



General Computer Science
Princeton University
Spring 2011

Kevin Wayne

www.princeton.edu/~cos126

Overview

What is COS 126? Broad, but technical, intro to computer science.

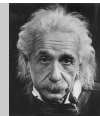
Goals.

- Demystify computer systems.
- Empower you to exploit available technology.
- Build awareness of substantial intellectual underpinnings.

Topics.

- **Programming** in Java.
- Machine architecture.
- Theory of computation.
- **Applications** to science, engineering, and commercial computing.

“Computers are incredibly fast, accurate, and stupid; humans are incredibly slow, inaccurate, and brilliant; together they are powerful beyond imagination.” – Albert Einstein



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The Basics

Lectures. [Kevin Wayne]

- Tuesdays and Thursdays, Frist 302.
- Same lecture at 10am and 11am.
- Office hours: M Th 2-3pm in CS 207. ← or just drop by

Precepts. [Donna Gabai (co-lead) · Keith Vertanen (co-lead) · 11 others]

- Tue+Thu or Wed+Fri.
- Tips on assignments, worked examples, clarify lecture material.

Computing laboratory. [Undergrad lab assistants]

- Sun 5-11pm, Mon-Fri 7-11pm, Sat 2-6pm in Friend 017.
- Help with debugging.

Full details and office hours. See www.princeton.edu/~cos126

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Grades

Course grades. No preset curve or quota.

9 programming assignments. 40%.

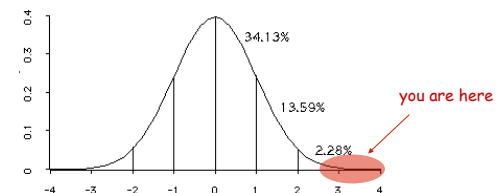
2 exams. 50%.

Final programming project. 10%.

Extra credit and staff discretion. Adjust borderline cases.

participation helps, frequent absences hurts

Check grades. Blackboard. [www.blackboard.princeton.edu]

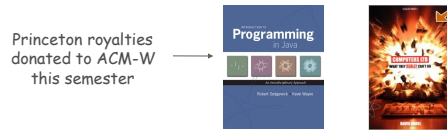


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Course website. [www.princeton.edu/~cos126]

- Programming assignments and checklists.
- Submit assignments.
- Lecture slides. ← print before lecture; annotate during lecture
- Exam archive.

Required readings. Sedgewick and Wayne. *Intro to Programming in Java: An Interdisciplinary Approach*. [Labyrinth] ← skim before lecture; read thoroughly afterwards



Recommended readings. Harel. *What computers can't do*. [Labyrinth]

Desiderata.

- Address an important scientific or commercial problem.
- Illustrate the importance of a fundamental CS concept.
- You solve problem from scratch!

Due. Mondays 11pm via Web submission.

Computing equipment.

- Your laptop. [OS X, Windows, Linux, iPhone, ...]
- OIT desktop. [Friend 016 and 017 labs]

What's Ahead?

Lecture 2. Intro to Java.

Precept 1. Meets today/tomorrow.

Precept 2. Meets Thu/Fri.

Not registered? Go to any precept now; officially register ASAP.

Change precepts? Use SCORE.

← for enrollment problems, see Colleen Kenny-McGinley in CS 210

Assignment 0. Due Monday, 11pm.

- Read Sections 1.1 and 1.2 in textbook.
- Install Java programming environment + a few exercises.
- Lots of help available, don't be bashful.

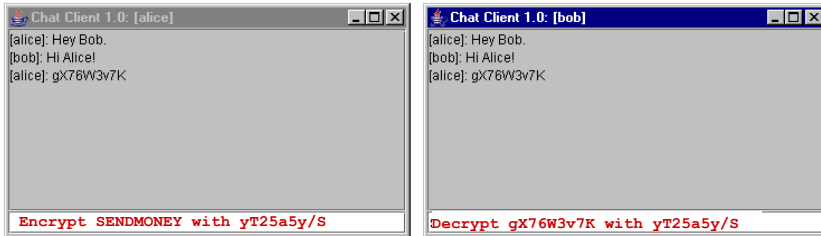
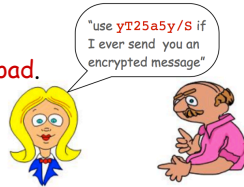
END OF ADMINISTRATIVE STUFF

0. Prologue: A Simple Machine

Secure Chat

Alice wants to send a secret message to Bob?

- Sometime in the past, they exchange a **one-time pad**.
- Alice uses the pad to encrypt the message.
- Bob uses the same pad to decrypt the message.



