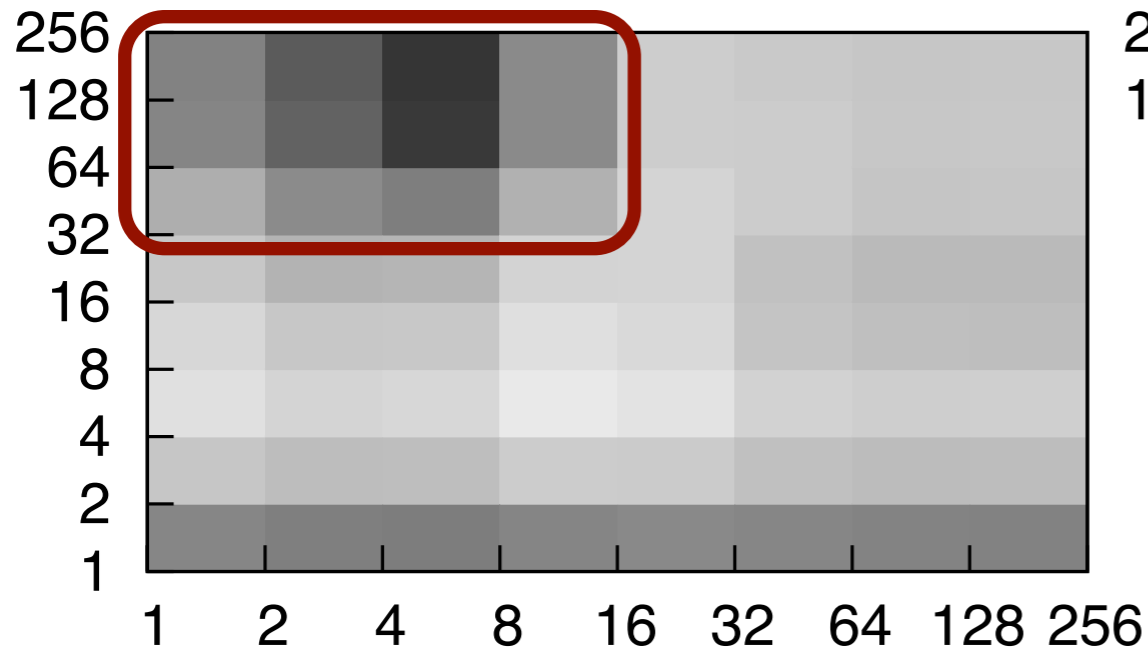
- ✓ Die-level parallelism, low latency IOPs
- ✗ Shared-controller and bus contention
- ✗ Erase-before-write overhead
- ✗ FTL and read-modify-write garbage collection

4 read/4 write tenants  
1:1 Pure Read/Pure Write

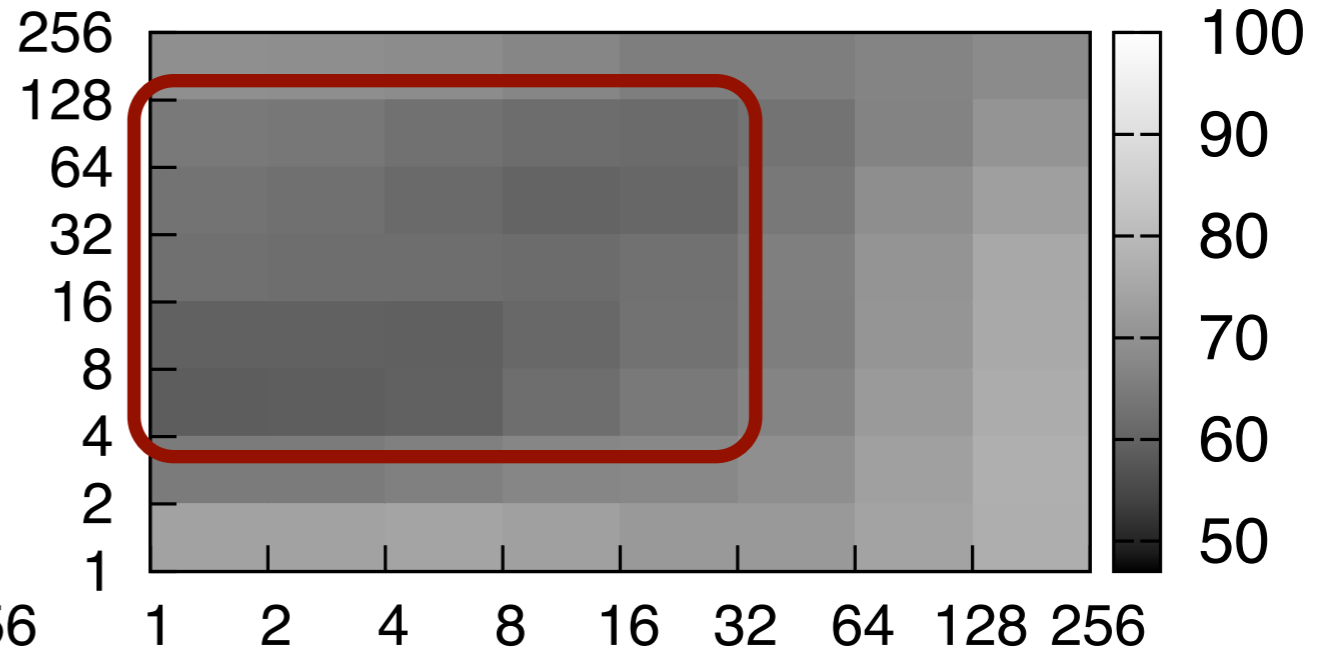


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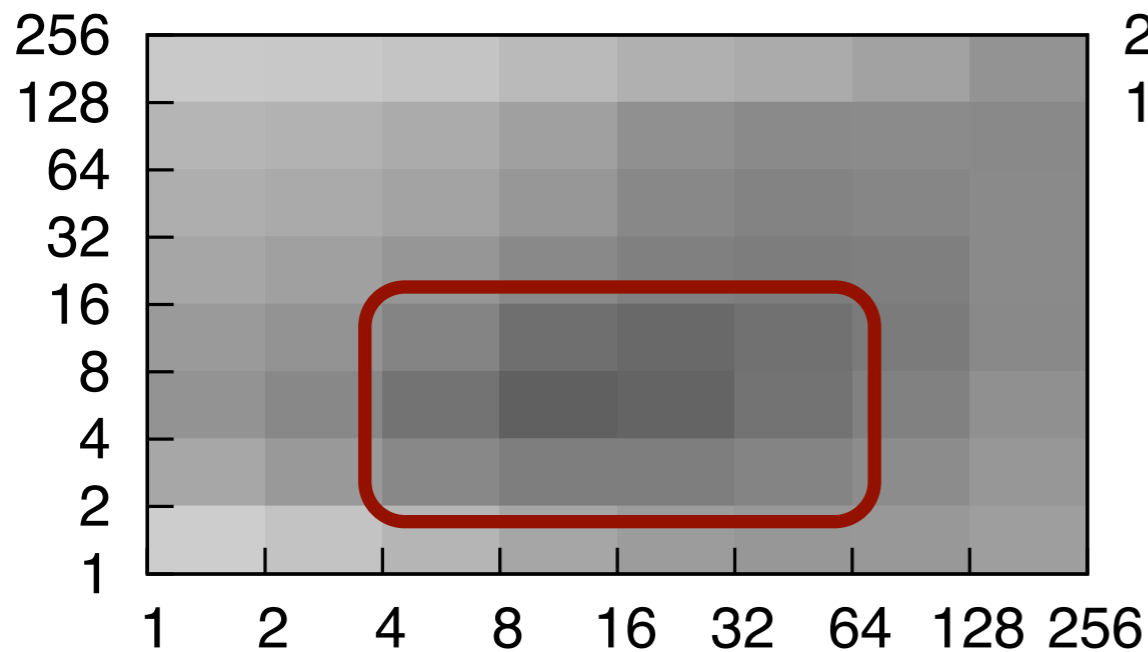
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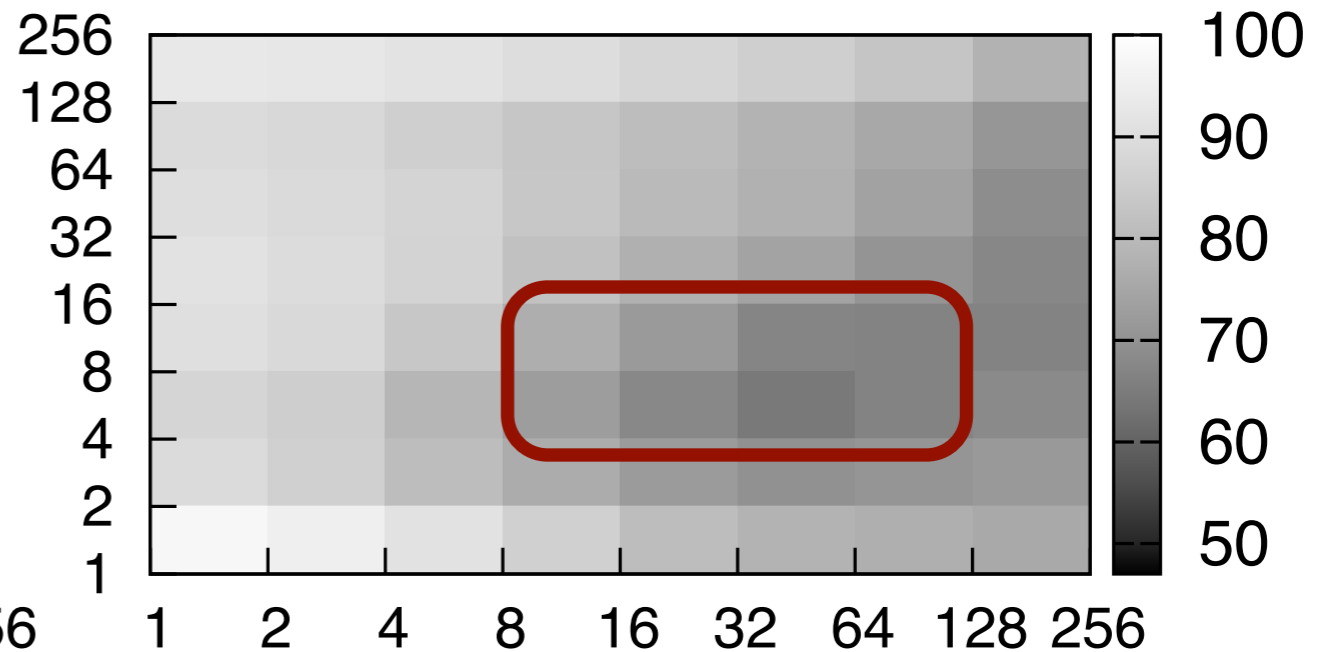
75:25 Read/Write Ratio



