

Simple Data

COS 326

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What is the single most important mathematical concept ever developed in human history?

What is the single most important mathematical concept ever developed in human history?

An answer: The mathematical variable

Why is the mathematical variable so important?

The mathematician says:

“Let x be some integer, we define a polynomial over x ...”

Why is the mathematical variable so important?

The mathematician says:

“Let x be some integer, we define a polynomial over x ...”

What is going on here? The mathematician has separated a *definition* (of x) from its *use* (in the polynomial). This is the most primitive kind of *abstraction*.

Abstraction is the key to controlling complexity and without it, modern mathematics, science, and computation would not exist.

O'CAML BASICS: LET DECLARATIONS

Abstraction

- Good programmers identify repeated patterns in their code and factor out the repetition into meaning components
- In O'Caml, the most basic technique for factoring your code is to use **let expressions**
- Instead of writing this expression:

```
(2 + 3) * (2 + 3)
```

Abstraction & Abbreviation

- Good programmers identify repeated patterns in their code and factor out the repetition into meaning components
- In O'Caml, the most basic technique for factoring your code is to use **let expressions**
- Instead of writing this expression:

```
(2 + 3) * (2 + 3)
```

- We write this one:

```
let x = 2 + 3 in  
x * x
```


A Few More Let Expressions

```
let x = 2 in  
let squared = x * x in  
let cubed = x * squared in  
squared * cubed
```

A Few More Let Expressions

```
let two = 2.0 in  
let zero = 0.0 in  
two *. zero
```

```
let x = 2 in  
let squared = x * x in  
let cubed = x * squared in  
squared * cubed
```

```
let a = "a" in  
let b = "b" in  
let as = a ^ a ^ a in  
let bs = b ^ b ^ b in  
as ^ bs
```

How do let expressions operate?

```
let x = 2 + 1 in x * x
```

How do let expressions operate?

```
let x = 2 + 1 in x * x
```

-->

```
let x = 3 in x * x
```

How do let expressions operate?

```
let x = 2 + 1 in x * x
```


-->

```
let x = 3 in x * x
```

-->

```
3 * 3
```

substitute
3 for x



How do let expressions operate?

```
let x = 2 + 1 in x * x
```

-->

```
let x = 3 in x * x
```


-->

```
3 * 3
```

-->

```
9
```

substitute
3 for x



How do let expressions operate?

```
let x = 2 + 1 in x * x
```

-->

```
let x = 3 in x * x
```


-->

```
3 * 3
```

-->

```
9
```

substitute
3 for x



Note: I write
 $e1 \rightarrow e2$
when $e1$ evaluates
to $e2$ in one step


Another Example

```
let x = 2 in  
let y = x + x in  
y * x
```


Another Example

```
let x = 2 in  
let y = x + x in  
y * x
```

substitute
2 for x




-->

```
let y = 2 + 2 in  
y * 2
```

Another Example

```
let x = 2 in  
let y = x + x in  
y * x
```

substitute
2 for x



-->

```
let y = 2 + 2 in  
y * 2
```

-->

```
let y = 4 in  
y * 2
```

Another Example

```
let x = 2 in  
let y = x + x in  
y * x
```

substitute
2 for x




-->

```
let y = 2 + 2 in  
y * 2
```

-->

```
let y = 4      in  
y * 2
```

substitute
4 for y



-->

```
4 * 2
```

Another Example

```
let x = 2 in  
let y = x + x in  
y * x
```

substitute
2 for x

-->

```
let y = 2 + 2 in  
y * 2
```

-->

```
let y = 4 in  
y * 2
```

substitute
4 for y

-->

```
4 * 2
```

-->


```
8
```

Moral: Let
operates by
substituting
computed values
for variables

Abstraction & Abbreviation

- Two kinds of let:

```
if tuesday() then
    let x = 2 + 3 in
    x + x
else
    0
;;
```




let ... in ... is an expression that declares a local variable for temporary use and produce a value

Abstraction & Abbreviation

- Two kinds of let:


```
if tuesday() then
    let x = 2 + 3 in
    x + x
else
    0
;;
```



let ... in ... is an *expression* that can appear inside any other *expression*

The scope of x does not extend outside the enclosing “in”

```
let x = 2 + 3 ;;
let y = x + 17 / x ;;
```