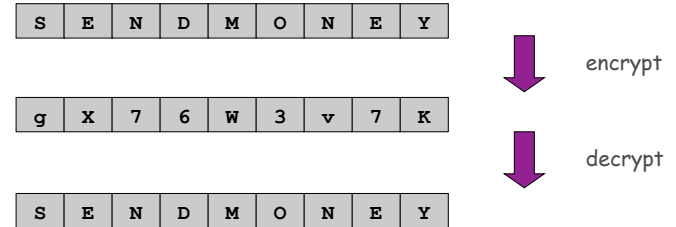
Key point. Without the pad, Eve cannot understand the message.



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Encryption Machine

Goal. Design a machine to encrypt and decrypt data.



Enigma encryption machine.

- "Unbreakable" German code during WWII.
- Broken by Turing bombe.
- One of first uses of computers.
- Helped win Battle of Atlantic by locating U-boats.



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A Digital World

Data is a sequence of bits. [bit = 0 or 1]

- **Text.**
- Programs, executables.
- Documents, pictures, sounds, movies, ...

File formats. txt, pdf, java, exe, docx, pptx, jpeg, mp3, divx, ...



computer with a lens



computer with earbuds



computer with a radio

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A Digital World

Data is a sequence of bits. [bit = 0 or 1]

- **Text.**
- Programs, executables.
- Documents, pictures, sounds, movies, ...

Base64 encoding. Use 6 bits to represent each alphanumeric symbol.

Binary Char	Binary Char	Binary Char	Binary Char	Binary Char	Binary Char
000000 A	001011 L	010110 W	100001 h	101100 s	110111 3
000001 B	001100 M	010111 X	100010 i	101101 t	111000 4
000010 C	001101 N	011000 Y	100011 j	101110 u	111001 5
000011 D	001110 O	011001 Z	100100 k	101111 v	111010 6
000100 E	001111 P	011010 a	100101 l	110000 w	111011 7
000101 F	010000 Q	011011 b	100110 m	110001 x	111100 8
000110 G	010001 R	011100 c	100111 n	110010 y	111101 9
000111 H	010010 S	011101 d	101000 o	110011 z	111110 +
001000 I	010011 T	011110 e	101001 p	110100 0	111111 /
001001 J	010100 U	011111 f	101010 q	110101 1	
001010 K	010101 V	100000 g	101011 r	110110 2	

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One-Time Pad Encryption

Encryption.

- Convert text message to N bits.

Base64 Encoding

char	dec	binary
A	0	000000
B	1	000001
...
M	12	001100
...

S	E	N	D	M	O	N	E	Y
010010	000100	001101	000011	001100	001110	001101	000100	011000

message

base64

One-Time Pad Encryption

Encryption.

- Convert text message to N bits.
- Generate N random bits (one-time pad).

S	E	N	D	M	O	N	E	Y
010010	000100	001101	000011	001100	001110	001101	000100	011000
110010	010011	110110	111001	011010	111001	100010	111111	010010

message

base64

random bits

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One-Time Pad Encryption

Encryption.

- Convert text message to N bits.
- Generate N random bits (one-time pad).
- Take bitwise XOR of two bitstrings.

XOR Truth Table

x	y	$x \oplus y$
0	0	0
0	1	1
1	0	1
1	1	0

sum corresponding pair of bits: 1 if sum is odd, 0 if even

S	E	N	D	M	O	N	E	Y
010010	000100	001101	000011	001100	001110	001101	000100	011000
110010	010011	110110	111001	011010	111001	100010	111111	010010
100000	010111	111011	111010	010110	110111	101111	111011	001010

message

base64

random bits

XOR

$0 \oplus 1 = 1$

One-Time Pad Encryption

Encryption.

- Convert text message to N bits.
- Generate N random bits (one-time pad).
- Take bitwise XOR of two bitstrings.
- Convert binary back into text.

Base64 Encoding

char	dec	binary
A	0	000000
B	1	000001
...
w	22	010110
...

S	E	N	D	M	O	N	E	Y
010010	000100	001101	000011	001100	001110	001101	000100	011000
110010	010011	110110	111001	011010	111001	100010	111111	010010
100000	010111	111011	111010	010110	110111	101111	111011	001010
g	x	7	6	w	3	v	7	K

message

base64

random bits

XOR

encrypted

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One-Time Pad Decryption

Decryption.

- Convert encrypted message to binary.

g	x	7	6	w	3	v	7	κ
---	---	---	---	---	---	---	---	---

encrypted

One-Time Pad Decryption

Decryption.

- Convert encrypted message to binary.

