

General Computer Science  
Princeton University  
Spring 2009

Kevin Wayne

Overview

What is COS 126? Broad, but technical, intro to CS.

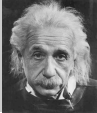
Goals.

- Demystify computer systems.
- Empower you to exploit available technology.
- Build awareness of substantial intellectual underpinnings.

Topics.

- **Programming** in Java.
- Machine architecture.
- Theory of computation.
- **Applications** to science, engineering, and commercial computing.

*“Computers are incredibly fast, accurate, and stupid; Humans are incredibly slow, inaccurate, and brilliant; together they are powerful beyond imagination.”*



The Basics

Lectures. [Kevin Wayne]

- Tuesdays and Thursdays, Frist 302.
- Same lecture at 10am and 11am.

Precepts. [Donna Gabai (lead) · Aditya Bhaskara · Will Clarkson · Rob Dockins · Michael Golightly · Thomas Mason · Chris Park · JP Singh]

- Tue+Thu or Wed+Fri.
- Tips on assignments, worked examples, clarify lecture material.

Friend 016/017 lab. [Undergrad lab assistants]

- Weekdays 7-11pm, some weekend hours.
- Full schedule on Web.

For full details: See [www.princeton.edu/~cos126](http://www.princeton.edu/~cos126)

Grades

Course grades. No preset curve or quota.

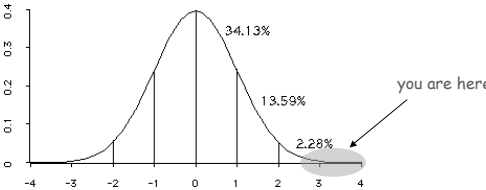
9 programming assignments. 40%.

2 exams. 50%.

Final programming project. 10%.

Extra credit and staff discretion. Adjust borderline cases.

participation helps, frequent absences hurts



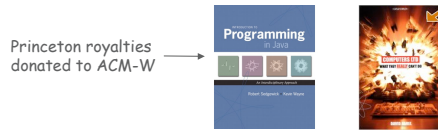
Course website. [www.princeton.edu/~cos126]

- Submit assignments, check grades.
- Programming assignments.
- Lecture slides.

print before lecture;  
annotate during lecture

skim before lecture;  
read thoroughly afterwards

Required readings. Sedgewick and Wayne. *Intro to Programming in Java: An Interdisciplinary Approach*. [Labyrinth Books]



Recommended readings. Harel. *What computers can't do*. [Labyrinth]

Desiderata.

- Address an important scientific or commercial problem.
- Illustrate the importance of a fundamental CS concept.

Examples.

- N-body simulation.
- Pluck a guitar string.
- DNA sequence alignment.
- Estimate Avogadro's number.

} you solve scientific problems from scratch!

Due. Mondays 11pm via Web submission.

Computing equipment.

- Your laptop. [OS X, Windows, Linux, iPhone, ... ]
- OIT desktop. [Friend 016 and 017 labs]

## What's Ahead?

Lecture 2. Intro to Java.

Precept 1. Meets today/tomorrow.

Precept 2. Meets Thu/Fri.

Not registered? Go to any precept now; officially register ASAP.

Change precepts? Use SCORE.

see Donna O'Leary in CS 410  
if the only precept you can attend is closed

Assignment 0. Due Monday, 11pm.

- Read Sections 1.1 and 1.2 in textbook.
- Install Java programming environment + a few exercises.
- Lots of help available, don't be bashful.

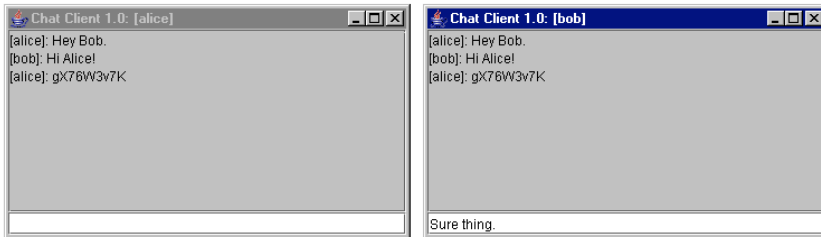
END OF ADMINISTRATIVE STUFF

## 0. Prologue: A Simple Machine

## Secure Chat

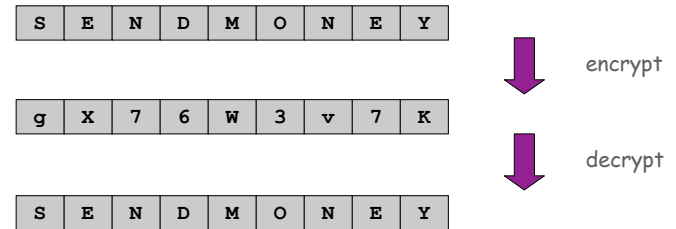
Alice wants to send a secret message to Bob?

