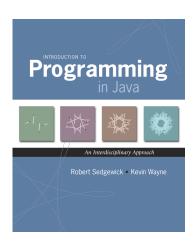
Symbol Table

4.4 Symbol Tables



Introduction to Programming in Java: An Interdisciplinary Approach · Robert Sedgewick and Kevin Wayne · Copyright © 2008 · April 8, 2009 9:51 AM

Symbol Table Applications

Application	Purpose	Key	Value
phone book	look up phone number	name	phone number
bank	process transaction	account number	transaction details
file share	find song to download	name of song	computer ID
file system	find file on disk	filename	location on disk
dictionary	look up word	word	definition
web search	find relevant documents	keyword	list of documents
book index	find relevant pages	keyword	list of pages
web cache	download	filename	file contents
genomics	find markers	DNA string	known positions
DNS	find IP address given URL	URL	IP address
reverse DNS	find URL given IP address	IP address	URL
compiler	find properties of variable	variable name	value and type
routing table	route Internet packets	destination	best route

Symbol table. Key-value pair abstraction.

- Insert a key with specified value.
- Given a key, search for the corresponding value.

Ex. [DNS lookup]

- Insert URL with specified IP address.
- Given URL, find corresponding IP address.

URL	IP address
www.cs.princeton.edu	128.112.136.11
www.princeton.edu	128.112.128.15
www.yale.edu	130.132.143.21
www.harvard.edu	128.103.060.55
www.simpsons.com	209.052.165.60
key	value

Symbol Table API

public class ST<Key extends Comparable<Key>, Value>

```
ST() create a symbol table

void put (Key key, Value v) put key-value pair into the table

Value get (Key key) return value paired with key, null if key not in table

boolean contains (Key key) is there a value paired with key?
```

3

Datasets

Frequency counter. [e.g., web traffic analysis, linguistic analysis]

- Read in a key.
- If key is in symbol table, increment counter by one;
 If key is not in symbol table, insert it with count = 1.

Zipf's Law

Linguistic analysis. Compute word frequencies in a piece of text.

```
% java Freq < mobydick.txt
4583 a
2 aback
2 abaft
3 abandon
7 abandoned
1 abandonedly
2 abandonment
2 abased
1 abasement
2 abashed
1 abate
...</pre>
```

```
% java Freq < mobydick.txt | sort -rn

13967 the

6415 of

6247 and

4583 a

4508 to

4037 in

2911 that

2481 his

2370 it

1940 i

1793 but
```

Zipf's law. In natural language, frequency of i^{th} most common word is inversely proportional to i.

e.g., most frequent word occurs about twice as often as second most frequent one

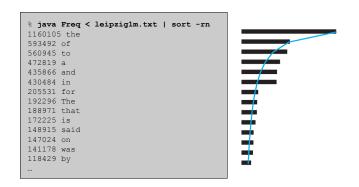
Linguistic analysis. Compute word frequencies in a piece of text.

File	Description	Words	Distinct
mobydick.txt	Melville's Moby Dick	210,028	16,834
leipzig100k.txt	100K random sentences	2,121,054	144,256
leipzig200k.txt	200K random sentences	4,238,435	215,515
leipzig1m.txt	1M random sentences	21,191,455	534,580

Reference: Wortschatz corpus, Univesität Leipzig
http://corpora.informatik.uni-leipzig.de

Zipf's Law

Linguistic analysis. Compute word frequencies in a piece of text.



Zipf's law. In natural language, frequency of i^{th} most common word is inversely proportional to i.

e.g., most frequent word occurs about twice as often as second most frequent one

Symbol Table: Elementary Implementations

Unordered array.

- Put: add key to the end (if not already there).
- Get: scan through all keys to find desired value.



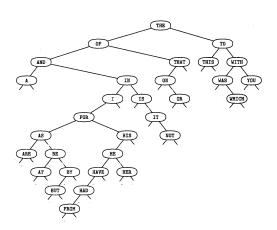
Ordered array.

- Put: find insertion point, and shift all larger keys right.
- Get: binary search to find desired key.



insert 28

Binary Search Trees



Reference: Knuth, The Art of Computer Programming

Symbol Table: Implementations Cost Summary

Unordered array. Hopelessly slow for large inputs.

Ordered array. Acceptable if many more searches than inserts; too slow if many inserts.

