

# Lecture 11. Quicksort

- Sort  $x[0..n-1]$  into increasing (or decreasing) order
- Quicksort is a well-known sorting algorithm: Recursion is natural and fast

To sort  $x[0..n-1]$ :

1. Pick a 'pivot' element
2. Rearrange  $x$  so that:  
 $x[k]$  holds this element,  $x[0..k-1] < x[k]$ , and  $x[k+1..n-1] > x[k]$
3. Sort  $x[0..k-1]$  and  $x[k+1..n-1]$  recursively

```
void quicksort(int x[], int l, int r) {
    if (r > l) {
        int k = partition(x, l, r);
        quicksort(x, l, k - 1);
        quicksort(x, k + 1, r);
    }
}
```

```
int main(void) {
    int n, array[1000];

    ...
    quicksort(array, 0, n - 1);
    ...
}
```

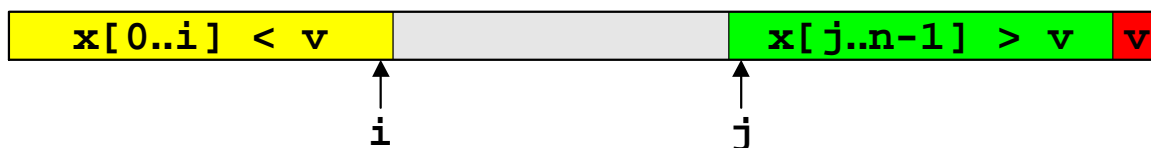
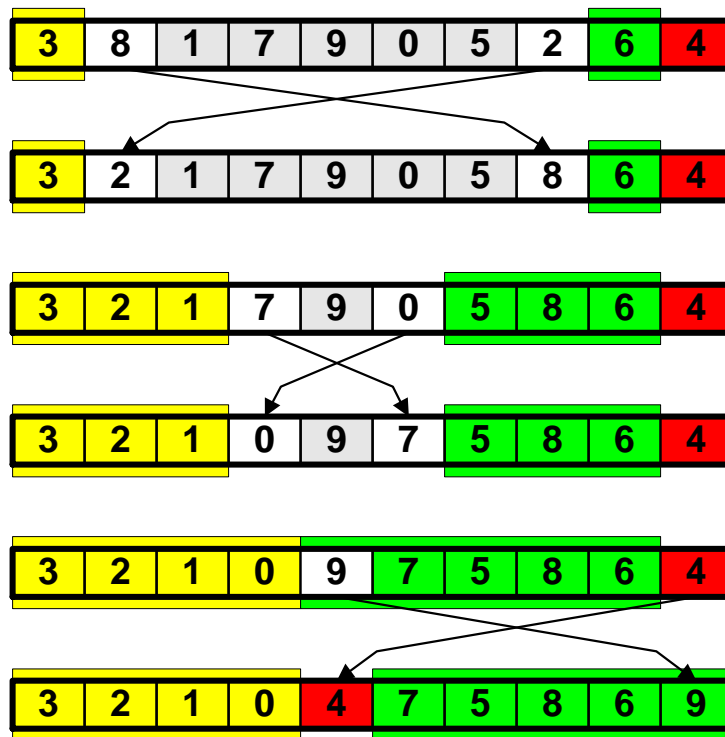
# Partitioning

```

int partition(int x[], int i, int j) {
    int k = j, v = x[k];

    i--;
    while (i < j) {
        while (x[++i] < v)
            ;
        while (--j > i && x[j] > v)
            ;
        if (i < j) {
            int t = x[i];
            x[i] = x[j];
            x[j] = t;
        }
    }
    x[k] = x[i];
    x[i] = v;
    return i;
}