50:50 Read/Write Ratio



25:75 Read/Write Ratio



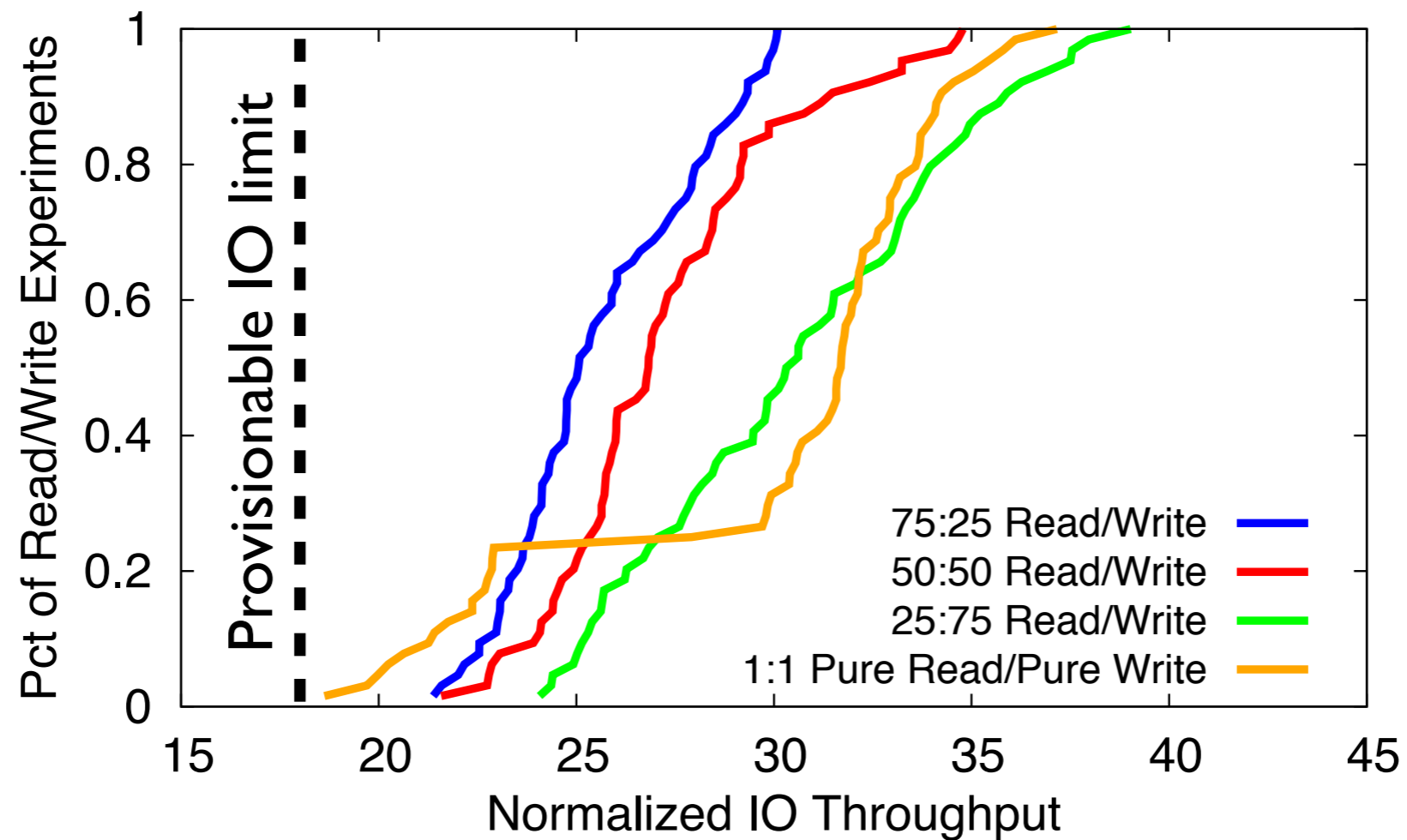
Write IOP Size (KB)

Pct of Ideal Throughput

Read IOP Size (KB)

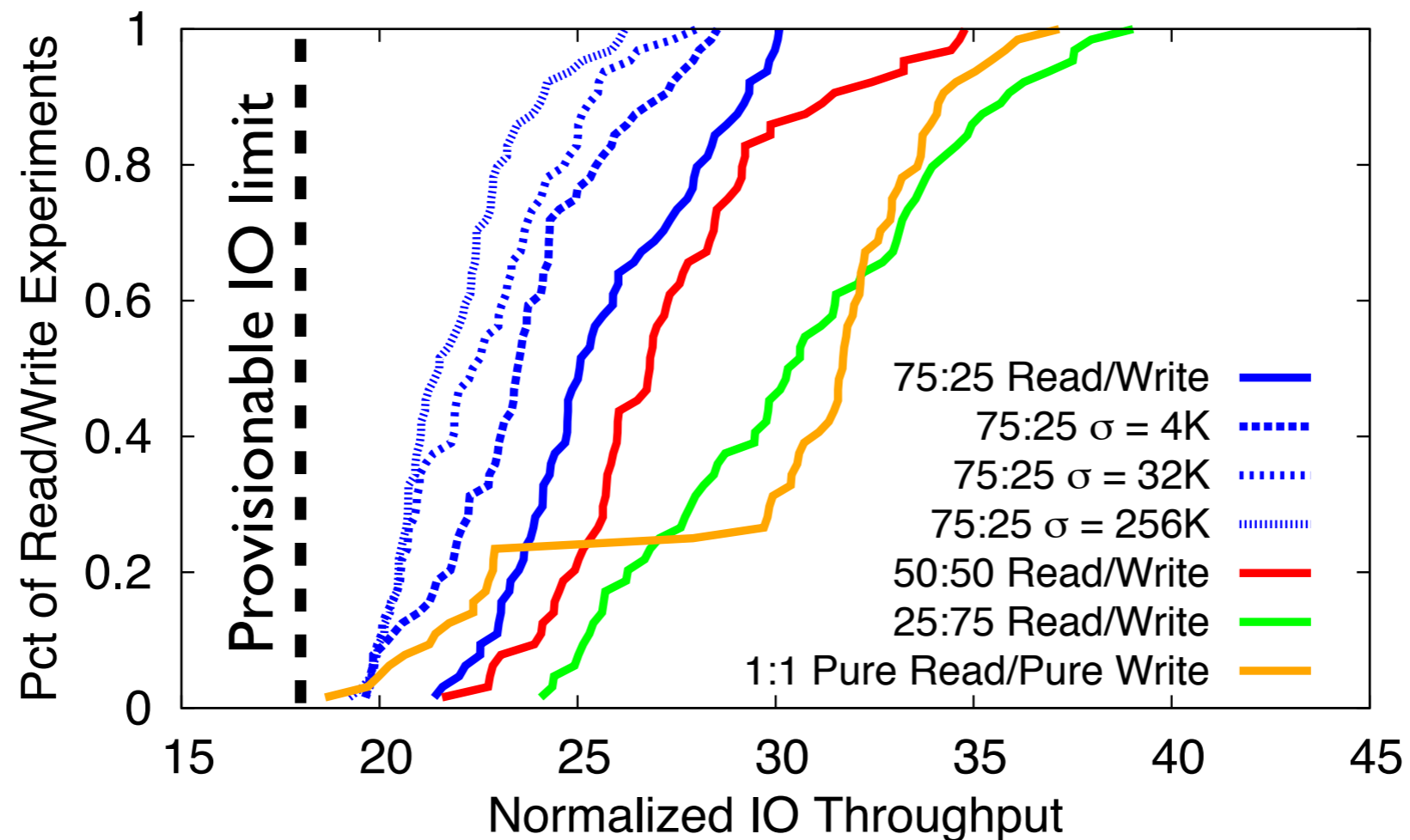
# Libra Underestimates IO Capacity to Ensure Provisionable Throughput

