

# **PATTERNS IN NETWORK ARCHITECTURE:**

## **MOBILITY**

# MOBILITY

## OUTLINE

- **Roadmap**

- **Mobility**

  - Definition of mobility**

  - Discussion of “The Design Space of Network Mobility”**

  - Interoperation and layering**

  - A composition theorem**

- **Unfinished Business**

  - VLANs, again**

  - Architecture and trust boundaries**

# **ROADMAP: NETWORKING REQUIREMENTS**

## **SEARCHING OR FINDING REQUIREMENTS**

**The initiator of a communication must find the entity it wishes to communicate with, or the information it wishes to access.**

**The result of the search is a name in some network.**

## **REACHABILITY OR PROGRESS REQUIREMENTS**

**Desired communication should succeed, with adequate performance, using resources efficiently, and despite failures.**

## **SECURITY OR SAFETY REQUIREMENTS**

**Networks should not be used in damaging or malicious ways.**

# ROADMAP: NETWORK ARCHITECTURE

## BASIC NETWORKING

- naming
- routing
- forwarding
- session services

## LAYERING OR VERTICAL COMPOSITION

## BRIDGING OR HORIZONTAL COMPOSITION

## OTHER CRUCIAL TOPICS

- mobility
- middleboxes
- the granularity of change in all structures

# ROADMAP: OBJECTIVES

- 1** To understand the relationships between requirements and architectures, so we can do a better job of meeting the requirements.
- 2** To formalize network architecture just enough so that solutions to problems, and how they compose, can be fully understood.
- 3** To encourage networking practice that exposes these formalized architectural elements, so that composition always works.

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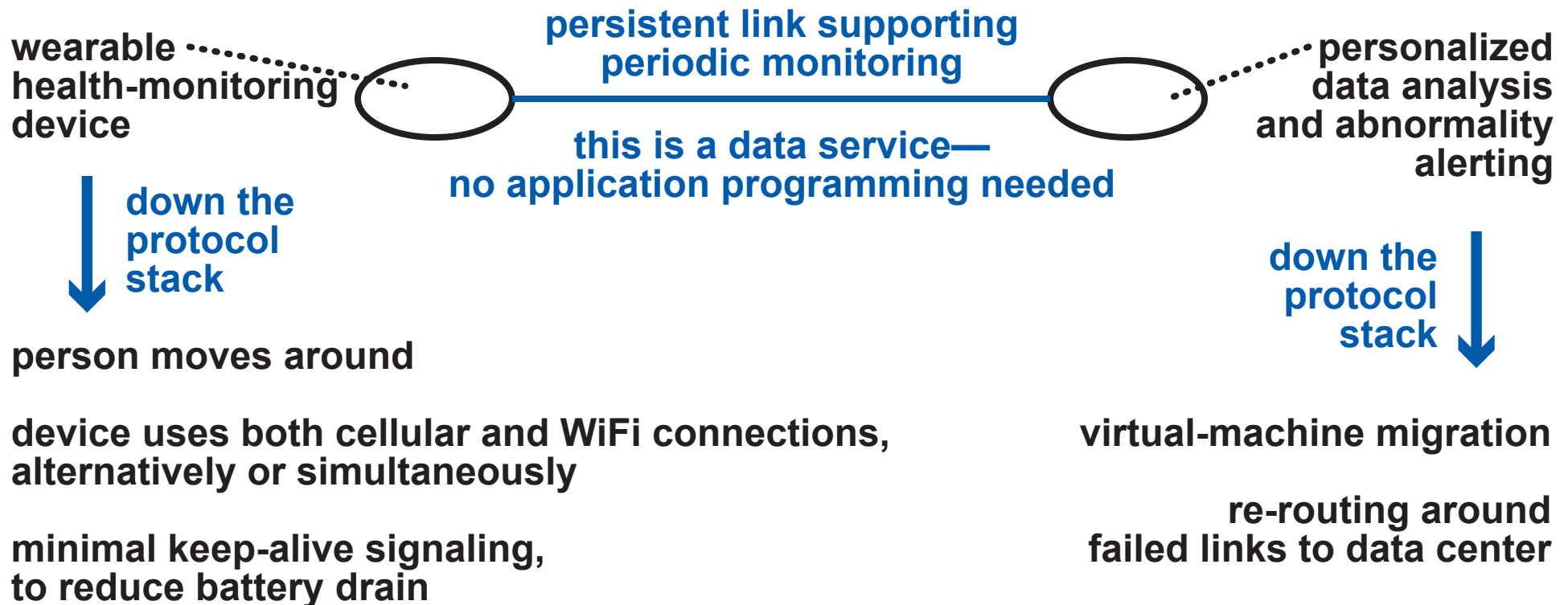
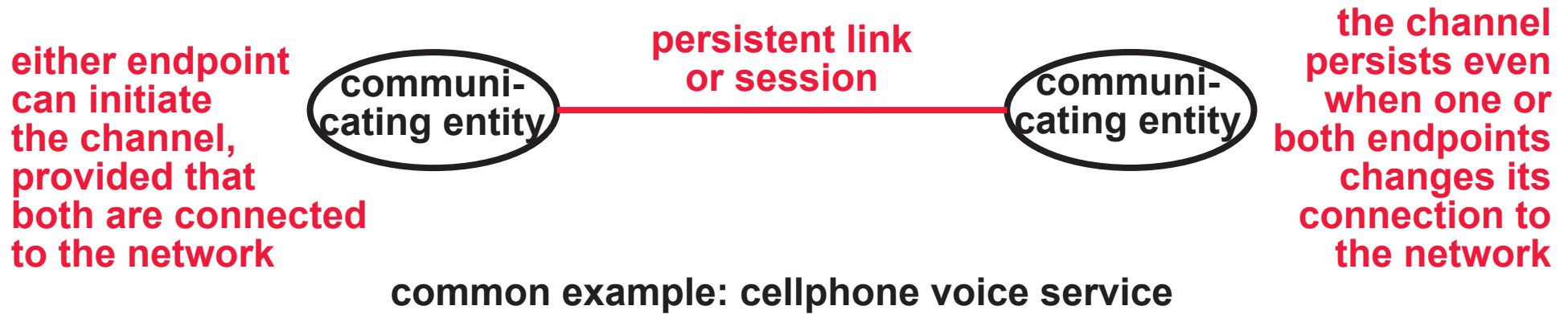
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# A DEFINITION OF MOBILITY



# WHY DIRECTORIES?

## MAP REAL-WORLD ENTITIES TO ROUTABLE NAMES

*e.g., telephone book, Domain Name System*

## MAP NAMES IN A NETWORK TO NAMES IN AN UNDERLAY

lookup at session setup, sometimes after a mobility event

## DIRECTORIES CAN BE CENTRALIZED OR DISTRIBUTED

IF DISTRIBUTED, THERE MUST BE SOME STRUCTURE FOR TARGETING LOOKUP QUERIES AND BALANCING THE LOAD

- hierarchical namespace

*Domain Name System, Mobile IP home agents*

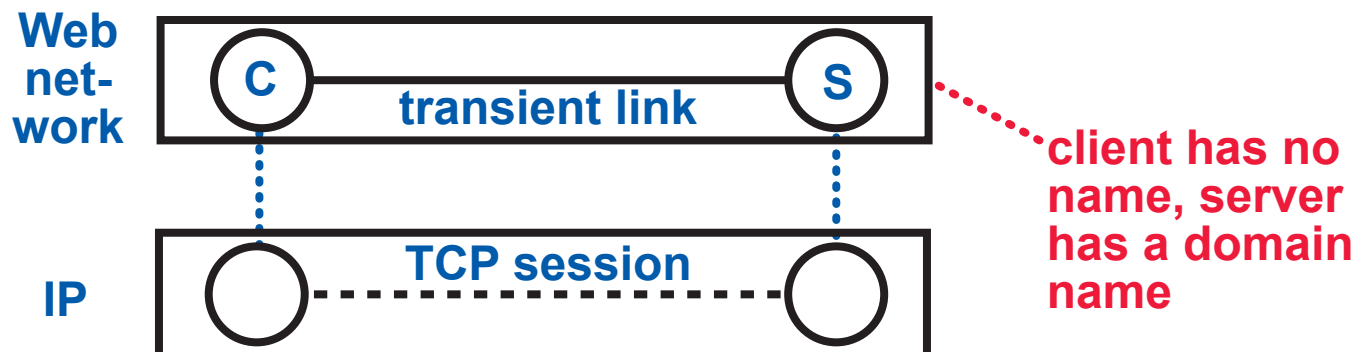
- numerical order on hash of name

*distributed hash table*

- simply limiting the size of the network

*in ARP flooding, each member is its own directory entry*

## IS THIS A USEFUL PERSPECTIVE?





**WHY IS MOBILITY DIFFICULT TO IMPLEMENT**

**ON THE INTERNET?**

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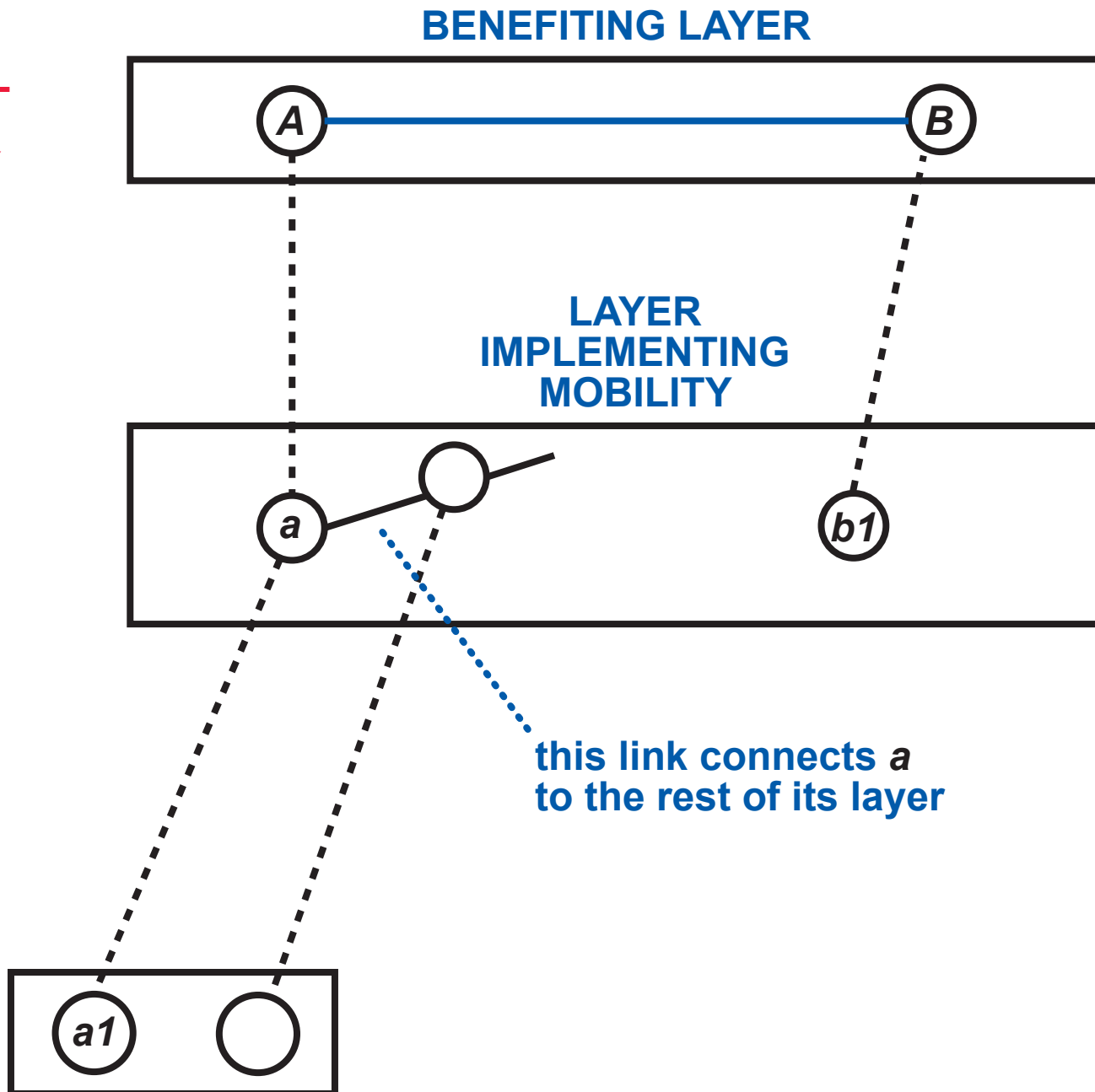
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# A PATTERN FOR IMPLEMENTING MOBILITY

