

1





Wireless Networks




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COS 461: Computer Networks
<http://www.cs.princeton.edu/courses/archive/spring14/cos461/>

2

Widespread Deployment

- Worldwide cellular subscribers
 - 1993: 34 million
 - 2005: more than 2 billion
 - 2012: 6.8 billion
 (2.1B with mobile broadband)
 >> 1.2B landline subscribers
- Wireless local area networks
 - Wireless adapters built into laptops, tablets, & phones
 - More ubiquitous than wired broadband? 700M in 2012

<http://mobithinking.com/mobile-marketing-tools/latest-mobile-stats/#subscribers>

3

Wireless Links

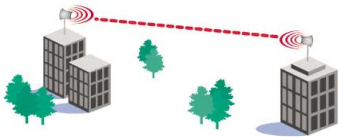
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Wireless Properties

- Interference / bit errors
 - More sources of corruption compared to wired
- Multipath propagation
 - Signal does not travel in a straight line
- Broadcast medium
 - All traffic to everyone
- Power trade-offs
 - Important for power constrained devices

Wireless Links: High Bit Error Rate

- **Decreasing signal strength**
 - Disperses as it travels greater distance
 - Attenuates as it passes through matter



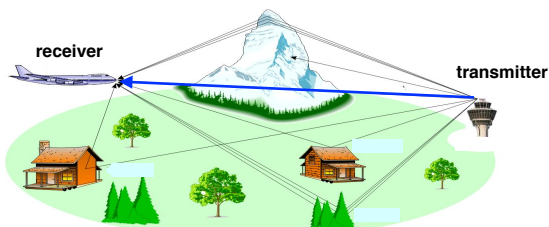
Wireless Links: High Bit Error Rate

- **Interference from other sources**
 - Radio sources in same frequency band
 - E.g., 2.4 GHz wireless phone interferes with 802.11b wireless LAN
 - Electromagnetic noise (e.g., microwave oven)



Wireless Links: High Bit Error Rate

- **Multi-path propagation**
 - Electromagnetic waves reflect off objects
 - Taking many paths of different lengths
 - Causing blurring of signal at the receiver

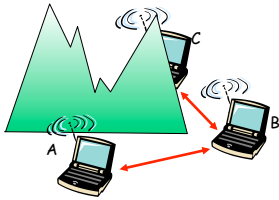


Dealing With Bit Errors

- **Wireless vs. wired links**
 - Wired: most loss is due to congestion
 - Wireless: higher, time-varying bit-error rate
- **Dealing with high bit-error rates**
 - Sender could increase transmission power
 - Requires more energy (bad for battery-powered hosts)
 - Creates more interference with other senders
 - Stronger error detection and recovery
 - More powerful error detection/correction codes
 - Link-layer retransmission of corrupted frames

Wireless Links: Broadcast Limitations

- **Wired broadcast links**
 - E.g., Ethernet bridging, in wired LANs
 - All nodes receive transmissions from all other nodes
- **Wireless broadcast: hidden terminal problem**

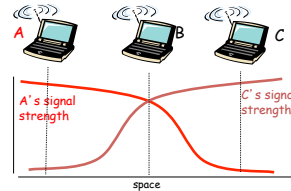


- A and B hear each other
- B and C hear each other
- But, A and C do not

So, A and C are unaware of their interference at B

Wireless Links: Broadcast Limitations

- **Wired broadcast links**
 - E.g., Ethernet bridging, in wired LANs
 - All nodes receive transmissions from all other nodes
- **Wireless broadcast: fading over distance**