Provisionable IO throughput = floor(workloads)  
(18 Kop/s)



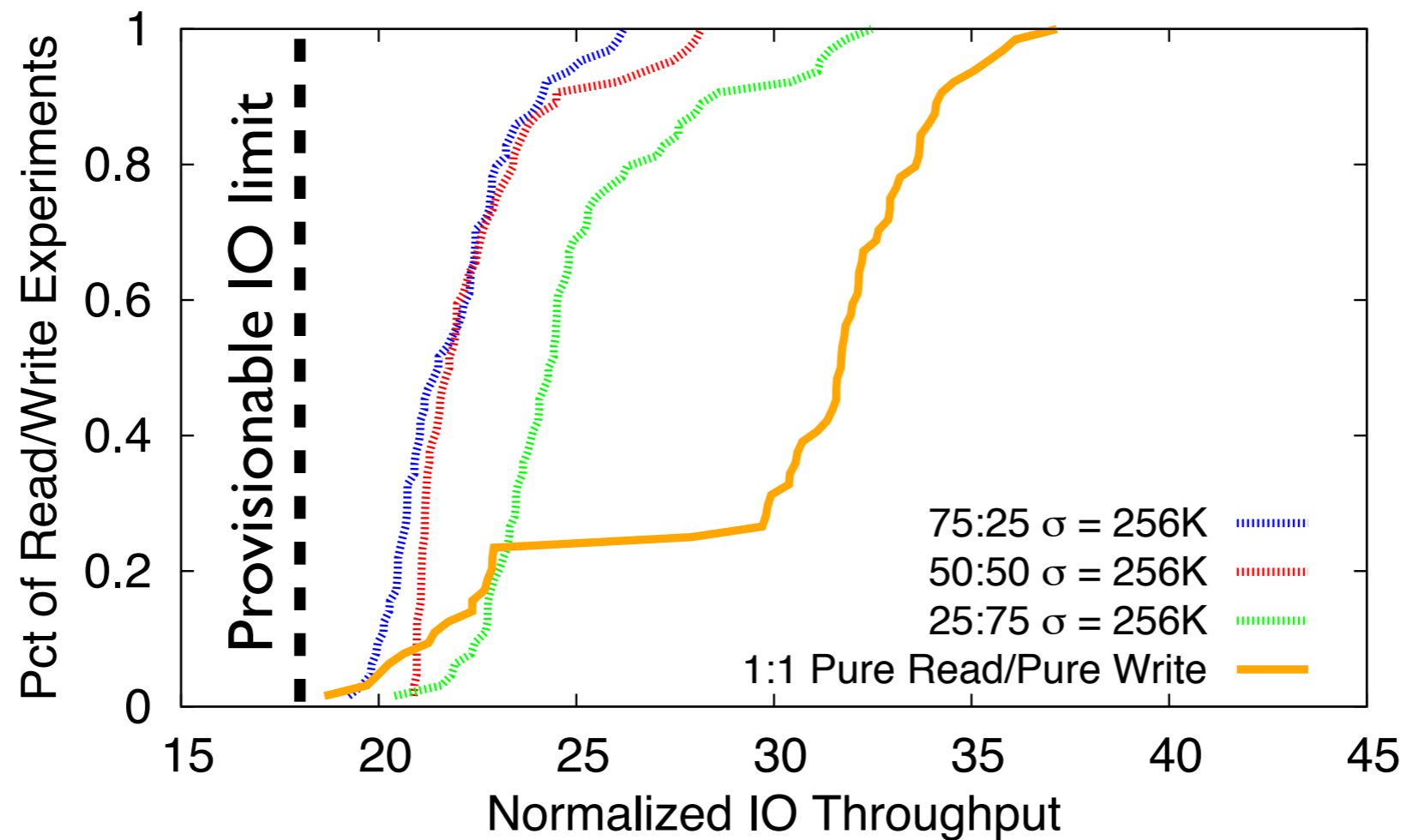
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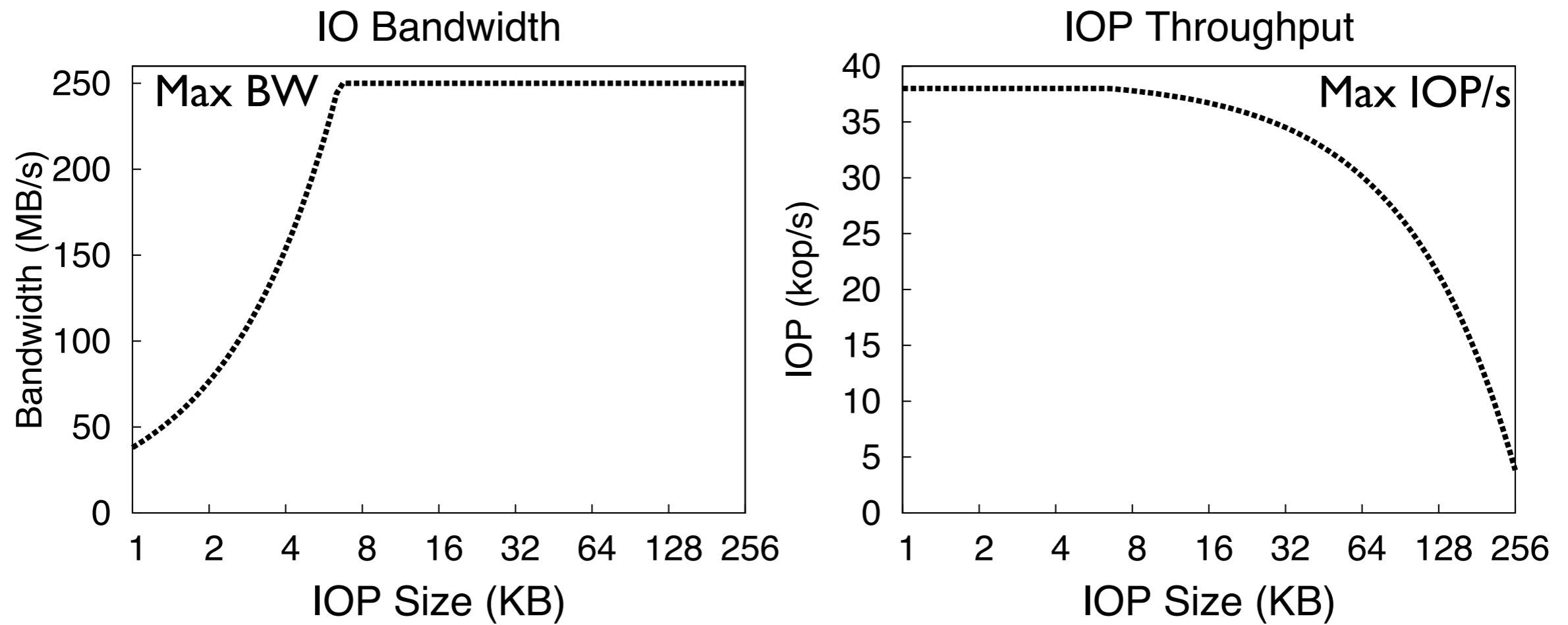
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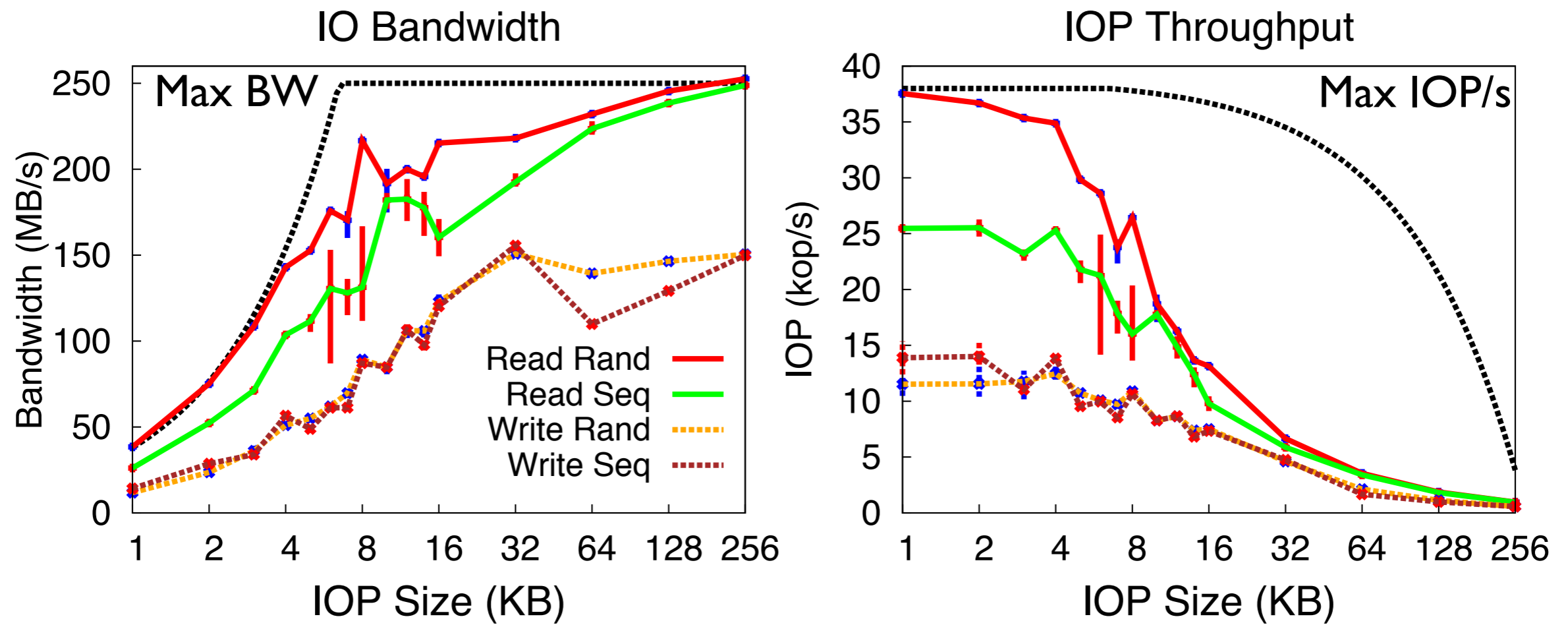