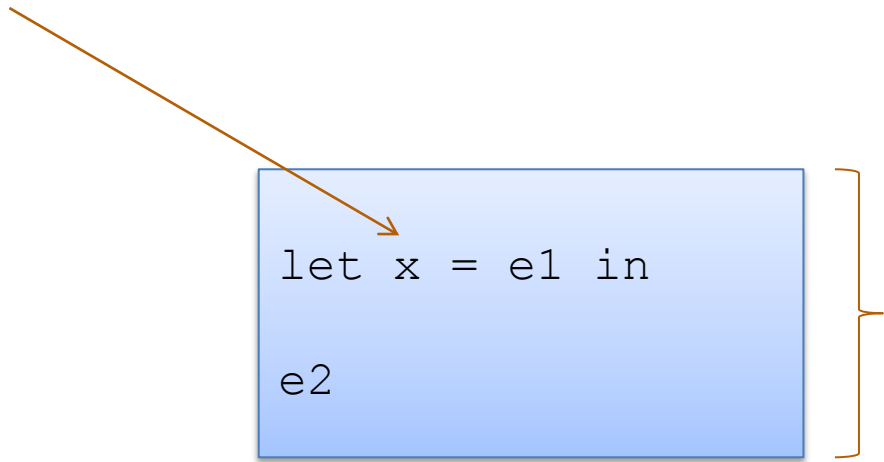


let ... ;; is a top-level *declaration* that appears at the top-level only.

Variables x and y may be exported; used by other modules

Typing Simple Let Expressions

x granted type of e1 for use in e2



let x = e1 in
e2

The diagram shows a blue rectangular box containing the text 'let x = e1 in' on the first line and 'e2' on the second line. An orange arrow points from the text 'x granted type of e1 for use in e2' to the variable 'x' in the first line. A brown curly bracket on the right side of the box encompasses both lines of code.

overall expression
takes on the type of e2

Typing Simple Let Expressions

x granted type of e1 for use in e2

```
let x = e1 in  
e2
```

overall expression
takes on the type of e2

x has type int
for use inside the
let body

```
let x = 3 + 4 in  
string_of_int x
```