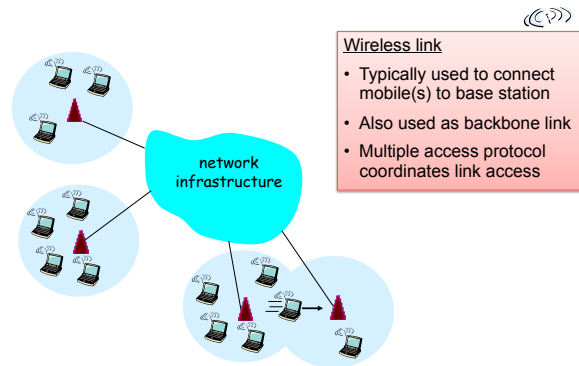
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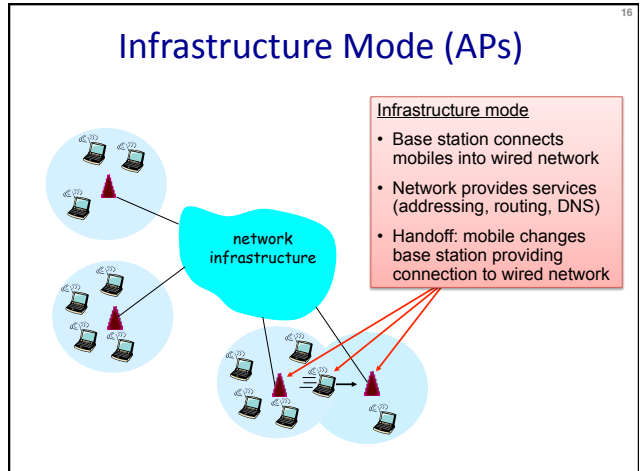
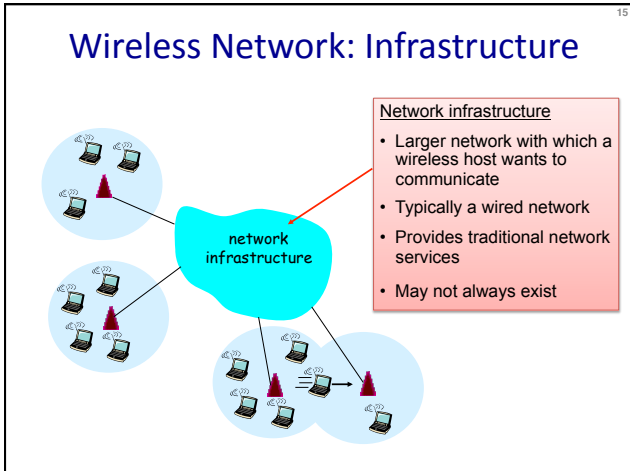
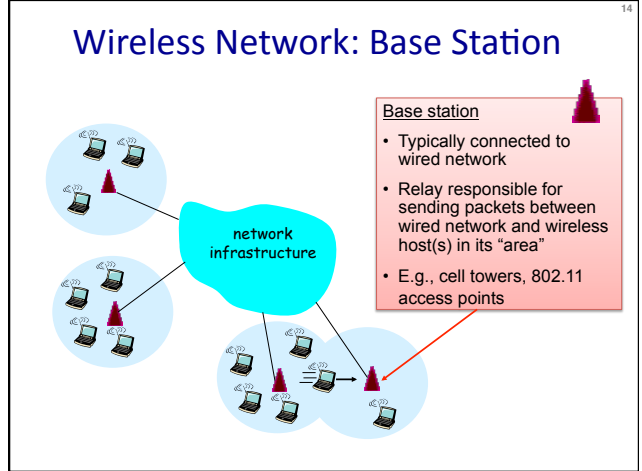
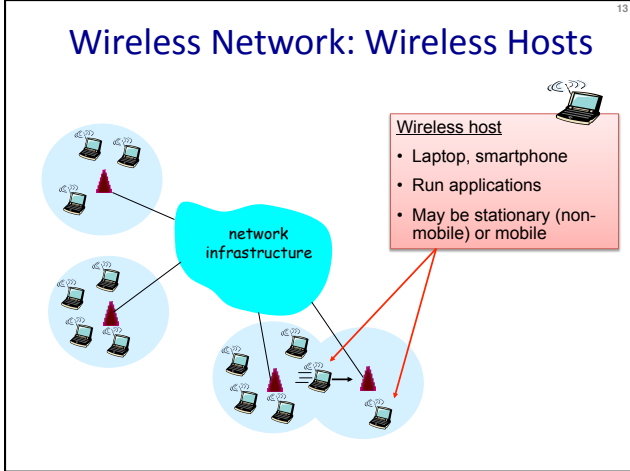
Example Wireless Link Technologies

- **Data networks**
 - 802.15.1 (Bluetooth): 2.1 Mbps – 10 m
 - 802.11b (WiFi): 5-11 Mbps – 100 m
 - 802.11a and g (WiFi): 54 Mbps – 100 m
 - 802.11n (WiFi): 200 Mbps – 100 m
 - 802.16 (WiMax): 70 Mbps – 10 km
- **Cellular networks, outdoors**
 - 2G: 56 Kbps
 - 3G: 384 Kbps
 - 3G enhanced (“4G”): 4 Mbps
 - LTE: 10-100 Mbps