Running Time Frequency Count

implementation	get	put	Moby	100K	200K	1M
unordered array	N	N	170 sec	4.1 hr	-	-
ordered array	log N	N	5.8 sec	5.8 min	15 min	2.1 hr

Challenge. Make all ops logarithmic.

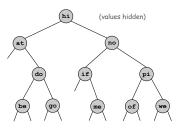
Binary Search Trees

Def. A binary search tree is a binary tree in symmetric order.

Binary tree is either:

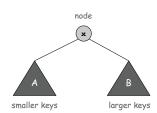
- Empty.
- A key-value pair and two binary trees.

we suppress values from figures



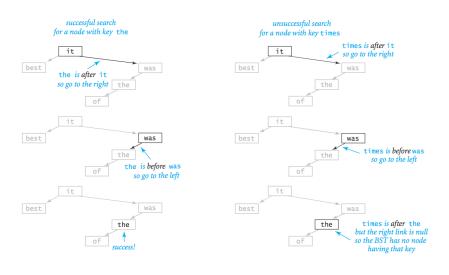
Symmetric order.

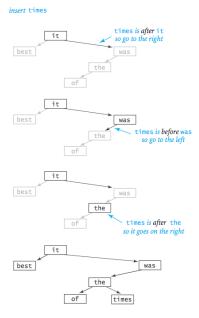
- Keys in left subtree are smaller than parent.
- Keys in right subtree are larger than parent.



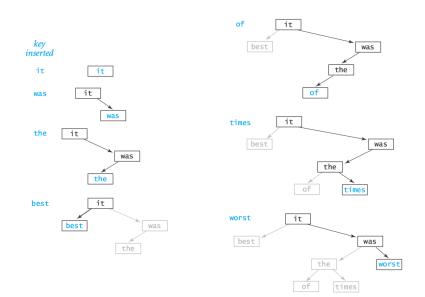
BST Search BST Insert

13





BST Construction



Binary Search Tree: Java Implementation

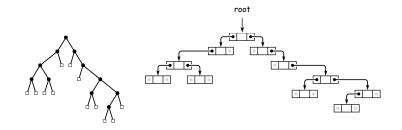
To implement: use two links per Node.

A Node is comprised of:

- A key.
- A value.
- A reference to the left subtree.
- A reference to the right subtree.

```
private class Node {
   private Key key;
   private Value val;
   private Node left;
   private Node right;
}
```

14



BST: Skeleton BST: Get

BST. Allow generic keys and values. requires Key to provide compareto() method; see book for details

```
public class BST<Key extends Comparable<Key>, Value> {
   private Node root; // root of the BST

private class Node {
   private Key key;
   private Value val;
   private Node left, right;

   private Node (Key key, Value val) {
      this.key = key;
      this.val = val;
   }
}

public void put(Key key, Value val) { ... }
   public Value get(Key key) { ... }
   public boolean contains(Key key) { ... }
```

17

BST: Put

Put. Associate val with key.

- Search, then insert.
- Concise (but tricky) recursive code.

```
public void put(Key key, Value val) {
   root = insert(root, key, val);
}

private Node put(Node x, Key key, Value val) {
   if (x == null) return new Node(key, val);
   int cmp = key.compareTo(x.key);
   if (cmp < 0) x.left = put(x.left, key, val);
   else if (cmp > 0) x.right = put(x.right, key, val);
   else x.val = val;
   return x;
   overwrite old value with new value
}
```

Get. Return val corresponding to given key, or null if no such key.

BST Implementation: Practice

Bottom line. Difference between a practical solution and no solution.

	Runnin	g lime	Frequency Count			
implementation	get	put	Moby	100K	200K	1M
unordered array	N	N	170 sec	4.1 hr	-	-
ordered array	log N	N	5.8 sec	5.8 min	15 min	2.1 hr
BST	?	?	.95 sec	7.1 sec	14 sec	69 sec

BST: Analysis

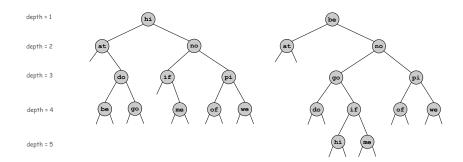
Running time per put/get.

- There are many BSTs that correspond to same set of keys.
- Cost is proportional to depth of node.

