



**What computers
just *cannot* do.**

COS 116: 3/1/2007

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Administrivia

- Do we want “precepts” at start of lab?
- Take-home midterm in midterms week (closed book, 3-hour test). Preferences?
- Couple review sessions before midterm in the evening.
- Handouts today: 2 articles; HW 2

“Prof, what’s with all the negative thinking?!?”



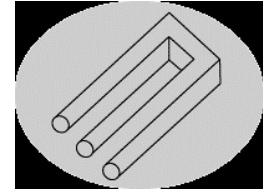
- An obvious motivation:
Understand the limits of technology



The power of negative thinking....

History-of-science perspective

Often, impossibility result \longrightarrow deep insight

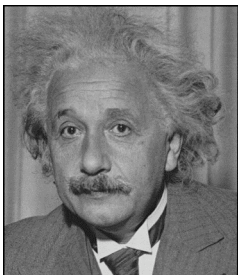


Examples

- Impossibility of trisecting angle with ruler and compass (Galois) \longrightarrow Group Theory and much of modern math



- Nothing travels faster than light \longrightarrow Relativity and modern physics



CAPTCHA & ESP

[Luis von Ahn]



Computer generated test that computers cannot pass easily.

Reminiscent of cryptography



Image labeling

Can mathematicians be replaced by machines?

[Hilbert, 1900]

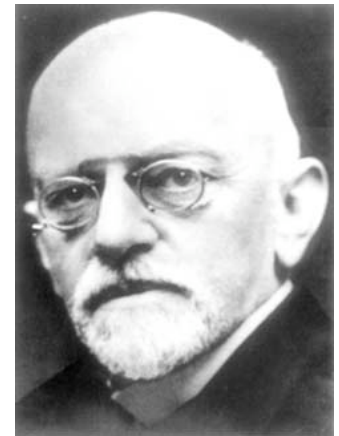
Math is axiomatic

Axioms – Set of statements

Derivation rules – finite set of rules for deriving
new statements from axioms

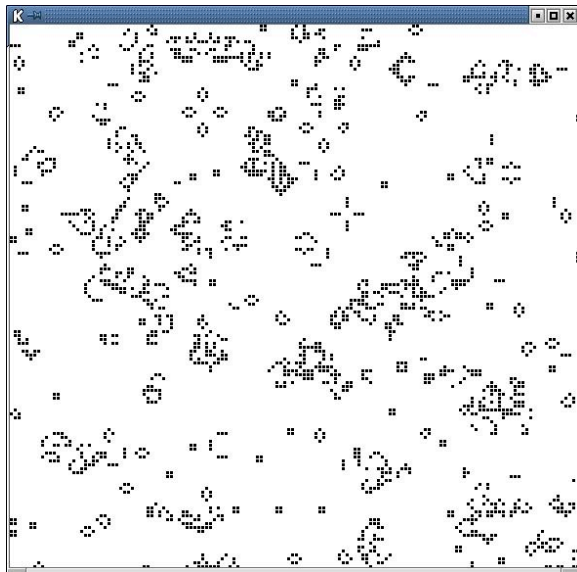
Theorems – Statements that *can* be derived from
axioms in a finite number of steps

Mathematician – Person who tries to determine
whether or not a statement is a theorem.



Can a simple set of mathematical equations
“solve” problems like:

Given starting configuration for the game of life, determine
whether or not cell (100,100) is ever occupied by a critter.



John Conway

Automated Checking of Software?



Windows XP: 40 million line program

Can computers check whether or not it will ever crash?

What is computation?

A formalization of an age-old notion

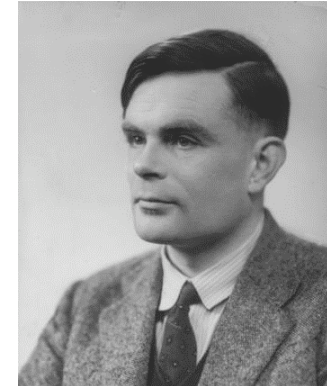
Basic Elements

- Scratch Pad
- Step-by-step description of what to do (“program”); should be finite!
- At each step:
 - Can only scan a fixed number of symbols
 - Can only write a fixed number of symbols



Turing's model

... 0 0 0 0 0 0 1 1 1 1 0 0 0 0 0 0 0 ...



- 1 dimensional unlimited scratchpad (“infinite”)
- Only symbols are 0/1 (tape initially has all 0s.)
- Can only scan/write one symbol per step
- Program looks like →

```
1. PRINT 0
2. GO LEFT
3. GO TO STEP 1 IF 1 SCANNED
4. PRINT 1
5. GO RIGHT
6. GO TO STEP 5 IF 1 SCANNED
7. PRINT 1
8. GO RIGHT
9. GO TO STEP 1 IF 1 SCANNED
10. STOP
```

The Doubling Program



What does this program do?