DYNAMIC-  
ROUTING  
MOBILITY



# A PATTERN FOR IMPLEMENTING MOBILITY

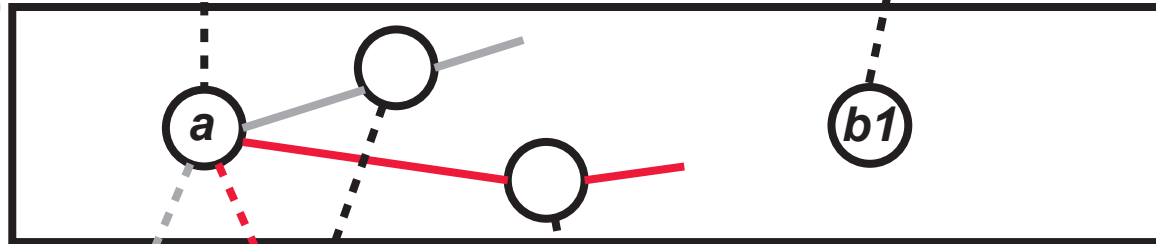
**DYNAMIC-ROUTING MOBILITY**

**BENEFITING LAYER**



as the attachment of a member changes, its links change, and the routing algorithm must find new routes to it

**LAYER IMPLEMENTING MOBILITY**

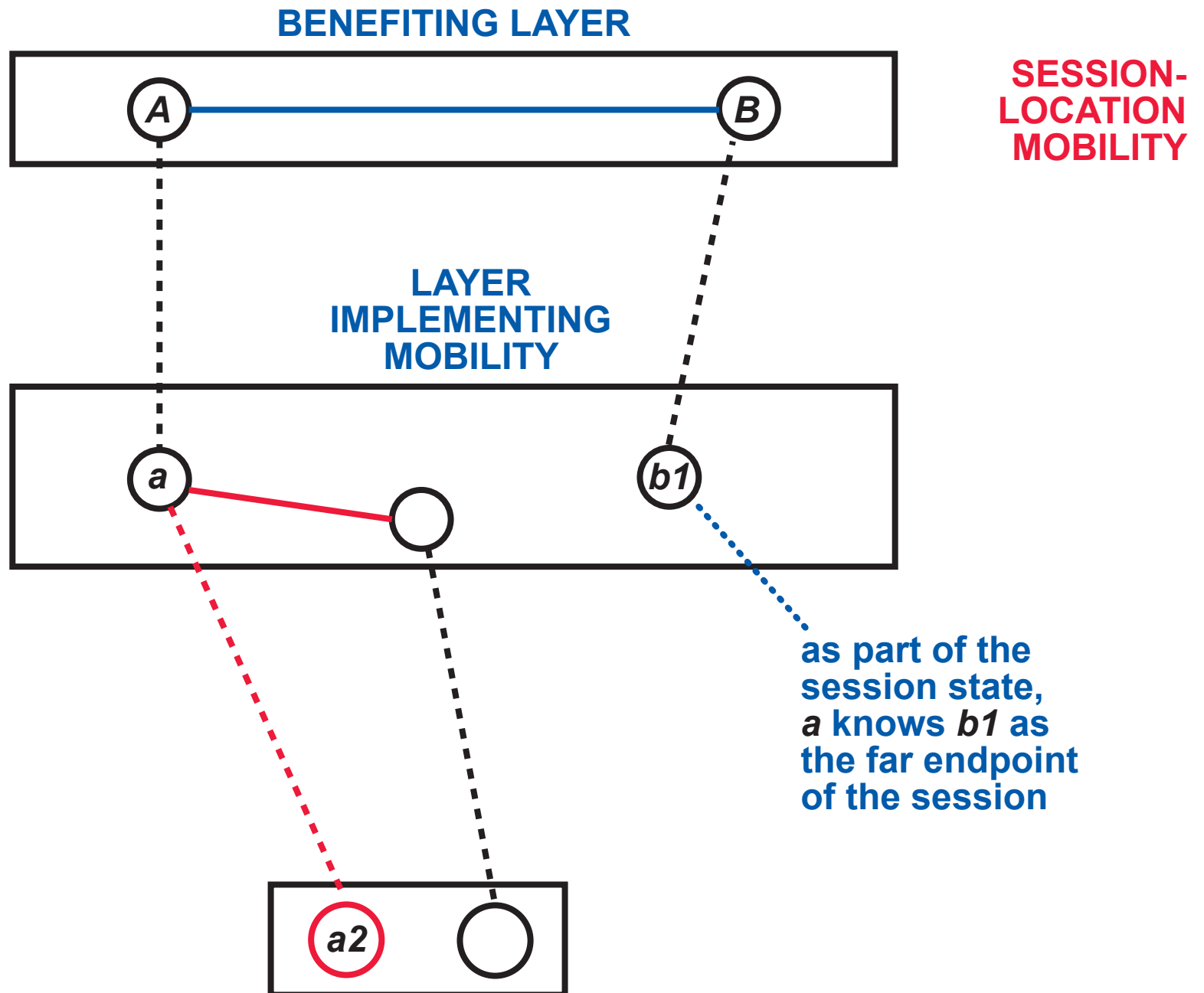


layer state components that change:

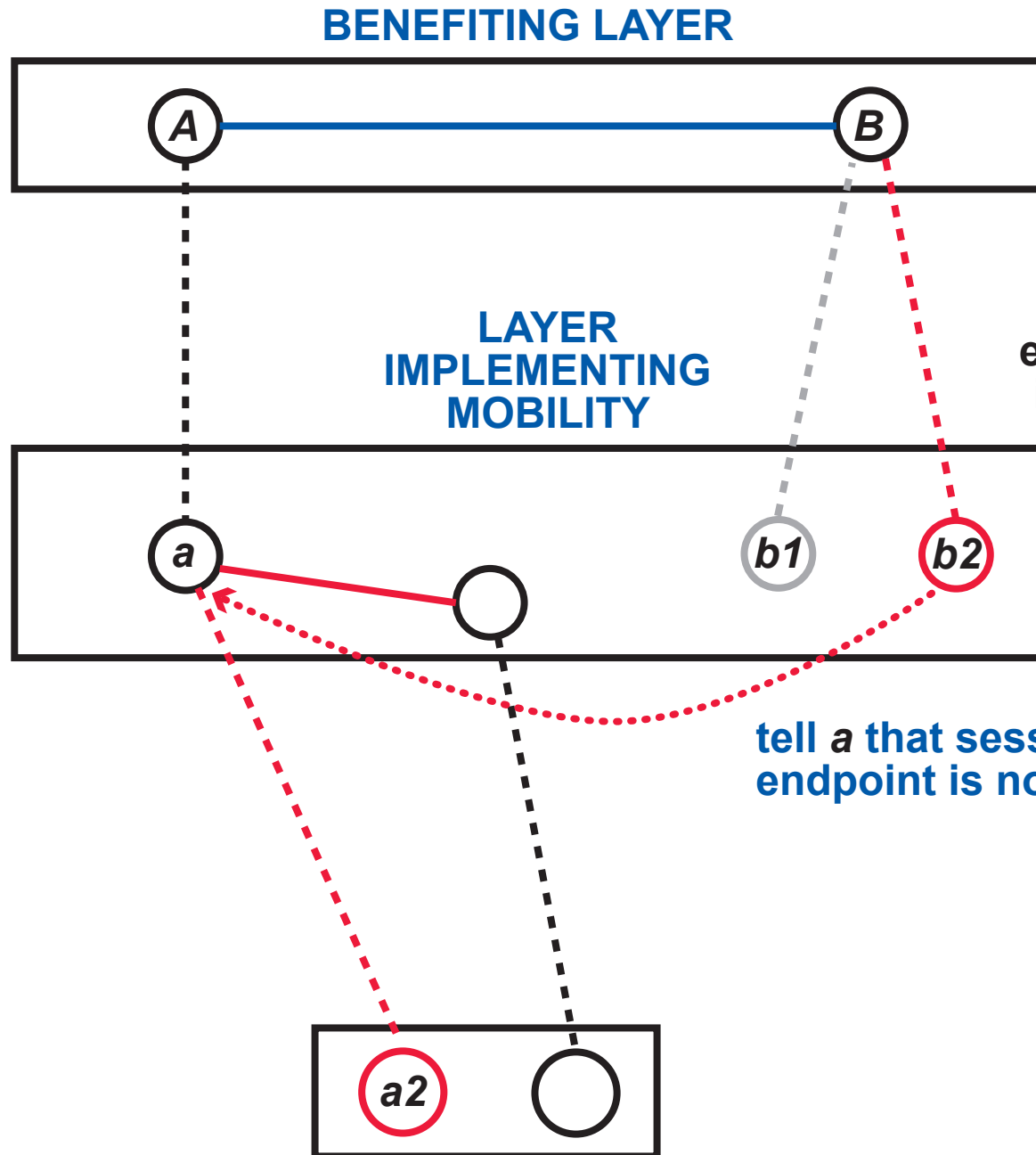
*locations*  
*links*  
*routes*



# ANOTHER PATTERN FOR IMPLEMENTING MOBILITY



# ANOTHER PATTERN FOR IMPLEMENTING MOBILITY



**SESSION-  
LOCATION  
MOBILITY**

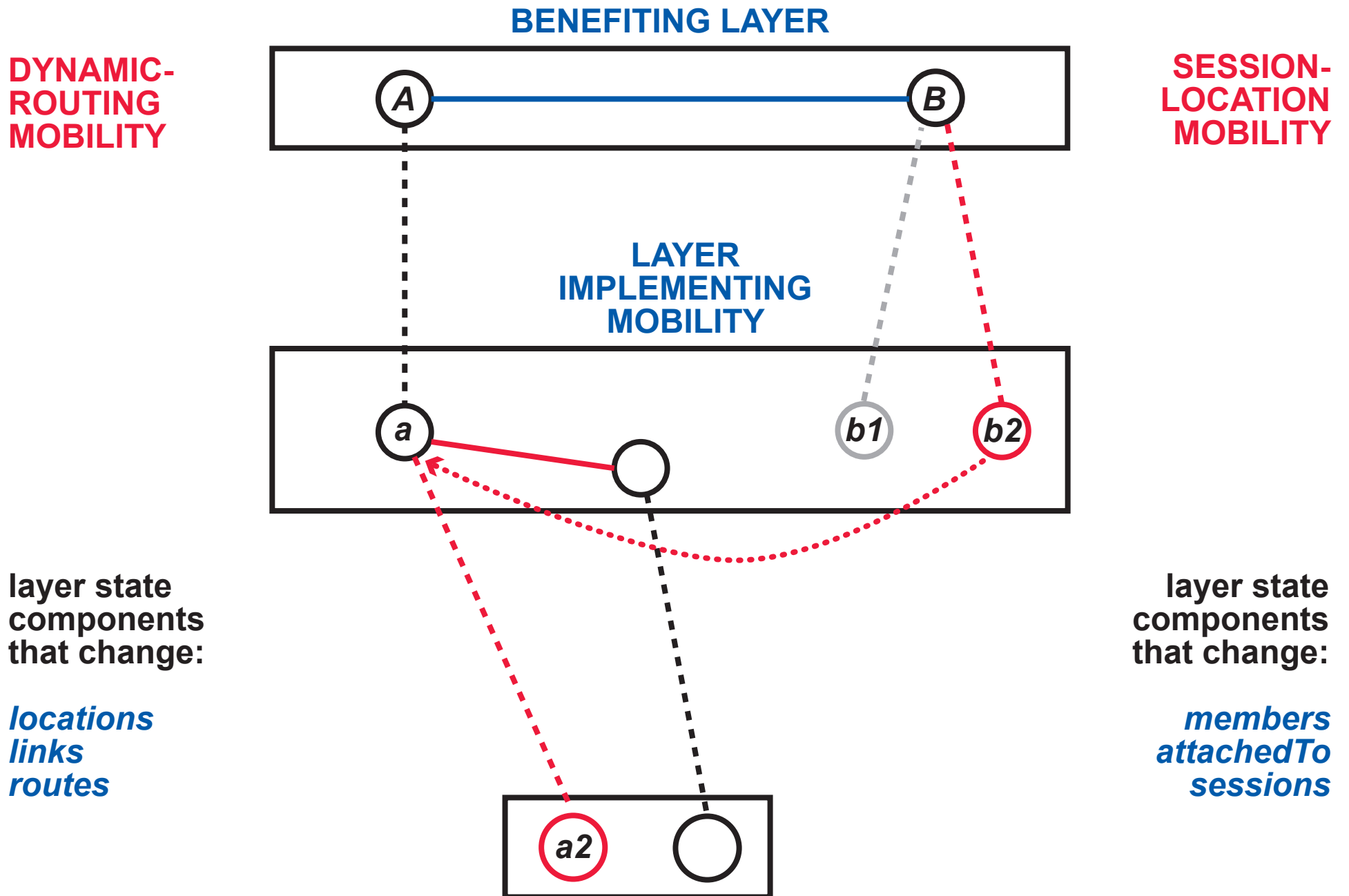
as the link  
endpoint changes  
its location in the  
implementing  
layer, the  
session state  
changes to  
match it