- Can you read the secret message gX76W3v7K ?
- But Bob can. How?



## Encryption Machine

Goal. Design a machine to encrypt and decrypt data.



Enigma encryption machine.

- "Unbreakable" German code during WWII.
- Broken by Turing bombe.
- One of first uses of computers.
- Helped win Battle of Atlantic by locating U-boats.



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## A Digital World

Data is a sequence of bits.

- Text.
- Documents, pictures, sounds, movies, ...
- Programs, executables.

File formats. txt, pdf, doc, ppt, jpeg, mp3, divx, java, exe, ...



computer with a lens



computer with earbuds



computer with a radio

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## A Digital World

Data is a sequence of bits.

- Text.
- Documents, pictures, sounds, movies, ...
- Programs, executables.

Base64 encoding. Use 6 bits to represent each alphanumeric symbol.

Binary Char	Binary Char	Binary Char	Binary Char	Binary Char	Binary Char
000000 A	001011 L	010110 W	100001 h	101100 s	110111 3
000001 B	001100 M	010111 X	100010 i	101101 t	111000 4
000010 C	001101 N	011000 Y	100011 j	101110 u	111001 5
000011 D	001110 O	011001 Z	100100 k	101111 v	111010 6
000100 E	001111 P	011010 a	100101 l	110000 w	111011 7
000101 F	010000 Q	011011 b	100110 m	110001 x	111100 8
000110 G	010001 R	011100 c	100111 n	110010 y	111101 9
000111 H	010010 S	011101 d	101000 o	110011 z	111110 +
001000 I	010011 T	011110 e	101001 p	110100 0	111111 /
001001 J	010100 U	011111 f	101010 q	110101 1	
001010 K	010101 V	100000 g	101011 r	110110 2	

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## One-Time Pad Encryption

### Encryption.

- Convert text message to N bits.

0 or 1

Base64 Encoding

char	dec	binary
A	0	000000
B	1	000001
...	...	...
Y	24	011000
...	...	...

S	E	N	D	M	O	N	E	Y
010010	000100	001101	000011	001100	001110	001101	000100	011000

message

base64

## One-Time Pad Encryption

### Encryption.

- Convert text message to N bits.
- Generate N random bits (one-time pad).

S	E	N	D	M	O	N	E	Y
010010	000100	001101	000011	001100	001110	001101	000100	011000
110010	010011	110110	111001	011010	111001	100010	111111	010010

message

base64

random bits

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## One-Time Pad Encryption

### Encryption.

- Convert text message to N bits.
- Generate N random bits (one-time pad).
- Take bitwise XOR of two bitstrings.

sum corresponding pair of bits: 1 if sum is odd, 0 if even

XOR Truth Table

x	y	$x \oplus y$
0	0	0
0	1	1
1	0	1
1	1	0

S	E	N	D	M	O	N	E	Y
010010	000100	001101	000011	001100	001110	001101	000100	011000
110010	010011	110110	111001	011010	111001	100010	111111	010010
100000	010111	111011	111010	010110	110111	101111	111011	001010

message

base64

random bits

XOR

$0 \oplus 1 = 1$

## One-Time Pad Encryption

### Encryption.

- Convert text message to N bits.
- Generate N random bits (one-time pad).
- Take bitwise XOR of two bitstrings.
- Convert binary back into text.

Base64 Encoding

char	dec	binary
A	0	000000
B	1	000001
...	...	...
X	23	010111
...	...	...

S	E	N	D	M	O	N	E	Y
010010	000100	001101	000011	001100	001110	001101	000100	011000
110010	010011	110110	111001	011010	111001	100010	111111	010010
100000	010111	111011	111010	010110	110111	101111	111011	001010
g	x	7	6	w	3	v	7	k

message

base64

random bits

XOR

encrypted

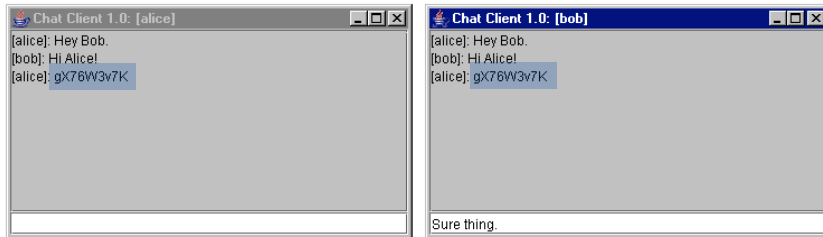
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## Secure Chat

Alice wants to send a secret message to Bob?