# Non-linear IO Performance



# Non-linear IO Performance

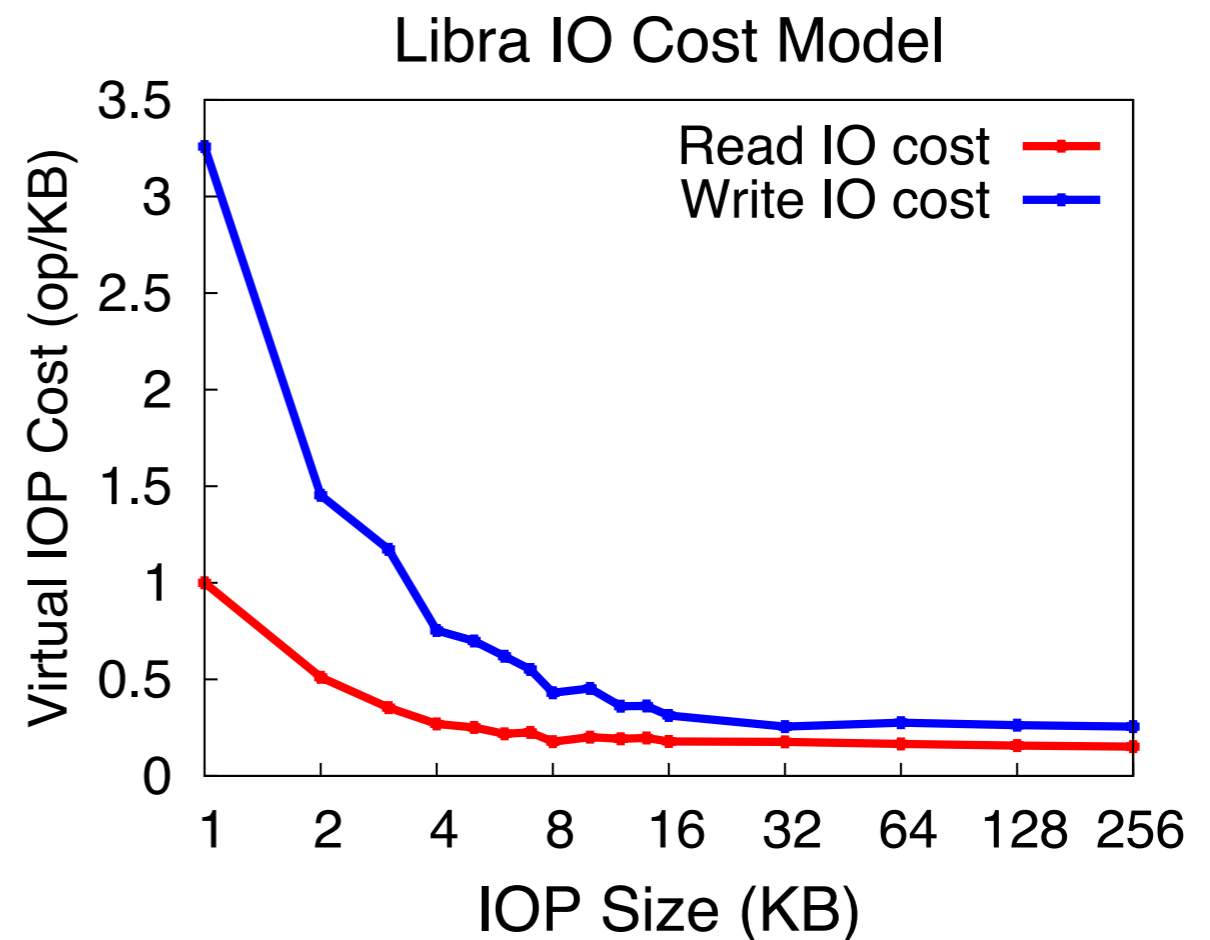


# Libra Uses Virtual IOPs to Model IO Resources

$$\text{VOP}_{\text{CPB}}(\text{IOP-size}) = \frac{\text{Max-IOP}}{\text{Achieved-IOP}(\text{IOP-size}) \times \text{IOP-size}}$$

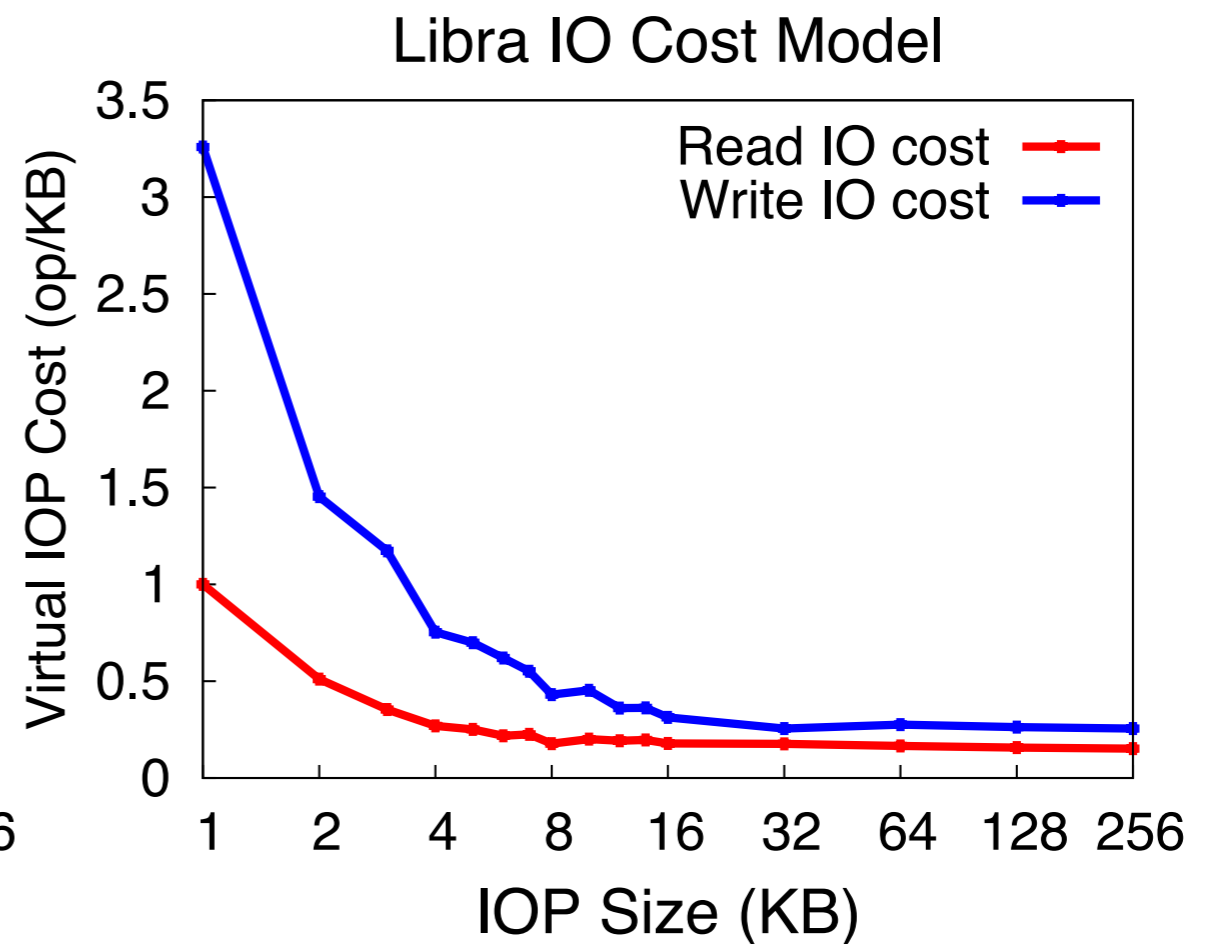
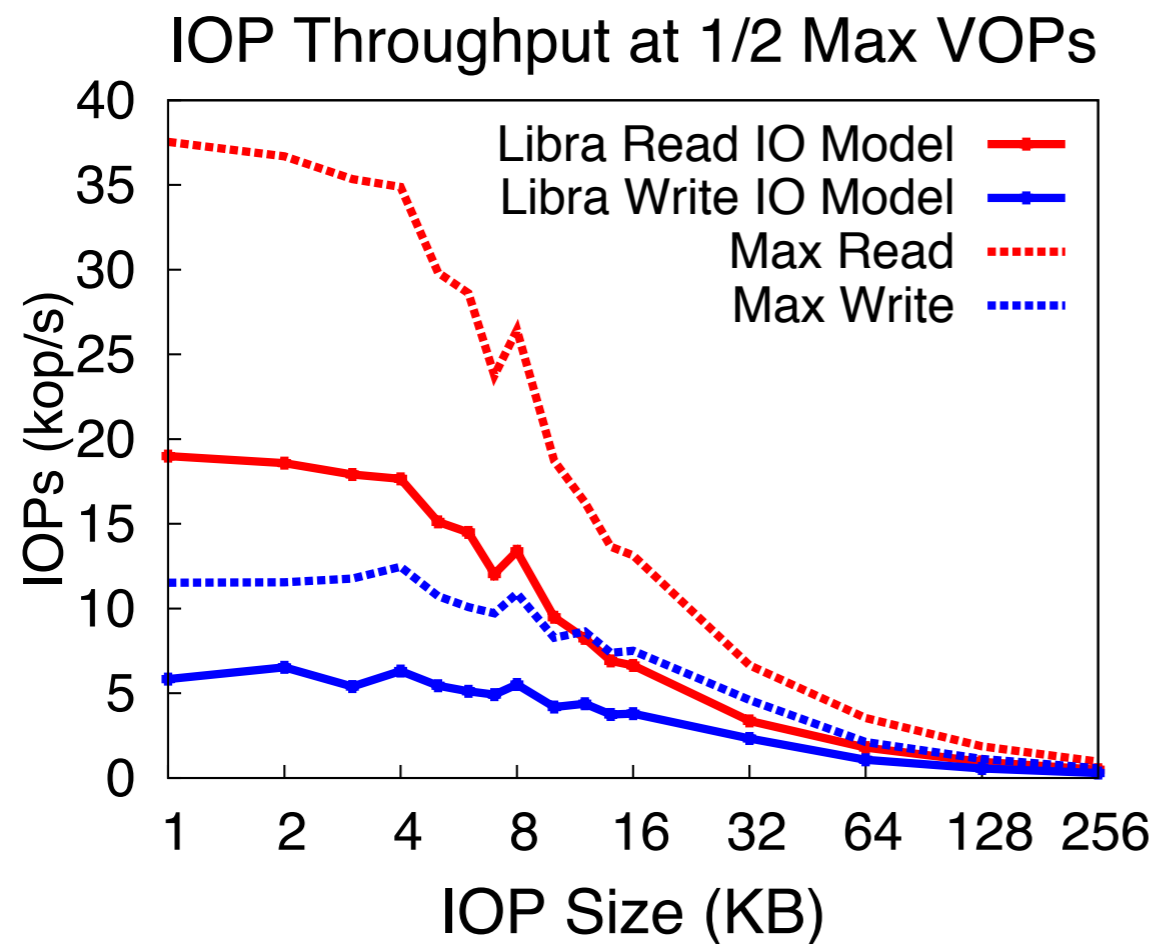
- ✓ Unifies IO cost into a single metric
- ✓ Captures non-linear IO performance
- ✓ Provides IO insulation

