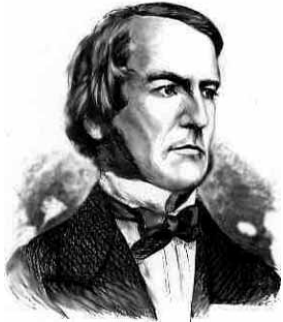


# 6.1 Combinational Circuits



George Boole (1815 - 1864)



Claude Shannon (1916 - 2001)

## Digital signals

- Binary (or "logical") values: 1 or 0, on or off, high or low voltage

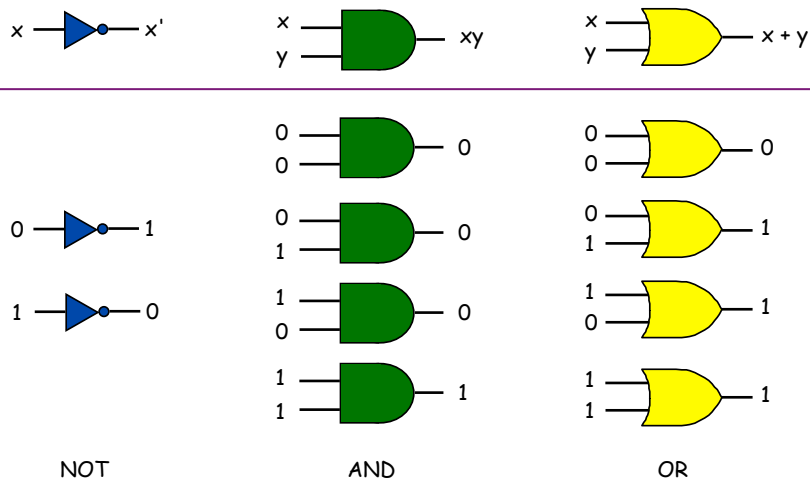
## Wires.

- Propagate digital signals from place to place.
- Signals "flow" from left to right.
  - A drawing convention, sometimes violated
  - Actually: flow from producer to consumer(s) of signal

## Logic Gates

### Logical gates.

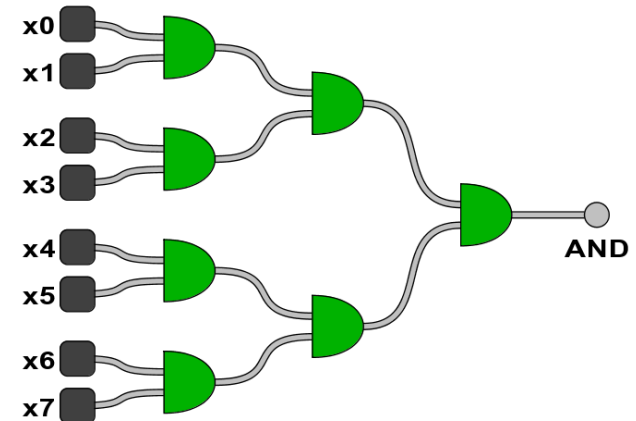
- Fundamental building blocks.



## Multiway AND Gates

### $AND(x_0, x_1, x_2, x_3, x_4, x_5, x_6, x_7)$

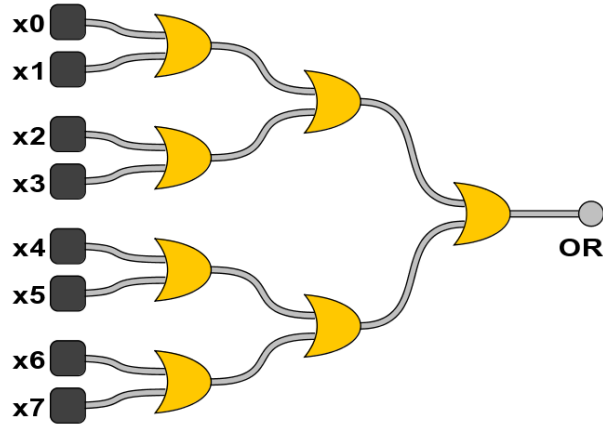
- 1 if all inputs are 1.
- 0 otherwise.



## Multiway OR Gates

$OR(x_0, x_1, x_2, x_3, x_4, x_5, x_6, x_7)$ .

- 1 if at least one input is 1.
- 0 otherwise.

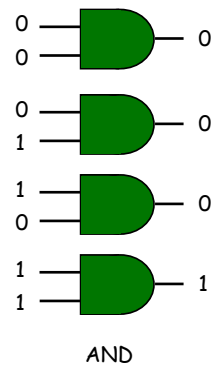


## Truth Table

Truth table.

- Systematic method to describe Boolean function.
- One row for each possible input combination.
- N inputs  $\Rightarrow 2^N$  rows.

AND Truth Table		
x	y	AND(x, y)
0	0	0
0	1	0
1	0	0
1	1	1



5

## Boolean Algebra

History.

- Developed by Boole to solve mathematical logic problems (1847).
- Shannon master's thesis applied it to digital circuits (1937).

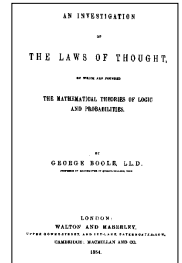
"possibly the most important, and also the most famous, master's thesis of the [20th] century" --Howard Gardner

Basics.