**tell *a* that session  
endpoint is now *b2***

layer state  
components  
that change:

*members  
attachedTo  
sessions*

# BOTH PATTERNS FOR IMPLEMENTING MOBILITY



# STRENGTHS AND WEAKNESSES OF THE PATTERNS

## DYNAMIC-ROUTING MOBILITY

### Strengths

Works well in a network with a smaller scope and a flat name space—usually dynamic routing for mobility is no different from “normal” routing.

*e.g., Ethernet*

Implemented by routers, which are more trusted.

### Weaknesses

In a larger network with a hierarchical name space, costs for dynamic routing to individual members are high.

*How many routers know where to find a mobile member?*

trade-off {  
if many, storage and update costs are high  
if few, path costs are high

## SESSION-LOCATION MOBILITY

### Strengths

Low storage and update costs.

No path costs.

### Weaknesses

Implemented by endpoints, which are less trusted, and harder to update with network software.

Packet losses during handover may be disruptive.



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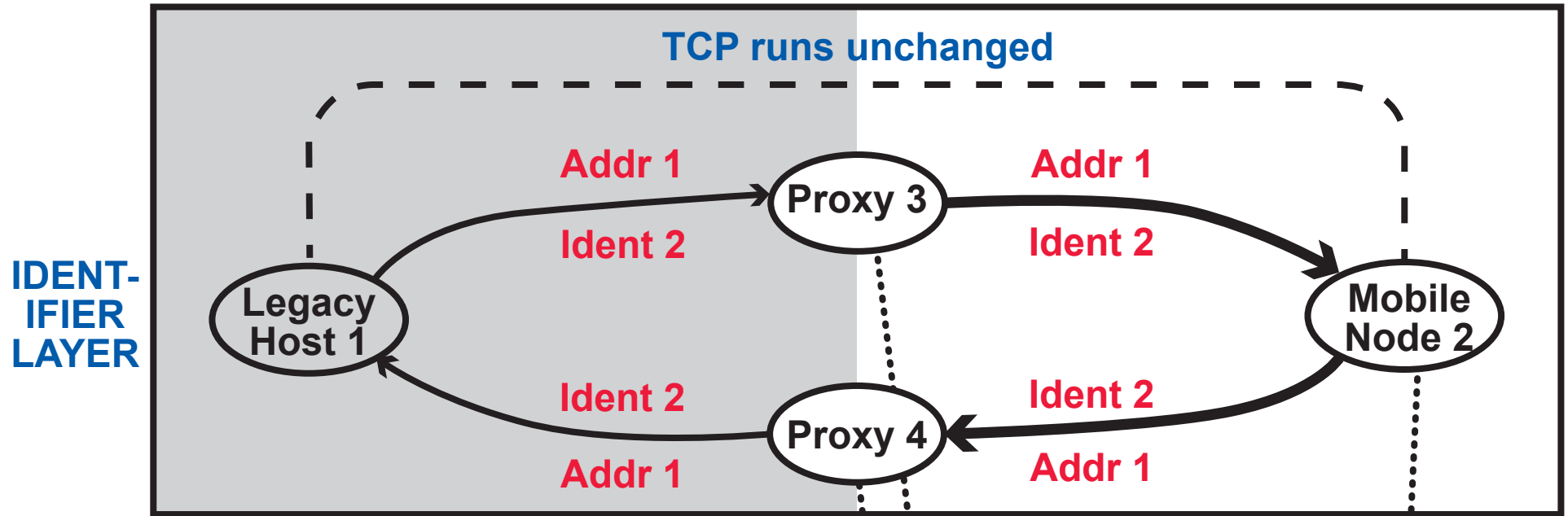
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# INTEROPERATION BRINGS THE TWO PATTERNS CLOSER

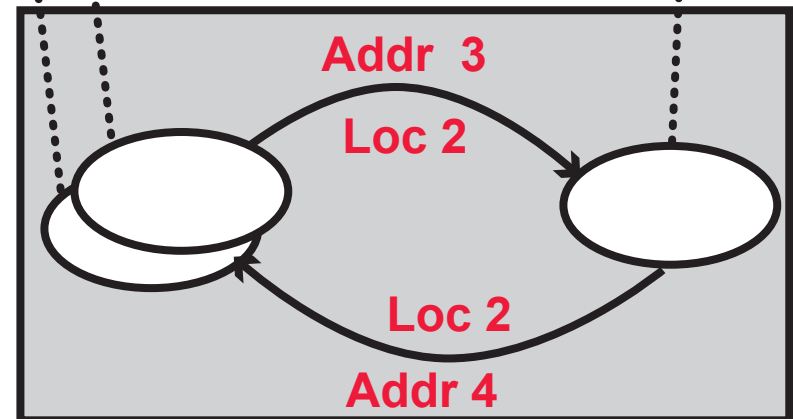


gray areas are normal IP

this picture describes either Mobile IP (which interoperates with IP by default) . . .

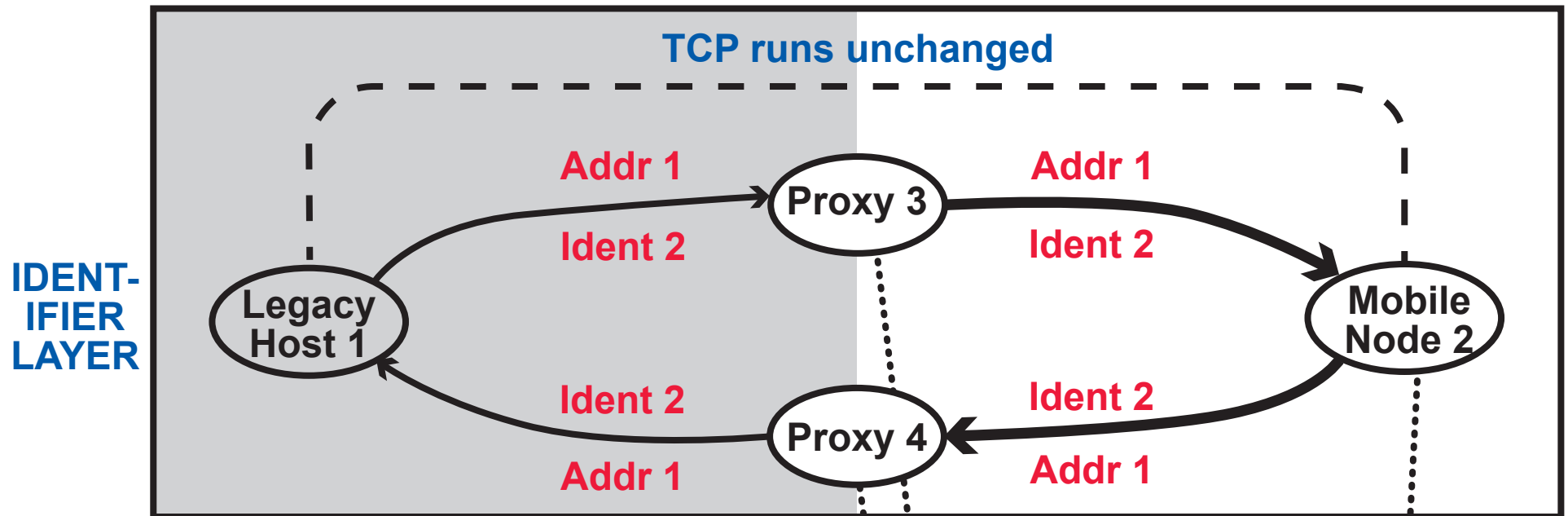
. . . or LISP Mobile Node interoperating with IP

LOCA-TOR LAYER



## Mobile IPv6 without route optimization:

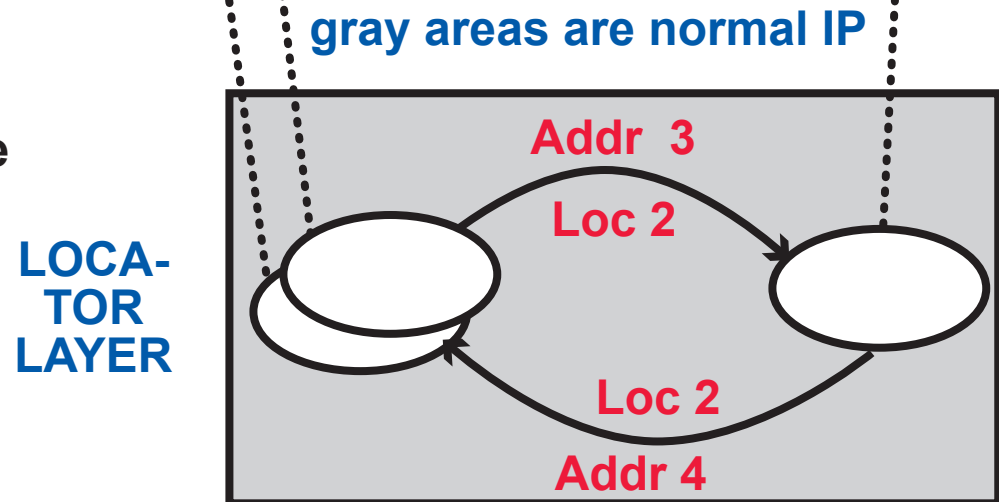
- Proxies 3 and 4 are both the Home Agent of Ident 2, in a fixed location
- HA routes to Ident 2 either locally or on a “special link”



## LISP Mobile Node:

- Proxies 3 and 4 can be anywhere
- routing Ident 2 to a proxy requires that LISP identifiers be aggregated into a prefix

**IN BOTH CASES, PROXIES HAVE LITTLE STATE**



# PURPOSES FOR LAYERING

TO BUILD A NETWORK WITH A LARGER SPAN OUT OF SMALLER, HETEROGENEOUS NETWORKS

IP over LANs and WANs

TO BUILD A NETWORK WITH BETTER PERFORMANCE OR RELIABILITY ON TOP OF AN EXISTING NETWORK

Resilient Overlay Networks over IP

TO SHARE THE RESOURCES OF AN EXISTING NETWORK IN A DISCIPLINED WAY

Virtual LANs over LANs

application networks over IP

TO BUILD A NETWORK WITH LINKS THAT OFFER A SUPERIOR COMMUNICATION SERVICE

Virtual Private Networks over IP

TO MAINTAIN A PERSISTENT, ROUTABLE IDENTIFIER FOR A MOBILE ENDPOINT

Mobile IP, LISP Mobile Node, and many other examples over IP

*how do these compose?*

A diagram illustrating the purposes for layering in network design. It features five distinct purposes, each with a corresponding example. Dotted blue lines connect the examples to a central question: "how do these compose?". The purposes and examples are: 1. Purpose: "TO BUILD A NETWORK WITH A LARGER SPAN OUT OF SMALLER, HETEROGENEOUS NETWORKS"; Example: "IP over LANs and WANs". 2. Purpose: "TO BUILD A NETWORK WITH BETTER PERFORMANCE OR RELIABILITY ON TOP OF AN EXISTING NETWORK"; Example: "Resilient Overlay Networks over IP". 3. Purpose: "TO SHARE THE RESOURCES OF AN EXISTING NETWORK IN A DISCIPLINED WAY"; Example: "Virtual LANs over LANs" and "application networks over IP". 4. Purpose: "TO BUILD A NETWORK WITH LINKS THAT OFFER A SUPERIOR COMMUNICATION SERVICE"; Example: "Virtual Private Networks over IP". 5. Purpose: "TO MAINTAIN A PERSISTENT, ROUTABLE IDENTIFIER FOR A MOBILE ENDPOINT"; Example: "Mobile IP, LISP Mobile Node, and many other examples over IP".

