

# COS 423: Projects 4

due Wed 05/16/2012

## Project 1: Generation of Simple Polygons

Design, analyze and implement an algorithm for constructing simple polygons. You may use a probabilistic algorithm. The sole input to your program is an integer  $n$ . The output should be a “complicated” simple polygon with  $n$  vertices. The quality of your program will depend on speed, but more important, on how winding, irregular, and cluster-free your polygons are. You should produce a number of printouts to display your polygons for various values of  $n$ .

## Project 2: Low-Cutting Paths

Implement an algorithm that given  $n$  points in the plane connects them in a simple path that no line can cut in more than  $O(\sqrt{n})$  points. Give graphical evidence that your program works.

## Project 3: Voronoi Diagrams

Learn about and implement Fortune’s Voronoi diagram algorithm (see lecture notes). Provide graphical evidence that your code works.

## Project 4: Minimum Enclosing Disk

Implement the minimum-enclosing disk algorithm and provide graphical evidence that your code works.

## Project 5: Linear Programming

Implement the randomized LP algorithm in fixed, arbitrary dimension, and provide graphical evidence that your code works.

## Project 6: Random Independent Sets

Consider an undirected, connected graph  $G = (V, E)$  and define the following Markov chain: a state is any independent set in  $G$  (ie, any subset of nodes with no two of them adjacent to each other); The chain has an edge from  $S$  to  $S'$  if they differ in exactly one element. The transition probabilities are inferred by the following process. Given the current state  $S_t$ , pick a random node  $v$  uniformly in  $V$ : (i) if  $v \in S_t$ , then set  $S_{t+1} = S_t \setminus \{v\}$ ; (ii) if  $v$  is neither in  $S_t$  nor adjacent to any node in  $S_t$ ,

then set  $S_{t+1} = S_t \cup \{v\}$ ; (iii) in all other cases, set  $S_{t+1} = S_t$ . Prove that the chain is ergodic (ie, irreducible and aperiodic). Find its stationary distribution.

It is often desired to achieve a given stationary distribution with a reversible chain. Here is the idea. Pick some  $\mu > 0$  and suppose that you want to sample a random independent set  $S$  with probability proportional to  $\mu^{|S|}$ . (The case  $\mu = 1$  gives the uniform distribution.) We modify the Markov chain as follows: Given the current state  $S_t$ , pick a random node  $v$  uniformly in  $V$ : (i) if  $v \in S_t$ , then set  $S_{t+1} = S_t \setminus \{v\}$  with probability  $\min\{1, 1/\lambda\}$ ; (ii) if  $v$  is neither in  $S_t$  nor adjacent to any node in  $S_t$ , then set  $S_{t+1} = S_t \cup \{v\}$  with probability  $\min\{1, \lambda\}$ ; (iii) in all other cases, set  $S_{t+1} = S_t$ . Prove that the chain is ergodic and time-reversible. Find its stationary distribution.

Pick a random graph  $G$  in  $\mathbf{G}(n, p)$ , ie, with  $n$  nodes and each pair of them connected at random, independently with probability  $p$ . Implement the first Markov chain above and estimate the average size of a random independent set as a function of  $p$  and  $n$ . Repeat the experiment with the second chain and study the dependency on  $\mu$ .