- Boolean variable: value is 0 or 1.
- Boolean function: function whose inputs and outputs are 0, 1.

Relationship to circuits.

- Boolean variables: signals.
- Boolean functions: circuits.



6

## Truth Table for Functions of 2 Variables

Truth table.

- 16 Boolean functions of 2 variables.
- every 4-bit value represents one

Truth Table for All Boolean Functions of 2 Variables									
x	y	ZERO	AND		x	y	XOR	OR	
0	0	0	0	0	0	0	0	0	0
0	1	0	0	0	0	1	1	1	1
1	0	0	0	1	1	0	0	1	1
1	1	0	1	0	1	0	1	0	1

Truth Table for All Boolean Functions of 2 Variables									
x	y	NOR	EQ	y'	x'		NAND	ONE	
0	0	1	1	1	1	1	1	1	1
0	1	0	0	0	0	1	1	1	1
1	0	0	0	1	1	0	0	1	1
1	1	0	1	0	1	0	1	0	1

7

8

## Truth Table for Functions of 3 Variables

Truth table.

- 16 Boolean functions of 2 variables.
  - every 4-bit value represents one
- 256 Boolean functions of 3 variables.
  - every 8-bit value represents one
- $2^{(2^N)}$  Boolean functions of N variables!

Some Functions of 3 Variables						
x	y	z	AND	OR	MAJ	ODD
0	0	0	0	0	0	0
0	0	1	0	1	0	1
0	1	0	0	1	0	1
0	1	1	0	1	1	0
1	0	0	0	1	0	1
1	0	1	0	1	1	0
1	1	0	0	1	1	0
1	1	1	1	1	1	1



## Sum-of-Products

Any Boolean function can be expressed using AND, OR, NOT.

- Sum-of-products is systematic procedure.
  - form AND term for each 1 in truth table of Boolean function
  - OR terms together

Expressing MAJ Using Sum-of-Products								
x	y	z	MAJ	$x'yz$	$xy'z$	$xyz'$	$xyz$	$x'yz + xy'z + xyz' + xyz$
0	0	0	0	0	0	0	0	0
0	0	1	0	0	0	0	0	0
0	1	0	0	0	0	0	0	0
0	1	1	1	1	0	0	0	1
1	0	0	0	0	0	0	0	0
1	0	1	1	0	1	0	0	1
1	1	0	1	0	0	1	0	1
1	1	1	1	0	0	0	1	1

## Universality of AND, OR, NOT

Any Boolean function can be expressed using AND, OR, NOT.

- "Universal"
- Example:  $XOR(x,y) = xy' + x'y$

Notation	Meaning
$x'$	NOT x
$x y$	x AND y
$x + y$	x OR y

Expressing XOR Using AND, OR, NOT							
x	y	$x'$	$y'$	$x'y$	$xy'$	$x'y + xy'$	XOR
0	0	1	1	0	0	0	0
0	1	1	0	1	0	1	1
1	0	0	1	0	1	1	1
1	1	0	0	0	0	0	0

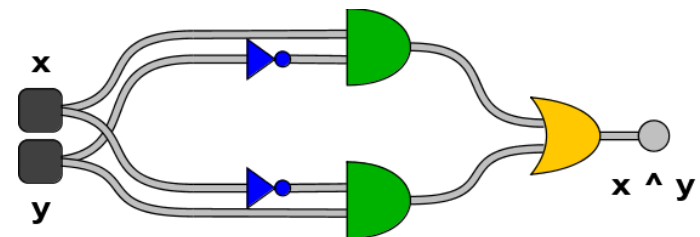
Exercise. Show {AND, NOT}, {OR, NOT}, {NAND}, {AND, XOR} are universal.

Hint. Use DeMorgan's Law:  $(xy)' = (x' + y')$  and  $(x + y)' = (x'y)'$

## Translate Boolean Formula to Boolean Circuit

Use sum-of-products form.

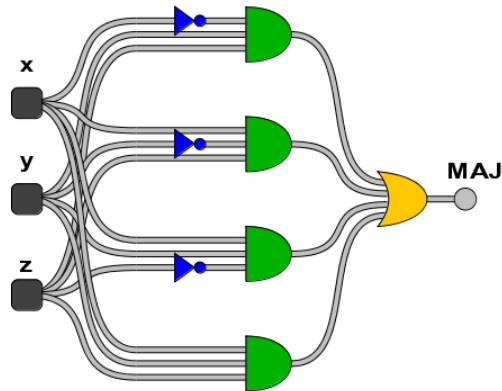
- $XOR(x, y) = xy' + x'y$ .



## Translate Boolean Formula to Boolean Circuit

Use sum-of-products form.

- MAJ(x, y, z) =  $x'yz + xy'z + xyz' + xyz$ .