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# COMPOSITIONAL NETWORK MOBILITY

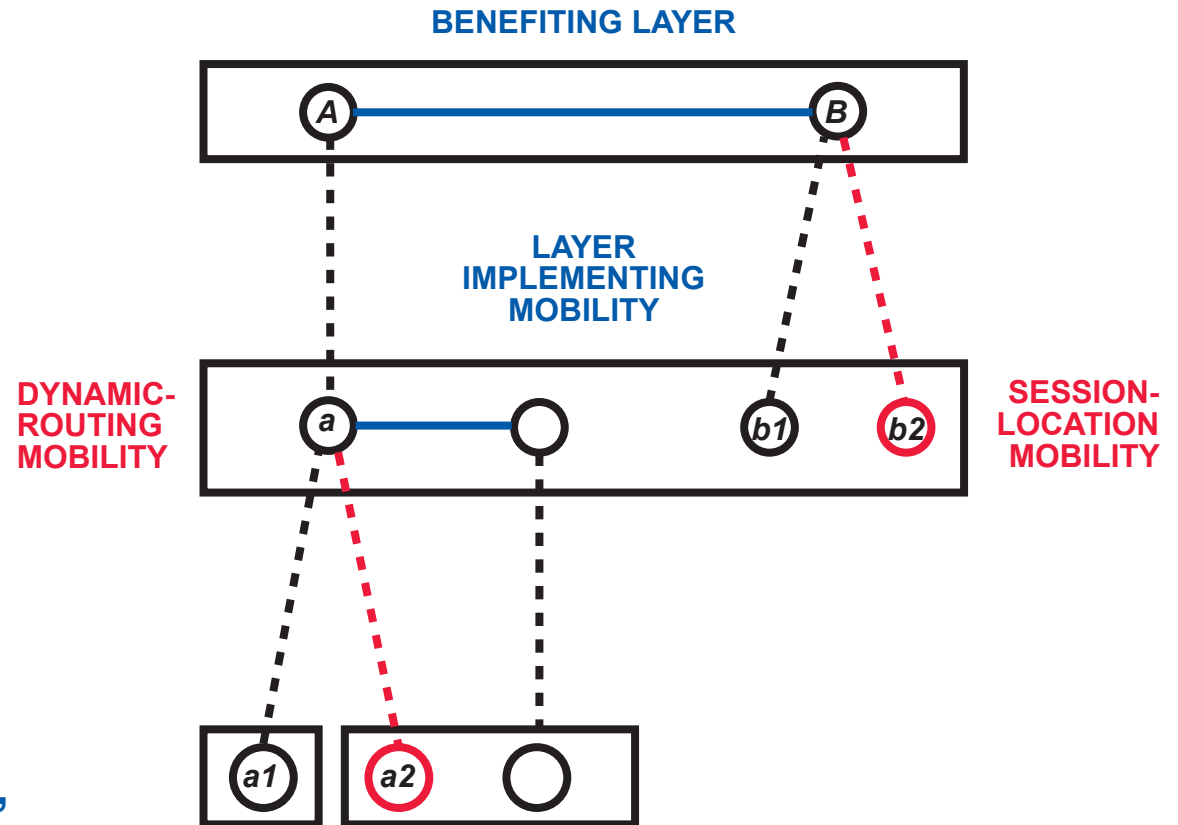
- every mobility mechanism specializes one of these patterns, or is a composition of the two

with enough design freedom, instances of mobility can be moved up and down the levels

- in principle, each instance of mobility could be handled with either of these patterns at any level below the benefiting layer —so mobility mechanisms could be everywhere

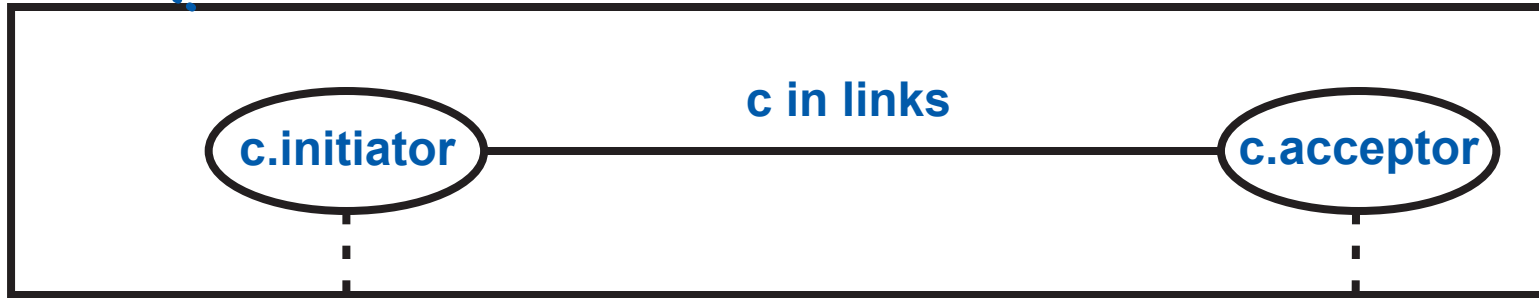
there is a large design space, much of it unexplored

- an interesting question: how do implementations of both patterns in the same layer compose?



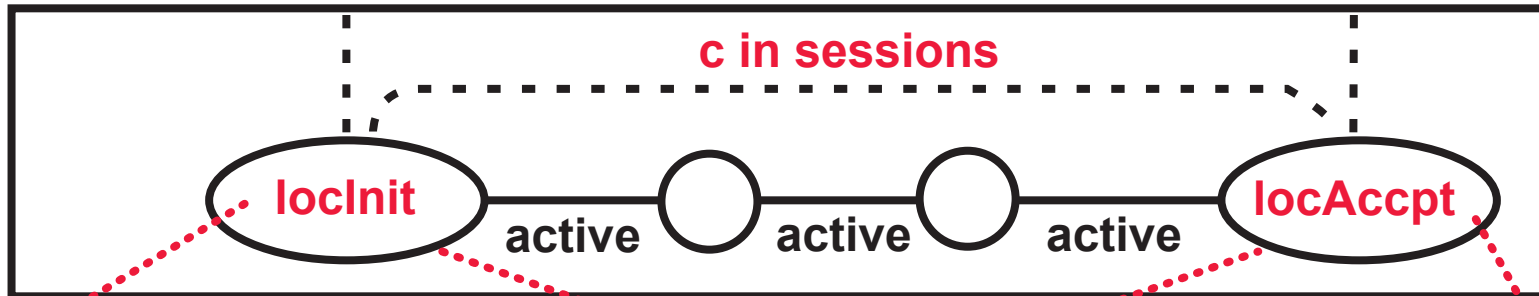
# AN ACTIVE IMPLEMENTED CHANNEL

c.userLayer



..... in attachments .....

c.implLayer



..... in locations .....

c.initFarLoc = locAccept

reachable

c.acceptFarLoc = locInit

# AN INACTIVE IMPLEMENTED CHANNEL

MODEL IMPLEMENTS BOTH PATTERNS IN EVERY LAYER

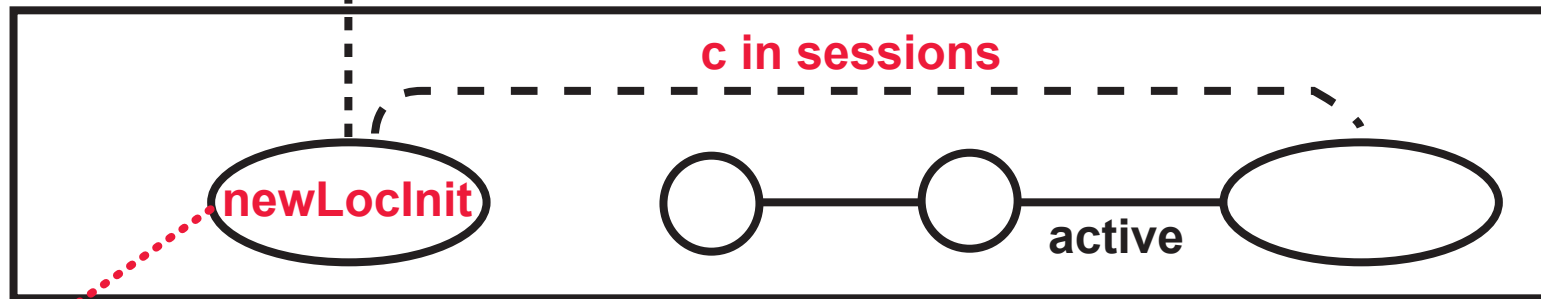
c.userLayer



..... in attachments

MOBILITY COULD DESTROY REGISTRATIONS

c.implLayer



..... in locations

MOBILITY COULD CAUSE FAR LOCATIONS IN SESSION STATE TO BE WRONG

MOBILITY COULD DESTROY OR INACTIVATE LINKS

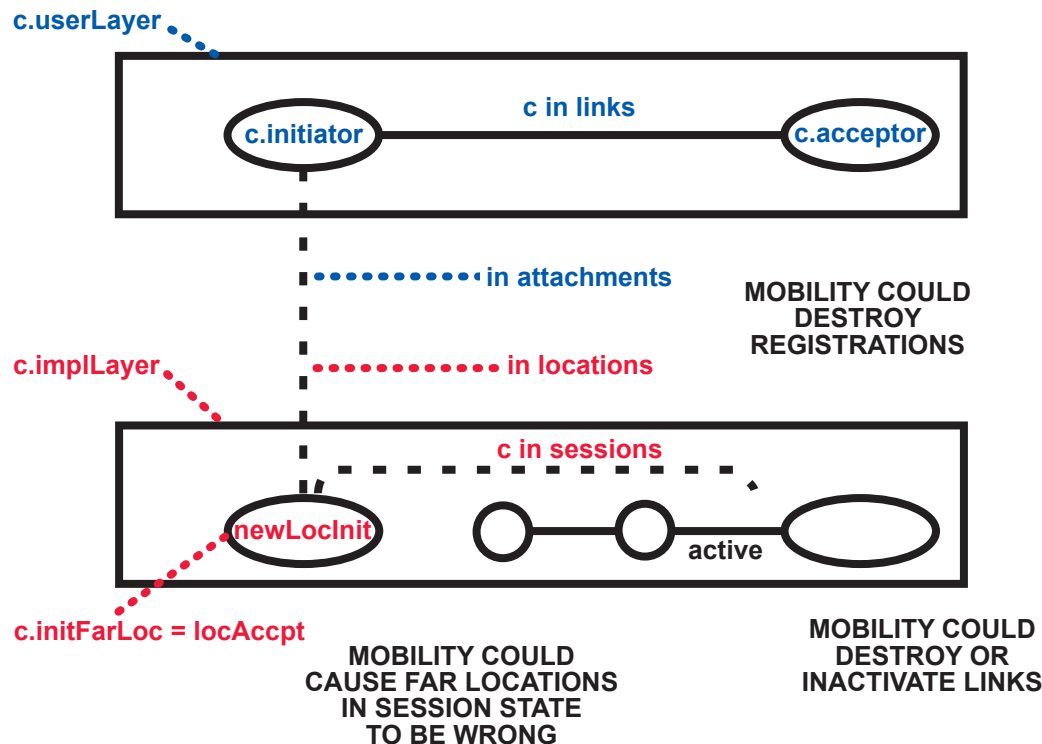
c.initFarLoc = locAccpt



# PROOF THAT MOBILITY MECHANISMS IN A LAYER COMPOSE WITHOUT INTERFERENCE

We cannot assume that mobile devices and network elements will perform all the requisite actions (to prove a true progress property).