- Can you read the secret message gX76W3v7K ?
- But Bob can. How?



## One-Time Pad Decryption

Decryption.

- Convert encrypted message to binary.

g X 7 6 W 3 v 7 K encrypted

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## One-Time Pad Decryption

Decryption.

- Convert encrypted message to binary.

Base64 Encoding

char	dec	binary
A	0	000000
B	1	000001
...	...	...
X	23	010111
...	...	...

g X 7 6 W 3 v 7 K encrypted  
 100000 010111 111011 111010 010110 110111 101111 111011 001010 base64

## One-Time Pad Decryption

Decryption.

- Convert encrypted message to binary.
- Use **same** N random bits (one-time pad).

g X 7 6 W 3 v 7 K encrypted  
 100000 010111 111011 111010 010110 110111 101111 111011 001010 base64  
 110010 010011 110110 111001 011010 111001 100010 111111 010010 random bits

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## One-Time Pad Decryption

### Decryption.

- Convert encrypted message to binary.
- Use same N random bits (one-time pad).
- Take bitwise XOR of two bitstrings.

XOR Truth Table

x	y	x ^ y
0	0	0
0	1	1
1	0	1
1	1	0

g	x	7	6	w	3	v	7	K	
100000	010111	111011	111010	010110	110111	101111	111011	001010	encrypted
110010	010011	110110	111001	011010	111001	100010	111111	010010	base64
010010	000100	001101	000011	001100	001110	001101	000100	011000	random bits
									XOR

$1 \wedge 1 = 0$

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## One-Time Pad Decryption

### Decryption.

- Convert encrypted message to binary.
- Use same N random bits (one-time pad).
- Take bitwise XOR of two bitstrings.
- Convert back into text.

Base64 Encoding

char	dec	binary
A	0	000000
B	1	000001
...	...	...
Y	24	011000
...	...	...

g	x	7	6	w	3	v	7	K	
100000	010111	111011	111010	010110	110111	101111	111011	001010	encrypted
110010	010011	110110	111001	011010	111001	100010	111111	010010	base64
010010	000100	001101	000011	001100	001110	001101	000100	011000	random bits
									XOR
S	E	N	D	M	O	N	E	Y	message

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## Why Does It Work?

Crucial property. Decrypted message = original message.

Notation	Meaning
a	original message bit
b	one-time pad bit
^	XOR operator
a ^ b	encrypted message bit
(a ^ b) ^ b	decrypted message bit

### Why is crucial property true?

- Use properties of XOR.
- $(a \wedge b) \wedge b = a \wedge (b \wedge b) = a \wedge 0 = a$

$\uparrow$  associativity of ^      $\uparrow$  always 0      $\uparrow$  identity

XOR Truth Table

x	y	x ^ y
0	0	0
0	1	1
1	0	1
1	1	0

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## One-Time Pad Decryption (with the wrong pad)

### Decryption.

- Convert encrypted message to binary.
- Use **wrong** N bits (bogus one-time pad).
- Take bitwise XOR of two bitstrings.
- Convert back into text: **Oops.**

g	x	7	6	w	3	v	7	K	
100000	010111	111011	111010	010110	110111	101111	111011	001010	encrypted
101000	011100	110101	101111	010010	111001	100101	101010	001010	base64
001000	001011	001110	010101	000100	001110	001010	010001	000000	wrong bits
									XOR
I	L	O	V	E	O	K	R	A	wrong message

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## Goods and Bads of One-Time Pads

### Good.

- Very simple encryption/decryption processes.
- Provably unbreakable if pad is truly random. [Shannon, 1940s]