overall expression
has type string

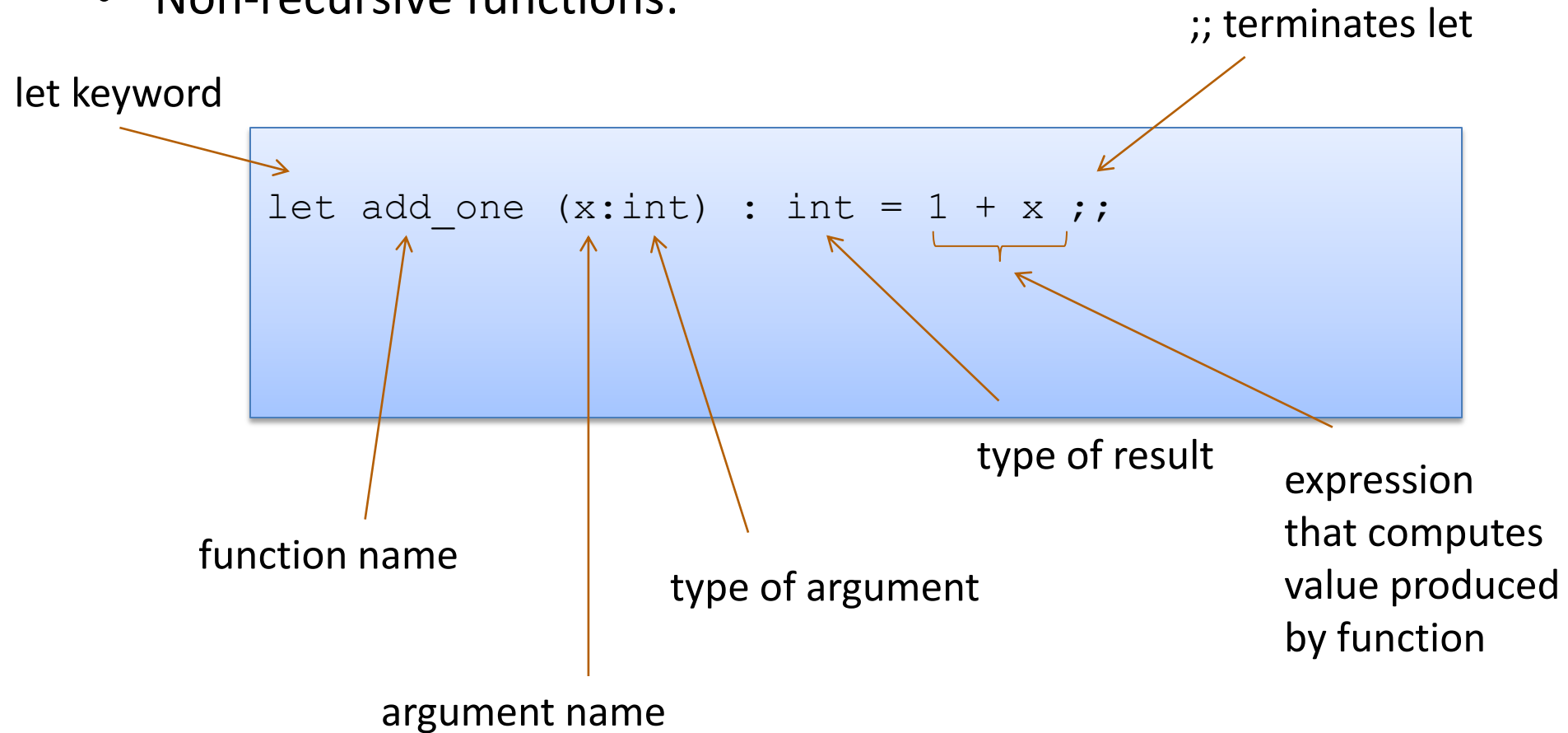
Defining functions

- Non-recursive functions:

```
let add_one (x:int) : int = 1 + x ;;
```

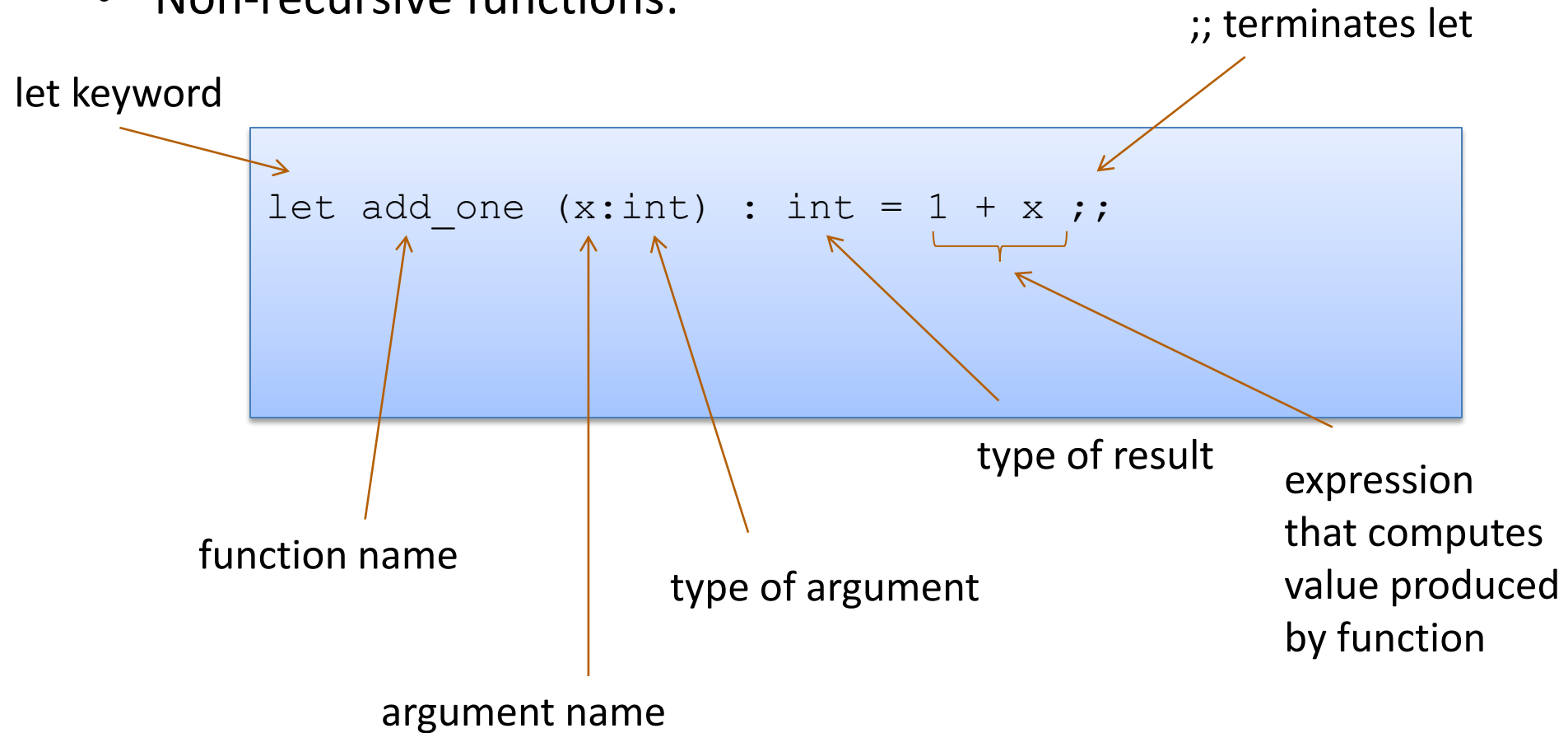
Defining functions

- Non-recursive functions:



