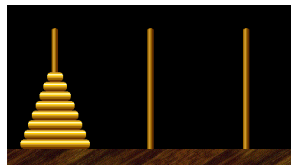
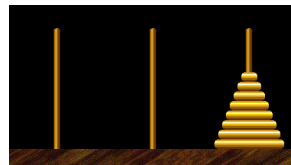


2.3 Recursion



start



finish

Greatest Common Divisor

Gcd. Find largest integer that evenly divides into p and q.

Ex. $\text{gcd}(4032, 1272) = 24$.

$$\begin{aligned} 4032 &= 2^6 \times 3^2 \times 7^1 \\ 1272 &= 2^3 \times 3^1 \times 53^1 \\ \text{gcd} &= 2^3 \times 3^1 = 24 \end{aligned}$$

Applications.

- Simplify fractions: $1272/4032 = 53/168$.
- RSA cryptosystem.

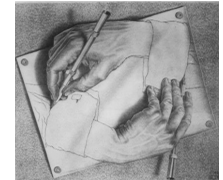
What is recursion? When one function calls **itself** directly or indirectly.

Why learn recursion?

- New mode of thinking.
- Powerful programming paradigm.

Many computations are naturally self-referential.

- Mergesort, FFT, gcd.
- Linked data structures.
- A folder contains files and other folders.



Reproductive Parts
M. C. Escher, 1948

Closely related to mathematical induction.

Greatest Common Divisor

Gcd. Find largest integer that evenly divides into p and q.

Euclid's algorithm. [Euclid 300 BCE]

$$\text{gcd}(p, q) = \begin{cases} p & \text{if } q = 0 \\ \text{gcd}(q, p \% q) & \text{otherwise} \end{cases}$$

← base case
← reduction step, converges to base case

$$\begin{aligned} \text{gcd}(4032, 1272) &= \text{gcd}(1272, 216) \\ &= \text{gcd}(216, 192) \\ &= \text{gcd}(192, 24) \\ &= \text{gcd}(24, 0) \\ &= 24. \end{aligned}$$

$4032 = 3 \times 1272 + 216$

Greatest Common Divisor

Gcd. Find largest integer d that evenly divides into p and q.

$$\text{gcd}(p, q) = \begin{cases} p & \text{if } q = 0 \\ \text{gcd}(q, p \% q) & \text{otherwise} \end{cases}$$

← base case
← reduction step, converges to base case

p							
q				q			
x	x	x	x	x	x	x	x

p = 8x
q = 3x
gcd(p, q) = x

↑
gcd

Recursive Graphics

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Greatest Common Divisor

Gcd. Find largest integer d that evenly divides into p and q.

$$\text{gcd}(p, q) = \begin{cases} p & \text{if } q = 0 \\ \text{gcd}(q, p \% q) & \text{otherwise} \end{cases}$$

← base case
← reduction step, converges to base case

Java implementation.

```
public static int gcd(int p, int q) {
    if (q == 0) return p;
    else return gcd(q, p % q);
}
```

