

# Graphs

- Definition
  - $G = (V, E)$
- Representation
  - Edge set
  - Adjacency matrix
  - Adjacency lists
- Graph API
  - create
  - `addEdge(v,w)`
  - `Adj(v)`
- Performance
  - Described in  $|E|$  and  $|V|$

# DFS and BFS

**DFS** (to visit a vertex  $v$ )

---

Mark  $v$  as visited.

Recursively visit all unmarked  
vertices  $w$  adjacent to  $v$ .

---

**BFS** (from source vertex  $s$ )

---

Put  $s$  onto a FIFO queue, and mark  $s$  as visited.

Repeat until the queue is empty:

- remove the least recently added vertex  $v$
  - add each of  $v$ 's unvisited neighbors to the queue,  
and mark them as visited.
-

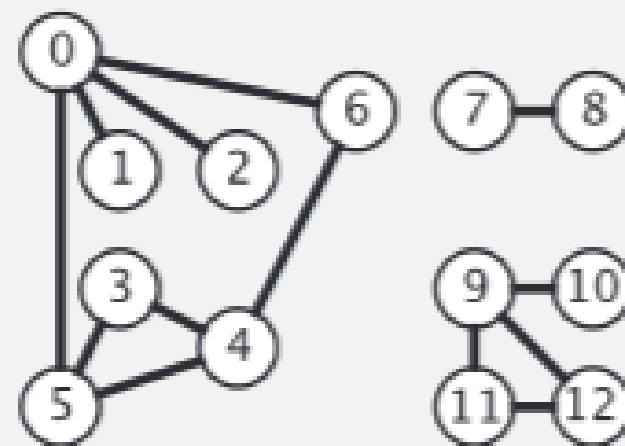
## Connected components

---

Initialize all vertices  $v$  as unmarked.

For each unmarked vertex  $v$ , run DFS to identify all vertices discovered as part of the same component.

---



# Challenge 1

- Determine if a graph is bipartite
- $G = (U, V, E)$ 
  - $U$  and  $V$  are two independent sets whose union is the vertex set
  - Each edge in  $E$  connects a vertex in  $U$  to a vertex in  $V$