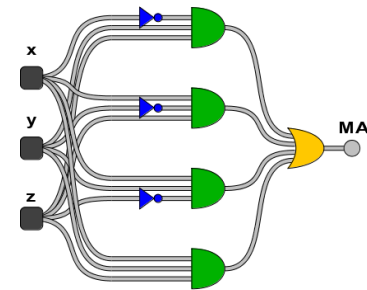


13

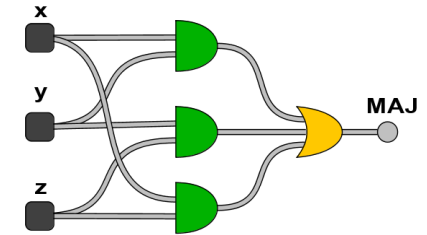
## Simplification Using Boolean Algebra

Many possible circuits for each Boolean function.

- Sum-of-products not necessarily optimal in:
  - number of gates (space)
  - depth of circuit (time)
- $MAJ(x, y, z) = x'yz + xy'z + xyz' + xyz = xy + yz + xz$ .



size = 8, depth = 3



size = 4, depth = 2

14

## Expressing a Boolean Function Using AND, OR, NOT

Ingredients.

- AND gates.
- OR gates.
- NOT gates.
- Wire.

Instructions.

- Step 1: represent input and output signals with Boolean variables.
- Step 2: construct truth table to carry out computation.
- Step 3: derive (simplified) Boolean expression using sum-of products.
- Step 4: transform Boolean expression into circuit.

15

## ODD Parity Circuit

ODD(x, y, z).

- 1 if odd number of inputs are 1.
- 0 otherwise.

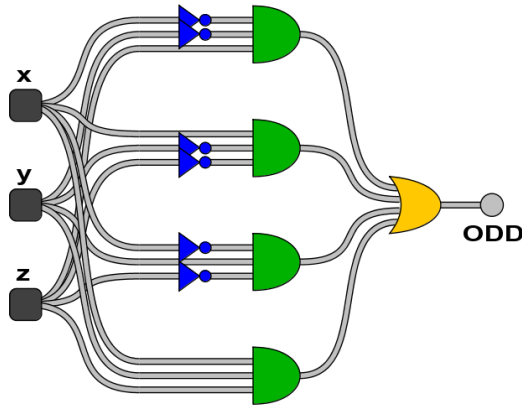
Expressing ODD Using Sum-of-Products								
x	y	z	ODD	$x'y'z$	$x'yz'$	$xy'z'$	$xyz$	$x'y'z + x'yz' + xy'z' + xyz$
0	0	0	0	0	0	0	0	0
0	0	1	1	1	0	0	0	1
0	1	0	1	0	1	0	0	1
0	1	1	0	0	0	0	0	0
1	0	0	1	0	0	1	0	1
1	0	1	0	0	0	0	0	0
1	1	0	0	0	0	0	0	0
1	1	1	1	0	0	0	1	1

16

## ODD Parity Circuit

ODD(x, y, z).

- 1 if odd number of inputs are 1.
- 0 otherwise.



17

## Let's Make an Adder Circuit

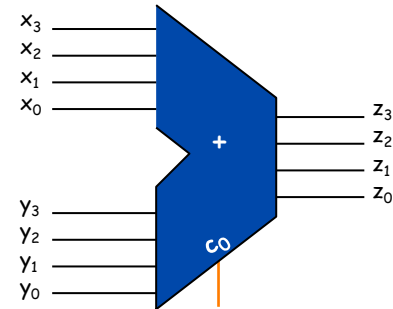
Goal:  $x + y = z$  for 4-bit integers.

- We build 4-bit adder: 9 inputs, 4 outputs.
- Same idea scales to 128-bit adder.
- Key computer component.

	1	1	1	0
	2	4	8	7
+	3	5	7	9
	6	0	6	6

Step 1.

- Represent input and output in binary.



	1	1	0	0
	0	0	1	0
+	0	1	1	1
	1	0	0	1

	$x_3$	$x_2$	$x_1$	$x_0$
+	$y_3$	$y_2$	$y_1$	$y_0$
	$z_3$	$z_2$	$z_1$	$z_0$

18

## Let's Make an Adder Circuit

Goal:  $x + y = z$  for 4-bit integers.

Step 2. (first attempt)

- Build truth table.
- Why is this a bad idea?
  - 128-bit adder:  $2^{256+1}$  rows > # electrons in universe!

					$c_0$
	$x_3$	$x_2$	$x_1$	$x_0$	
+	$y_3$	$y_2$	$y_1$	$y_0$	
	$z_3$	$z_2$	$z_1$	$z_0$	

4-Bit Adder Truth Table												
$c_0$	$x_3$	$x_2$	$x_1$	$x_0$	$y_3$	$y_2$	$y_1$	$y_0$	$z_3$	$z_2$	$z_1$	$z_0$
0	0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	1	0	0	0
0	0	0	0	0	0	0	0	1	0	0	0	1
0	0	0	0	0	0	0	0	1	0	0	1	0
0	0	0	0	0	0	0	1	1	0	0	1	1
0	0	0	0	0	0	1	0	0	0	1	0	0
0	0	0	0	0	0	1	0	1	0	1	0	1
.	.	.	.	.	.	.	.	.	.	.	.	.
1	1	1	1	1	1	1	1	1	1	1	1	1

$2^{8+1} = 512$  rows!

19

## Let's Make an Adder Circuit

Goal:  $x + y = z$  for 4-bit integers.

Step 2. (do one bit at a time)

- Build truth table for carry bit.
- Build truth table for summand bit.