```



- For more, read R. Sedgewick, *Algorithms in C*, Addison-Wesley, 1990

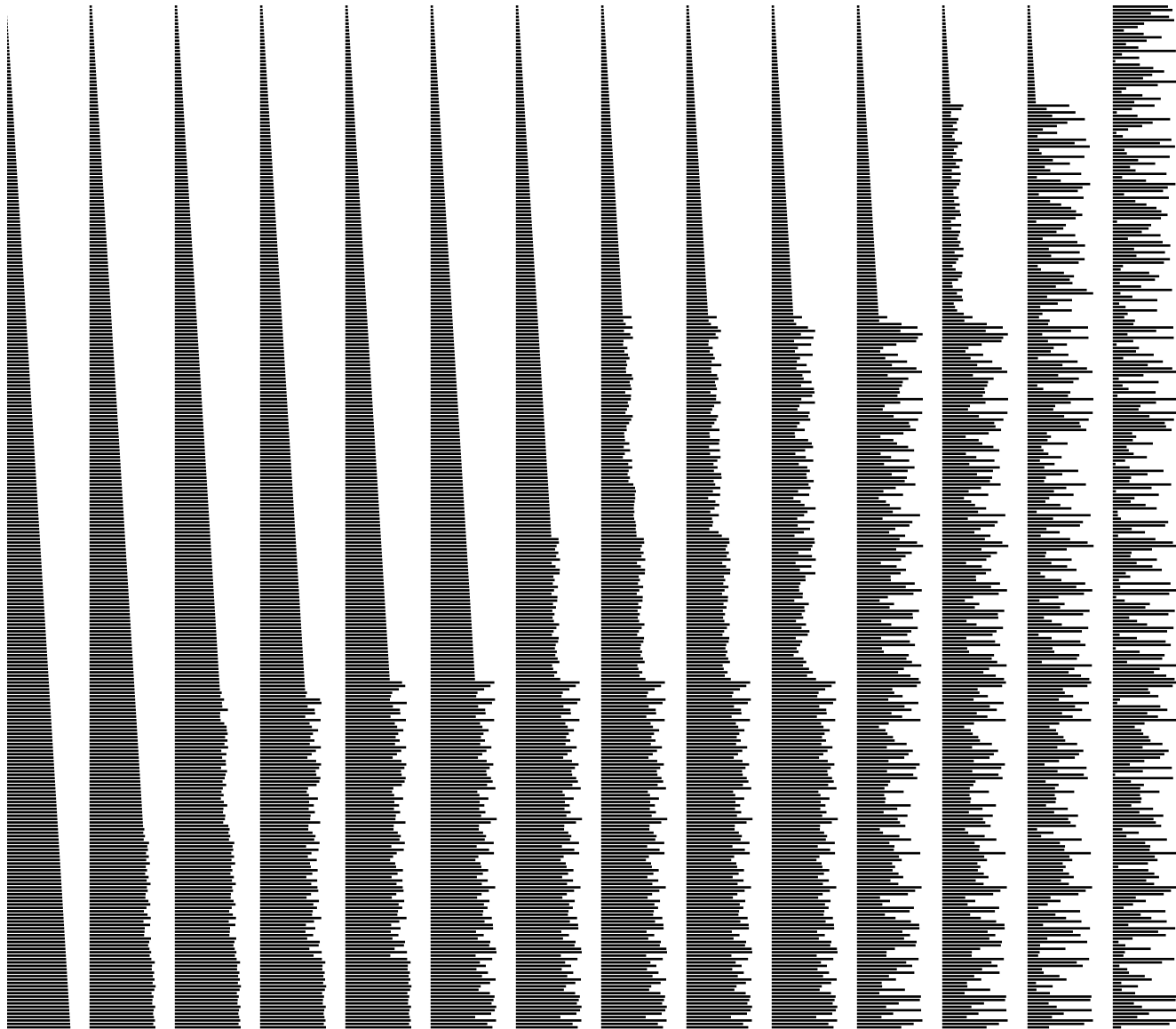
## Quicksort in Action

```

quicksort(x, 0, 9)      3  8  1  7  9  0  5  2  6  4
                       3  2  1  7  9  0  5  8  6  4
                       3  2  1  0  9  7  5  8  6  4
                       3  2  1  0  4 7  5  8  6  9
quicksort(x, 0, 3)     3  2  1  0 4  7  5  8  6  9
                       0 2  1  3  4  7  5  8  6  9
quicksort(x, 0, -1)
quicksort(x, 1, 3)     0  2  1  3 4  7  5  8  6  9
                       0  2  1  3 4  7  5  8  6  9
quicksort(x, 1, 2)     0  2  1 3  4  7  5  8  6  9
                       0  1  2  3  4  7  5  8  6  9
quicksort(x, 1, 0)
quicksort(x, 2, 2)
quicksort(x, 4, 3)
quicksort(x, 5, 9)     0  1  2  3  4  7  5  8  6  9
                       0  1  2  3  4  7  5  8  6  9
quicksort(x, 5, 8)     0  1  2  3  4  7  5  8  6  9
                       0  1  2  3  4  5  7  8  6  9
                       0  1  2  3  4  5  6  8  7  9
quicksort(x, 5, 5)
quicksort(x, 7, 8)     0  1  2  3  4  5  6  8  7  9
                       0  1  2  3  4  5  6  7  8  9
quicksort(x, 7, 6)
quicksort(x, 8, 8)
quicksort(x, 10, 9)

