

5. The TOY Machine II



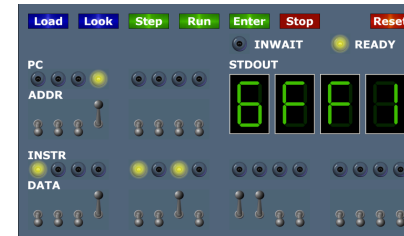
Laboratory Instrument Computer (LINC)

TOY machine.

- Box with switches and lights.
- 16-bit memory locations, 16-bit registers, 8-bit pc.
- $4,328 \text{ bits} = (255 \times 16) + (15 \times 16) + (8) = 541 \text{ bytes!}$
- von Neumann architecture.

TOY programming.

- TOY instruction set architecture: 16 instruction types.
- Variables, arithmetic, loops.



What We Do Today

Data representation. Negative numbers.

Input and output. Standard input, standard output.

Manipulate addresses. References (pointers) and arrays.

TOY simulator in Java.

Data Representation



Data is a sequence of bits. (interpreted in different ways)

- Integers, real numbers, characters, strings, ...
- Documents, pictures, sounds, movies, Java programs, ...

Ex. 01110101

- As binary integer: $1 + 4 + 16 + 32 + 64 = 117_{10}$.
- As character: 117th Unicode character = 'u'.
- As music: 117/256 position of speaker.
- As grayscale value: 45.7% black.

Decimal and binary addition.

carries

1	1 1
013	0 0 0 0 1 1 0 1
+ 092	+ 0 1 0 1 1 1 0 0
105	0 1 1 0 1 0 0 1

Subtraction. Add a negative integer.

e.g., $6 - 4 = 6 + (-4)$

Q. How to represent negative integers?

Representing Negative Integers

TOY words are 16 bits each.

- We could use 16 bits to represent 0 to $2^{16} - 1$.
- We want negative integers too.
- Reserving half the possible bit-patterns for negative seems fair.

Highly desirable property. If x is an integer, then the representation of -x, when added to x, is zero.

$$\begin{array}{r}
 x \quad \quad 0 \ 0 \ 1 \ 1 \ 0 \ 1 \ 0 \ 0 \\
 + (-x) \quad + \ ? \ ? \ ? \ ? \ ? \ ? \ ? \\
 \hline
 0 \quad \quad 0 \ 0 \ 0 \ 0 \ 0 \ 0 \ 0 \ 0
 \end{array}$$

flip bits and add 1

$$\begin{array}{r}
 x \quad \quad 0 \ 0 \ 1 \ 1 \ 0 \ 1 \ 0 \ 0 \\
 + (-x) \quad + \ 1 \ 1 \ 0 \ 0 \ 1 \ 0 \ 1 \ 1 \\
 \hline
 0 \quad \quad 0 \ 0 \ 0 \ 0 \ 0 \ 0 \ 0 \ 0
 \end{array}$$

Two's Complement Integers

To compute -x from x:

- Start with x.

+4

0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	0	0
---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---

- Flip bits.

-5

1	1	1	1	1	1	1	1	1	1	1	1	1	1	0	1	1
---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---

- Add one.

-4

1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	0	0
---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---

Two's Complement Integers

		15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
dec	hex	binary															
+32767	7FFF	0	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1
...																	
+4	0004	0	0	0	0	0	0	0	0	0	0	0	0	0	1	0	0
+3	0003	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	1
+2	0002	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	0
+1	0001	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1
+0	0000	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
-1	FFFF	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1
-2	FFFE	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	0
-3	FFFD	1	1	1	1	1	1	1	1	1	1	1	1	1	1	0	1
-4	FFFC	1	1	1	1	1	1	1	1	1	1	1	1	1	1	0	0
...																	
-32768	8000	1	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0

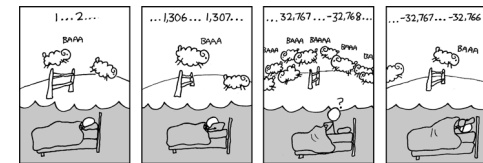
Properties of Two's Complement Integers

Properties.

- Leading bit (bit 15) signifies sign.
- Addition and subtraction are easy.
- 0000000000000000 represents zero.
- Checking for arithmetic overflow is easy.
- Negative integer $-x$ represented by $2^{16} - x$.
- Not symmetric: can represent -32768 but not 32768 .

Java. Java's `int` data type is a 32-bit two's complement integer.

Ex. $2147483647 + 1$ equals -2147483648 .



<http://xkcd.com/571>

9

10

Representing Other Primitive Data Types in TOY

Bigger integers. Use two 16-bit TOY words per 32-bit Java `int`.