1. PRINT 0
2. GO RIGHT
3. GO TO STEP 1 if 1 SCANNED
4. GO TO STEP 1 if 0 SCANNED



Turing –Church Thesis

This model exactly captures what computation is.

It can **simulate** every other computational model that can be physically built.

“Code” for a program

= Binary Representation



Many conventions possible (e.g., ASCII)

Davis’s convention:

Code	Instruction
000	PRINT 0
001	PRINT 1
010	GO LEFT
011	GO RIGHT
1010...01	GO TO STEP <i>i</i> IF 0 IS SCANNED
1101...10	GO TO STEP <i>i</i> IF 1 IS SCANNED
100	STOP

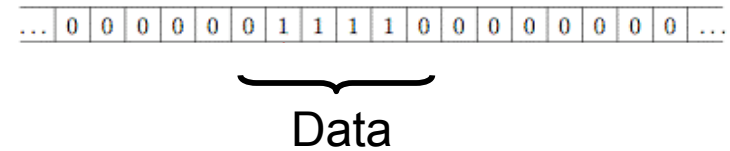
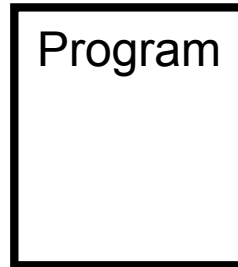
P ↔ Code (P)



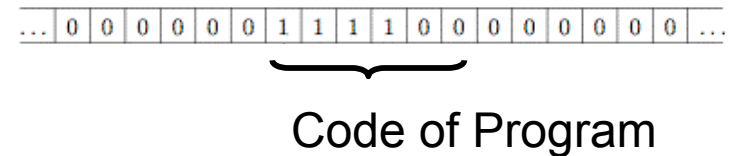
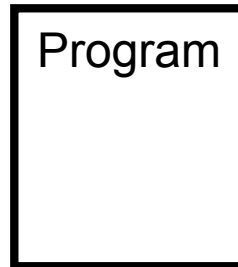
Programs and Data

A False Dichotomy

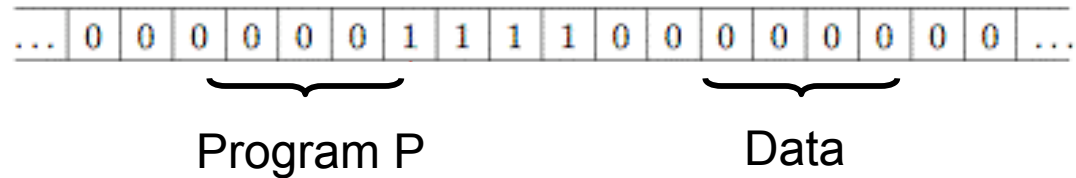
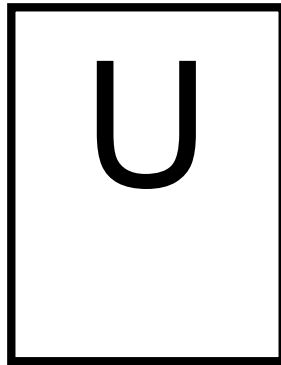
Usual viewpoint -



But can have -



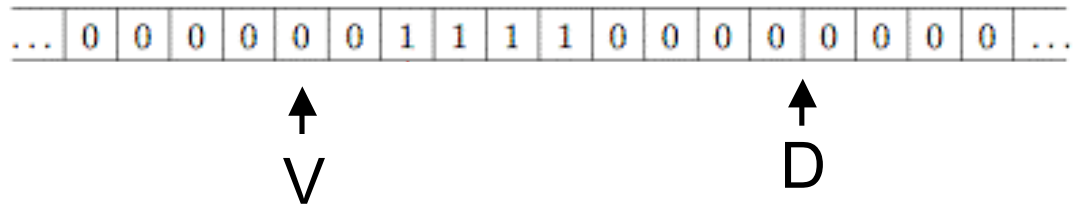
Universal Program U



- U “simulates” what P would do on that data

Automated Bug Checking Revisited

Halting Problem



Let P = program such that $\text{code}(P) = V$.
Does P halt on data D ?

Trivial Idea: Simulate P using universal program U .
If P halts, will eventually detect.

Problem: But if P never halts, neither does the simulation.



Next Time: Halting Problem is unsolvable by another program

Also, some class discussion of the readings.

Need to understand notion of Turing-Post program (e.g., doubling program) and what a universal program is.