Defining functions

- Non-recursive functions:



Note: recursive functions with begin with **"let rec"**

Defining functions

- Non-recursive functions:

```
let add_one (x:int) : int = 1 + x ;;
```

```
let add_two (x:int) : int = add_one (add_one x) ;;
```

Defining functions

- Non-recursive functions:

```
let add_one (x:int) : int = 1 + x ;;  
let add_two (x:int) : int = add_one (add_one x) ;;
```

- With a local definition:

```
let add_two' (x:int) : int =  
  let add_one x = 1 + x in  
  add_one (add_one x)  
;;
```

local function definition
hidden from clients

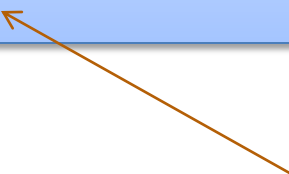
I left off the types.
O'Caml figures them out

Good style: types on
top-level definitions

Types for Functions

- Some functions:

```
let add_one (x:int) : int = 1 + x ;;  
let add_two (x:int) : int = add_one (add_one x) ;;  
let add (x:int) (y:int) : int = x + y ;;
```



- Types for functions: function with two arguments

```
add_one : int -> int  
add_two : int -> int  
add : int -> int -> int
```

Rule for type-checking functions

General Rule:

If a function $f : T1 \rightarrow T2$
and an argument $e : T1$
then $f e : T2$

Example:

```
add_one : int -> int
```

```
3 + 4 : int
```

```
add_one (3 + 4) : int
```

Rule for type-checking functions

- Recall the type of add:

Definition:

```
let add (x:int) (y:int) : int =  
  x + y  
;;
```

Type:

```
add : int -> int -> int
```


Rule for type-checking functions

- Recall the type of add:

Definition:

```
let add (x:int) (y:int) : int =  
  x + y  
;;
```

Type:

```
add : int -> int -> int
```

Same as:

```
add : int -> (int -> int)
```

Rule for type-checking functions

General Rule:

If a function $f : T1 \rightarrow T2$
and an argument $e : T1$
then $f e : T2$

Remember:

$A \rightarrow B \rightarrow C$
is the same as
 $A \rightarrow (B \rightarrow C)$

Example:

```
add : int -> int -> int
```

```
3 + 4 : int
```

```
add (3 + 4) : ???
```

Rule for type-checking functions

General Rule:

If a function $f : T1 \rightarrow T2$
and an argument $e : T1$
then $f e : T2$

Remember:

$A \rightarrow B \rightarrow C$
is the same as
 $A \rightarrow (B \rightarrow C)$

Example:

```
add : int -> (int -> int)
```

```
3 + 4 : int
```

```
add (3 + 4) :
```

Rule for type-checking functions

General Rule:

If a function $f : T1 \rightarrow T2$
and an argument $e : T1$
then $f e : T2$

Remember:

$A \rightarrow B \rightarrow C$
is the same as
 $A \rightarrow (B \rightarrow C)$

Example:

```
add : int -> (int -> int)
3 + 4 : int
add (3 + 4) : int -> int
```

Rule for type-checking functions

General Rule:

If a function $f : T1 \rightarrow T2$
and an argument $e : T1$
then $f e : T2$

Remember:

$A \rightarrow B \rightarrow C$
is the same as
 $A \rightarrow (B \rightarrow C)$


Example:

```
add : int -> int -> int
```

```
3 + 4 : int
```

```
add (3 + 4) : int -> int
```

```
(add (3 + 4)) 7 : int
```



Rule for type-checking functions

General Rule:

If a function $f : T1 \rightarrow T2$
and an argument $e : T1$
then $f e : T2$

Remember:

$A \rightarrow B \rightarrow C$
is the same as
 $A \rightarrow (B \rightarrow C)$

Example:

```
add : int -> int -> int
```

```
3 + 4 : int
```

```
add (3 + 4) : int -> int
```

```
add (3 + 4) 7 : int
```

Rule for type-checking functions

Example:

```
let munge (b:bool) (x:int) : ?? =  
  if not b then  
    string_of_int x  
  else  
    "hello"  
;;  
  
let y = 17;;
```

```
munge (y > 17) : ??
```

```
munge true (f (munge false 3)) : ??  
f : ??
```

```
munge true (g munge) : ??  
g : ??
```

Rule for type-checking functions

Example:

```
let munge (b:bool) (x:int) : ?? =  
  if not b then  
    string_of_int x  
  else  
    "hello"  
;;  
  
let y = 17;;
```

```
munge (y > 17) : ??
```

```
munge true (f (munge false 3)) : ??  
  f : string -> int
```

```
munge true (g munge) : ??  
  g : (bool -> int -> string) -> int
```