Base64 Encoding

char	dec	binary
A	0	000000
B	1	000001
...
W	22	010110
...

g	x	7	6	w	3	v	7	κ
---	---	---	---	---	---	---	---	---

encrypted

100000	010111	111011	111010	010110	110111	101111	111011	001010
--------	--------	--------	--------	--------	--------	--------	--------	--------

base64

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One-Time Pad Decryption

Decryption.

- Convert encrypted message to binary.
- Use **same** N random bits (one-time pad).

g	x	7	6	w	3	v	7	κ
---	---	---	---	---	---	---	---	---

encrypted

100000	010111	111011	111010	010110	110111	101111	111011	001010
--------	--------	--------	--------	--------	--------	--------	--------	--------

base64

110010	010011	110110	111001	011010	111001	100010	111111	010010
--------	--------	--------	--------	--------	--------	--------	--------	--------

random bits

One-Time Pad Decryption

Decryption.

- Convert encrypted message to binary.
- Use same N random bits (one-time pad).
- Take bitwise XOR of two bitstrings.

XOR Truth Table

x	y	x ^ y
0	0	0
0	1	1
1	0	1
1	1	0

g	x	7	6	w	3	v	7	κ
---	---	---	---	---	---	---	---	---

encrypted

100000	010111	111011	111010	010110	110111	101111	111011	001010
--------	--------	--------	--------	--------	--------	--------	--------	--------

base64

110010	010011	110110	111001	011010	111001	100010	111111	010010
--------	--------	--------	--------	--------	--------	--------	--------	--------

random bits

010010	000100	001101	000011	001100	001110	001101	000100	011000
--------	--------	--------	--------	--------	--------	--------	--------	--------

XOR

1 ^ 1 = 0

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One-Time Pad Decryption

Decryption.

- Convert encrypted message to binary.
- Use same N random bits (one-time pad).
- Take bitwise XOR of two bitstrings.
- Convert back into text.

Base64 Encoding

char	dec	binary
A	0	000000
B	1	000001
...
M	12	001100
...

g	X	7	6	W	3	v	7	K	encrypted
100000	010111	111011	111010	010110	110111	101111	111011	001010	base64
110010	010011	110110	111001	011010	111001	100010	111111	010010	random bits
010010	000100	001101	000011	001100	001110	001101	000100	011000	XOR
S	E	N	D	M	O	N	E	Y	message

Why Does It Work?

Crucial property. Decrypted message = original message.

Notation	Meaning
a	original message bit
b	one-time pad bit
^	XOR operator
a ^ b	encrypted message bit
(a ^ b) ^ b	decrypted message bit

Why is crucial property true?

- Use properties of XOR.
 - $(a \oplus b) \oplus b = a \oplus (b \oplus b) = a \oplus 0 = a$
- ↙ associativity of ^ ↙ always 0 ↙ identity

XOR Truth Table

x	y	x ^ y
0	0	0
0	1	1
1	0	1
1	1	0

One-Time Pad Decryption (with the wrong pad)

Decryption.

- Convert encrypted message to binary.
- Use **wrong** N bits (bogus one-time pad).
- Take bitwise XOR of two bitstrings.
- Convert back into text: **Oops.**

g	X	7	6	W	3	v	7	K	encrypted
100000	010111	111011	111010	010110	110111	101111	111011	001010	base64
101000	011100	110101	101111	010010	111001	100101	101010	001010	wrong bits
001000	001011	001110	010101	000100	001110	001010	010001	000000	XOR
I	L	O	V	E	O	K	R	A	wrong message

Goods and Bads of One-Time Pads

Good.

- Easily computed by hand.
- Very simple encryption/decryption processes.
- Provably unbreakable if bits are truly random. [Shannon, 1940s]

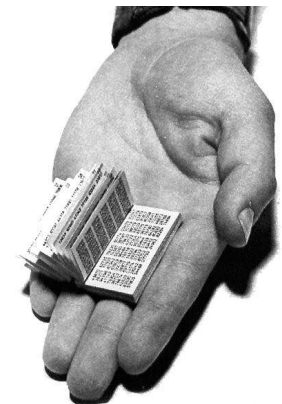
← eavesdropper Eve sees only random bits

Bad.

- Easily breakable if pad is re-used.
- Pad must be as long as the message.
- Truly random bits are very hard to come by.
- Pad must be distributed securely.

"one time" means one time only

← impractical for Web commerce



a Russian one-time pad

Pseudo-Random Bit Generator

Practical middle-ground.

- Let's make a "random" bit generator gadget.
- Alice and Bob each get identical small gadgets.

