

## 1 From Last Time

$H$  is finite. Then with probability  $\geq 1 - \delta$ , if  $h \in H$  is consistent then

$$err(h) \leq \frac{\ln|H| + \ln\frac{1}{\delta}}{m}. \quad (1)$$

We saw last time that this bound works if  $|H|$  is finite. What if  $|H|$  is infinite?

## 2 Intuition and examples

Even if we have infinitely many possible hypotheses, learning is possible from a finite sample.

### Example 1:

Let's say we have 3 examples. Then there are infinitely many possible hypotheses but only four possible labelings. Labelings are also called *behaviors* or *dichotomies*. In Fig. 1, all the possible labelings for the possible hypotheses are shown.

In such a case if we have  $m$  samples, there are  $m + 1$  possible labelings.

**Example 2 - Learning Intervals:** In this case, there are  $\frac{m(m-1)}{2} + m + 1 = \binom{m}{2} + m + 1$  possible labelings, where  $+m$  is for the intervals having just single points. As it can be seen, the number of labelings is  $O(m^2)$  for this example.

## 3 An upper bound for $err(h)$ when $|H|$ is not finite

### 3.1 Notation

The following notation was introduced:

$$\begin{aligned} S &= \langle x_1, x_2, \dots, x_m \rangle, \\ \Pi_H(S) &= \{ \langle h(x_1), h(x_2), \dots, h(x_m) \rangle : h \in H \}, \\ \Pi_H(m) &= \max_{|S|=m} |\Pi_H(S)| \leq 2^m. \end{aligned}$$

Note:  $\Pi_H(S)$  is the set of all possible labelings for all possible hypotheses and  $\Pi_H(m)$  is the number we computed in the above examples.

### 3.2 Finding the upper bound

For any  $H$ , there are 2 possible cases:

1. Either  $\forall m, \Pi_H(m) = 2^m$ , which is the worst case,

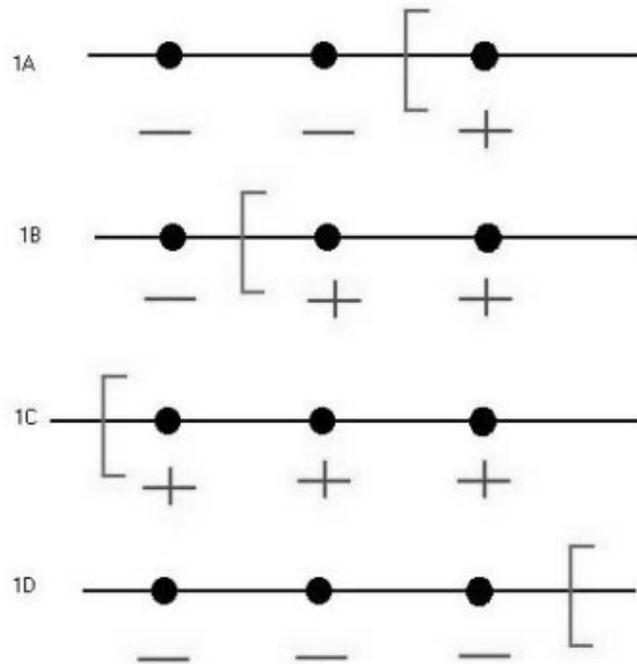


Figure 1: Possible labelings when we have 3 samples

2. or,  $\Pi_H(m) = O(m^d)$ , which is a really nice case. Here,  $d$  is the VC-dimension of  $H$  where  $VC$  stands for Vapnik-Chervonenkis. VC-dimension will be defined in next lecture.

**Step 1:** Derive an error bound in which  $\ln |H|$  is replaced by  $\ln |O(m^d)|$  so, we will get a result analogous to Occam's razor result.

**Theorem:** With probability  $\geq 1 - \delta$ ,  $\forall h \in H$ , if  $h$  is consistent, then

$$err(h) \leq O\left(\frac{\ln \Pi_H(2m) + \ln(\frac{1}{\delta})}{m}\right) \quad (2)$$

**Proof:** First, we will try to show that with probability  $\geq 1 - \delta$

$$(\forall h \in H : h \text{ is consistent, } err(h) \leq \epsilon). \quad (3)$$

Let's define event  $B$  and  $\Pr[B]$  as follows:

$$\Pr_S[\underbrace{\exists h \in H : h \text{ is consistent on } S \text{ but } err(h) > \epsilon}_{event B}].$$

Note that event  $B$  is the negation of the event defined in (3). We are trying to bound  $\Pr_S[B]$ . Because, if  $\Pr_S[B] < \delta$ , then  $\Pr[\text{event defined in (3)}] \geq 1 - \delta$  which is what we want to show.

*Trick:* Replace the error with error on another sample. In this new sample, there will be finitely many errors we need to consider. Let

$S' =$  second sample of  $m$  examples.

The data is independent identically distributed. We will argue that it is unlikely to see many errors on one sample, and no errors on the second sample.