One key thing to remember

- If you have a function f with a type like this:

$A \rightarrow B \rightarrow C \rightarrow D \rightarrow E \rightarrow F$

- Then each time you add an argument, you can get the type of the result by knocking off the first type in the series

$f\ a1 : B \rightarrow C \rightarrow D \rightarrow E \rightarrow F$ (if $a1 : A$)

$f\ a1\ a2 : C \rightarrow D \rightarrow E \rightarrow F$ (if $a2 : B$)

$f\ a1\ a2\ a3 : D \rightarrow E \rightarrow F$ (if $a3 : C$)


$f\ a1\ a2\ a3\ a4\ a5 : F$ (if $a4 : D$ and $a5 : E$)

Binding Variables to Values

- Each O'Caml variable is *bound* to 1 value
- The value to which a variable is bound to never changes!

```
let x = 3 ;;
```

```
let add_three (y:int) : int = y + x ;;
```




Binding Variables to Values


- Each O'Caml variable is *bound* to 1 value
- The value to which a variable is bound to never changes!

```
let x = 3 ;;
```

```
let add_three (y:int) : int = y + x ;;
```



*It does not
matter what
I write next.
add_three
will always
add 3!*



Binding Variables to Values

- Each O'Caml variable is bound to 1 value
- The value a variable is bound to never changes!

a distinct
variable that
"happens to
be spelled the
same"

```
let x = 3 ;;
```

```
let add_three (y:int) : int = y + x ;;
```

```
let x = 4 ;;
```

```
let add_four (y:int) : int = y + x ;;
```

Binding Variables to Values

- Since the 2 variables (both happened to be named x) are actually different, unconnected things, we can rename one of them

rename x
to zzz
if you want
to, replacing
its uses

```
let x = 3 ;;
```

```
let add_three (y:int) : int = y + x ;;
```

```
let zzz = 4 ;;
```

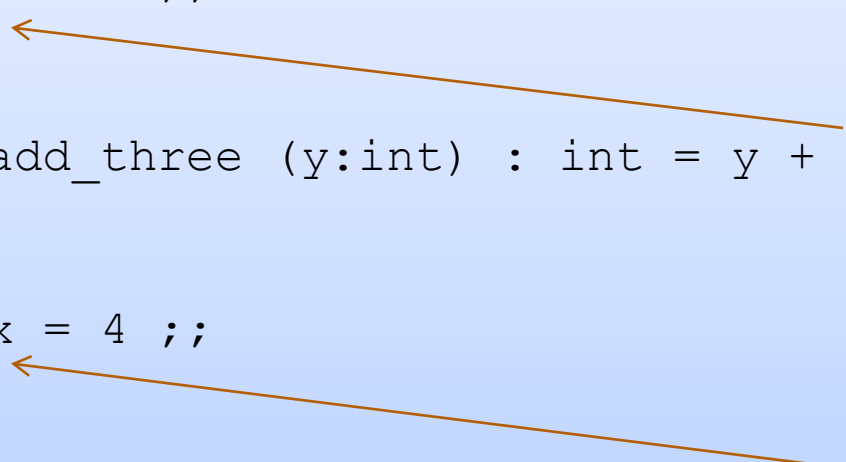
```
let add_four (y:int) : int = y + zzz ;;
```

```
let add_seven (y:int) : int =  
  add_three (add_four y)  
;;
```

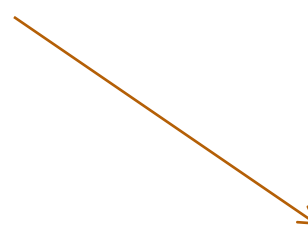
Binding Variables to Values

- Each O'Caml variable is bound to 1 value
- O'Caml is a **statically scoped** language

```
let x = 3 ;;  
  
let add_three (y:int) : int = y + x ;;  
  
let x = 4 ;;  
  
let add_four (y:int) : int = y + x ;;  
  
let add_seven (y:int) : int =  
  add_three (add_four y)  
;;
```



we can use
add_three
without worrying
about the second
definition of x



OUR FIRST* COMPLEX DATA STRUCTURE! THE TUPLE

* it is really our second complex data structure since functions are data structures too!