Real numbers.

- Use IEEE floating point (like scientific notation).
- Use four 16-bit TOY words per 64-bit Java `double`.

Characters.

- Use Unicode (16 bits per char).
- Use one 16-bit TOY word per 16-bit Java `double`.

	0	1	2	3	4	5	6	7	8	9	A	B	C	D	E	F
0	NUL	SOH	STX	ETX	END	CR	LF	VT	FF	CR	SO	SI	DEL			
1	SP	DEL	DC1	DC2	DC3	DC4	NAK	SYN	ETB	CAN	EM	END	ESC	FS	CS	RS
2	SP	!	"	#	\$	%	&	'	()	*	+	,	-	.	/
3	0	1	2	3	4	5	6	7	8	9	:	<	=	>	?	
4	@	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O
5	P	Q	R	S	T	U	V	W	X	Y	Z	[\]	^	_
6	`	a	b	c	d	e	f	g	h	i	j	k	l	m	n	o
7	p	q	r	s	t	u	v	w	x	y	z	{		}	~	---

Hexadecimal-to-ASCII conversion table

Note. Real microprocessors add hardware support for `int` and `double`.

Standard Input and Output

Standard Output

Standard output.

- Writing to memory location `FF` sends one word to TOY stdout.
- Ex. `9AFF` writes the integer in register `A` to stdout.

```

00: 0000  0
01: 0001  1

10: 8A00  RA ← mem[00]      a = 0
11: 8B01  RB ← mem[01]      b = 1
12: 9AFF  write RA to stdout  print a
13: 1AAB  RA ← RA + RB          a = a + b
14: 2BAB  RB ← RA - RB          b = a - b
15: DA12  if (RA > 0) goto 12  } while (a > 0)
16: 0000  halt
    
```

fibonacci.toy

```

0000
0001
0001
0002
0003
0005
0008
0008
000D
0015
0022
0037
0059
0090
0090
00E9
0179
0262
03DB
063D
0A18
1055
1A6D
2AC2
452F
6FF1
    
```

13

Standard Input

Standard input.

- Loading from memory address `FF` loads one word from TOY stdin.
- Ex. `8AFF` reads an integer from stdin and store it in register `A`.

Ex: read in a sequence of integers and print their sum.

- In Java, stop reading when EOF.
- In TOY, stop reading when user enters 0000.

```

while (!StdIn.isEmpty()) {
    a = StdIn.readInt();
    sum = sum + a;
}
StdOut.println(sum);
    
```

```

00: 0000  0
10: 8C00  RC ← mem[00]
11: 8AFF  read RA from stdin
12: CA15  if (RA == 0) pc ← 15
13: 1CCA  RC ← RC + RA
14: C011  pc ← 11
15: 9CFF  write RC
16: 0000  halt
    
```

```

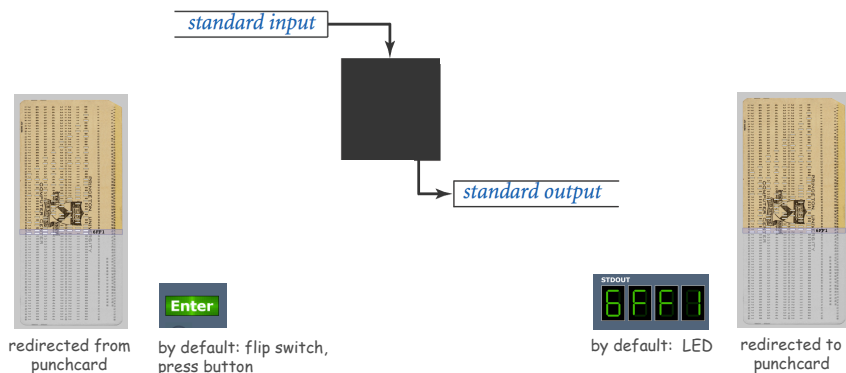
00AE
0046
0003
0000
00F7
    
```

14

Standard Input and Output: Implications

Standard input and output enable you to:

- Get information out of machine.
- Put information from real world into machine.
- Process more information than fits in memory.
- Interact with the computer while it is running.



15

Pointers

16

Load Address (a.k.a. Load Constant)

Load address. [opcode 7]

- Loads an 8-bit integer into a register.
- 7A30 means load the value 30 into register a.

Applications.

- Load a small **constant** into a register.
- Load a 8-bit **memory address** into a register.

```
a = 0x30;
```

Java code

register stores "pointer" to a memory cell

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	1	1	1	0	1	0	0	0	1	1	0	0	0	0
7 ₁₆				A ₁₆				3 ₁₆				0 ₁₆			
opcode				dest d				addr							

17

TOY Implementation of Reverse

TOY implementation of reverse.