$S = \langle x_1, x_2, \dots, x_m \rangle$  all i.i.d,  
 $S' = \langle x'_1, x'_2, \dots, x'_m \rangle$  all i.i.d,  
 $S; S'$  has  $2m$  samples.

**NOTATION:**

$M(h) = |\{i : h(x'_i) \neq c(x'_i)\}|$ . (number of mistakes)

$B' \equiv \exists h \in H : h$  is consistent on  $S$  and  $M(h) \geq \frac{m\epsilon}{2}$ . (We have  $m$  samples and probability of making error for each sample is  $\epsilon$ .)

*Claim:*  $\Pr[B'|B] \geq \frac{1}{2}$  i.e. if you are in bad case  $B$ , the probability that you are in case  $B'$  is  $\geq \frac{1}{2}$ .

If you know  $B$  happens, i.e., if  $h$  is consistent on  $S$  and  $\text{err}(h) > \epsilon$ , then  $M(h) \geq \frac{m\epsilon}{2}$  with probability  $\geq \frac{1}{2}$  which implies  $\Pr[B'|B] \geq \frac{1}{2}$ . (This will be proven later.)

$$\begin{aligned} \Pr[B'] &\geq \Pr[B' \wedge B] \\ &= \Pr[B] \cdot \Pr[B'|B] \\ &\geq \frac{1}{2} \Pr[B] \end{aligned} \tag{4}$$

(4) implies  $\Pr[B] \leq 2\Pr[B']$ . So, if probability of event  $B'$  happening is small, then the probability of event  $B$  happening is also small. Thus, instead of bounding probability of event  $B$ , we can start working with event  $B'$  and bound its probability.

*Experiment I:* Draw  $S$  at random and then draw  $S'$  at random.

*Experiment II:* Draw  $S, S'$ . With probability  $1/2$  interchange  $x_i$  and  $x'_i$  and with probability  $1/2$  leave them as they are. Doing this will not change the sample distribution.

As Experiment I and Experiment II will give the same distribution of examples, we can work with experiment II. So,

FIX  $h, S, S'$ . We will try to bound  $\Pr[B'|S, S']$ .  
 Recall,  $B' \equiv \exists h \in H : h$  is consistent on  $S$  and  $M(h) \geq \frac{m\epsilon}{2}$ .

$$\begin{array}{rcccccc}
& & x_1 & x_2 & \dots & & \\
S & : & 0 & 1 & 0 & 1 & 0 & 0 \\
S' & : & 1 & 1 & 0 & 0 & 1 & 1 \\
& & x_{1'} & x_{2'} & \dots & & & 
\end{array}$$

$S : 0 0 0 0 \dots$  means  $h$  is consistent with  $S$ .

If  $\exists i$  such that both  $x_i$  and  $x'_i$  are 1, then there is no way we can have all zeros in  $S$ . So,

$$Pr[h \text{ is consistent on } S \text{ and } M(h) \geq \frac{m\epsilon}{2}] = 0. \quad (5)$$

If there are  $M(h)$   $i$ 's where exactly one of  $x_i$  or  $x'_i$  is 1, then,

$$Pr[h \text{ is consistent on } S] \leq 2^{-M(h)} \text{ (this is the probability of all the 1s ending up in } S'). \quad (6)$$

We can think of this as follows: If  $x_1$  is 0 and  $x'_1$  is 1, w.p  $1/2$   $x_1$  will remain 0. If  $x_2$  is 0 and  $x'_2$  is 0, w.p  $1$   $x_2$  will remain 0 ...etc. So; the probability of all  $x_i$ 's being 0 is:

$$\frac{1}{2} \cdot 1 \cdot \frac{1}{2} \dots = \left(\frac{1}{2}\right)^{\text{number of } i\text{'s for which only one of } x_i \text{ or } x'_i \text{ is 1}}$$

unless there is an  $i$  for which both  $x_i$  and  $x'_i$  are 1 in which case the probability is zero.

Let  $H'(S) =$  one representative from  $H$  for each dichotomy in  $S$ . Then;

$$B' \equiv \left( \underbrace{\exists h \in H'(S; S') : h \text{ is consistent on } S \text{ and } M(h) \geq \frac{m\epsilon}{2}}_{e(h)} \right)$$

$$\begin{aligned}
Pr[B'|S, S'] &= Pr[\exists h \in H'(S; S') : e(h)|S, S'] \\
&= Pr[e(h_1) \vee e(h_2) \vee \dots \vee e(h_N)|S, S'] \\
&\leq \sum_{i=1}^N Pr[e(h_i)|S, S'] \\
&\leq |H'(S, S')| 2^{-m\epsilon/2} \\
&= |\Pi_H(S, S')| 2^{-m\epsilon/2}
\end{aligned}$$

The last equality comes from the fact that there is one representative for each labeling. So; number of representatives is equal to number of labelings.

In the next lecture;  $Pr[B']$  will be written as an expectation and the bound, found above, for  $Pr[B'|S, S']$  will be used to bound  $Pr[B']$  which will in turn give a bound for  $Pr[B]$ . Because  $Pr[B] \leq 2Pr[B']$ .