Tuples

- A tuple is a fixed, finite, ordered collection of values
- Some examples with their types:

```
(1, 2) : int * int
```

```
("hello", 7 + 3, true) : string * int * bool
```

```
('a', ("hello", "goodbye")) : char * (string * string)
```


Tuples

- To use a tuple, we extract its components
- General case:

```
let (id1, id2, ..., idn) = e1 in e2
```

- An example:

```
let (x, y) = (2, 4) in x + x + y
```

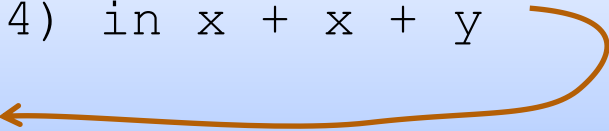
Tuples

- To use a tuple, we extract its components
- General case:

```
let (id1, id2, ..., idn) = e1 in e2
```

- An example:

```
let (x, y) = (2, 4) in x + x + y  
--> 2 + 2 + 4
```



substitute!

Tuples

- To use a tuple, we extract its components
- General case:

```
let (id1, id2, ..., idn) = e1 in e2
```

- An example:

```
let (x, y) = (2, 4) in x + x + y  
--> 2 + 2 + 4  
--> 8
```

Rules for Typing Tuples

if $e1 : t1$ and $e2 : t2$
then $(e1, e2) : t1 * t2$

Rules for Typing Tuples

if $e1 : t1$ and $e2 : t2$
then $(e1, e2) : t1 * t2$

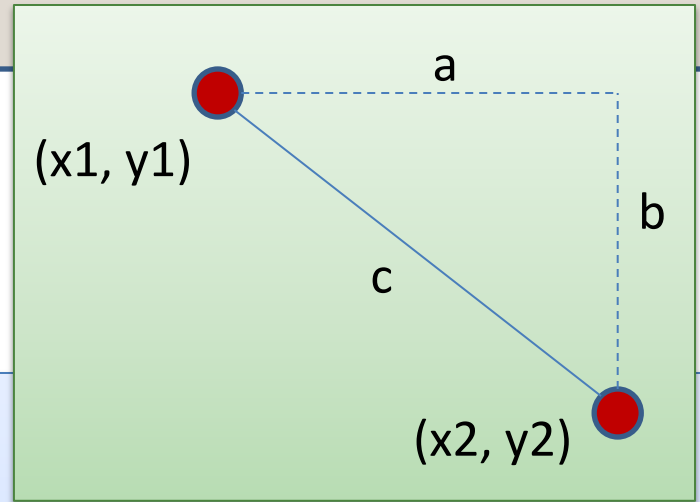
if $e1 : t1 * t2$ then
 $x1 : t1$ and $x2 : t2$
inside the expression $e2$

let $(x1, x2) = e1$ in
 $e2$

overall expression
takes on the type of $e2$

Distance between two points

$$c^2 = a^2 + b^2$$



Problem:

- A point is represented as a pair of floating point values.
- Write a function that takes in two points as arguments and returns the distance between them as a floating point number

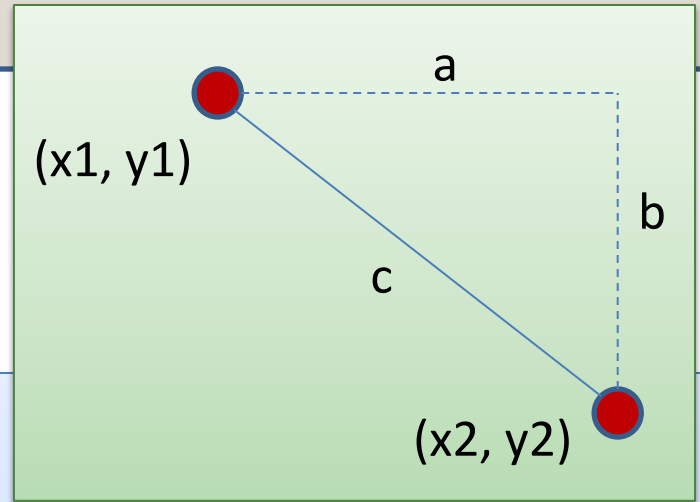
Writing Functions Over Typed Data

- Steps to writing functions over typed data:
 1. **Write down** the function and argument **names**
 2. **Write down** argument and result **types**
 3. **Write down** some examples (in a comment)
 4. **Deconstruct** input data structures
 - the argument types may suggest how to do it
 5. **Build** new output values
 - the result type may suggest how you do it
 6. **Clean up** by identifying repeated patterns
 - define and reuse helper functions
 - your code should be elegant and easy to read