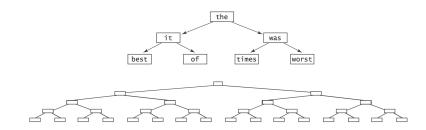
number of nodes on path from root to node

21

23

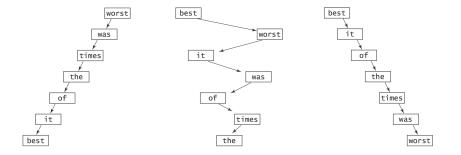


Best case. If tree is perfectly balanced, depth is at most $\lg N$.



BST: Analysis

Worst case. If tree is unbalanced, depth is N.

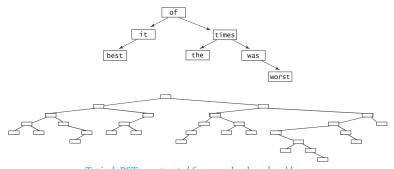


BST: Analysis

22

Average case. If keys are inserted in random order, average depth is $2 \ln N$.

requires proof (see COS 226)



Typical BSTs constructed from randomly ordered keys

Symbol Table: Implementations Cost Summary

BST. Logarithmic time ops if keys inserted in random order.

Running Time	Frequency Count

implementation	get	put	Moby	100K	200K	1M
unordered array	N	N	170 sec	4.1 hr	-	-
ordered array	log N	N	5.8 sec	5.8 min	15 min	2.1 hr
BST	log N †	log N †	.95 sec	7.1 sec	14 sec	69 sec

† assumes keys inserted in random order

Q. Can we guarantee logarithmic performance?

Red-Black Tree

Red-black tree. A clever BST variant that guarantees depth $\leq 2 \lg N$.

see COS 226

Running	Time	Frequency	Count
Running	THIC	1 1 equency	Courti

implementation	get	put	Moby	100K	200K	1M
unordered array	N	N	170 sec	4.1 hr	-	-
ordered array	$\log N$	N	5.8 sec	5.8 min	15 min	2.1 hr
BST	$\log N^{\dagger}$	$\log N^{\dagger}$.95 sec	7.1 sec	14 sec	69 sec
red-black	log N	log N	.95 sec	7.0 sec	14 sec	74 sec

† assumes keys inserted in random order

Red-Black Tree

Red-black tree. A clever BST variant that guarantees depth $\leq 2 \lg N.$ $$_{\rm see COS \, 226}$$

Iteration

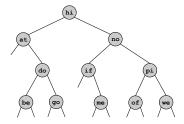
Inorder Traversal

Enhanced For Loop

Enhanced for loop. Enable client to iterate over items in a collection.

Inorder traversal.

- Recursively visit left subtree.
- Visit node.
- Recursively visit right subtree.



inorder: at be do go hi if me no of pi we

```
public inorder() { inorder(root); }

private void inorder(Node x) {
   if (x == null) return;
   inorder(x.left);
   StdOut.println(x.key);
   inorder(x.right);
}
```



ST<String, Integer> st = new ST<String, Integer>(); ... for (String s : st) { StdOut.println(st.get(s) + " " + s); }

3

Enhanced For Loop with BST

BST. Add following code to support enhanced for loop.

see COS 226 for details

```
import java.util.Iterator;
import java.util.NoSuchElementException;
public class BST<Key extends Comparable<Key>, Value> implements Iterable<Key> {
  private Node root;
  private class Node { ... }
  public void put(Key key, Value val) { ... }
  public Value get(Key key)
public boolean contains(Key key)
   public Iterator<Key> iterator() { return new Inorder(); }
    orivate class Inorder implements Iterator<Key> {
       Inorder() { pushLeft(root); }
       public void remove() { throw new UnsupportedOperationException(); }
public boolean hasNext() { return !stack.isEmpty(); }
       public Key next() {
          if (!hasNext()) throw new NoSuchElementException();
           Node x = stack.pop();
           pushLeft(x.right);
            return x.key;
       public void pushLeft(Node x) {
   while (x != null) {
              stack.push(x);
              x = x.left;
```

Symbol Table: Summary

Symbol table. Quintessential database lookup data type.

Choices. Ordered array, unordered array, BST, red-black, hash,

- Different performance characteristics.
- Java libraries: TreeMap, HashMap.

Remark. Better symbol table implementation improves all clients.

31