← base case
← reduction step

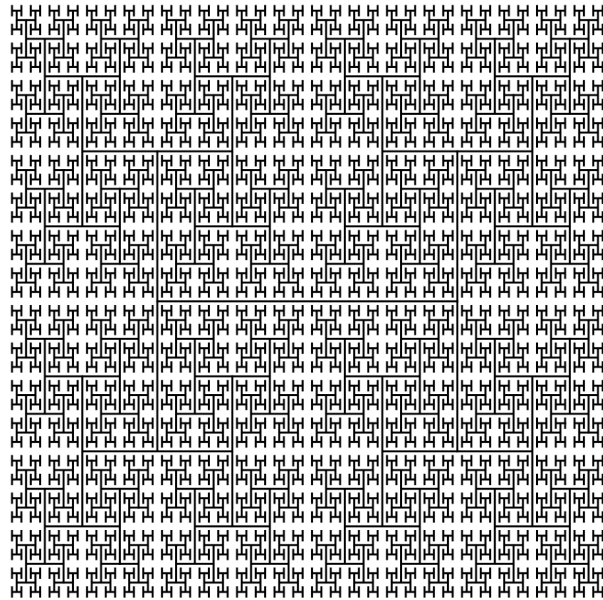


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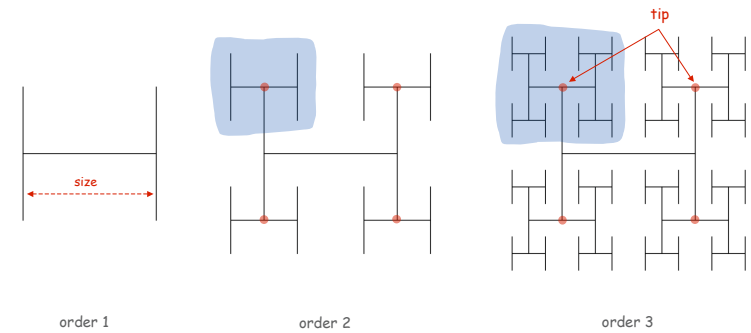


Htree

H-tree of order n.

- Draw an H.
- Recursively draw 4 H-trees of order n-1, one connected to each tip.

and half the size



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Htree in Java

```
public class Htree {
    public static void draw(int n, double size, double x, double y) {
        if (n == 0) return;
        double x0 = x - size/2, x1 = x + size/2;
        double y0 = y - size/2, y1 = y + size/2;

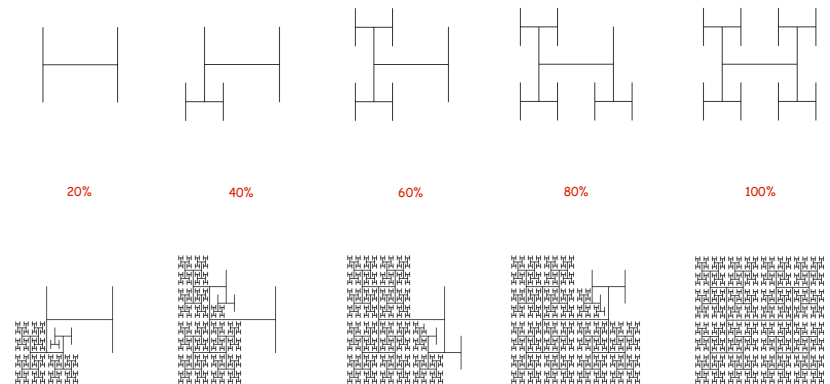
        StdDraw.line(x0, y, x1, y); ← draw the H, centered on (x, y)
        StdDraw.line(x0, y0, x0, y1);
        StdDraw.line(x1, y0, x1, y1);

        draw(n-1, size/2, x0, y0); ← recursively draw 4 half-size Hs
        draw(n-1, size/2, x0, y1);
        draw(n-1, size/2, x1, y0);
        draw(n-1, size/2, x1, y1);
    }

    public static void main(String[] args) {
        int n = Integer.parseInt(args[0]);
        draw(n, .5, .5, .5);
    }
}
```

Animated H-tree

Animated H-tree. Pause for 1 second after drawing each H.



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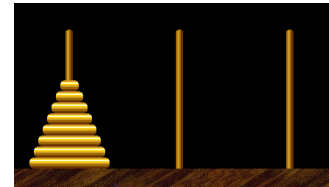
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Towers of Hanoi

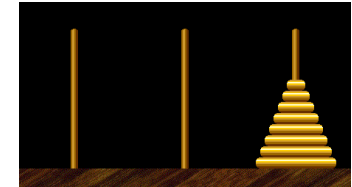
Towers of Hanoi

Move all the discs from the leftmost peg to the rightmost one.