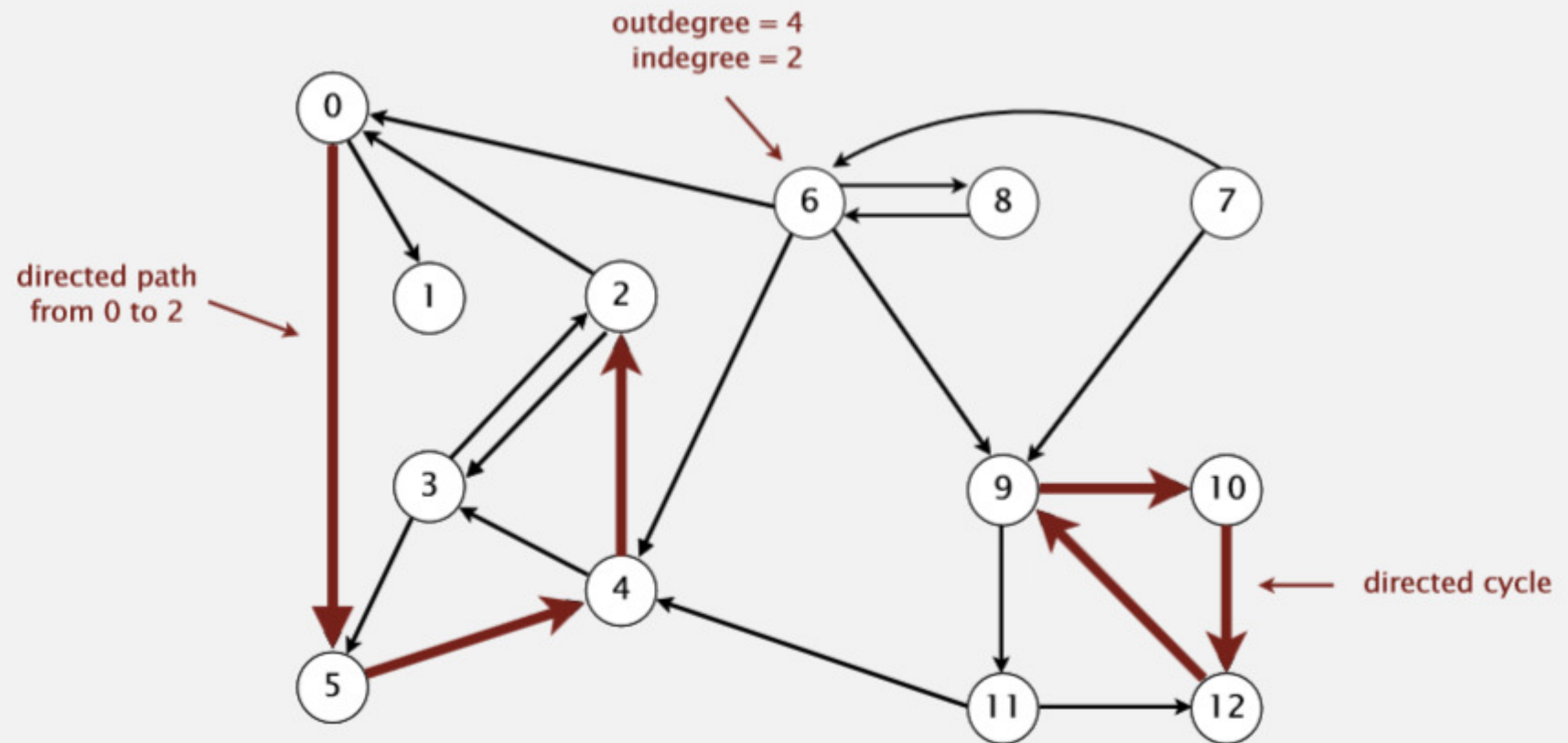
# Challenges

- Find a cycle in a graph (if any)
- Find a cycle in a graph that uses every edge exactly once (Euler path)
- Find a cycle that uses every vertex exactly once (Hamilton cycle)
- Are two graphs isomorphic?
- Layout a graph on a plane w/o crossing edges (Tarjan, 1970)

