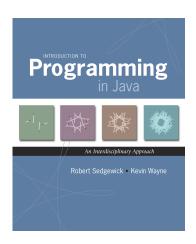
2.3 Recursion



 $Introduction \ to \ Programming \ in \ Java: \ An \ Interdisciplinary \ Approach \\ \qquad \cdot \quad Robert \ Sedgewick \ and \ Kevin \ Wayne \\ \qquad \cdot \quad Copyright \ @ \ 2008 \\ \qquad \cdot \quad * \ *$

Greatest Common Divisor

Gcd. Find largest integer that evenly divides into p and q.

Ex. gcd(4032, 1272) = 24.

$$4032 = 2^6 \times 3^2 \times 7^1$$

 $1272 = 2^3 \times 3^1 \times 53^1$
 $9cd = 2^3 \times 3^1 = 24$

Applications.

- Simplify fractions: 1272/4032 = 53/168.
- RSA cryptosystem.

Overview

What is recursion? When one function calls itself directly or indirectly.

Why learn recursion?

- New mode of thinking.
- Powerful programming paradigm.

Many computations are naturally self-referential.

- Mergesort, FFT, gcd.
- Linked data structures.
- A folder contains files and other folders.

Closely related to mathematical induction.





Reproductive Parts M. C. Escher, 1948

Greatest Common Divisor

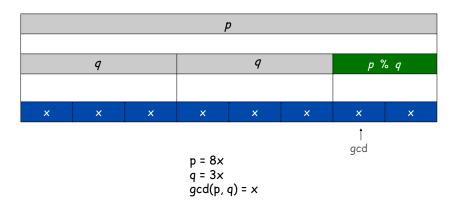
Gcd. Find largest integer that evenly divides into p and q.

Euclid's algorithm. [Euclid 300 BCE]

$$\gcd(p,q) = \begin{cases} p & \text{if } q = 0 \\ \gcd(q, p \% q) & \text{otherwise} \end{cases} \quad \begin{array}{l} \longleftarrow \text{ base case} \\ \longleftarrow \text{ reduction step,} \\ \text{converges to base case} \\ \end{array}$$

Gcd. Find largest integer d that evenly divides into p and q.

$$\gcd(p,q) = \begin{cases} p & \text{if } q = 0 \\ \gcd(q,p \% q) & \text{otherwise} \end{cases} \quad \begin{array}{l} \longleftarrow & \text{base case} \\ \longleftarrow & \text{reduction step,} \\ \text{converges to base case} \\ \end{array}$$



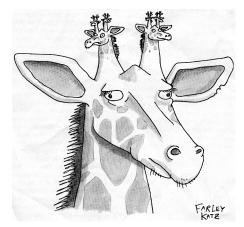
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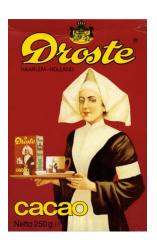
Java implementation.



Recursive Graphics



New Yorker Magazine, August 11, 2008





Htree in Java

```
public class Htree {
   public static void draw(int n, double sz, double x, double y) {
      if (n == 0) return;
      double x0 = x - sz/2, x1 = x + sz/2;
      double y0 = y - sz/2, y1 = y + sz/2;
                                       draw the H, centered on (x, y)
      StdDraw.line(x0, y, x1, y);
      StdDraw.line(x0, y0, x0, y1);
      StdDraw.line(x1, y0, x1, y1);
                                        ← recursively draw 4 half-size Hs
      draw(n-1, sz/2, x0, y0);
      draw(n-1, sz/2, x0, y1);
      draw(n-1, sz/2, x1, y0);
      draw(n-1, sz/2, x1, y1);
                                                       \phi(x_0, y_1)
   public static void main(String[] args) {
      int n = Integer.parseInt(args[0]);
      draw(n, .5, .5, .5);
```

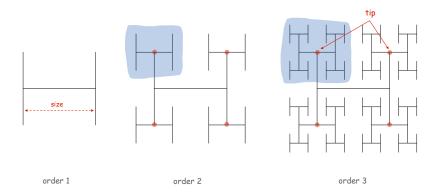
Htree

H-tree of order n.

and half the size

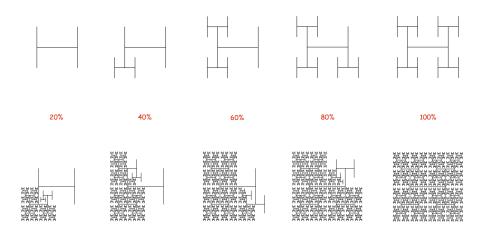
Draw an H.

• Recursively draw 4 H-trees of order n-1, one connected to each tip.



Animated H-tree

Animated H-tree. Pause for 1 second after drawing each H.



n .

Towers of Hanoi



http://en.wikipedia.org/wiki/Image:Hanoiklein.jpg

Towers of Hanoi: Recursive Solution



Move n-1 smallest discs right.



Move n-1 smallest discs right.



Move largest disc left.

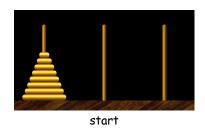
cyclic wrap-around

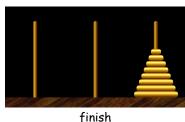


Towers of Hanoi

Move all the discs from the leftmost peg to the rightmost one.

- Only one disc may be moved at a time.
- A disc can be placed either on empty peg or on top of a larger disc.