```

# Quicksort in Action, cont'd



# Implementing Recursive Functions

- Consider `sum(10)`: each call must have its own argument `n` and its return address
- Use a stack to hold arguments, local variables, and the return address

```

sum(n=10) calls
  sum(9)
    sum(8)
      sum(7)
        sum(6)
          sum(5)
            sum(4)
              sum(3)
                sum(2)
                  sum(1)
                    sum(0)
                      returns 0
                    returns 1
                  returns 3
                returns 6
              returns 10
            returns 15
          returns 21
        returns 28
      returns 36
    returns 45
  returns 55

```

ret. addr.
n=0
ret. addr.
n=1
ret. addr.
n=2
ret. addr.
n=3
ret. addr.
n=4
ret. addr.
n=5
ret. addr.
n=6
ret. addr.
n=7
ret. addr.
n=8
ret. addr.
n=9
ret. addr.
n=10

## Implementing Recursive Functions, cont'd

- Use conventions for the stack and for how arguments, etc. are 'pushed'

Use R<sub>7</sub> as the 'stack pointer:' it holds the address of the top element

Stack starts at FF<sub>16</sub> and grows 'down' — toward lower addresses

Push the arguments onto the stack before calling a function; push the return address upon entering a function

<b>30:</b>	B201	R2 ← 1	push the return address
<b>31:</b>	2772	R7 ← R7 - R2 = R7 - 1	
<b>32:</b>	A670	M[R7+0] ← R6	
<b>33:</b>	9171	R1 ← M[R7+1]	R1 ← n
<b>34:</b>	2312	R3 ← R1 - R2 = R1 - 1	R3 ← n - 1
<b>35:</b>	633D	jump to <b>3D</b> if R3 < 0	if (n == 0) return 0
<b>36:</b>	2772	R7 ← R7 - R2 = R7 - 1	push n - 1
<b>37:</b>	A370	M[R7+0] ← R3	
<b>38:</b>	8630	R6 ← PC, PC ← 30	call sum
<b>39:</b>	B201	R2 ← 1	pop n - 1
<b>3A:</b>	1772	R7 ← R7 + R2 = R7 + 1	
<b>3B:</b>	9271	R2 ← M[R7+1]	R2 ← n
<b>3C:</b>	1112	R1 ← R1 + R2	R1 ← sum(n-1) + n
<b>3D:</b>	9670	R6 ← M[R7+0]	pop return address
<b>3E:</b>	B201	R2 ← 1	
<b>3F:</b>	1772	R7 ← R7 + R2 = R7 + 1	
<b>40:</b>	7600	PC ← R6	return

## Implementing Recursive Functions, cont'd

- Main program makes the first call

00:	B000	R0 ← 0	R0 holds 0
01:	B7FF	R7 ← FF	initialize stack pointer
02:	B210	R2 ← 50	R2 ← address of n
03:	9220	R2 ← M[R2+0]	R2 ← n
04:	B101	R1 ← 1	push n
05:	2771	R7 ← R7 - R1 = R7 - 1	
06:	A270	M[R7+0] ← R2	
07:	86 <u>30</u>	R6 ← PC, PC ← 30	call sum
08:	B201	R2 ← 1	pop n
09:	1772	R7 ← R7 + R2 = R7 + 1	
0A:	4102	print R1	print sum(n)
0B:	0000	halt	
50:	0000		n