- Only one disc may be moved at a time.
- A disc can be placed either on empty peg or on top of a larger disc.



start



finish



Towers of Hanoi demo

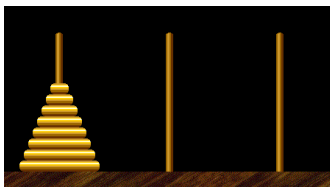


Edouard Lucas (1883)

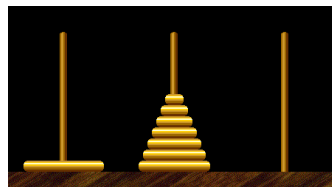
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Towers of Hanoi: Recursive Solution

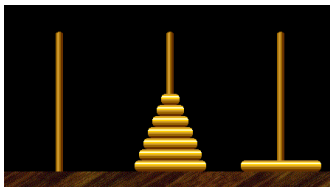


Move n-1 smallest discs right.

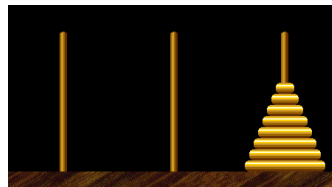


Move largest disc left.

← cyclic wrap-around



Move n-1 smallest discs right.



Towers of Hanoi Legend

- Q. Is world going to end (according to legend)?
 - 64 golden discs on 3 diamond pegs.
 - World ends when certain group of monks accomplish task.
- Q. Will computer algorithms help?

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Towers of Hanoi: Recursive Solution

```
public class TowersOfHanoi {

    public static void moves(int n, boolean left) {
        if (n == 0) return;
        moves(n-1, !left);
        if (left) System.out.println(n + " left");
        else System.out.println(n + " right");
        moves(n-1, !left);
    }

    public static void main(String[] args) {
        int N = Integer.parseInt(args[0]);
        moves(N, true);
    }
}
```

moves(n, true) : move discs 1 to n one pole to the left
 moves(n, false): move discs 1 to n one pole to the right

Towers of Hanoi: Recursive Solution

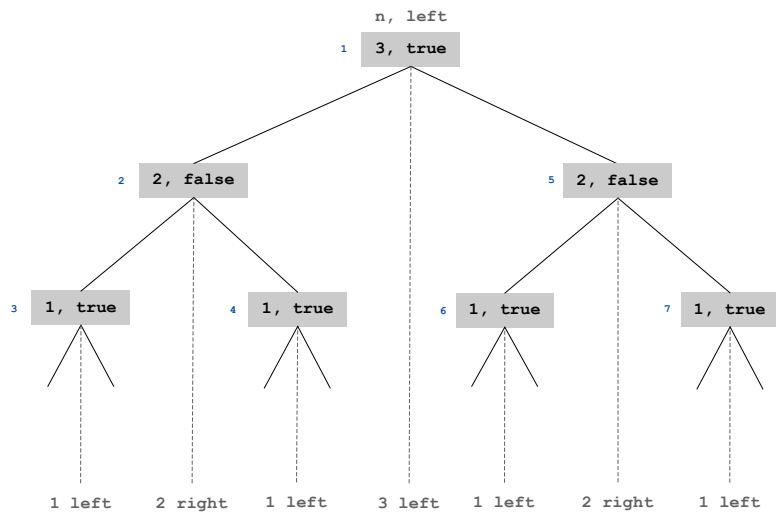
```
% java TowersOfHanoi 3
1 left
2 right
1 left
3 left
1 left
2 right
1 left

% java TowersOfHanoi 4
1 right
2 left
1 right
3 right
1 right
2 left
1 right
4 left
1 right
2 left
1 right
3 right
1 right
2 left
1 right
```

every other move is smallest disc

subdivisions of ruler

Towers of Hanoi: Recursion Tree



Properties of Towers of Hanoi Solution

Remarkable properties of recursive solution.

- Takes $2^n - 1$ steps to solve n disc problem.
- Sequence of discs is same as subdivisions of ruler.
- Smallest disc always moves in same direction.

Recursive algorithm yields non-recursive solution!

- Alternate between two moves:
 - move smallest disc to right if n is even
 - make only legal move not involving smallest disc
- to left if n is odd

Recursive algorithm may reveal fate of world.

- Takes 585 billion years for n = 64 (at rate of 1 disc per second).
- Reassuring fact: any solution takes at least this long!