	$c_3$	$c_2$	$c_1$	$c_0 = 0$
	$x_3$	$x_2$	$x_1$	$x_0$
+	$y_3$	$y_2$	$y_1$	$y_0$
	$z_3$	$z_2$	$z_1$	$z_0$

Carry Bit			
$x_i$	$y_i$	$c_i$	$c_{i+1}$
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	1
1	0	0	0
1	0	1	1
1	1	0	1
1	1	1	1

Summand Bit			
$x_i$	$y_i$	$c_i$	$z_i$
0	0	0	0
0	0	1	1
0	1	0	1
0	1	1	0
1	0	0	1
1	0	1	0
1	1	0	0
1	1	1	1

20

## Let's Make an Adder Circuit

Goal:  $x + y = z$  for 4-bit integers.

Step 3.

- Derive (simplified) Boolean expression.

	$c_3$	$c_2$	$c_1$	$c_0 = 0$
$x_3$	$x_2$	$x_1$	$x_0$	
$+$	$y_3$	$y_2$	$y_1$	$y_0$
	$z_3$	$z_2$	$z_1$	$z_0$

Carry Bit				
$x_i$	$y_i$	$c_i$	$c_{i+1}$	MAJ
0	0	0	0	0
0	0	1	0	0
0	1	0	0	0
0	1	1	1	1
1	0	0	0	0
1	0	1	1	1
1	1	0	1	1
1	1	1	1	1

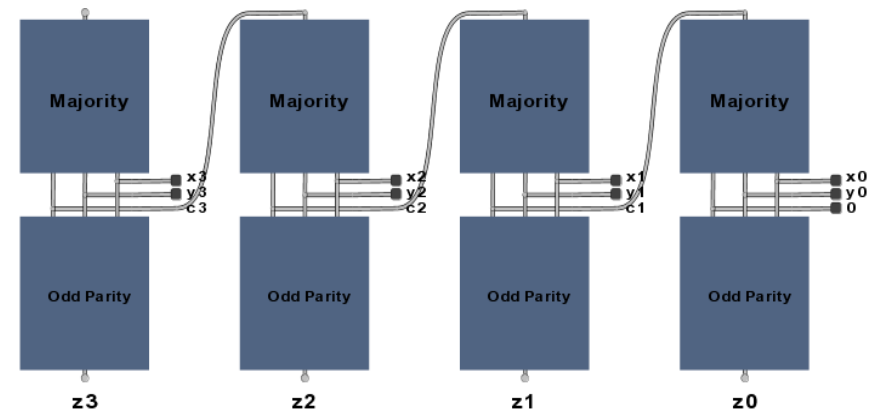
Summand Bit				
$x_i$	$y_i$	$c_i$	$z_i$	ODD
0	0	0	0	0
0	0	1	1	1
0	1	0	1	1
0	1	1	0	0
1	0	0	1	1
1	0	1	0	0
1	1	0	0	0
1	1	1	1	1

## Let's Make an Adder Circuit

Goal:  $x + y = z$  for 4-bit integers.

Step 4.

- Transform Boolean expression into circuit.
- Chain together 1-bit adders.



21

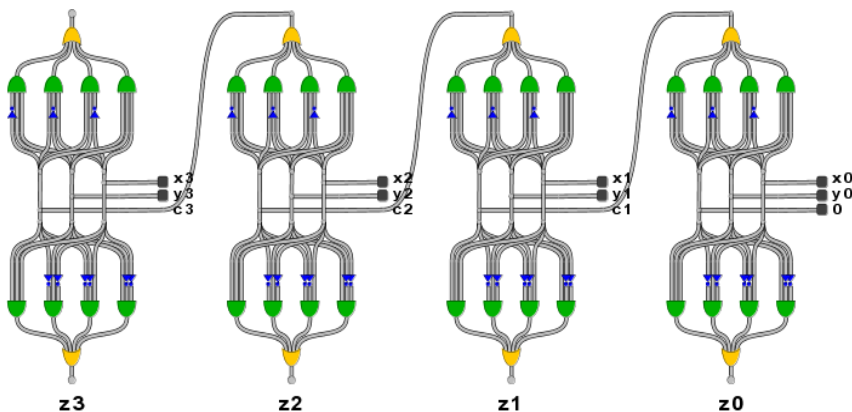
22

## Let's Make an Adder Circuit

Goal:  $x + y = z$  for 4-bit integers.

Step 4.

- Transform Boolean expression into circuit.
- Chain together 1-bit adders.

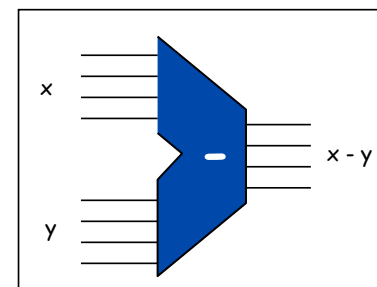


23

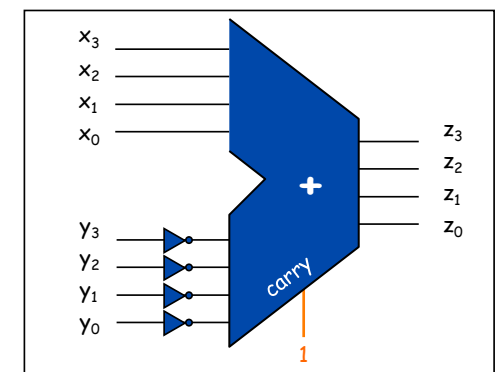
## Subtractor

Subtractor circuit:  $z = x - y$ .