Edouard Lucas (1883)

Towers of Hanoi Legend

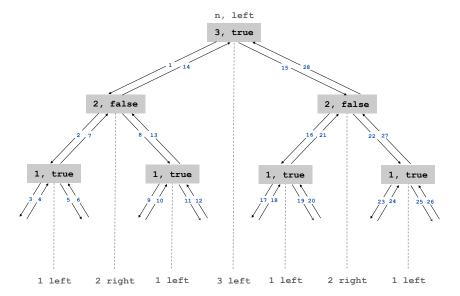
- Q. Is world going to end (according to legend)?
- 64 golden discs on 3 diamond pegs.
- World ends when certain group of monks accomplish task.
- Q. Will computer algorithms help?

moves (n, true): move discs 1 to n one pole to the left moves (n, false): move discs 1 to n one pole to the right

% java TowersOfHanoi 4 % java TowersOfHanoi 3 1 left 1 right 2 left 2 right 1 left 1 right 3 left 3 right 1 left 1 right 2 right 2 left 1 left _1 right 4 left √1 right 2 left →1 right 3 right every other move is smallest disc 1 right 2 left 1 right subdivisions of ruler

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Towers of Hanoi: Recursion Tree



Towers of Hanoi: Properties of Solution

Remarkable properties of recursive solution.

- Takes 2ⁿ 1 moves to solve n disc problem.
- Sequence of discs is same as subdivisions of ruler.
- Every other move involves smallest disc.

Recursive algorithm yields non-recursive solution!

- Alternate between two moves:
- to lett it n is odd
- move smallest disc to right if n is even
- make only legal move not involving smallest disc

Recursive algorithm may reveal fate of world.

- Takes 585 billion years for n = 64 (at rate of 1 disc per second).
- Reassuring fact: any solution takes at least this long!

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Divide-and-Conquer

Divide-and-conquer paradigm.

- Break up problem into smaller subproblems of same structure.
- Solve subproblems recursively using same method.
- Combine results to produce solution to original problem.

Divide et impera. Veni, vidi, vici. - Julius Caesar

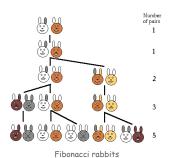
Many important problems succumb to divide-and-conquer.

- FFT for signal processing.
- Parsers for programming languages.
- Multigrid methods for solving PDEs.
- Quicksort and mergesort for sorting.
- Hilbert curve for domain decomposition.
- Quad-tree for efficient N-body simulation.
- Midpoint displacement method for fractional Brownian motion.

Fibonacci Numbers

Fibonacci numbers. 0, 1, 1, 2, 3, 5, 8, 13, 21, 34, ...

$$F(n) = \begin{cases} 0 & \text{if } n = 0\\ 1 & \text{if } n = 1\\ F(n-1) + F(n-2) & \text{otherwise} \end{cases}$$





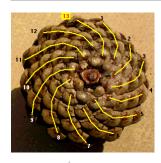
L. P. Fibonacci (1170 - 1250)

Fibonacci Numbers

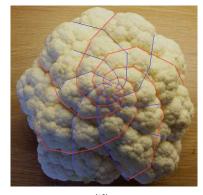
Fibonacci Numbers and Nature

Fibonacci numbers. 0, 1, 1, 2, 3, 5, 8, 13, 21, 34, ...

$$F(n) = \begin{cases} 0 & \text{if } n = 0\\ 1 & \text{if } n = 1\\ F(n-1) + F(n-2) & \text{otherwise} \end{cases}$$



pinecone



cauliflower

23

21

Fibonacci numbers. 0, 1, 1, 2, 3, 5, 8, 13, 21, 34, ...

```
if n = 0
                          if n = 1
F(n) =
       F(n-1) + F(n-2) otherwise
```

A natural for recursion?

```
public static long F(int n) {
  if (n == 0) return 0;
  if (n == 1) return 1;
  return F(n-1) + F(n-2);
}
```

Recursion Challenge 2 (easy and also important)

Q. Is this an efficient way to compute F(50)?

```
public static long(int n) {
   long[] F = new long[n+1];
   F[0] = 0; F[1] = 1;
   for (int i = 2; i \le n; i++)
      F[i] = F[i-1] + F[i-2];
   return F[n];
}
```

FYI: classic math $= |\phi^n/\sqrt{5}|$ ϕ = golden ratio ≈ 1.618 25

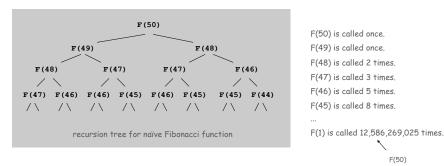
A. Yes. This code does it with 50 additions. Lesson. Don't use recursion to engage in exponential waste.

Context. This is a special case of an important programming technique known as dynamic programming (stay tuned).

Q. Is this an efficient way to compute F(50)?

```
public static long F(int n) {
  if (n == 0) return 0;
  if (n == 1) return 1;
  return F(n-1) + F(n-2);
```

A. No, no, no! This code is spectacularly inefficient.



Summary

How to write simple recursive programs?

- Base case, reduction step.
- Trace the execution of a recursive program.
- Use pictures.

Towers of Hanoi by W. A. Schloss.

F(50)

Why learn recursion?

- New mode of thinking.
- Powerful programming tool.

Divide-and-conquer. Elegant solution to many important problems.

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