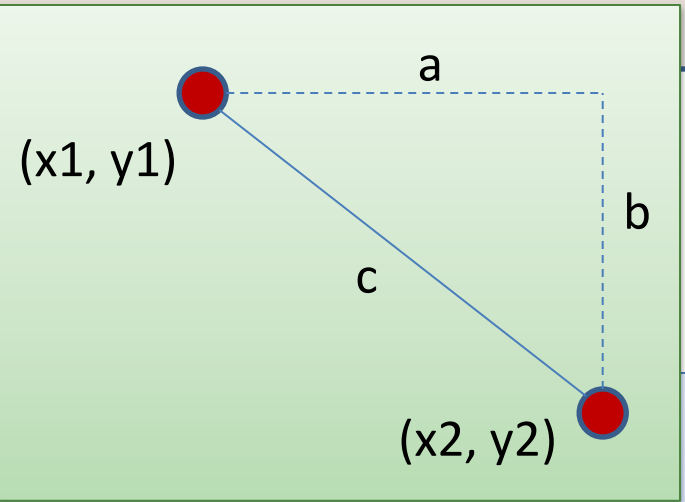
Distance between two points

a type abbreviation

```
type point = float * float
```



Distance between two points



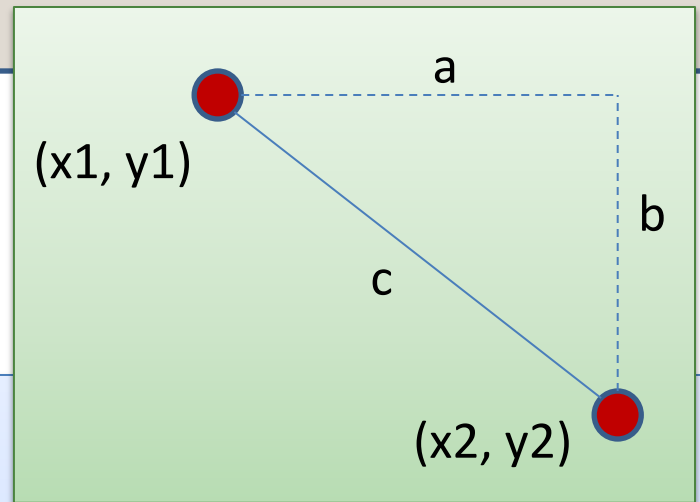
```
type point = float * float
```

```
let distance (p1:point) (p2:point) : float =
```

```
;;
```

write down function name
argument names and types

Distance between two points



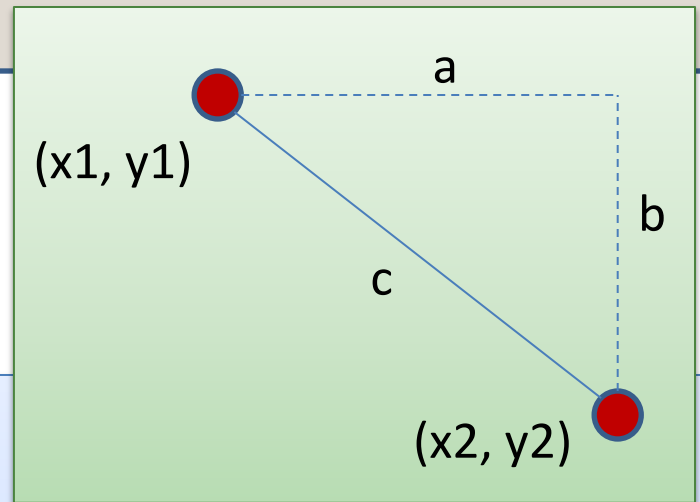
examples

```
type point = float * float
```

```
(* distance (0.0,0.0) (0.0,1.0) == 1.0
 * distance (0.0,0.0) (1.0,1.0) == sqrt(1.0 + 1.0)
 *
 * from the picture:
 * distance (x1,y1) (x2,y2) == sqrt(a^2 + b^2)
 *)
```

```
let distance (p1:point) (p2:point) : float =
```

Distance between two points



```
type point = float * float
```

```
let distance (p1:point) (p2:point) : float =
```

```
  let (x1,y1) = p1 in
```

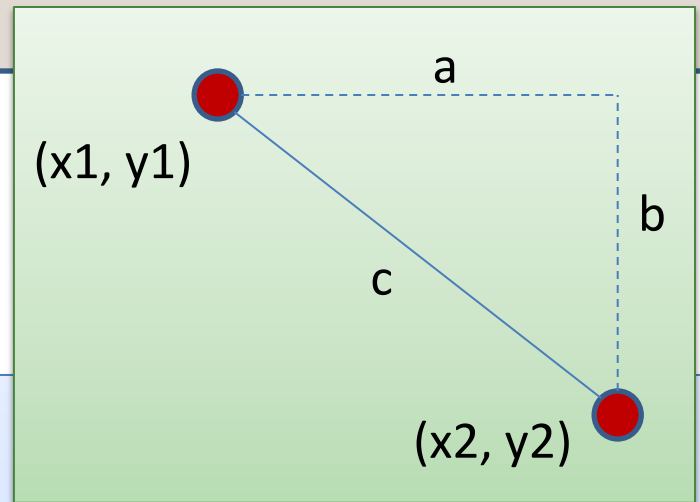
```
  let (x2,y2) = p2 in
```

```
  ...
```

```
;;
```

deconstruct
function inputs

Distance between two points



```
type point = float * float
```

```
let distance (p1:point) (p2:point) : float =
```

```
  let (x1,y1) = p1 in
```

```
  let (x2,y2) = p2 in
```