- One approach: design like adder circuit.
- Better idea: reuse adder circuit.
  - 2's complement: to negate an integer, flip bits, then add 1



4-Bit Subtractor Interface



4-Bit Subtractor Implementation

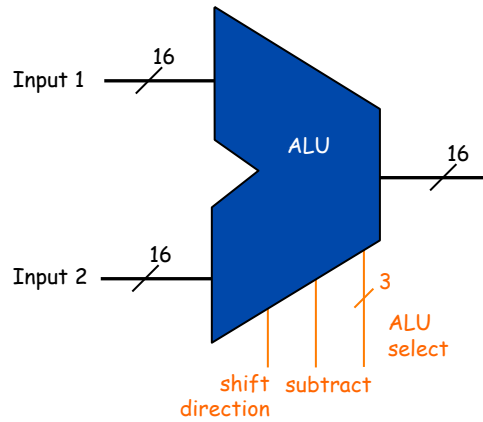
24

## Arithmetic Logic Unit: Interface

### ALU Interface.

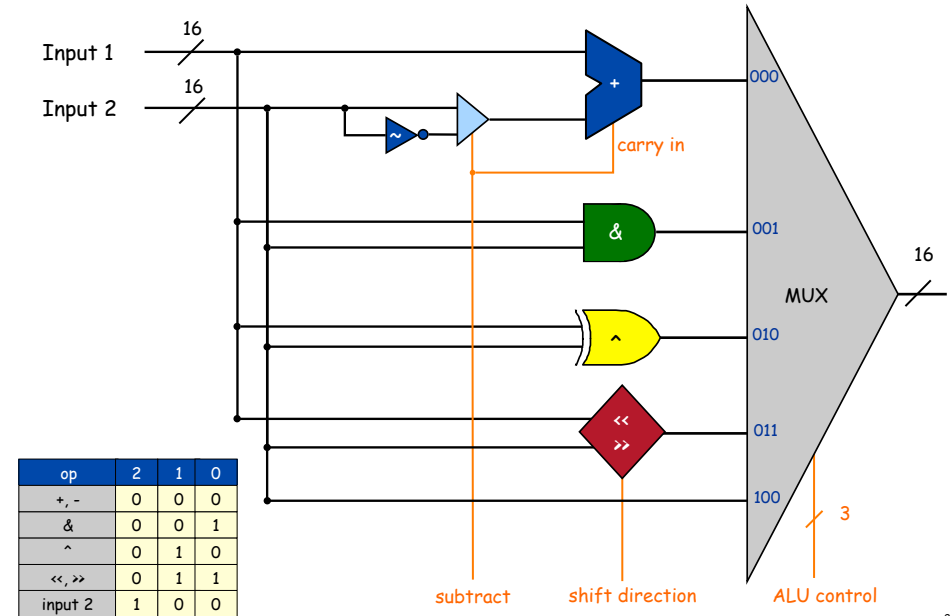
- Add, subtract, bitwise and, bitwise xor, shift left, shift right, copy.
- Associate 3-bit integer with 5 primary ALU operations.
  - ALU performs operations in parallel
  - control wires select which result ALU outputs

op	2	1	0
+, -	0	0	0
&	0	0	1
^	0	1	0
<<, >>	0	1	1
input 2	1	0	0



25

## Arithmetic Logic Unit: Implementation



op	2	1	0
+, -	0	0	0
&	0	0	1
^	0	1	0
<<, >>	0	1	1
input 2	1	0	0

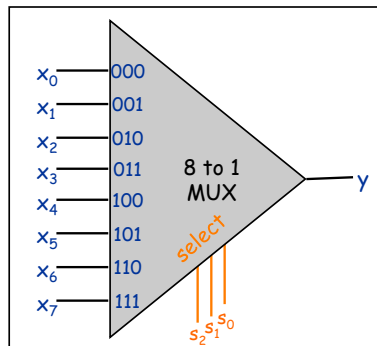
26

## 2<sup>n</sup>-to-1 Multiplexer

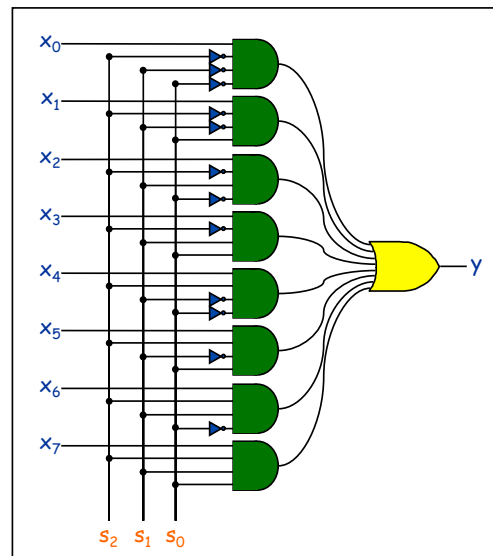
← n = 8 for main memory

### 2<sup>n</sup>-to-1 multiplexer.

- n select inputs, 2<sup>n</sup> data inputs, 1 output.
- Copies "selected" data input bit to output.