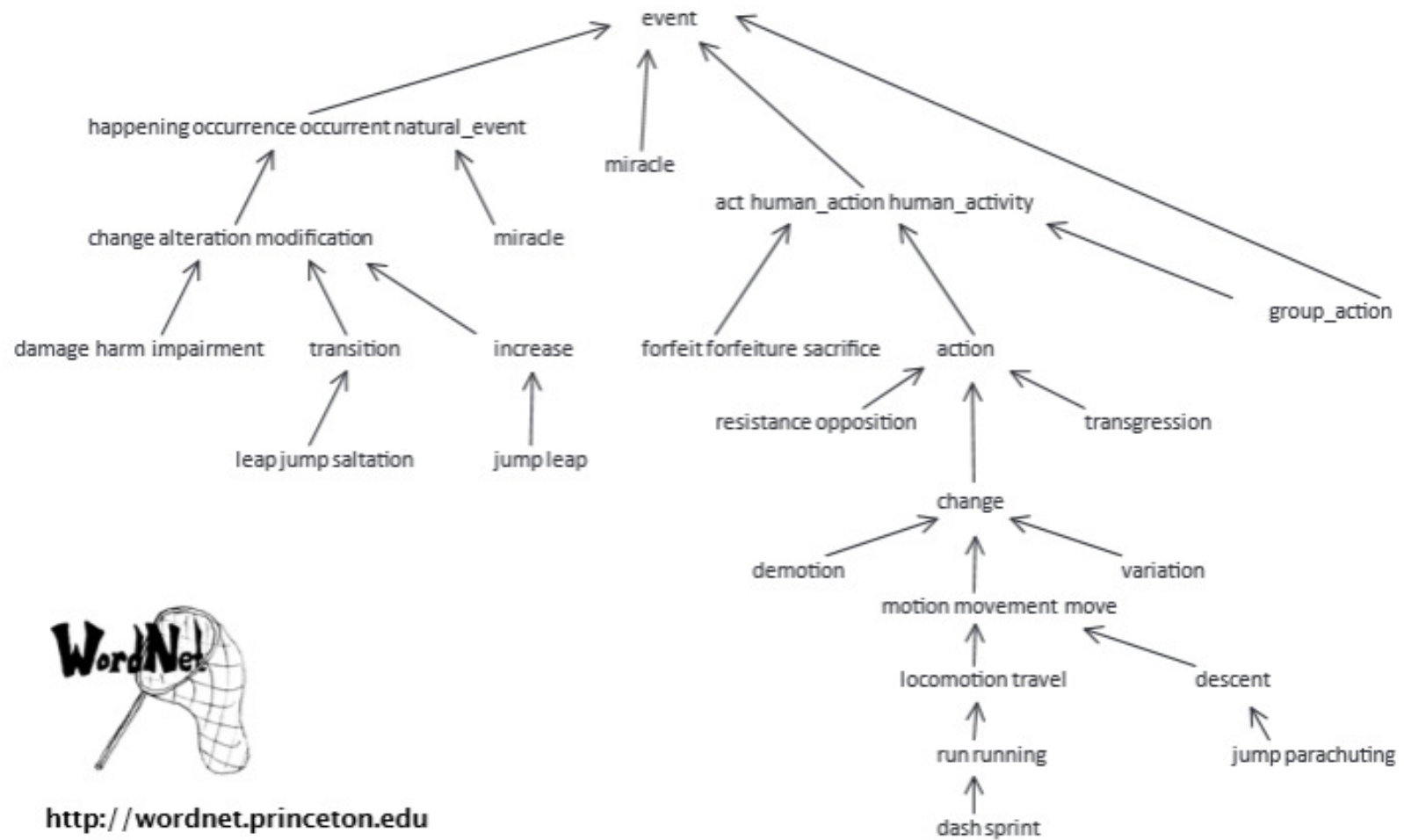
# What we know

problem	BFS	DFS	time
path between s and t	✓	✓	$E + V$
shortest path between s and t	✓		$E + V$
connected components	✓	✓	$E + V$
biconnected components		✓	$E + V$
cycle	✓	✓	$E + V$
Euler cycle		✓	$E + V$
Hamilton cycle			$2^{1.657 V}$
bipartiteness	✓	✓	$E + V$
planarity		✓	$E + V$
graph isomorphism			$2^{c\sqrt{V \log V}}$

# Directed Graphs



Vertex = synset; edge = hypernym relationship.

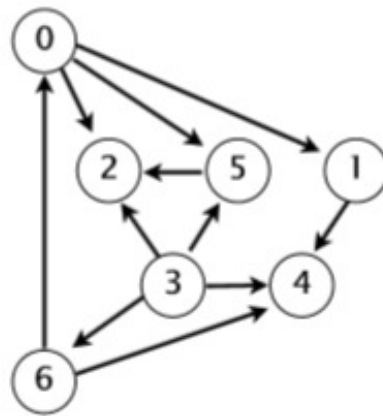




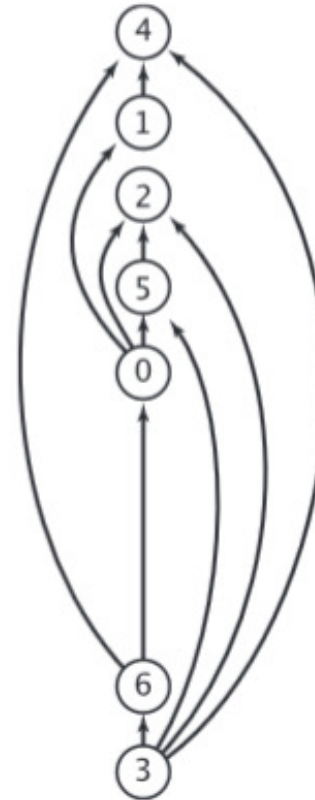
# representation

representation	space	insert edge from $v$ to $w$	edge from $v$ to $w$ ?	iterate over vertices pointing from $v$ ?
list of edges	$E$	1	$E$	$E$
adjacency matrix	$V^2$	1†	1	$V$
adjacency lists	$E + V$	1	$outdegree(v)$	$outdegree(v)$

# Topological sort



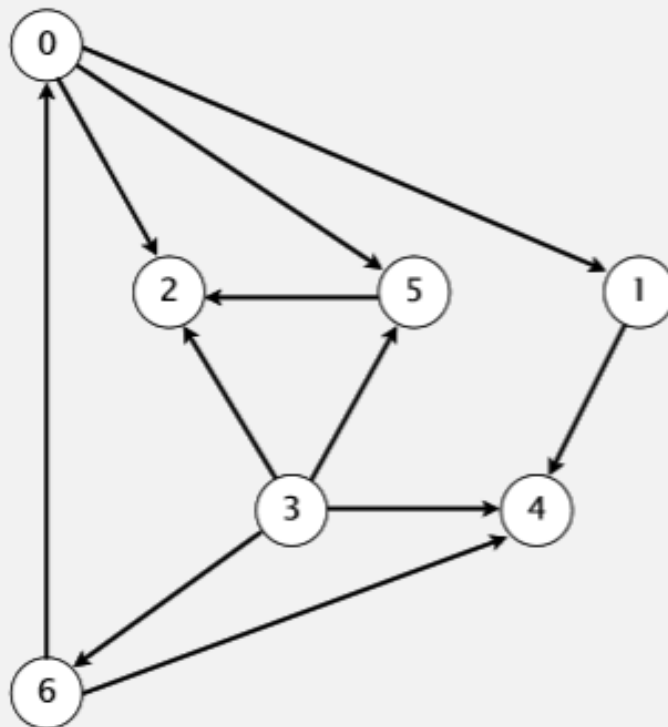
**precedence constraint graph**



**feasible schedule**

# Find a topological sort

- Run depth-first search.
- Return vertices in reverse postorder.



**postorder**

4 1 2 5 0 6 3

**topological order**

3 6 0 5 2 1 4

# DFS orderings

## Orderings.

- Preorder: order in which `dfs()` is called.
- Postorder: order in which `dfs()` returns.
- Reverse postorder: reverse order in which `dfs()` returns.

# What is the relation between these two graphs?