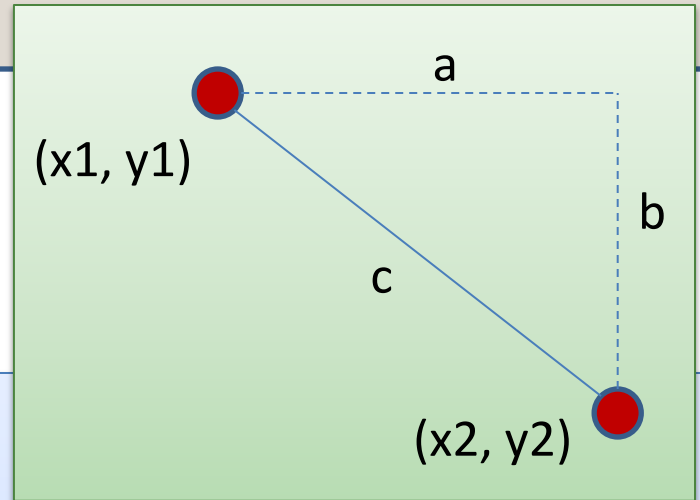
```
  sqrt ((x2 -. x1) *. (x2 -. x1) +.  
        (y2 -. y1) *. (y2 -. y1))
```

```
;;
```

} compute
function
results

notice operators on
floats have a "." in them

Distance between two points



```
type point = float * float
```

```
let distance (p1:point) (p2:point) : float =
```

```
  let square x = x *. x in
```

```
  let (x1,y1) = p1 in
```

```
  let (x2,y2) = p2 in
```

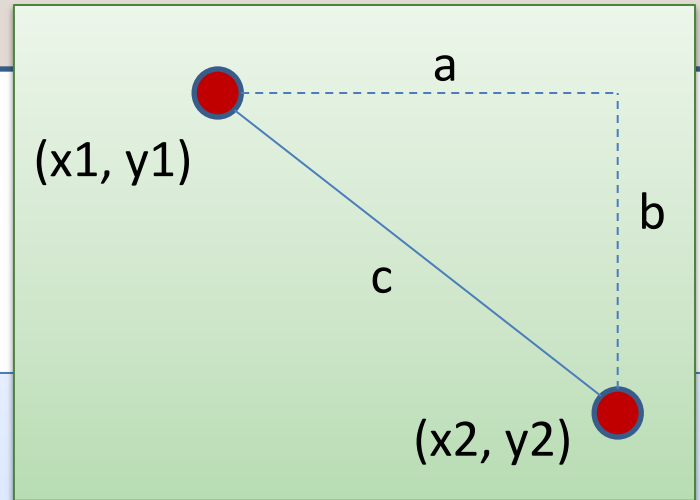
```
  sqrt (square (x2 -. x1)) +.
```

```
        square (y2 -. y1))
```

```
;;
```

define helper functions to
avoid repeated code

Distance between two points



```
type point = float * float

let distance (p1:point) (p2:point) : float =
  let square x = x *. x in
  let (x1,y1) = p1 in
  let (x2,y2) = p2 in
  sqrt (square (x2 -. x1) +. square (y2 -. y1))
;;

let pt1 = (2.0,3.0);;
let pt2 = (0.0,1.0);;
let dist12 = distance pt1 pt2;;
```

testing

SUMMARY:

BASIC FUNCTIONAL PROGRAMMING

Writing Functions Over Typed Data

- Steps to writing functions over typed data:
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 - define and reuse helper functions
 - your code should be elegant and easy to read

Writing Functions Over Typed Data

- Steps to writing functions over typed data:
 1. Write down the function and argument names
 2. Write down argument and result types
 3. Write down some examples (in a comment)
 4. **Deconstruct** input data structures
 5. **Build** new output values
 6. Clean up by identifying repeated patterns
- For tuples:
 - when the **input** has type $t1 * t2$
 - use `let (x,y) = ...` to **deconstruct**
 - when the **output** has type $t1 * t2$
 - use `(e1, e2)` to **construct**
- We will see this paradigm repeat itself over and over

END