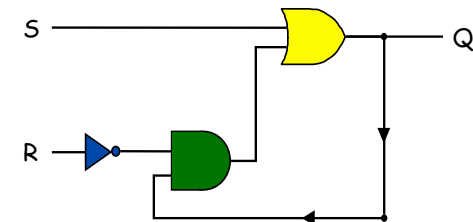
8-to-1 Mux Interface



8-to-1 Mux Implementation

27

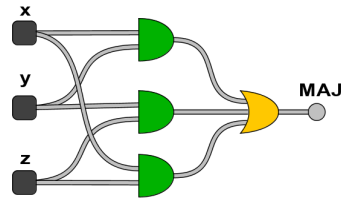
## 6.2: Sequential Circuits



## Sequential vs. Combinational Circuits

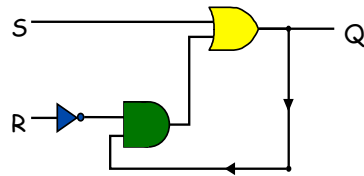
### Combinational circuits.

- Output determined solely by inputs.
- Can draw solely with left-to-right signal paths.



### Sequential circuits.

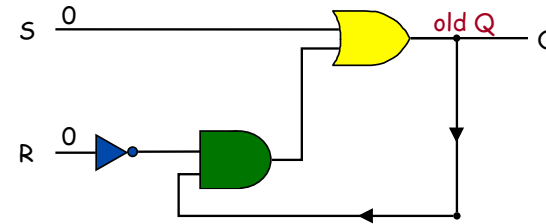
- Output determined by inputs AND previous outputs.
- Feedback loop.



## SR Flip-Flop

What is the value of Q if:

- $S = 1$  and  $R = 0$  ?  $\Rightarrow$  Q is surely 1.
- $S = 0$  and  $R = 1$  ?  $\Rightarrow$  Q is surely 0.
- $S = 0$  and  $R = 0$  ?  $\Rightarrow$  Q is possibly 0 . . . or possibly 1.



While  $S = R = 0$ , Q remembers what it was the last time S or R was 1.

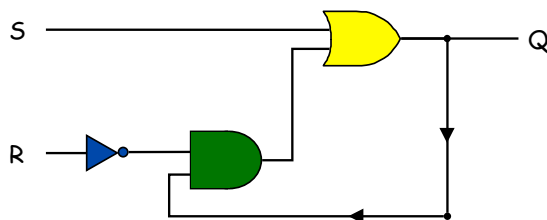
29

30

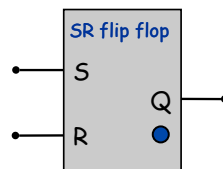
## SR Flip-Flop

### SR Flip-Flop.

- $S = 1, R = 0$  (set)  $\Rightarrow$  "Flips" bit on.
- $S = 0, R = 1$  (reset)  $\Rightarrow$  "Flips" bit off.
- $S = R = 0$   $\Rightarrow$  Status quo.
- $S = R = 1$   $\Rightarrow$  Not allowed.



Implementation

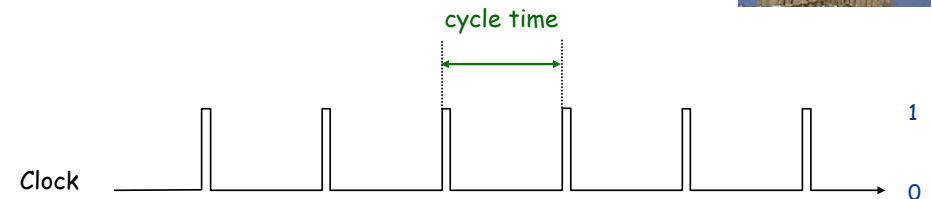


Interface

## Clock

### Clock.

- Fundamental abstraction.
  - regular on-off pulse
- External analog device.
- Synchronizes operations of different circuit elements.
- 1 GHz clock means 1 billion pulses per second.



31

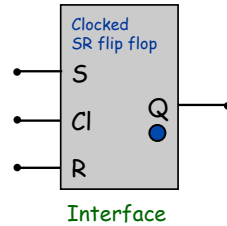
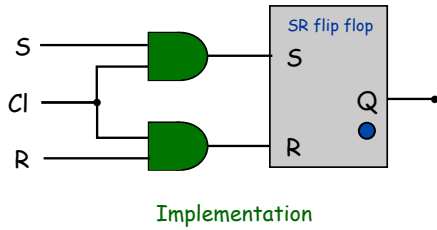
32



## Clocked SR Flip-Flop

### Clocked SR Flip-Flop.

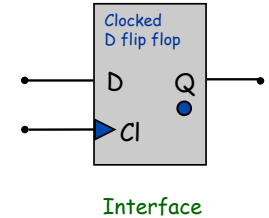
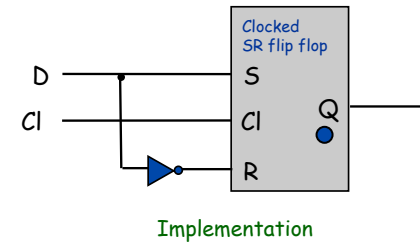
- Same as SR flip-flop except S and R only active when clock is 1.