2 equal-allocation tenants  
IO Insulation = 1/2 Max Read/Write



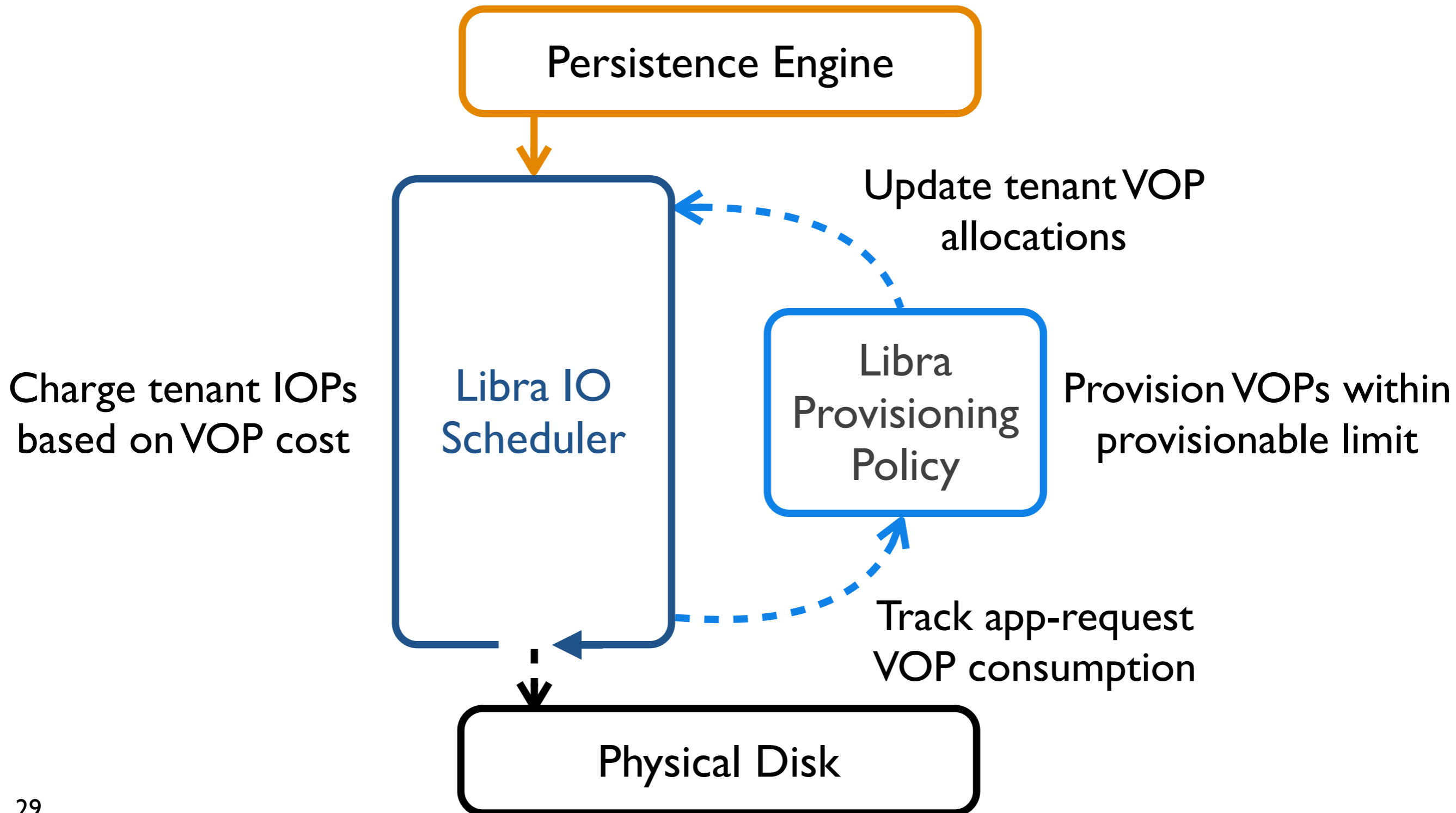
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2 equal-allocation tenants  
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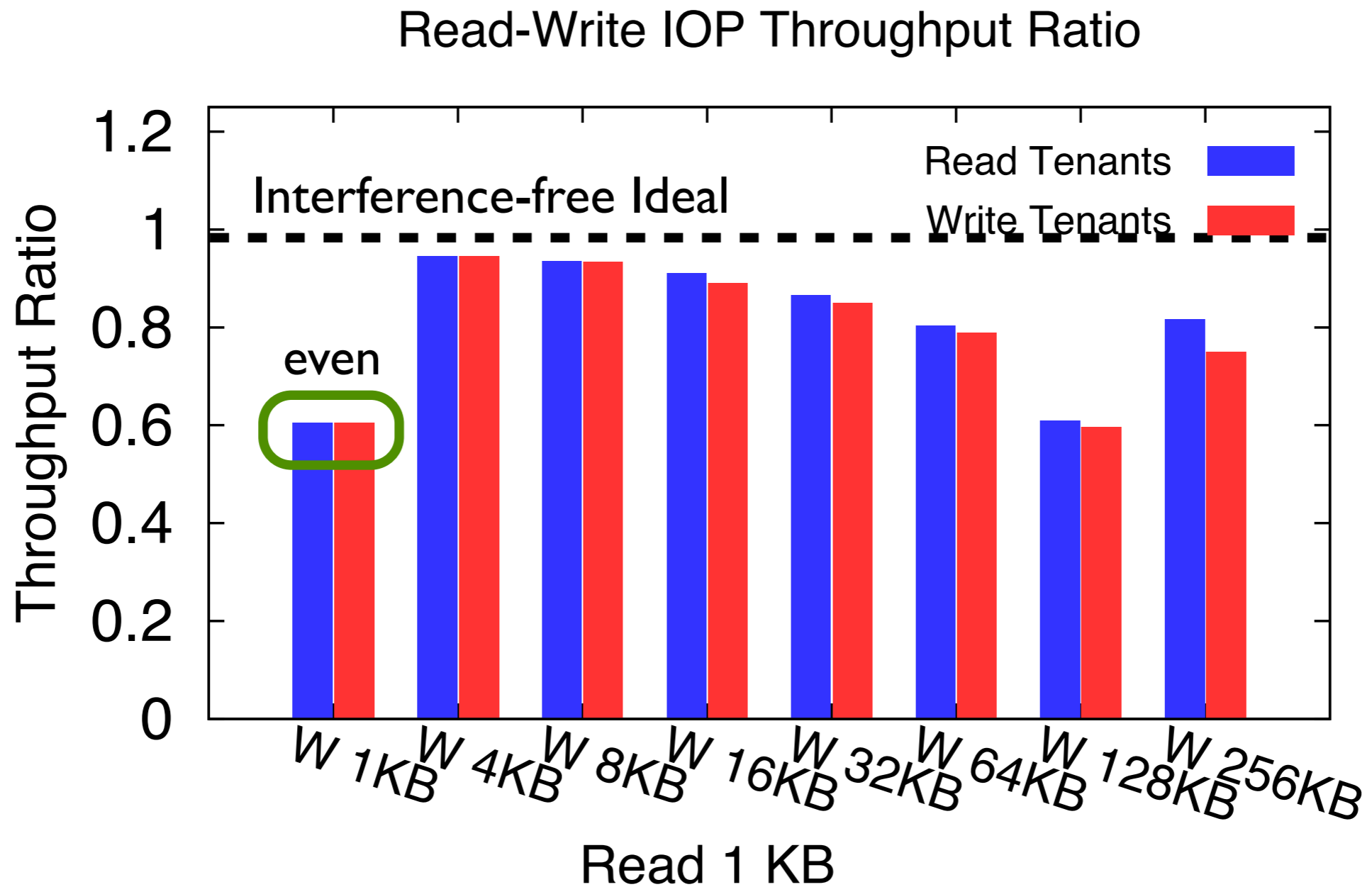
# Libra Design