↑ eavesdropper Eve sees only random bits

### Bad.

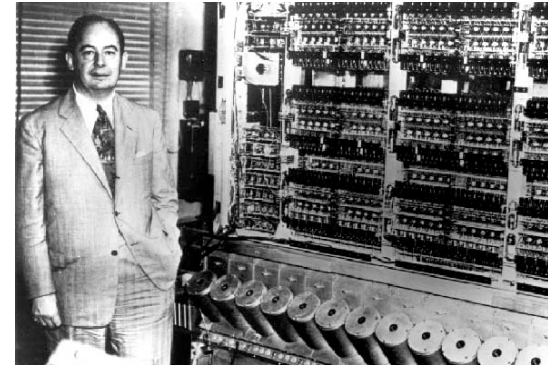
- Easily breakable if pad is re-used.
- Pad must be as long as the message.
- Truly random bits are very hard to come by.
- **Pad must be distributed securely.**

↑ impractical for Web commerce

"one time" means one time only

## Random Numbers

*"Anyone who considers arithmetical methods of producing random digits is, of course, in a state of sin."*



Jon von Neumann (left), ENIAC (right)

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## Pseudo-Random Bit Generator

### Practical middle-ground.

- Let's make a **pseudo**-random bit generator gadget.
- Alice and Bob each get identical small gadgets.

↑ instead of identical large one-time pads

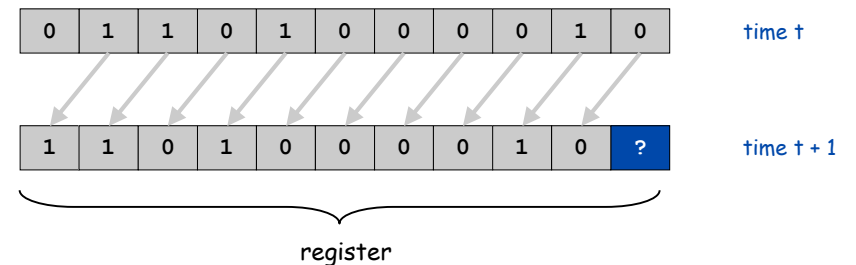
### How to make small gadget that produces "random" numbers.

- **Linear feedback shift register.**
- Linear congruential generator.
- Blum-Blum-Shub generator.
- ...

## Shift Register

### Shift register terminology.

- Bit: 0 or 1.
- Cell: storage element that holds one bit.
- Register: sequence of cells.
- Seed: initial sequence of bits.
- Shift register: when clock ticks, bits propagate one position to left.



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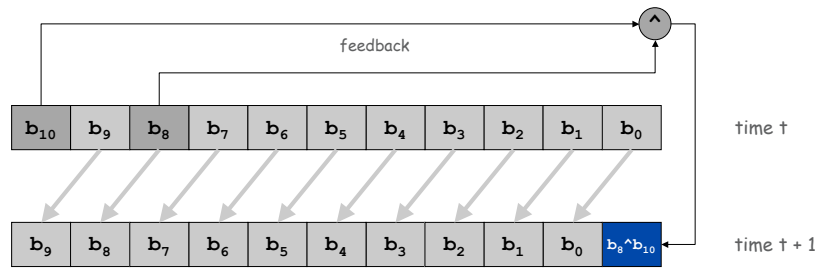
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{8, 10} linear feedback shift register.

- 11 bit shift register.
- New output bit 0 is XOR of previous bits 8 and 10.
- Output bit = bit 0.



LFSR demo



LFSR Challenge 1

Goal. Decrypt/encrypt 1,000 characters. Can we use an 11-bit LFSR?

- A. Yes, no problem.
- B. No, the bits it produces are not truly random.
  - True, but that's beside the point.
- C. No, need a longer LFSR.
  - Only  $2^{11}$  bit patterns for register.
  - "Random" bits cycle after  $2^{11} - 1 = 2047$  steps.

Lesson. LFSR are scalable: 20 cells for 1 million bits; 30 cells for 1 billion. (but need theory of finite groups to know where to put taps)