Wireless Network: Wireless Link

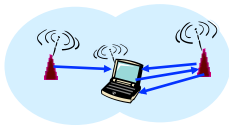
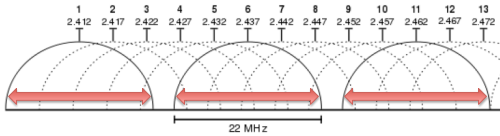


- Wireless link**
- Typically used to connect mobile(s) to base station
 - Also used as backbone link
 - Multiple access protocol coordinates link access



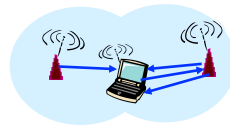
Channels and Association

- **Multiple channels at different frequencies**
 - Network administrator chooses frequency for AP
 - Interference if channel is same as neighboring AP



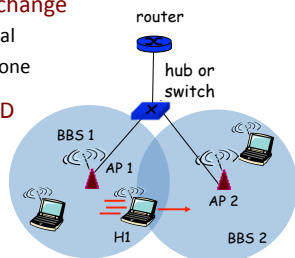
Channels and Association

- **Multiple channels at different frequencies**
 - Network administrator chooses frequency for AP
 - Interference if channel is same as neighboring AP
- **Access points send periodic beacon frames**
 - Containing AP's name (SSID) and MAC address
 - Host scans channels, listening for beacon frames
 - Host selects an access point: association request/response protocol between host and AP



Mobility Within the Same Subnet

- **H1 remains in same IP subnet**
 - IP address of the host can remain same
 - Ongoing data transfers can continue uninterrupted
- **H1 recognizes the need to change**
 - H1 detects a weakening signal
 - Starts scanning for stronger one
- **Changes APs with same SSID**
 - H1 disassociates from one
 - And associates with other
- **Switch learns new location**
 - Self-learning mechanism



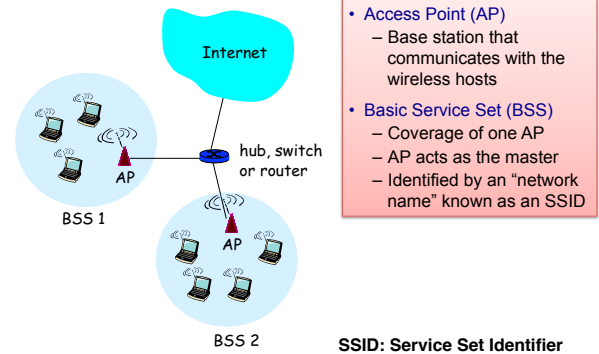
Questions

- **Loss is primary caused by bit errors**
 - (A) Ethernet (Wired)
 - (B) 802.11 (Wireless)
 - (C) Both
 - (D) Neither
- **All hosts on subnet see all communication**
 - (A) Ethernet (Wired)
 - (B) 802.11 (Wireless)
 - (C) Both
 - (D) Neither

WiFi: 802.11 Wireless LANs

21

802.11 LAN Architecture

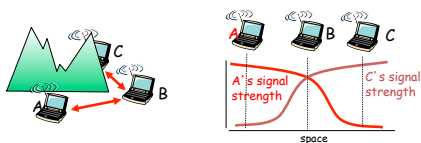


22

CSMA: Carrier Sense, Multiple Access

23

- **Multiple access: channel is shared medium**
 - Station: wireless host or access point
 - Multiple stations may want to transmit at same time
- **Carrier sense: sense channel before sending**
 - Station doesn't send when channel is busy
 - To prevent collisions with ongoing transfers
 - But, detecting ongoing transfers isn't always possible



CA: Collision Avoidance, Not Detection

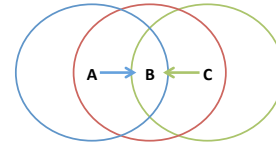
24

- **Collision detection in wired Ethernet**
 - Station listens while transmitting
 - Detects collision with other transmission
 - Aborts transmission and tries sending again
- **Problem #1: cannot detect all collisions**
 - Hidden terminal problem
 - Fading

CA: Collision Avoidance, Not Detection

- **Collision detection in wired Ethernet**
 - Station listens while transmitting
 - Detects collision with other transmission
 - Aborts transmission and tries sending again
- **Problem #1: cannot detect all collisions**
 - Hidden terminal problem
 - Fading
- **Problem #2: listening while sending**
 - Strength of received signal is much smaller
 - Expensive to build hardware that detects collisions
- **So, 802.11 does collision avoidance, not detection**