## Clocked D Flip-Flop

### Clocked D Flip-Flop.

- Output follows D input while clock is 1.
- Output is remembered while clock is 0.



33

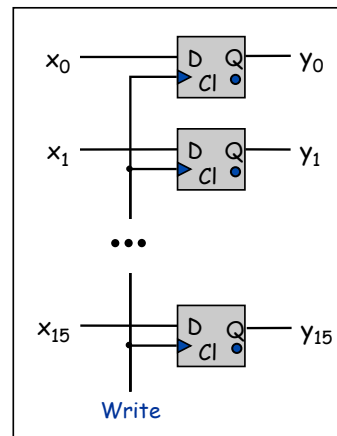
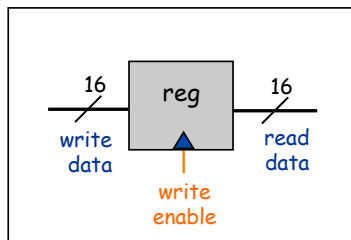
34

## Stand-Alone Register

### k-bit register.

- Stores k bits.
- Register contents always available on output.
- If write enable is asserted, k input bits get copied into register.

Ex: Program Counter, 16 TOY registers, 256 TOY memory locations.

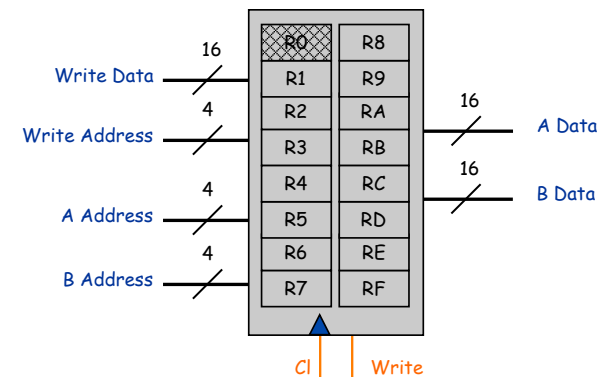


35

## Registers

### TOY registers: fancy 16 x 16-bit register file.

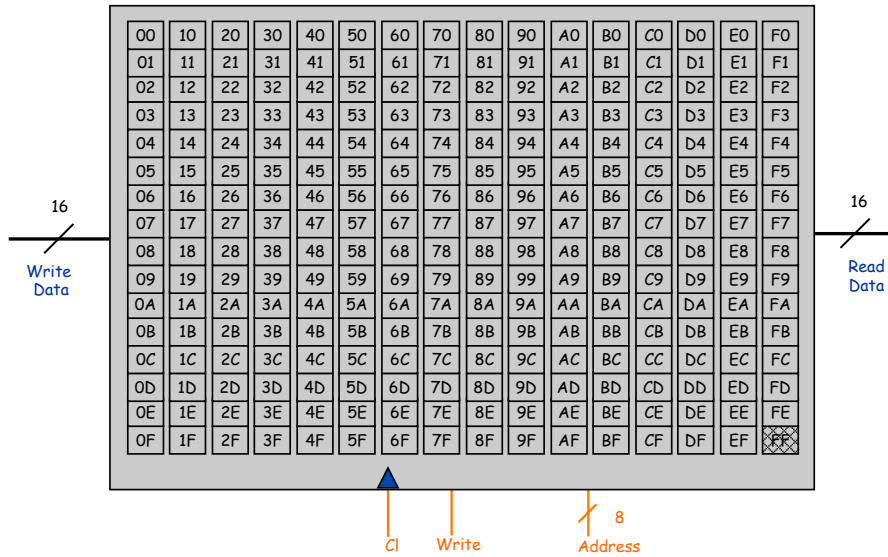
- Want to be able to read two registers, and write to a third in the same instructions:  $R1 \leftarrow R2 + R3$ .
- 3 address inputs, 1 data input, 2 data outputs.
- Add decoders and muxes for additional ports.



36

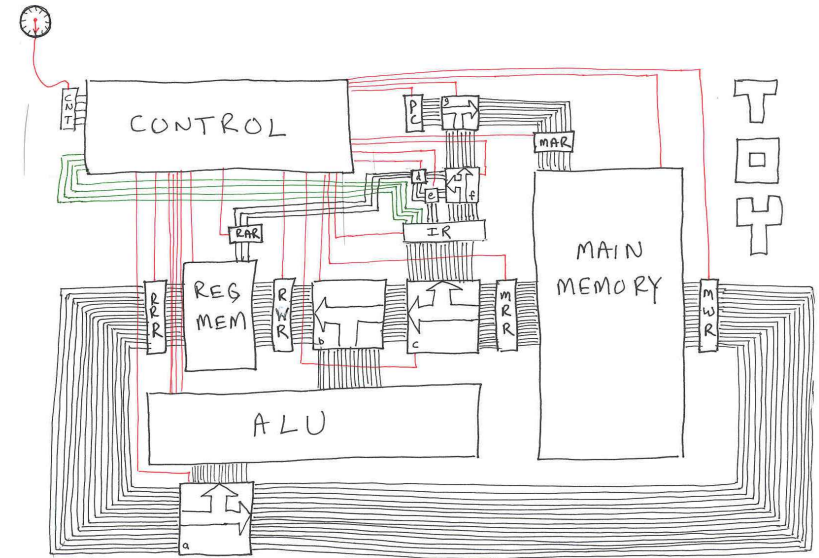
## Main Memory

TOY main memory: 256 x 16-bit register file.