# Evaluation

- Does Libra's IO resource model achieve accurate resource allocations?
- Does Libra's IO threshold make an acceptable tradeoff of performance for predictability in a real storage stack?
- Can Libra ensure per-tenant app-request reservations while achieving high utilization?

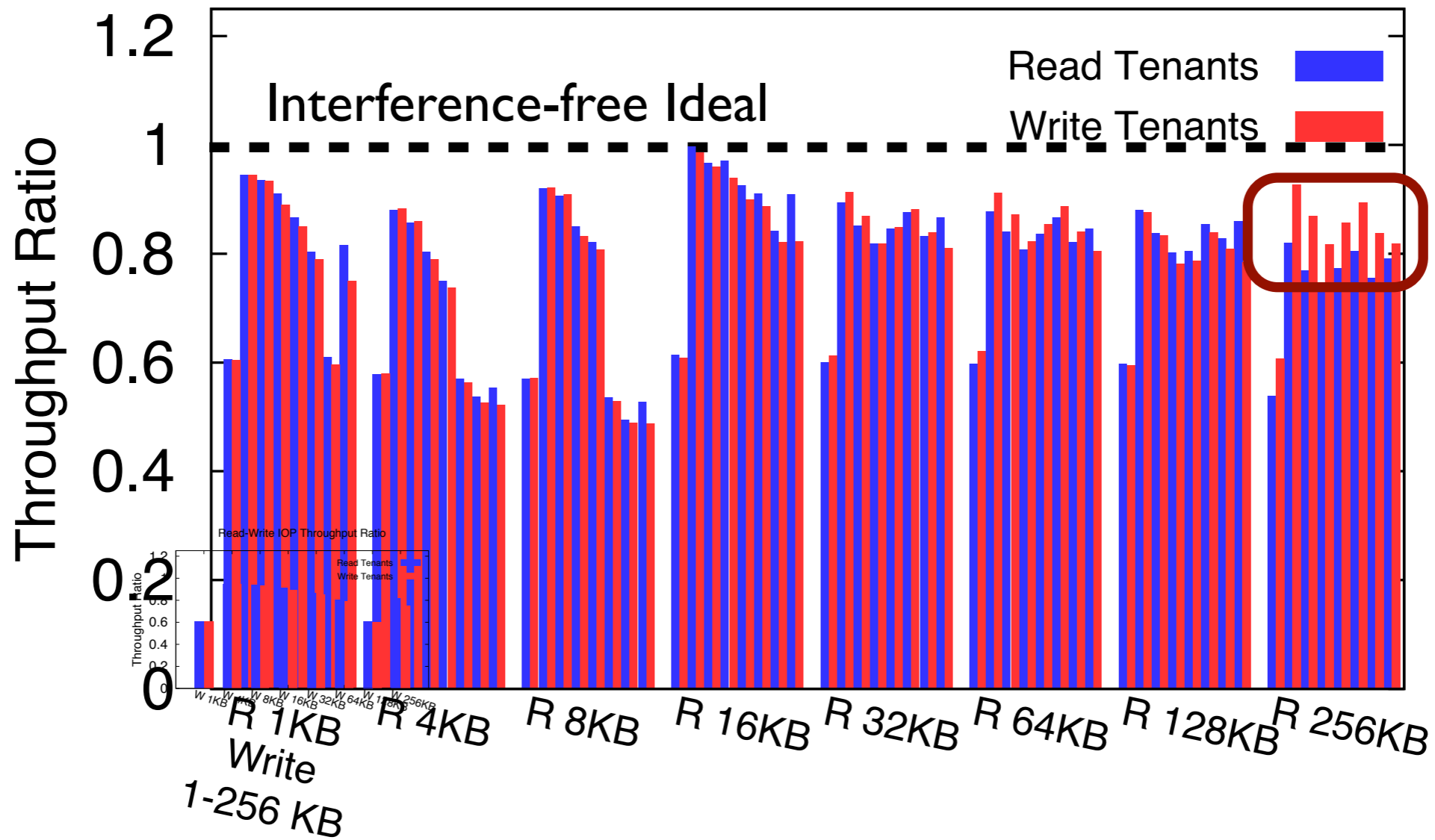
# Libra Achieves Accurate IO Allocations



Throughput Ratio = Actual / Expected (IO Insulation)

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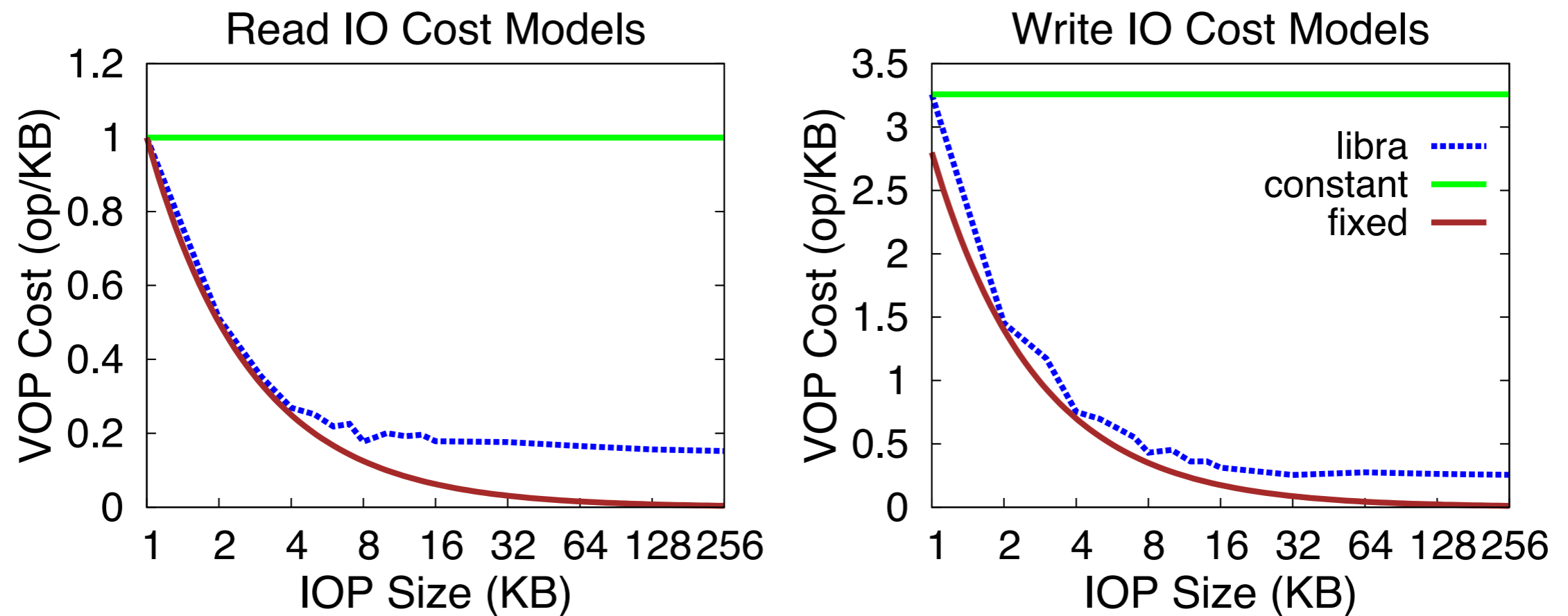
Read-Write IOP Throughput Ratio