- Read in a sequence of integers and store in memory 30, 31, 32, ...
- Stop reading if 0000.
- Print sequence in reverse order.

```

10: 7101 R1 ← 0001      constant 1
11: 7A30 RA ← 0030      a[]
12: 7B00 RB ← 0000      n

13: 8CFF read RC          while(true) {
14: CC19 if (RC == 0) goto 19    c = StdIn.readInt();
15: 16AB R6 ← RA + RB          if (c == 0) break;
16: BC06 mem[R6] ← RC          memory address of a[n]
17: 1BB1 RB ← RB + R1          a[n] = c;
18: C013 goto 13              n++;
                                }

```

read in the data

19

Arrays in TOY

TOY main memory is a giant array.

- Can access memory cell 30 using load and store.
- 8C30 means load mem[30] into register c.
- Goal: access memory cell *i* where *i* is a variable.

...	...
30	0000
31	0001
32	0001
33	0002
34	0003
35	0005
36	0008
37	000D
...	...

TOY memory

Load indirect. [opcode A] ← a variable index

- AC06 means load mem[R6] into register c.

Store indirect. [opcode B] ← a variable index

- BC06 means store contents of register c into mem[R6].

```

for (int i = 0; i < N; i++)
    a[i] = StdIn.readInt();

for (int i = 0; i < N; i++)
    StdOut.println(a[N-i-1]);

```

18

TOY Implementation of Reverse

TOY implementation of reverse.

- Read in a sequence of integers and store in memory 30, 31, 32, ...
- Stop reading if 0000.
- Print sequence in reverse order.

```

19: CB20 if (RB == 0) goto 20      while (n > 0) {
1A: 16AB R6 ← RA + RB              address of a[n]
1B: 2661 R6 ← R6 - R1              address of a[n-1]
1C: AC06 RC ← mem[R6]              c = a[n-1];
1D: 9CFF write RC                  StdOut.println(c);
1E: 2BB1 RB ← RB - R1              n--;
1F: C019 goto 19                    }
20: 0000 halt

```

print in reverse order

20

Unsafe Code at any Speed

Q. What happens if we make array start at 00 instead of 30?

A. Self modifying program; can overflow buffer and run arbitrary code!

```

10: 7101 R1 ← 0001      constant 1
11: 7A00 RA ← 0000      a[]
12: 7B00 RB ← 0000      n

13: 8CFF read RC          while(true) {
14: CC19 if (RC == 0) goto 19   c = StdIn.readInt();
15: 16AB R6 ← RA + RB         if (c == 0) break;
16: BC06 mem[R6] ← RC        address of a[n]
17: 1BB1 RB ← RB + R1        a[n] = c;
18: C013 goto 13             n++;
                                }

```

```

% more crazy8.txt
1 1 1 1 1 1 1 1
1 1 1 1 1 1 1 1
8888 8810
98FF C011

```

21

What Can Happen When We Lose Control (in C or C++)?

Buffer overflow.

- Array `buffer[]` has size 100.
- User might enter 200 characters.
- Might lose control of machine behavior.

```

#include <stdio.h>
int main(void) {
    char buffer[100];
    scanf("%s", buffer);
    printf("%s\n", buffer);
    return 0;
}

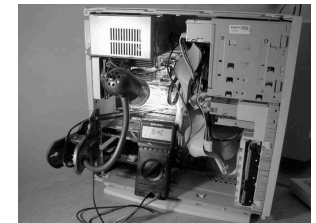
```

unsafe C program

Consequences. Viruses and worms.

Java enforces security.

- Type safety.
- Array bounds checking.
- Not foolproof.



shine 50W bulb at DRAM
[Appel-Govindavajhala '03]

22

Buffer Overflow Example: JPEG of Death

Stuxnet worm. [July 2010]

- Step 1. Natanz centrifuge fuel-refining plant employee plugs in USB flash drive.
- Step 2. Machine is Owned; data becomes code by exploiting Windows buffer overflow.
- Step 3. Uranium enrichment in Iran stalled.



Buffer overflow attacks. Morris worm, Code Red, SQL Slammer, iPhone unlocking, Xbox softmod, GDI+ library for JPEG, ...

Moral.

- Not easy to write error-free software.
- Embrace Java security features.
- Don't try to maintain several copies of the same file.
- Keep your OS patched.