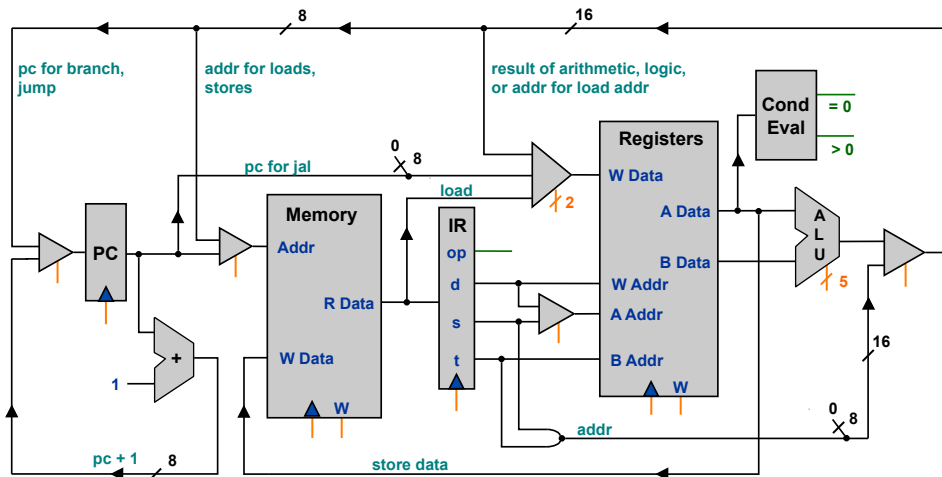


37

## 6.3: Machine Architecture



## The TOY Datapath

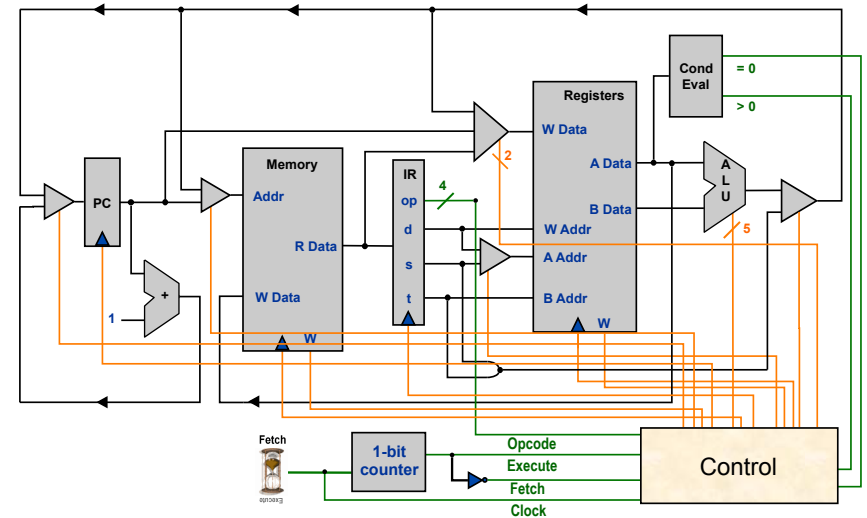


39

## Control

Control: controls components, enables connections.

- Input: opcode, clock, conditional evaluation. (green)
- Output: control wires. (orange)



40

## Summary

### Combinational circuits: how to compute things

- And, Or, Not primitives sufficient for any Boolean function
- Systematic method: truth tables and sum-of-products
- Examples
  - Majority
  - Binary adder
  - Multiplexor

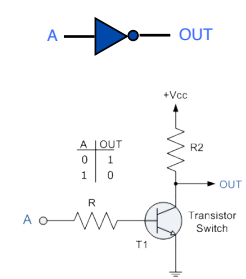
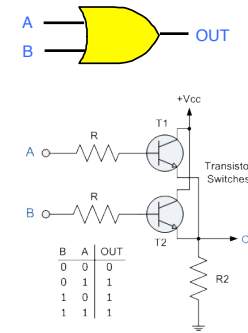
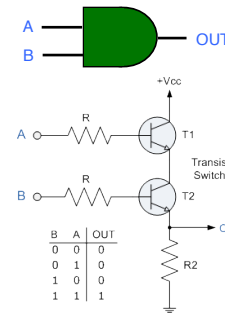
### Sequential circuits: where to put things

- Flip flop primitive holds one bit
- Many flip flops make a *register* (16 for TOY)
- Many registers make a *register file*
- Lots and lot of registers make a *memory* (256 for TOY)

### A whole computer

- Uses *combinational circuits* to perform computations
- Uses *sequential circuits* to store results
- Uses a little of each for control

## The final secret



All three of our logic primitives can be made using a *single\** type of electronic primitive: the *transistor*!