We do assume that a mobile device can always become a member of a layer of its choice.



**Theorem:**

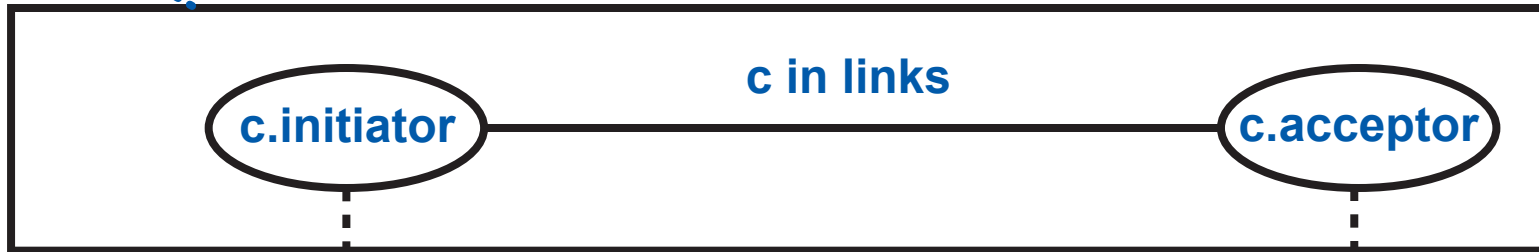
In any state in which an implemented link is inactive, some event is enabled whose execution will make progress toward making the link active (a safety property).

**Proof at one level:**

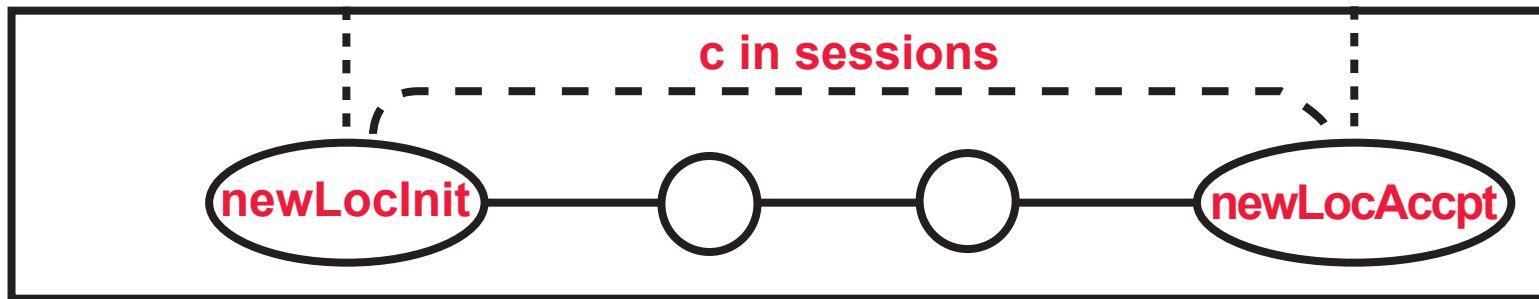
Manual enumeration of possible event sequences, automated checking of their preconditions with the Alloy Analyzer (verification over small domains).

# WHAT COULD GO WRONG?

c.userLayer



c.implLayer



c.initFarLoc = locAcpt

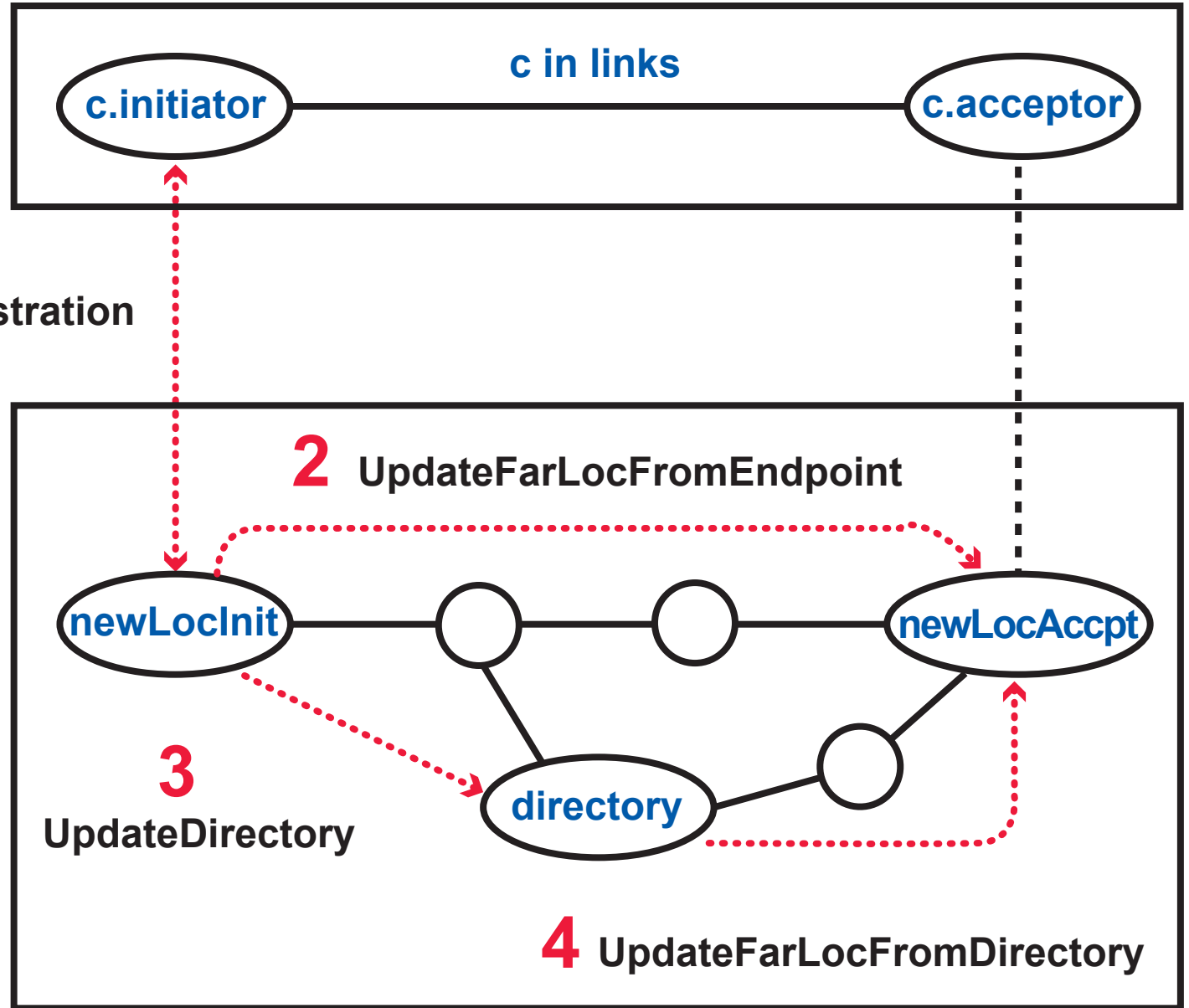
both endpoints  
have moved

c.acptFarLoc = locInit

both endpoints have  
the wrong far location

neither can send an update  
message to the other

# SOME EVENT SEQUENCES



in the double-handoff scenario,  
1, 2, 3 and 1, 3, 2  
do not work, but  
1, 3, 4 does

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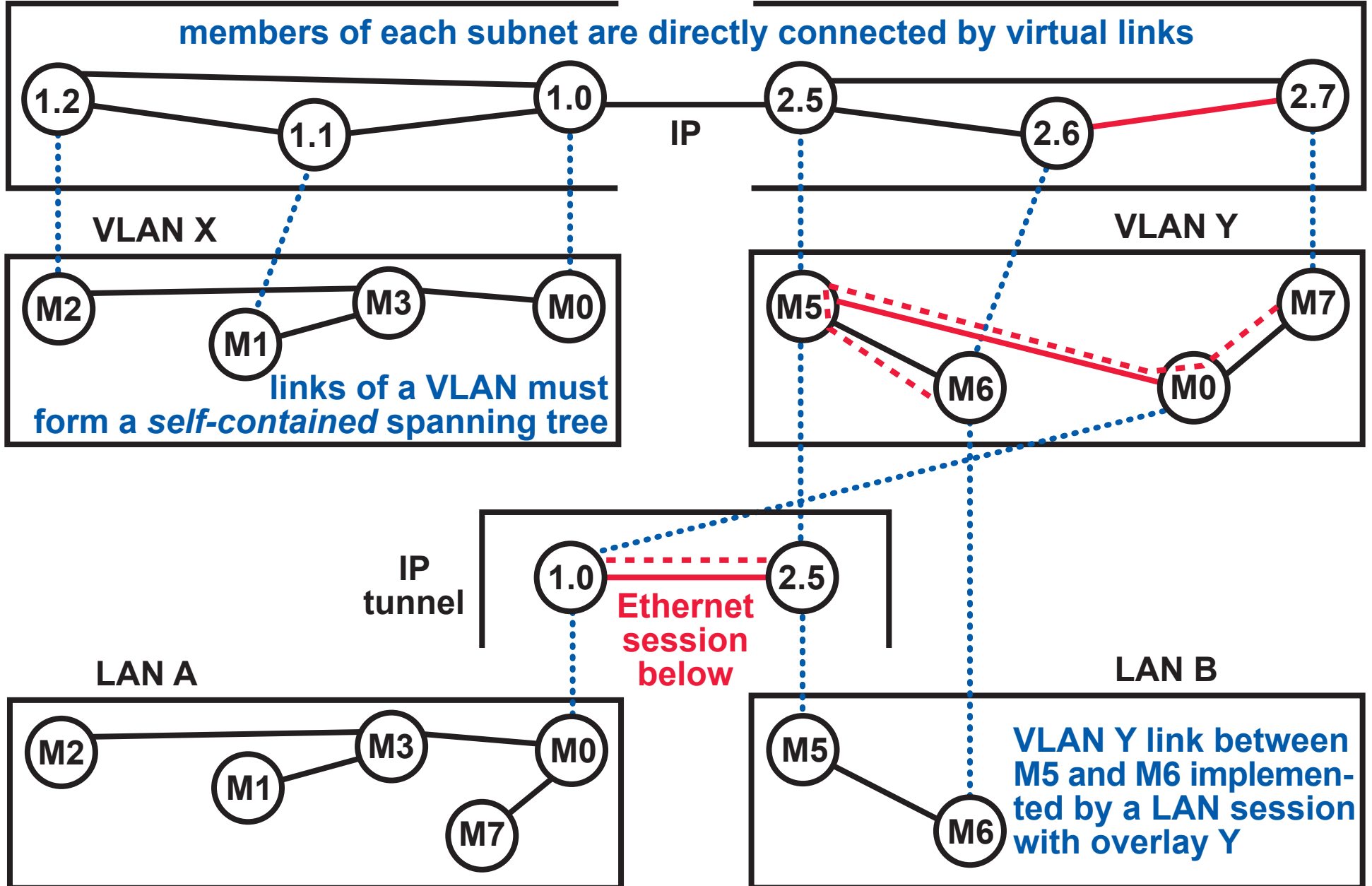
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# VLAN TECHNOLOGY