Q. Are these 2000 numbers random? If not, what is the pattern?

```

1100100100111101011100101101011100110001011111101001000010011010010111100110010011111
10111000001010110001000011101010011010000111001001100111011111110101000001000010001010
01010100001000001011110001001001101010111000110100110111001110101110010001001110101
0111010000010100100010001101010101110000000101100000100110001011101010010101100110000
11111100110000011111000110000110111001110100111010011101001110010011101101101010101000
0000000100000001010000001000100001010101001000000101000001101000001100100011011101011010100
010100001010001001000101011010100001100001001110010111001110010111011100100101011011
0000101011100100001011101001001010011010001111011101100101010111000000100110000101111
10010010001110110101101101100011000111011101010100110000110011100111110111100001010
011001000111110101100001000111001010110110000110101100111000111101101100010110111010
01101010011110000111001100110111111101000000010010000010110100010011001010111100001
00001100101001111100011100011011101101101101101101010110110000011011000011010110100011
011001011101110010101001110000011101100011010111011000101010101000000110010000111110
100110001001111101011100010001010101010011000000111100001100011001111011111000100000
111000100110101011101100010010110101100101000111100010110011010011111100111000011110
11001100101111111001000000111010000110100100110011011011110101010001000000101010000
10000010010100010110001010011101000111010001011010011001100111111111000000001100000000
111100000110011000111111101100000010110000100101100101100111100111100111100011100111
0110011110111100010100011010001011100101001011100011001011100111001101000111100010110
00110011101101110101101001000110011010111110001000001101010001110000101101100100110
11110111101001010010001100010111110111010001010100101000001100010001111010101001000001
1110100011001001011111011001000010111101010010010000110110100111011001110101111010000100
0100101010101100000000111000000110110000110111001101010111100000100011000101011101010
    
```

A. No. This is output of an 11 bit LFSR!

LFSR Challenge 2

Goal. Decrypt/encrypt 1 gigabyte movie. How big an LFSR?

- A. 30 bits should be enough.
  - Probably not; Eve can try all  $2^{30}$  possibilities and see which one results in a movie.
- B. 100 bits is safe.
  - Maybe; Eve would need  $10^{12}$  centuries to try all  $2^{100}$  possibilities.
- C. 1000 bits makes it sufficiently secure.
  - Experts have cracked LFSR.
  - More complicated machines needed.



What else can we do with a LFSR?

- DVD encryption with CSS.
- DVD decryption with DeCSS!
- Subroutine in military cryptosystems.

```

/*  efdtt.c  Author: Charles M. Hannum <root@ihack.net>  */
/*  Usage is: cat title-key scrambled.vob | efdtt >clear.vob  */

#define m(i) (x[i]^s[i+84])<<

    unsigned char x[5]          ,y,s[2048];main(
n){for( read(0,x,5          );read(0,s ,n=2048
); write(1 ,s,n          )if(s
[y=s [13]%8+20] /16%4 ==1 ) (int
i=m( 1)17 ^256 +m(0) 8,k ==m(2)
0,j= m(4) 17^ m(3) 9^k* 2-k%8
^8,a =0,c =26;for (s[y] -=16;
--c;j *=2)a= a*2^i& 1,i=i /2^j&1
<<24;for(j= 127; ++j<n;c=c>
y)
c
+
+=y*i/8^i>>4^i>>12,
i=i>>8^y<<17,a^=a>>14,y=a^a*8^a<<6,a=a
>>8^y<<9,k=s[j],k
&7]+2^i^cr3sfw6v;*k+>n." [k>>4]*2^k*257/
8,s[j]=k^(k&k*2&34)*6^c+y
;}}
    
```

<http://www.cs.cmu.edu/~dst/DeCSS/Gallery>

Important properties.

- Built from simple components.
- Scales to handle huge problems.
- Requires a deep understanding to use effectively.

Basic Component	LFSR	Computer
control	start, stop, load	same
clock	regular pulse	2.8 GHz pulse
memory	11 bits	1 GB
input	seed	sequence of bits
computation	shift, XOR	logic, arithmetic, ...
output	pseudo-random bits	Sequence of bits

**Critical difference.** General purpose machine can be programmed to simulate ANY abstract machine.

A Profound Idea

**Programming.** Can write a Java program to simulate the operations of any abstract machine.

- Basis for theoretical understanding of computation. [stay tuned]
- Basis for bootstrapping real machines into existence. [stay tuned]

**Stay tuned.** See Assignment 5.

```

public class LFSR {
    private int fill[];
    private int tap;
    private int N;

    public LFSR(String fill, int tap) { ... }

    public int step() { ... }

    public static void main(String[] args) {
        LFSR lfsr = new LFSR("01101000010", 8);
        for (int i = 0; i < 2000; i++)
            StdOut.println(lfsr.step());
    }
}
    
```

```

% java LFSR
11001001001111011011100101101
01110011000101111110100100001
00110100101111001100100111...
    
```

A Profound Question

Q. What is a random number?

LFSR does not produce random numbers.

- It is a very simple deterministic machine.
- But it is hard to distinguish the bits it produces from random ones.

Q. Are random processes found in nature?

- Motion of cosmic rays or subatomic particles?
- Mutations in DNA?

Q. Is the natural world a (not-so-simple) deterministic machine?

“ God does not play dice. ” – Albert Einstein