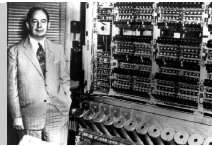
← instead of identical large one-time pads

How to make small gadget that produces pseudo-random numbers.

- Enigma.
- Linear feedback shift register.
- Linear congruential generator.
- Blum-Blum-Shub generator.
- ...

"Anyone who considers arithmetical methods of producing random digits is, of course, in a state of sin."

– Jon von Neumann (left)
– ENIAC (right)

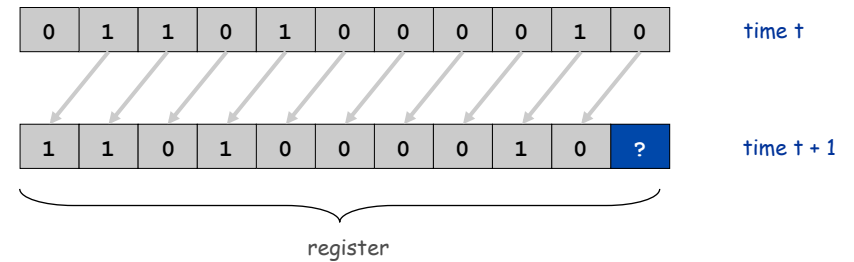


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Shift Register

Shift register terminology.

- Bit: 0 or 1.
- Cell: storage element that holds one bit.
- Register: sequence of cells.
- Seed: initial sequence of bits.
- Shift register: when clock ticks, bits propagate one position to left.



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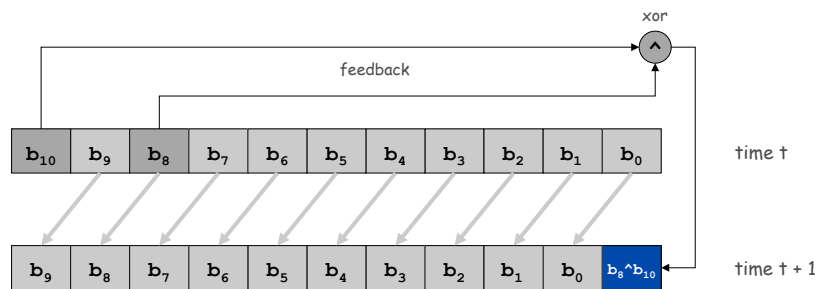
Linear Feedback Shift Register (LFSR)

{8, 10} linear feedback shift register.

- Shift register with 11 cells.
- Bit b_0 is XOR of previous bits b_8 and b_{10} .
- Pseudo-random bit = b_0 .



LFSR demo



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Random Numbers

Q. Are these 2000 numbers random? If not, what is the pattern?

```
110010010011110110111001011011011100110001011111101001000010011010010111100110010011111
10111000001010110001000011101010011010000111001001100110111111010100001000010001010
0101010001100001011110001001001101011011100010100110111001110101110010001001110101
011101000010100100010001101010101110000001011000010011100010111011010010101100110000
11111100110000111110001100001011100111010011101001110100110110110101010101010000
000000100000001010000010001000010101010010000010100000110010001101101011101010100
010100001010001001001010110101000011000010011100101110011001011101100100101011011
000101011100100010111010010010011011000111101110110010101011100000010011000010111
10010010001110101010101000110011101101010010110000110011100111110111100001010
0110010001111101011000100011100101011011000011010110011100011110110110001011011010
011010100111000011100110011011111101000000100100000101101000100110010101111100001
000011001010011110001110001101101101101101010110100000101110000110110110110100011
01100101110111001010011100000110110001010111011000101010110000001100000011110
1001100010011110101100010001010101010011000000111100001100011001111011111001010000
1110001001101101111011000100101110101100101000111000101100111100111000011110
11001100101111100100000011010000101001001100110111011101010001000000101010000
10000010010100010100010100111010001101001010100110011001111111100000000110000000
11110000011001100011111101100000010111000010010110011100111001111001111001111001
011001111011110001010001011000101100101001011100110010110111100110100011110010110
0011100111011011101011010010001001101011111000100000101010001110000101101100100110
111101111010010100100010111101101101000101010010100000110001000111101011001000001
11010001100100101111011001000101110101001000001010100111010011101011101000100
01001010101011000000001110000010110000110110001101100010101111000010001010111010
```

A. No. This is output of {8, 10} LFSR with seed 01101000010!

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LFSR Encryption

LFSR encryption.

- Convert text message to N bits.
- Initialize LFSR with small seed.
- Generate N bits **with LFSR**.
- Take bitwise XOR of two bitstrings.
- Convert binary back into text.

Base64 Encoding

char	dec	binary
A	0	000000
B	1	000001
...
w	22	010110
...