IP SUBNETWORK, PREFIX 192.168.1/24

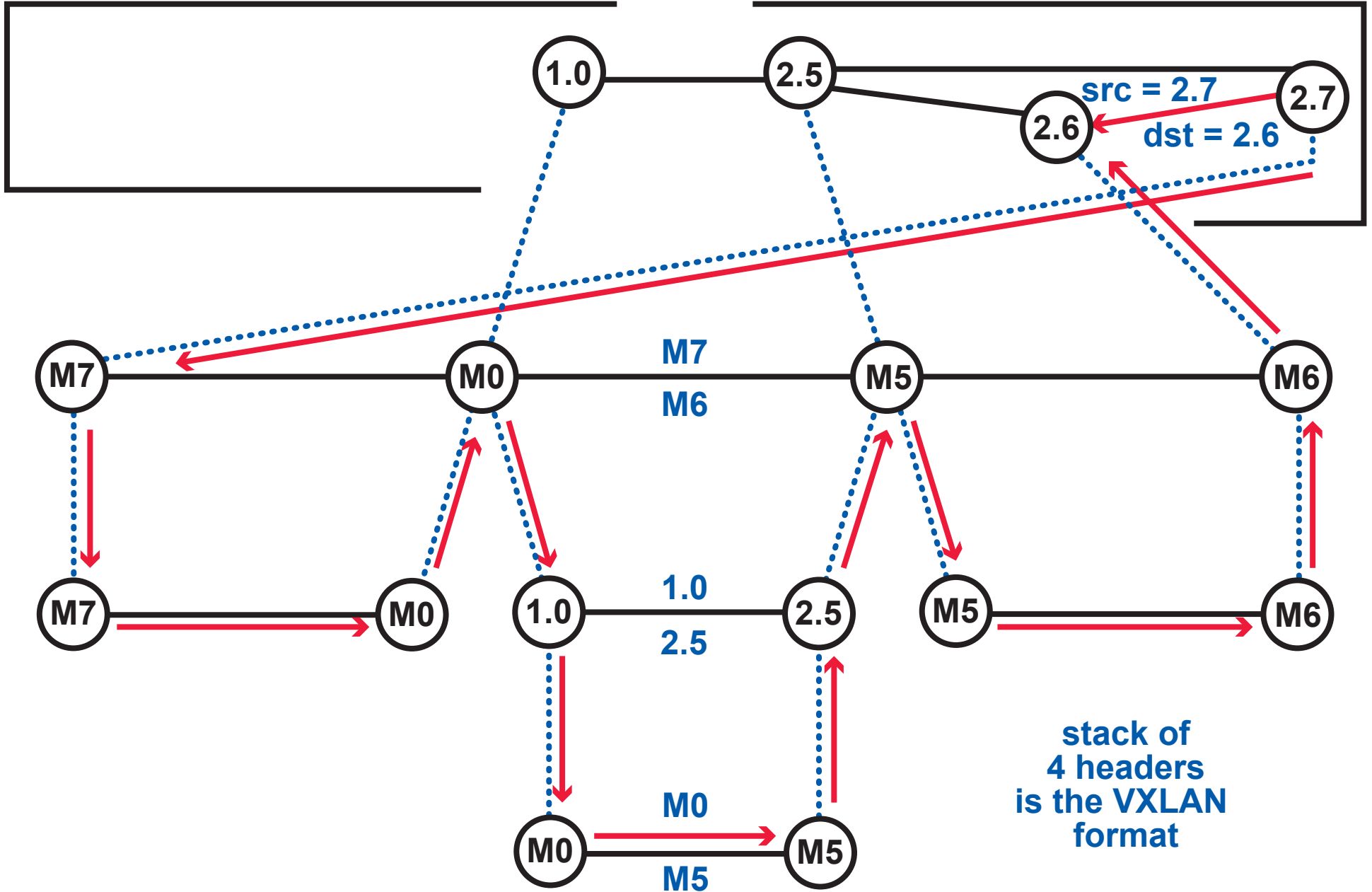
IP SUBNETWORK, PREFIX 192.168.2/24



# VXLAN TECHNOLOGY

IP SUBNETWORK, PREFIX 192.168.1/24

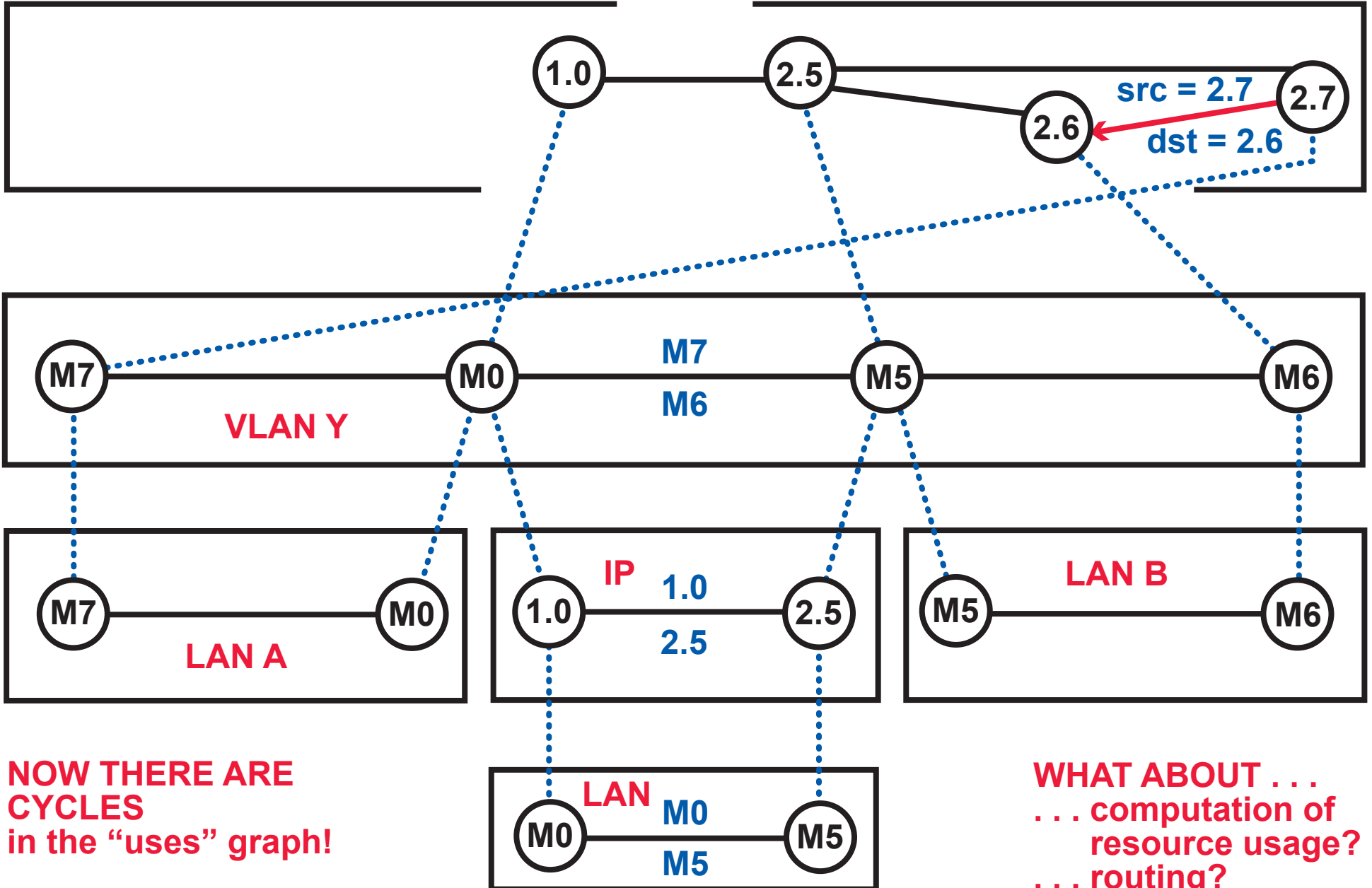
IP SUBNETWORK, PREFIX 192.168.2/24



# UNDERSTANDING THE COMPOSITION

IP SUBNETWORK, PREFIX 192.168.1/24

IP SUBNETWORK, PREFIX 192.168.2/24



NOW THERE ARE CYCLES in the "uses" graph!

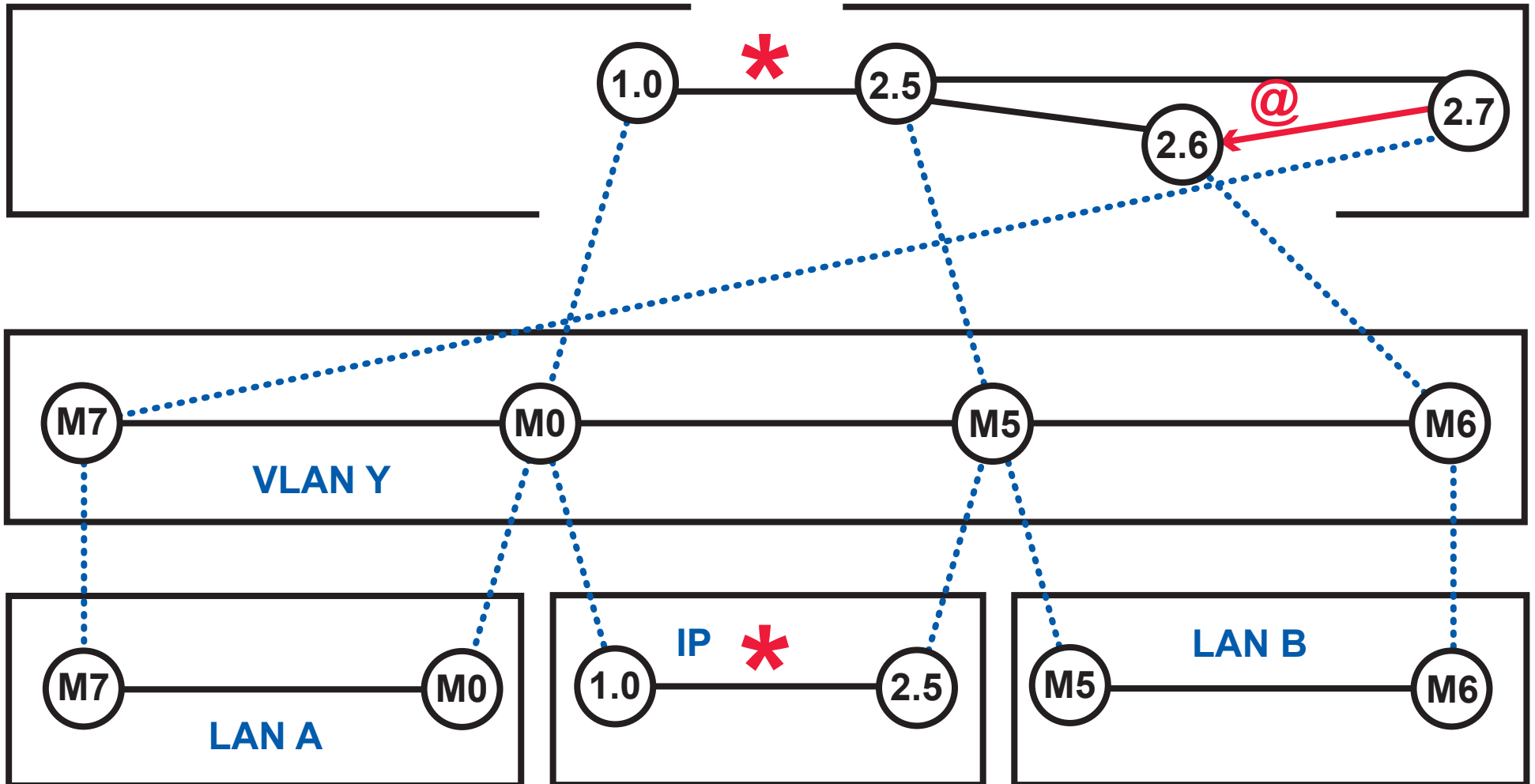
WHAT ABOUT ...  
... computation of resource usage?  
... routing?