Hidden Terminal Problem

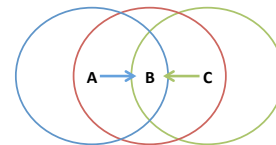


- A and C can't see each other, both send to B
- Occurs b/c 802.11 relies on physical carrier sensing, which is susceptible to hidden terminal problem

Virtual carrier sensing

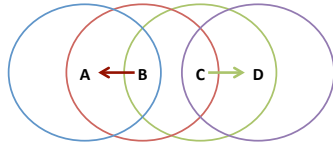
- **First exchange control frames before transmitting data**
 - Sender issues "Request to Send" (RTS), incl. length of data
 - Receiver responds with "Clear to Send" (CTS)
- If sender sees CTS, transmits data (of specified length)
- If other node sees CTS, will idle for specified period
- If other node sees RTS but not CTS, free to send

Hidden Terminal Problem



- A and C can't see each other, both send to B
- **RTS/CTS can help**
 - Both A and C would send RTS that B would see first
 - B only responds with one CTS (say, echoing A's RTS)
 - C detects that CTS doesn't match and won't send

Exposed Terminal Problem



- B sending to A, C wants to send to D
- As C receives packets, carrier sense would prevent it from sending to D, even though wouldn't interfere
- RTS/CTS can help
 - C hears RTS from B, but not CTS from A
 - C knows it's transmission will not interfere with A
 - C is safe to transmit to D

Impact on Higher-Layer Protocols

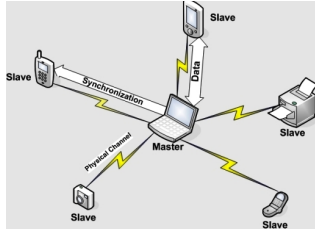
- **Wireless and mobility change path properties**
 - Wireless: higher packet loss, not from congestion
 - Mobility: transient disruptions, and changes in RTT
- **Logically, impact should be minimal ...**
 - Best-effort service model remains unchanged
 - TCP and UDP can (and do) run over wireless, mobile
- **But, performance definitely *is* affected**
 - TCP treats packet loss as a sign of congestion
 - TCP tries to estimate the RTT to drive retransmissions
 - TCP does not perform well under out-of-order packets
- **Internet not designed with these issues in mind**

Questions

- **RTS/CTS more like:**
 - A. Statistical multiplexing
 - B. Time-division multiplexing
 - C. Frequency-division multiplexing
- **Which of following is NOT true?**
 - A. Collisions are minimized when RTS/CTS used.
 - B. Sender can always detect a collision without feedback from receiver.
 - C. TCP congestion control works poorly in wireless without link-layer retransmission.
 - D. Wireless generally has higher loss rates than wired.

Bluetooth: 802.15.1
“personal-area-networks”

Bluetooth piconets



- Up to 7 “slave devices and 225 “parked” devices
- Operates on unlicensed wireless spectrum
 - How to prevent interference?

PHY: Spread Spectrum – Frequency Hopping

- Nodes rapidly jump between frequencies
- Sender and receiver coordinated in jumps
 - How coordinate? Pseudorandom number generator, with shared input known to sender/receiver
- If randomly collide with other transmitted, only for short period before jump again
- Bluetooth
 - 79 frequencies, on each frequency for 625 microseconds
 - Each channel also uses TDMA, with each frame taking 1/3/5 consecutive slots.
 - Only master can start in odd slot, slave only in response

Infrastructure vs. Ad Hoc

- Infrastructure mode
 - Wireless hosts are associated with a base station
 - Traditional services provided by the connected network
 - E.g., address assignment, routing, and DNS resolution
- Ad hoc networks
 - Wireless hosts have no infrastructure to connect to
 - Hosts themselves must provide network services
- Similar in spirit to the difference between
 - Client-server communication
 - Peer-to-peer communication

Conclusions

- Wireless
 - Already a major way people connect to the Internet
 - Gradually becoming more than just an access network
- Mobility (not discussed)
 - Today’s users tolerate disruptions as they move
 - ... and applications try to hide the effects
 - Tomorrow’s users expect seamless mobility
- Challenges the design of network protocols
 - Wireless breaks the abstraction of a link, and the assumption that packet loss implies congestion
 - Mobility breaks association of address and location
 - Higher-layer protocols don’t perform as well

The Upside to Interference

- Some systems leverage interference
- If packets collide once, likely will again
 - Can use both collisions to construct original packets
 - Reduce effective error rate significantly
- If two hosts send to each other through an AP, and collide, AP can broadcast collision to both
 - Both know what they sent, can “subtract” that from collision to get the other
 - Improves throughput of system!