Divide-and-Conquer

Divide-and-conquer paradigm.

- Break up problem into smaller subproblems of same structure.
- Solve subproblems recursively using same method.
- Combine results to produce solution to original problem.

Divide et impera. Veni, vidi, vici. - Julius Caesar

Many important problems succumb to divide-and-conquer.

- FFT for signal processing.
- Multigrid methods for solving PDEs.
- Adaptive quadrature for integration.
- Quicksort and mergesort for sorting.
- Hilbert curve for domain decomposition.
- Quad-tree for efficient N-body simulation.
- Midpoint displacement method for fractional Brownian motion.

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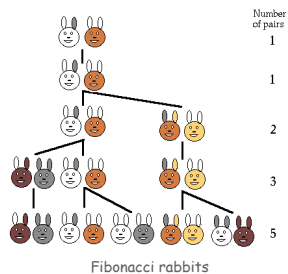
Pitfalls

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Fibonacci Numbers

Fibonacci numbers. 0, 1, 1, 2, 3, 5, 8, 13, 21, 34, ...

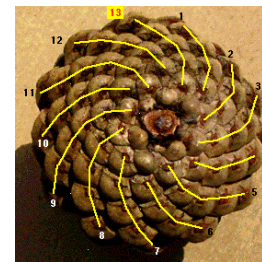
$$F_n = \begin{cases} 0 & \text{if } n = 0 \\ 1 & \text{if } n = 1 \\ F_{n-1} + F_{n-2} & \text{otherwise} \end{cases}$$



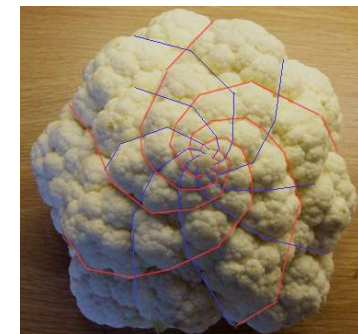
L. P. Fibonacci
(1170 - 1250)

23

Fibonacci Numbers and Nature



pinecone



cauliflower

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Possible Pitfalls With Recursion

Fibonacci numbers. 0, 1, 1, 2, 3, 5, 8, 13, 21, 34, ...

$$F_n = \begin{cases} 0 & \text{if } n=0 \\ 1 & \text{if } n=1 \\ F_{n-1} + F_{n-2} & \text{otherwise} \end{cases}$$

A natural for recursion?

```
public static long F(int n) {
    if (n == 0) return 0;
    if (n == 1) return 1;
    return F(n-1) + F(n-2);
}
```

spectacularly inefficient code

Observation. It takes a really long time to compute $F(50)$.

Summary

How to write simple recursive programs?

- Base case, reduction step.
- Trace the execution of a recursive program.
- Use pictures.



Towers of Hanoi by W. A. Schloss.

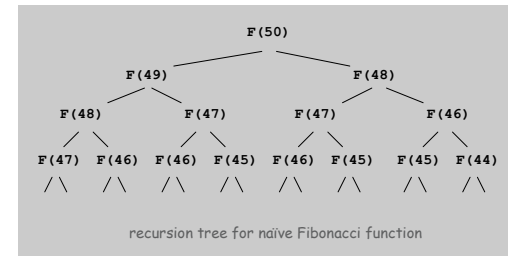
Why learn recursion?

- New mode of thinking.
- Powerful programming tool.

Divide-and-conquer. Elegant solution to many important problems.

Possible Pitfalls With Recursion

Caveat. Can easily write remarkably inefficient programs.



F(50) is called once.
 F(49) is called once.
 F(48) is called 2 times.
 F(47) is called 3 times.
 F(46) is called 5 times.
 F(45) is called 8 times.
 ...
 F(1) is called 12,586,269,025 times.

↖
F(50)

Binet's formula.

$$F(n) = \frac{\phi^n - (1-\phi)^n}{\sqrt{5}}$$

$$= \left\lfloor \frac{\phi^n}{\sqrt{5}} \right\rfloor$$

↖
 $\phi = \text{golden ration} = 1.618$