## COS 226 Midterm Exam, Spring 2010

This test is 10 questions, weighted as indicated. The exam is closed book, except that you are allowed to use a one page cheatsheet. No calculators or other electronic devices are permitted. Give your answers and show your work in the space provided. *Put your name, login ID, and precept number on this page (now)*, and write out and sign the Honor Code pledge before turning in the test. You have 80 minutes to complete the test.

"I pledge my honor that I have not violated the Honor Code during this examination."

1	/5
2	/5
3	/10
4	/5
5	/5
6	/10
7	/10
8	/10
9	/10
10	/20
11	/10
TOTAL	/100

1. **Partitioning** (5 points). Give the result of partitioning the array with standard Quicksort partitioning (taking the N at the left as the partitioning element).



2. Estimating running time (5 points). Suppose that you run the code fragment below (generate and then Mergesort an array of random double values) for N = 10,000,000 and observe that it takes 5.3 seconds.

```
Double[] a = new Double[N];
for (int i = 0; i < N; i++)
a[i] = Math.random();
Merge.sort(a);
```

Assuming you have enough memory which of the following is a reasonable prediction of its running time (in seconds) for N = 1,000,000,000?

- A. 53 seconds.
- B. 340 seconds
- C. 530 seconds.
- D. 680 seconds
- E. 1060 seconds
- F. 5300 seconds

3. **Social networking** (10 points). Suppose that a social networking website FRIENDS needs to support two operations: (*i*) declare A and B to be friends (thus making all of As friends and all of Bs friends friends of each other); and (*ii*) determine whether A and B are friends.

Which APIs should FRIENDS use to support these operations (circle two)?

- A. Queue.
- B. Union-find.
- C. Stack.
- D. Priority queue.
- E. Symbol table.
- F. Randomized queue.

Give the worst case order of growth of the running time that FRIENDS can guarantee for *M* operations, where *N* is the number of people listed on the website (circle one).

- G. N log M.
- H. Mlog N.
- I. N log N.
- J. M.
- K. N log\* M.
- L.  $M \log^* N$ .

In *one or two sentences*, justify your answer (describe how FRIENDS should implement the two operations).

Use a symbol table to map names to ids. For (i), union(A, B). For (ii), find(A, B).

4. **Sorting algorithms** (5 points). Match each of the types of input files described at right below with the *most appropriate sorting algorithm* (as presented in lecture and in the book) by writing the letter corresponding to an algorithm in the blank to the left of the corresponding file type. You should use each letter only once (and leave two letters unused).



- G. Shellsort
- 5. **Random sort** (5 points). Operating under court order, a certain computer company recently decided to randomly assign the order of browsers for customers to choose by using a system sort with the following broken compareTo() implementation.

```
public int compareTo(Browser b)
{ if (Math.random() < 0.5) return -1; else return +1; }</pre>
```

Assume that (since the list of browsers is short) the system uses our version of insertion sort for the task. Where would you prefer that your company's browser be in the list given as input to the sort?

- A. At the beginning.
- B. Second from the beginning.
- C. Doesn't matter, since the sort randomizes the array.
- D. Next to last.
- E. At the end.
- F. Either at the beginning or second from the beginning.

6. Mergesort (10 points). Consider the following implementation of recursive mergesort:

```
public class Merge
  public static void sort(Comparable[] a, Comparable[] aux, int lo, int hi)
     if (hi <= lo) return;
     int mid = lo + (hi - lo) / 2;
     sort(a, aux, lo, mid);
     sort(a, aux, mid+1, hi);
     merge(a, aux, lo, mid, hi); // merges 2 sorted subarrays into a[lo..hi].
     System.out.print(lo + " " + hi + " ");
     for (int i = lo; i <= hi; i++)
       System.out.print(a[i] + " ");
     System.out.println();
  }
 public static void sort(Comparable[] a)
     int N = a.length;
     Comparable[] aux = new Comparable[N];
     sort(a, aux, 0, N-1);
  }
}
```

Note that the last three lines of the recursive method have been instrumented to print the values of the indices and the contents of the array. The output produced by these methods when invoked by the following code appears below in scrambled order:

```
Character[] a = { 'z', 'y', 'x', 'w', 'v', 'u', 't', 's', 'r' };
     Merge.sort(a);
Α.
    0 2 x y z
B.
    34 v w
C.
    0 4 v w x y z
D.
    58rstu
E.
    0 8 r s t u v w x y z
F.
    0 1 y z
G.
   78rs
H.
    56tu
```

Give the order in which these lines actually appear in the output by writing one letter in each of the blanks below (the last one is filled in for you).



7. **LLRB insertion** (10 points). The following diagram shows a left-leaning red-black tree Thick lines are red links.



A. (2 points) Draw the tree that results after E is inserted.



B. (8 points) Draw the tree that results after F is inserted into your tree from Part A. *Hint*: You might find it easiest to convert to the 2-3 tree representation, then do the insertion, then convert back to the red-black tree representation.