# COMPUTATION OF RESOURCE USAGE

There can be usage cycles among networks,  
but not among links!

IP SUBNETWORK, PREFIX 192.168.1/24

IP SUBNETWORK, PREFIX 192.168.2/24



@ USES \*

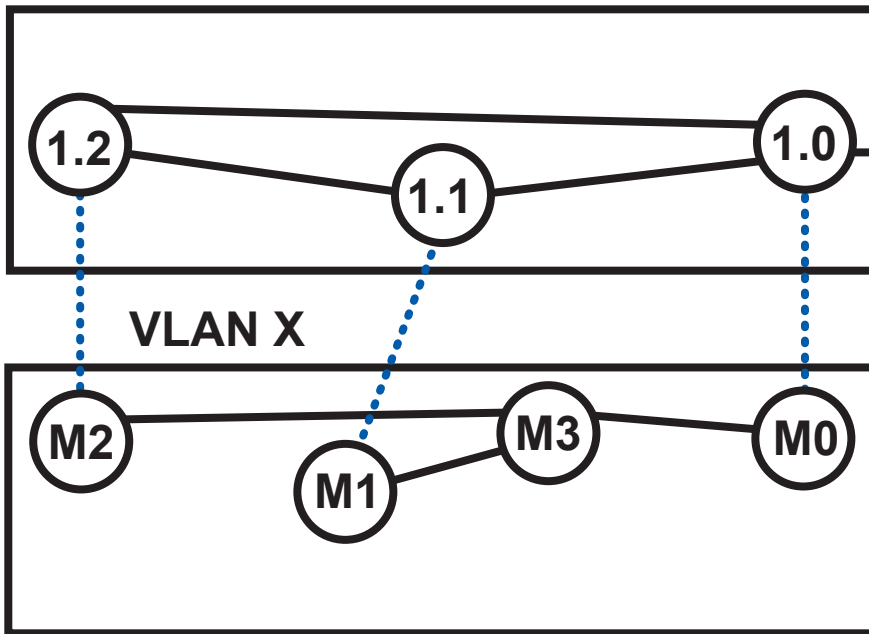
\* does not use anything but an isolated Ethernet link



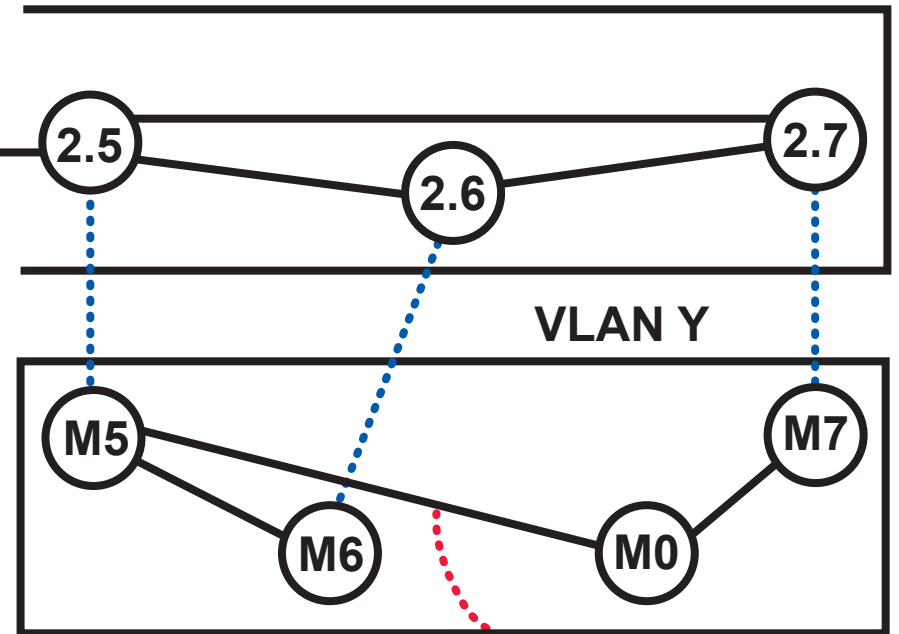
# ROUTING

WITH THE VXLAN ARCHITECTURE,  
ROUTING IN EVERY NETWORK IS NORMAL

IP SUBNETWORK, PREFIX 192.168.1/24



IP SUBNETWORK, PREFIX 192.168.2/24



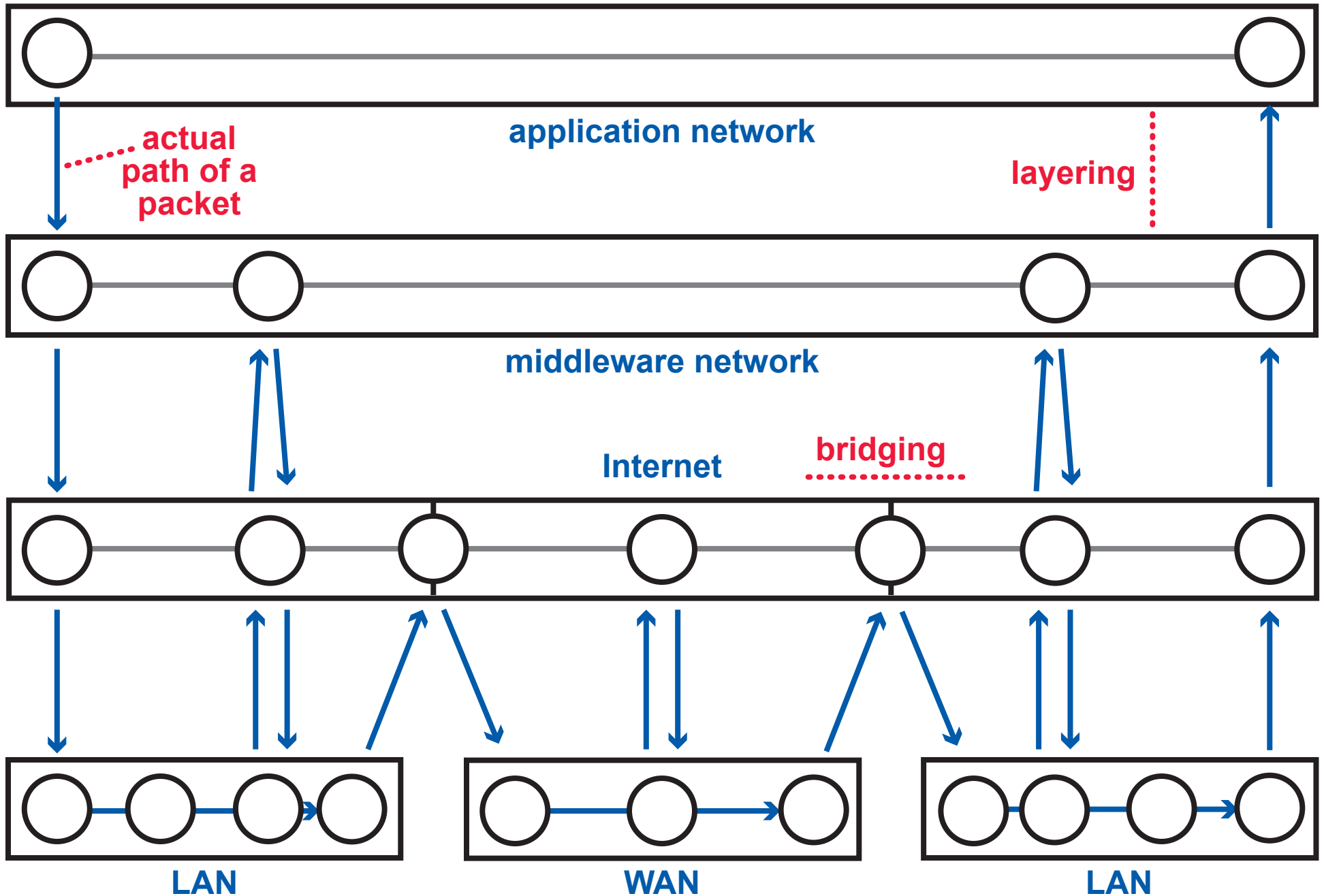
IP

note that switches  
and links can  
belong to multiple  
VLANs, although  
hosts cannot

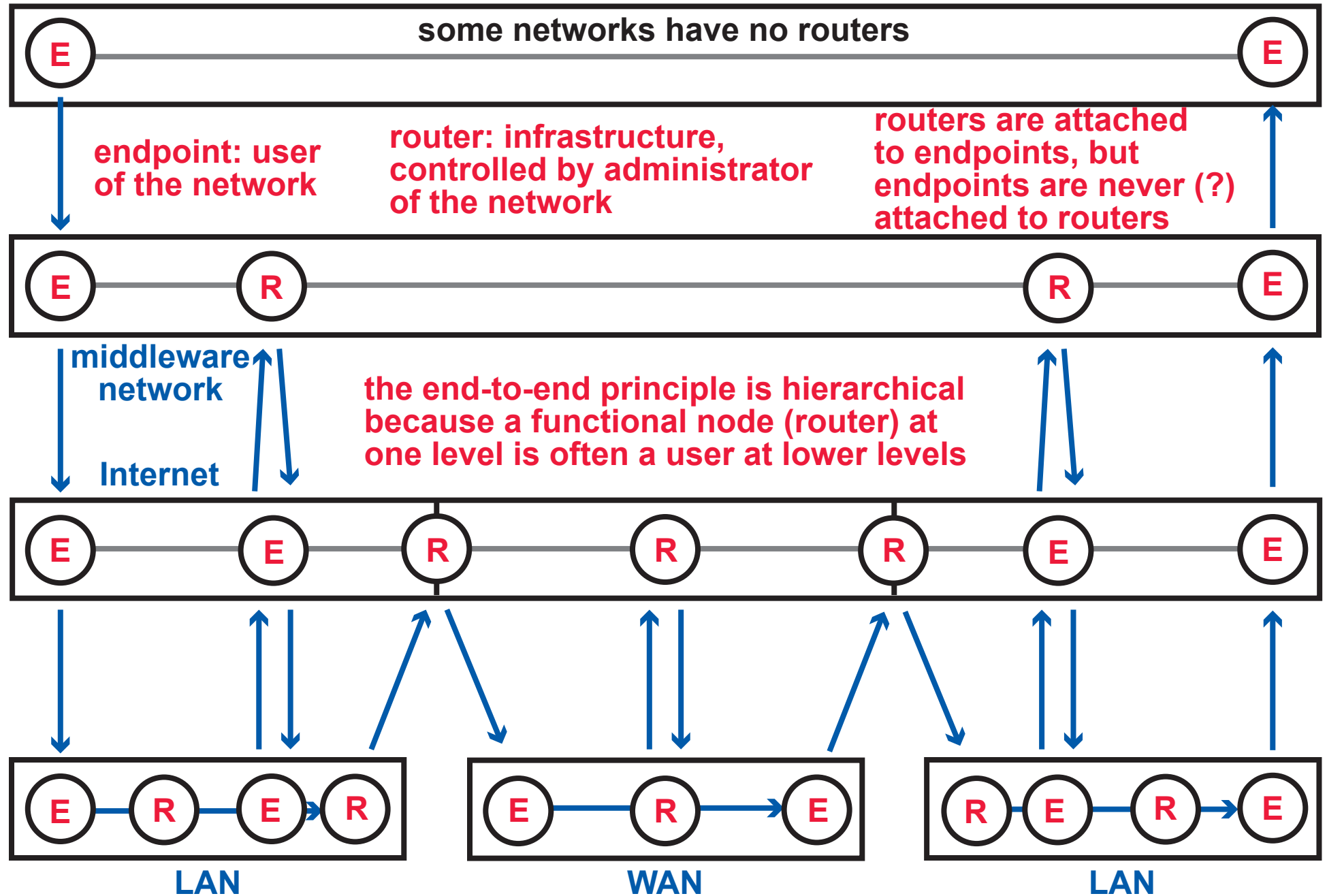
this causes no confusion  
because VLANs are isolated,  
both for security and to  
limit broadcast domains

it does not matter to  
VLAN routing that  
M0 and M5 are  
physically connected  
through the IP backbone,  
which is used here.....