S	E	N	D	M	O	N	E	Y	message
010010	000100	001101	000011	001100	001110	001101	000100	011000	base64
110010	010011	110110	111001	011010	111001	100010	111111	010010	LFSR bits
100000	010111	111011	111010	010110	110111	101111	111011	001010	XOR
g	x	7	6	w	3	v	7	K	encrypted

LFSR Decryption

LFSR Decryption.

- Convert encrypted message to binary.
- Initialize identical LFSR with same seed.
- Generate N bits **with LFSR**.
- Take bitwise XOR of two bitstrings.
- Convert binary back into text.

Base64 Encoding

char	dec	binary
A	0	000000
B	1	000001
...
M	12	001100
...

g	x	7	6	w	3	v	7	K	encrypted
100000	010111	111011	111010	010110	110111	101111	111011	001010	base64
110010	010011	110110	111001	011010	111001	100010	111111	010010	LFSR bits
010010	000100	001101	000011	001100	001110	001101	000100	011000	XOR
S	E	N	D	M	O	N	E	Y	message

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Goods and Bads of LFSR Encryption

Goods.

- Easily computed with simple machine.
- Very simple encryption / decryption process.
- Scalable: 20 cells for 1 million bits; 30 cells for 1 billion bits.
[but need theory of finite groups to know where to put taps]



a commercially available LFSR

Bads.

- Still need secure, independent way to distribute LFSR seed.
- The bits are not truly random.
[bits in our 11-bit LFSR cycle after $2^{11} - 1 = 2047$ steps]
- Experts have cracked LFSR.
[more complicated machines needed]

Other LFSR Applications

What else can we do with a LFSR?

- DVD encryption with CSS.
- DVD decryption with DeCSS!
- Subroutine in military cryptosystems.

```

/*  efdtt.c  Author: Charles M. Hannum <root@ihack.net>  */
/*  Usage is: cat title-key scrambled.vob | efdtt >clear.vob  */

#define m(i) (x[i]^s[i+84])<<

unsigned char x[5], y, s[2048]; main(
n){for( read(0,x,5 );read(0,s ,n=2048
); write(1 ,s,n )if(s
[y=s [13]%8+20] /16%4 ==1 )int
i=( 1)17 ^256 +m(0) 8,k =m(2)
0,j= m(4) 17^ m(3) 9*k* 2-k%8
^8,a =0,c =26;for (s[y] --=16;
--c;j *=2)a= a*2^i% 1,i=i /2^j%1
<<24;for(j= 127; ++j<n;c=c>
y)
c
+=y^i^i/8^i>>4^i>>12,
i=i>>8^y<<17,a^=a>>14,y=a^a*8^a<<6,a=a
>>8^y<<9,k=s[j],k =*7Wo-'G _216"[k
&7]+2^"cs3sfw6v;*k+>/n." [k>>4]*2^k*257/
8,s[j]=k^(k&k*2&34)*6^c+-y
);}
    
```

<http://www.cs.cmu.edu/~dst/DeCSS/Gallery>

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Important properties.

- Built from simple components.
- Scales to handle huge problems.
- Requires a deep understanding to use effectively.

Basic Component	LFSR	Computer
control	start, stop, load	same
clock	regular pulse	2.8 GHz pulse
memory	11 bits	1 GB
input	seed	sequence of bits
computation	shift, XOR	logic, arithmetic, ...
output	pseudo-random bits	Sequence of bits

Critical difference. General purpose machine can be programmed to simulate ANY abstract machine.

Programming. Can write a Java program to simulate the operations of any abstract machine.

- Basis for theoretical understanding of computation. [stay tuned]
- Basis for bootstrapping real machines into existence. [stay tuned]

Stay tuned. See Assignment 5.

```
public class LFSR {
    private int seed[];
    private final int tap;
    private final int N;

    public LFSR(String seed, int tap) { ... }

    public int step() { ... }

    public static void main(String[] args) {
        LFSR lfsr = new LFSR("01101000010", 8);
        for (int i = 0; i < 2000; i++)
            StdOut.println(lfsr.step());
    }
}
```

```
% java LFSR
11001001001111011011100101101
01110011000101111110100100001
00110100101111001100100111...
```

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A Profound Question

Q. What is a random number?

LFSR does not produce random numbers.

- It is a very simple deterministic machine.
- But not obvious how to distinguish the bits it produces from random.

Q. Are random processes found in nature?

- Motion of cosmic rays or subatomic particles?
- Mutations in DNA?

Q. Or, is the natural world a (not-so-simple) deterministic machine?

*"God does not play dice."
– Albert Einstein*



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