23

Dumping

Q. Work all day to develop operating system. How to save it?

A. Write short program `dump.toy` and run it to dump contents of memory onto tape.

```

00: 7001 R1 ← 0001
01: 7210 R2 ← 0010      i = 10
02: 73FF R3 ← 00FF

03: AA02 RA ← mem[R2]   do {
04: 9AFF write RA       a = mem[i]
05: 1221 R2 ← R2 + R1   print a
06: 2432 R4 ← R3 - R2   i++
07: D403 if (R4 > 0) goto 03 } while (i < 255)
08: 0000 halt

```

dump.toy

25

Booting

Q. How do you get it back?

A. Write short program `boot.toy` and run it to read contents of memory from tape.

```

00: 7001  R1 ← 0001
01: 7210  R2 ← 0010          i = 10
02: 73FF  R3 ← 00FF

                                do {
03: 8AFF  read RA              read a
04: BA02  mem[R2] ← RA        mem[i] = a
05: 1221  R2 ← R2 + R1        i++
06: 2432  R4 ← R3 - R2
07: D403  if (R4 > 0) goto 03 } while (i < 255)
08: 0000  halt
    
```

boot.toy

26

TOY Simulator

TOY Simulator

Goal. Write a program to "simulate" the behavior of the TOY machine.

- ➔ ■ TOY simulator in Java.
- TOY simulator in TOY!

```

public class TOY {
    public static void main(String[] args) {
        int pc = 0x10; // program counter
        int[] R = new int[16]; // registers
        int[] mem = new int[256]; // main memory

        // READ IN .toy FILE

        while (true) {
            // FETCH INSTRUCTION and DECODE
            ...
            // EXECUTE
            ...
        }
    }
}
    
```

```

% java TOY add-stdin.toy
A012 ← standard input
002B ← standard output
A03D ←
    
```

28

TOY Simulator: Fetch

Fetch. Extract destination register of `1cAB` by shifting and masking.

0	0	0	1	1	1	0	0	1	0	1	0	1	0	1	1	inst
1 ₁₆				C ₁₆				A ₁₆				B ₁₆				
0	0	0	0	0	0	0	0	0	0	0	1	1	1	0	0	inst >> 8
0 ₁₆				0 ₁₆				1				C ₁₆				
0	0	0	0	0	0	0	0	0	0	0	0	1	1	1	1	15
0 ₁₆				0 ₁₆				0 ₁₆				F ₁₆				
0	0	0	0	0	0	0	0	0	0	0	0	1	1	0	0	(inst >> 8) & 15
0 ₁₆				0 ₁₆				0				C ₁₆				

```

int inst = mem[pc++]; // fetch and increment
int op = (inst >> 12) & 15; // opcode (bits 12-15)
int d = (inst >> 8) & 15; // dest d (bits 08-11)
int s = (inst >> 4) & 15; // source s (bits 04-07)
int t = (inst >> 0) & 15; // source t (bits 00-03)
int addr = (inst >> 0) & 255; // addr (bits 00-07)
    
```

29

```

if (op == 0) break;          // halt

switch (op) {
case 1: R[d] = R[s] + R[t];  break;
case 2: R[d] = R[s] - R[t];  break;
case 3: R[d] = R[s] & R[t];  break;
case 4: R[d] = R[s] ^ R[t];  break;
case 5: R[d] = R[s] << R[t]; break;
case 6: R[d] = R[s] >> R[t]; break;
case 7: R[d] = addr;        break;
case 8: R[d] = mem[addr];    break;
case 9: mem[addr] = R[d];    break;
case 10: R[d] = mem[R[t]];   break;
case 11: mem[R[t]] = R[d];   break;
case 12: if (R[d] == 0) pc = addr; break;
case 13: if (R[d] > 0) pc = addr; break;
case 14: pc = R[d]; pc = addr; break;
case 15: R[d] = pc; pc = addr; break;
}

```

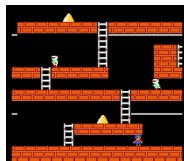
Simulation

Consequences of simulation.

- Test out new machine or microprocessor using simulator.
(cheaper and faster than building actual machine)
- Easy to add new functionality to simulator.
(trace, single-step, breakpoint debugging)
- Reuse software from old machines.

Ancient programs still running on modern computers.

- Ticketron.
- Lode Runner on Apple IIe.



Omitted details.

- Register 0 is always 0.
- reset `R[0]=0` after each fetch-execute step
- Standard input and output.
- if `addr` is `FF` and opcode is load (indirect) then read in data
- if `addr` is `FF` and opcode is store (indirect) then write out data
- TOY registers are 16-bit integers; program counter is 8-bit.
- Java `int` is 32-bit; Java `short` is 16-bit
- use casts and bit-whacking

Complete implementation. See `toy.java` on booksite.