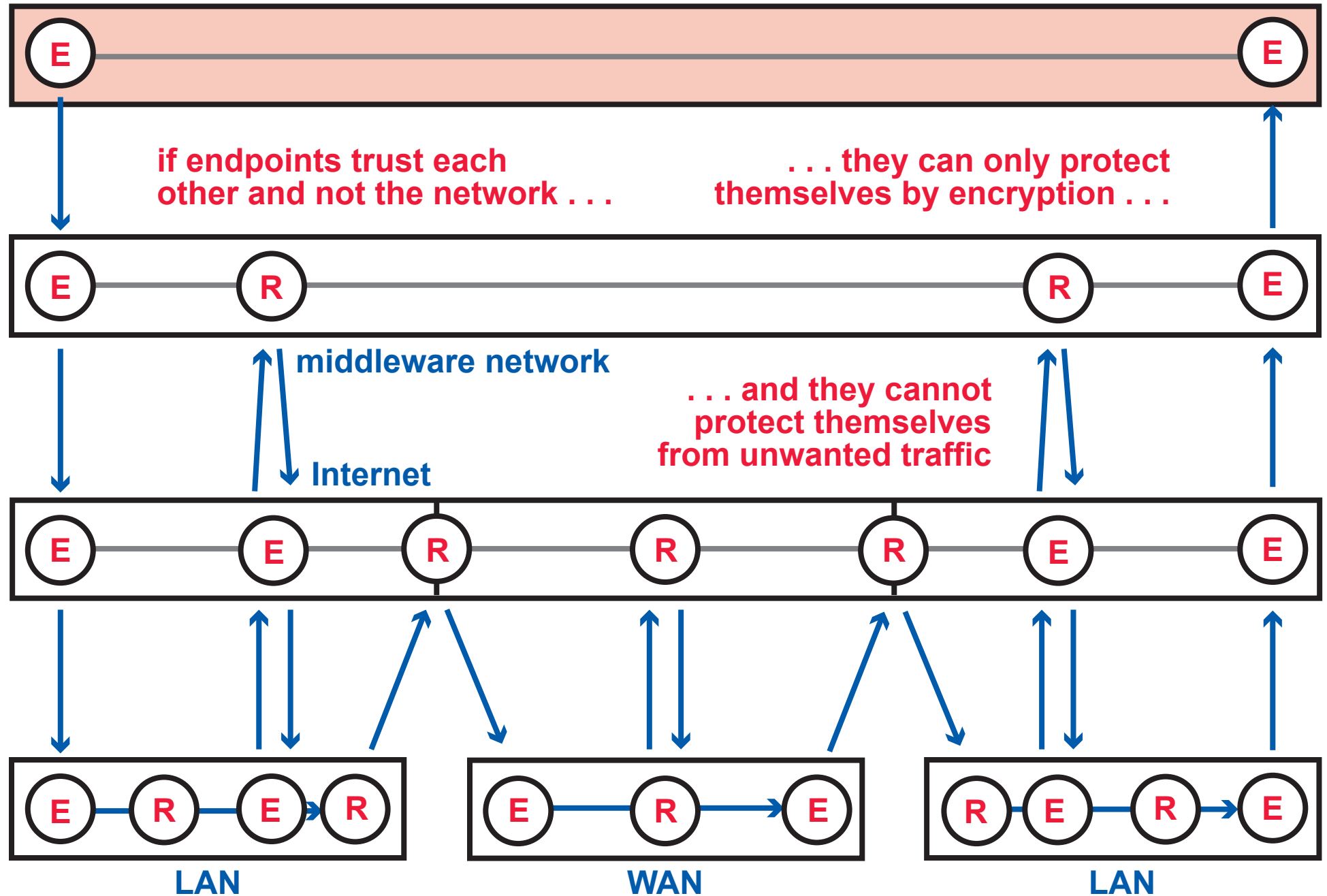
# ARCHITECTURE AND TRUST



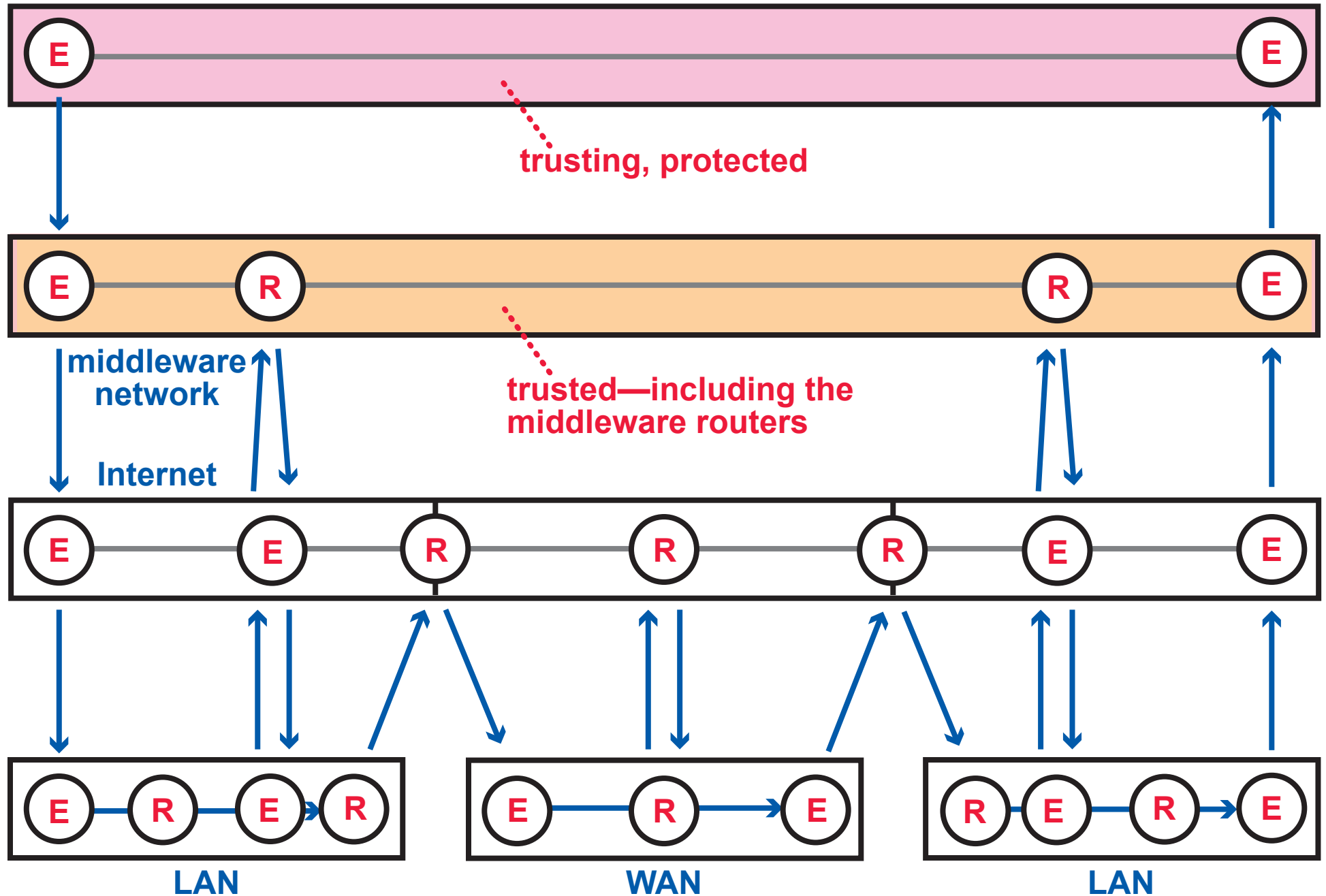
# ARCHITECTURE AND TRUST: ENDPOINTS AND ROUTERS



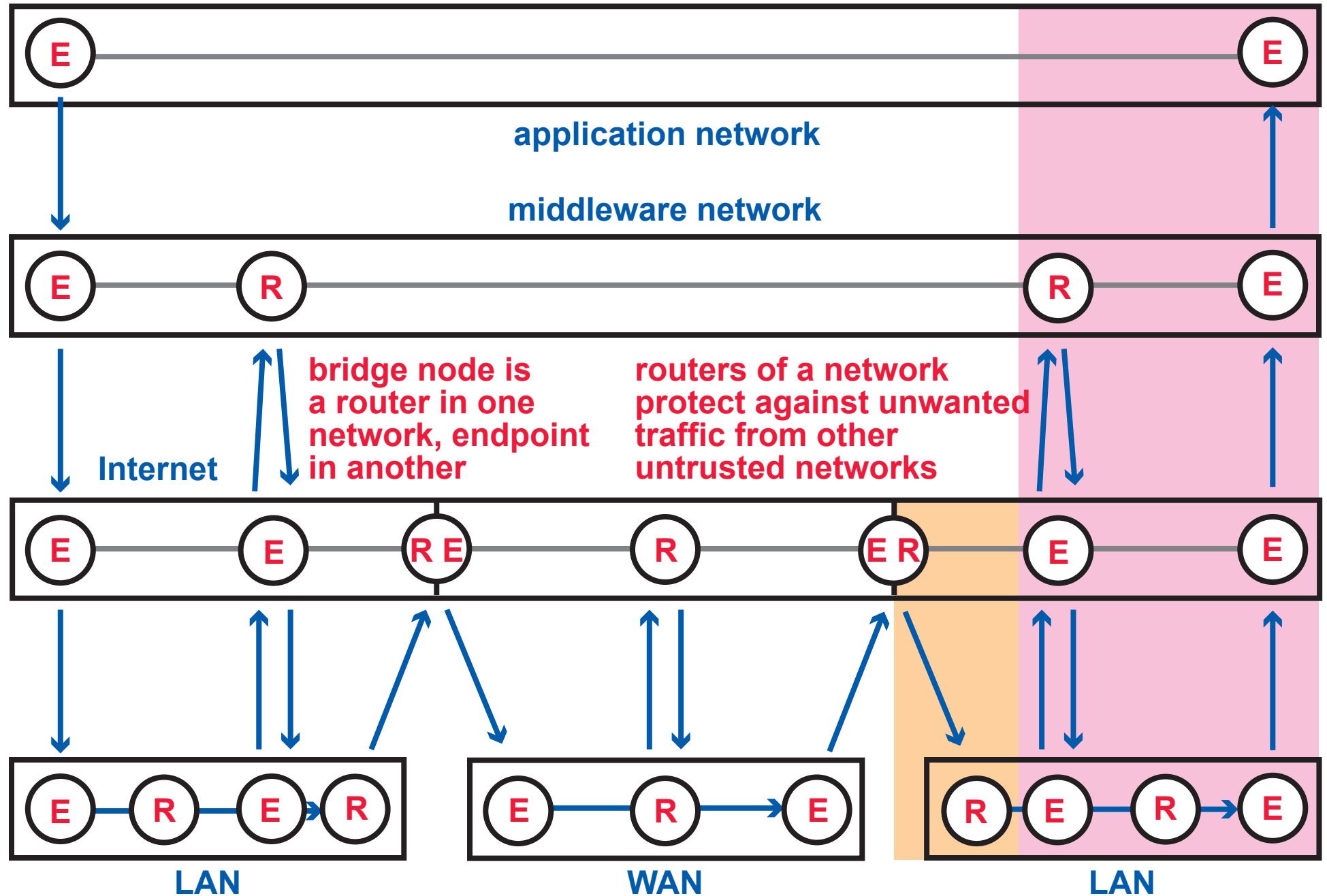
# ARCHITECTURE AND TRUST: ENDPOINT TRUST



# ARCHITECTURE AND TRUST: ENDPOINT TRUST



# ARCHITECTURE AND TRUST: NETWORK VS. NETWORK



# ARCHITECTURE AND TRUST: ROUTERS VS. ENDPOINTS