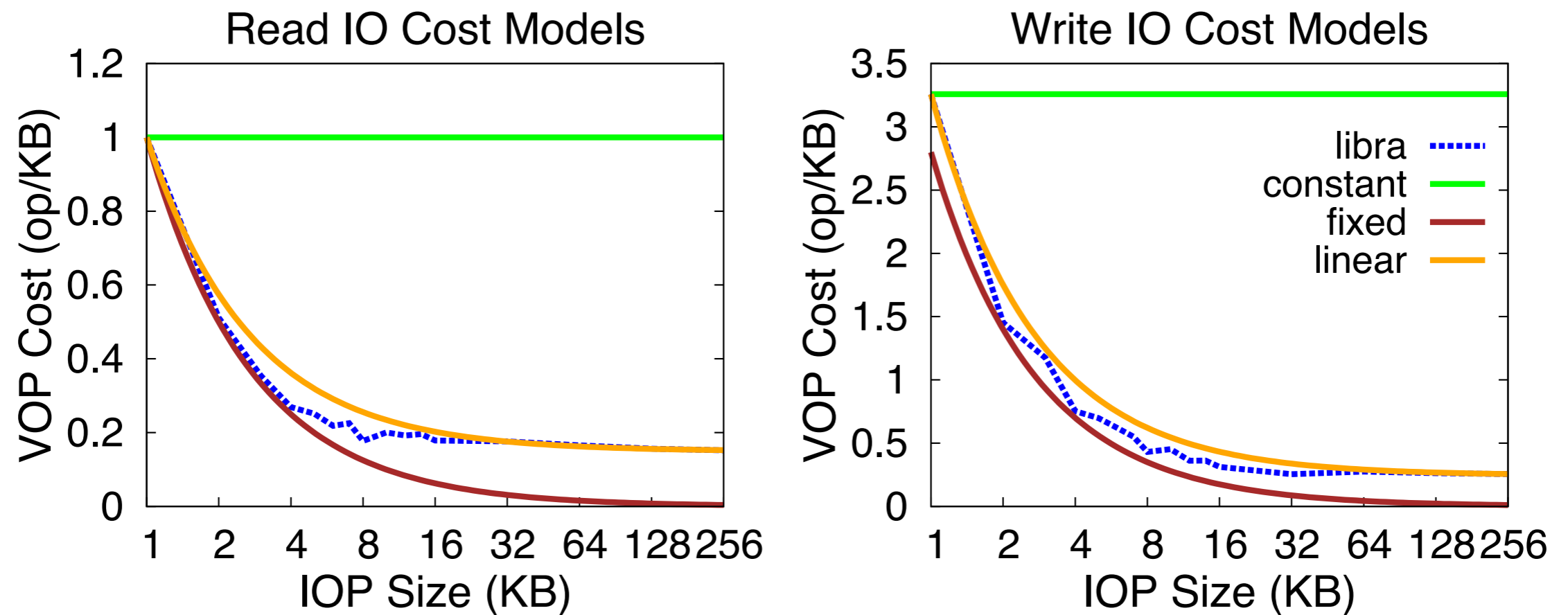
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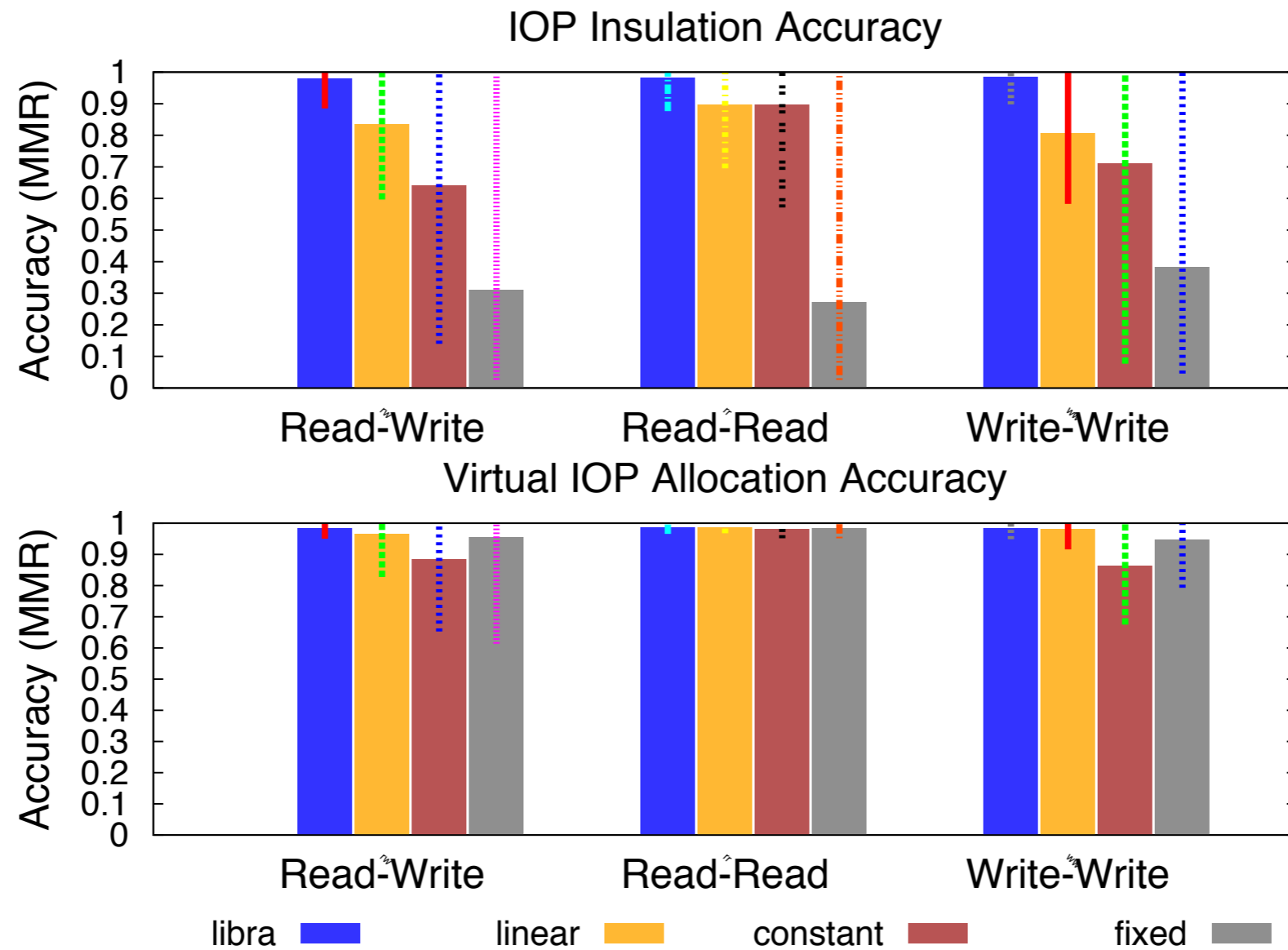


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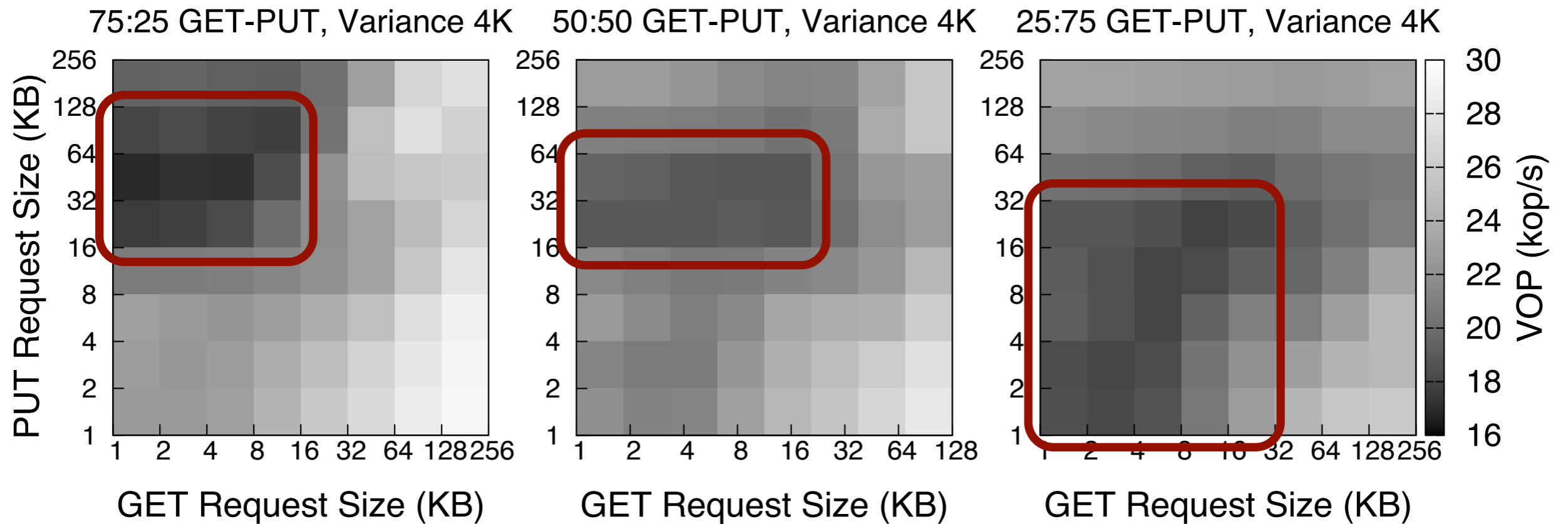