8. Heap operations (10 points). Consider the following max-heap:



A. Draw the result of inserting Z.



B. Draw the result of deleting the maximum from the original max-heap shown above (before Z has been inserted).



9. Linear probing (10 points). Give the result of inserting the following keys P R O B I N G into an empty linear probing hash table of size M = 7, using the hash function f(x) = i % 7, where x is the ith letter of the alphabet.

	Ν	0	Ρ	В	R	Ι	G
35	0	1	2	3	4	5	6
i f(i)	16 2	18 4	15 1	2 9 2 2	14 0	7 0	
x	P	R 	0	B I	N	G	

10. 7 sorting algorithms (20 points). The leftmost column is the original input of strings to be sorted, and the rightmost column is the sorted result. The other columns are the contents at some intermediate step during one of the 7 sorting algorithms listed below. Match up each algorithm by writing its letter under the corresponding column. Use each letter exactly once.

rugh	abba	hlue	abba	fivy	abba	nevo	zman	ahha
korn	acdc	cars	hlue	inxs	acdc	korn	vani	acdc
fivv	hlup	devo	card	korn	heck	fivv	vovo	heck
inve	bruc beck	enva	devo	rugh	blue	inya	tatu	blue
card	card	fivv	dido	card	cake	card	etvv	cake
onva	cake	fupl	enva	devo	care	enva	ween	Card
devo	devo	inve	fivv	enva	cher	devo	goal	cher
fuel	enmd	korn	fuol	fuol	devo	fuol	long	devo
tatu	cher	mohy	inve	hlue	dido	long	kiee	dido
atuv	inve	ruch	korn	mohy	doom	mime	nofv	doom
blue	dido	ctuy	mohy	atuv	anua	hluo	nrag	000111
mohy	fupl	tatu	milee	tatu	enmd	mohy	ruch	enmd
abba	doom	abba	rugh	ahha	rugh	abba	nevo	fivv
muco	kiga		rusii	dido	mugo		muco	fual
acol	00172	gool	atuv	mildo	gool	ahor	mima	inva
dido	long	dido	totu	acol	totu	dido	fuol	kiga
hogk	fivy	hogk	lalu Dada	ada	fivy	hogk	hogk	korn
Deck	LIXX	beck	hoal	hode		beck	j nura	long
KISS add	korn	KISS add	doom	beck	kiss	KISS add	anda	nima
acuc	mohu	acuc		KISS	KOIII	acuc	acuc	mehr
yanı nofu	mugo	yanı Dofu	KISS	yanı doom	yanı Dofu	epilia	learn	mugo
NIOIX	muse	A DOLX	nolx	0000	nolx	NOLX	KOrn	muse
aoom	pras	doom	pras	nolx	SLYX	doom		neyo
pras	mims	pras	yanı	pras	pras	pras	brue	noix
уоуо	seal	уоуо	уоуо	yoyo	уоуо	саке	moby	pras
ween	notx	ween	ween	cake	ween	rush	İlXX	rush
zman	tatu	zman	zman	neyo	zman	zman	abba	sea⊥
neyo	rush	neyo	neyo	ween	neyo	ween	enya	styx
cake	yani	cake	cake	zman	inxs	уоуо	cake	tatu
epmd	ween	epmd	epmd	cher	moby	yani	epmd	ween
cher	zman	cher	cher	epmd	fuel	seal	cher	yani
mims	styx	mims	mims	lons	mims	styx	devo	уоуо
lons	уоуо	lons	lons	mims	lons	tatu	dido	zman

\_B\_\_\_C\_\_F\_\_A\_\_E\_\_D\_\_G\_

- A. Bottom-up mergesort
- B. Shellsort
- C. Insertion sort
- D. Quicksort (with no random shuffle)
- E. Selection sort
- F. Top-down mergesort
- G. Heapsort

11. Interval clusters (10 points). Consider the following data type, for intervals on the line:

```
public class Interval implements Comparable<Interval>
{
    private final int left;
    private final int right;
    Interval(int left, int right)
    { this.left = left; this.right = right; }
    public int compareTo(Interval b)
    { return this.left - b.left; }
}
```

For a particular application, *clusters* of intervals are of importance. To find clusters, replace any pair of intervals that intersects (by even an endpoint) by the union of the two intervals, continuing until all intervals do not intersect. For example, the following set of intervals has 3 clusters:



Note that you are guaranteed to have Intervals with non-negative numbers. Given an array of intervals, how many clusters are there? The brute-force algorithm is quadratic, but an enterprising COS226 student figured out a way to find the number of clusters in an array of intervals in *linearithmic* time, with the following code to be added to Interval.

```
public static int count(Interval[] a)
{
    Arrays.sort(a);
    int cnt = 1;
    int max = a[0].right;
    for (int i = 1; i < a.length; i++)
    {
        // "Missing" line of code
        if (a[i].left > max) cnt++;
        if (a[i].right > max) max = a[i].right;
    }
    return cnt;
}
```

In the space below, write the *one line* of code that is missing.