# Libra Achieves Accurate IO Allocations

Min-Max Ratio =  
 Min Throughput Ratio / Max Throughput Ratio



# Libra Trades-off Nominal IO Throughput For Predictability



# Libra Trades-off Nominal IO Throughput For Predictability

Unprovisionable Throughput As a Percentage of Total Throughput

Workload	Percentile			
	10th	50th	80th	All
99:1	1.6%	30.5%	40.5%	45.8%
25:75	1.4%	14.9%	25.0%	34.7%
1:99	0.7%	12.2%	19.5%	28.1%

< 10th percentile covered by SLA and higher-level policies

# Libra Achieves App-request Reservations

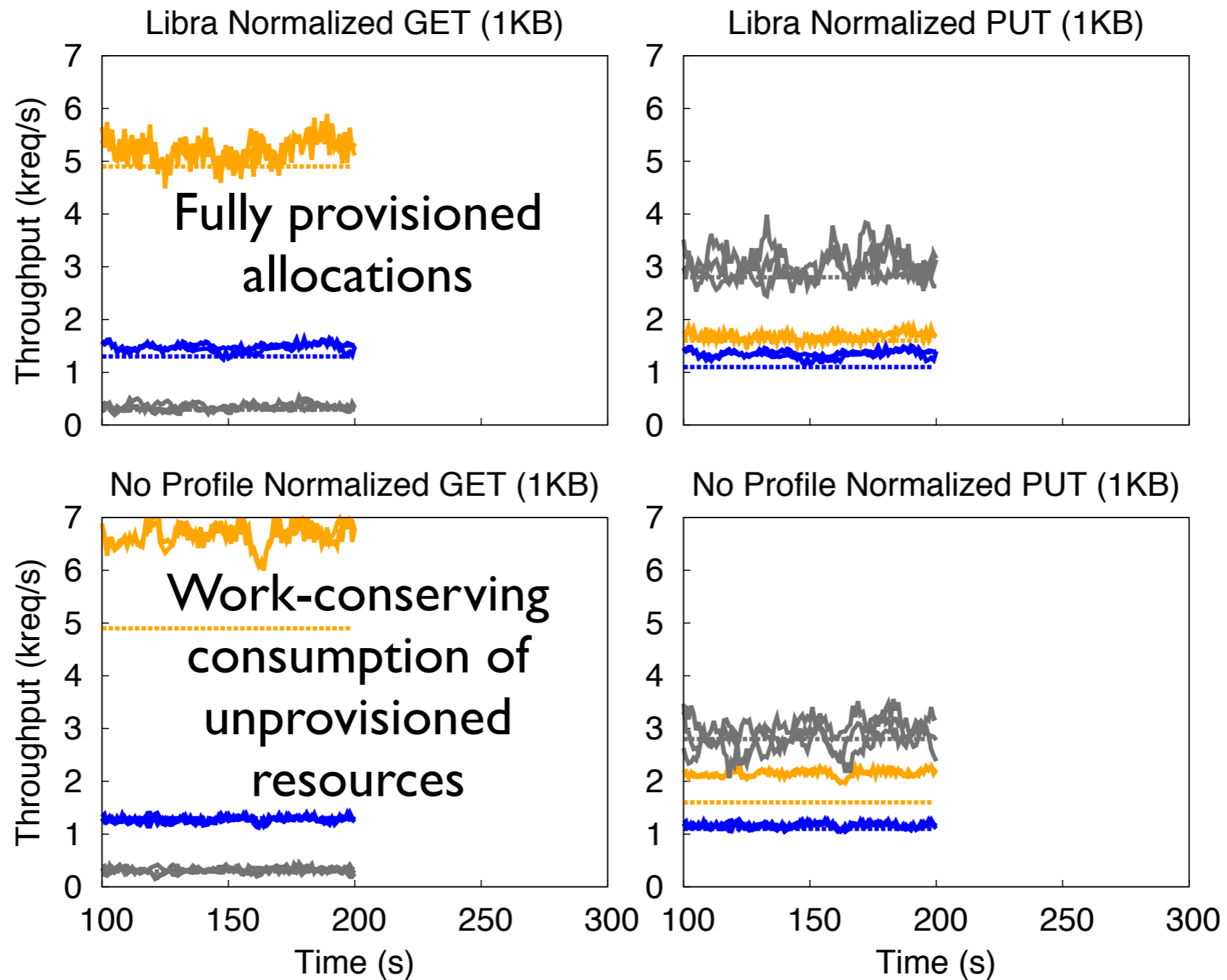
0.5x

1.5x

Read Heavy

Mixed

Write Heavy

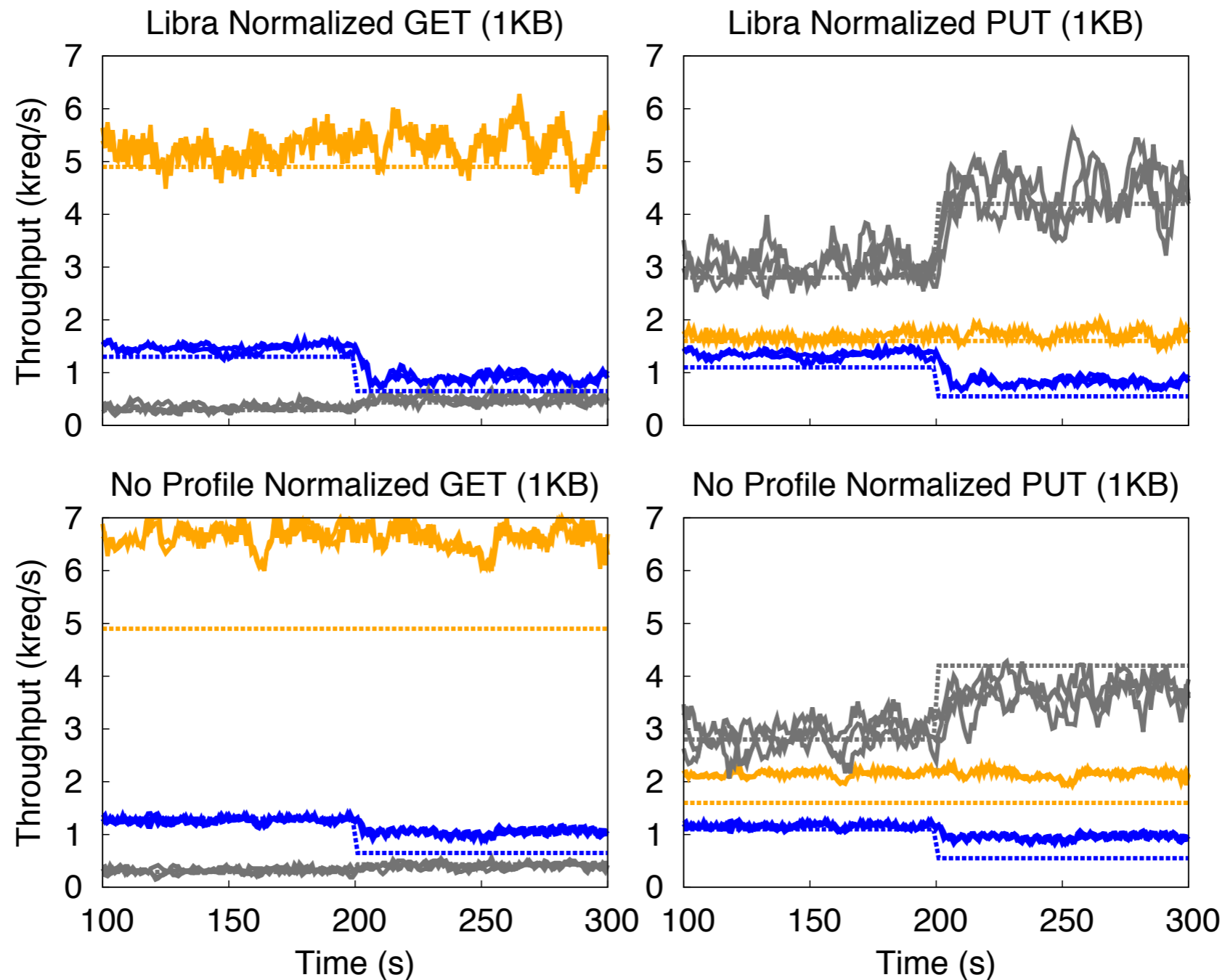


# Libra Achieves App-request Reservations

Read Heavy

Mixed

Write Heavy



# Conclusion

- **Libra IO Scheduler**

- Provisions *IO allocations* for *app-request reservations* w/ high utilization.
- Supports arbitrary object distributions and workloads.

- **2 key mechanisms**

- Track per-tenant app-request resource profiles.
- Model IO resources with Virtual IOPs.

- **Evaluation**

- Achieves accurate low-level IO allocations.
- Provisions the majority of IO resources over a wide range of workloads
- Satisfies app-request reservations w/ high utilization.