

## Lecture 17. Analysis of Algorithms

- An ***algorithm*** is a ‘method’ for solving a problem that is ***independent*** of a specific computer or programming language
- ***Design***: Finding a way to solve the problem
- ***Analysis***: Determining the algorithm’s cost in machine-independent terms, e.g.  $\lg N$
- Need to make a program faster?

Get a new machine

Costs \$\$\$ or more

Makes ‘everything’ run faster

But, it may — or ***may not*** — have much impact on a specific problem

Get a new algorithm

Costs ¢ or less

Can make or break a specific problem by allowing it to be solved at all

But, it may have ***little or no*** impact on ‘everything’



## A Simple Brute-Force Solution

- Try all possible sublists of  $n$  integers:  $x[lb..ub]$  for all  $lb, ub$  from 0 to  $n$

```

void sublist(int x[], int n) {
    int lb, ub, l, r, max = 0;

    for (lb = 0; lb < n; lb++)
        for (ub = lb; ub < n; ub++) {
            int i, sum = 0;
            for (i = lb; i <= ub; i++)
                sum += x[i];
            if (sum > max) {
                max = sum;
                l = lb;
                r = ub;
            }
        }
    printf("x[%d..%d] = %d\n", l, r, max);
}

% lcc -I/u/cs126/include sublistn3.c /u/cs126/lib/libmisc.a
% echo 31 -41 59 26 -53 58 97 -93 -23 84 | a.out
x[2..6] = 187

```

## Profiling

- Program profiles help understand execution frequencies; use `lcc -b` and `bprint`

```
% lcc -b -I/u/cs126/include sublistn3.c /u/cs126/lib/libmisc.a
% echo 31 -41 59 26 -53 58 97 -93 -23 84 | a.out
x[2..6] = 187
% bprint
```

...

```

1 for (<1>lb = 0; <11>lb < n; <10>lb++)
    2 for (<10>ub = lb; <65>ub < n; <55>ub++) {
        int i, sum = <55>0;
        3 for (<55>i = lb; <275>i <= ub; <220>i++)
            <220>sum += x[i];
            if (<55>sum > max) {
                <6>max = sum;
                <6>l = lb;
                <6>r = ub;
            }
    }
    <1>printf("x[%d..%d] = %d\n", l, r, max);
```

- For  $N = 10$

Loop	<b>1</b>	is executed	$11 \approx 10^1$	times
	<b>2</b>		$65 \approx 10^2/2$	
	<b>3</b>		$275 \approx 10^3/3$	

Execution time  $\approx N^3$ , can't solve  $N = 10,000$ , since  $10^{12}$  microseconds  $\approx 11$  days

## A Better Algorithm

- 
**Don't recompute the whole sum every time**

$$x[lb] + x[lb+1] + \dots + x[ub] = (x[lb] + \dots + x[ub-1]) + x[ub]$$

```
void sublist(int x[], int n) {
    int lb, ub, l, r, max = 0;

    for (lb = 0; lb < n; lb++) {
        int sum = 0;
        for (ub = lb; ub < n; ub++) {
            sum += x[ub];
            if (sum > max) {
                max = sum;
                l = lb;
                r = ub;
            }
        }
    }

    printf("x[%d..%d] =
}
```

	31	-41	59	26	-53	58	97	-93	-23	84
	31	-10	49	75	22	80	177	84	61	145
		-41	18	44	-9	49	146	53	30	114
			59	85	32	90	187	94	71	155
				26	-27	31	128	35	12	96
					-53	5	102	9	-14	70
						58	155	62	39	123
							97	4	-19	65
								-93	-116	-32
									-23	61
										84

## Profiling the Better Algorithm

```

1 for (<1>lb = 0; <11>lb < n; <10>lb++) {
    int sum = <10>0;
    2 for (<10>ub = lb; <65>ub < n; <55>ub++) {
        <55>sum += x[ub];
        if (<55>sum > max) {
            <6>max = sum;
            <6>l = lb;
            <6>r = ub;
        }
    }
}
<1>printf("x[%d..%d] = %d\n", l, r, max);

```


- For  $N = 10$

Loop **1** is executed  $11 \approx 10^1$  times  
**2**  $65 \approx 10^2/2$

Execution time  $\approx N^2$ , but can't solve  $N = 1,000,000$ , because  $10^{12}$  microseconds  $\approx$  11 days

- There is a divide-and-conquer algorithm that takes  $\approx N \lg N$ , but there's even a better way

## The Optimal Algorithm

- 
 Keep track of the maximum sum so far and the sum of the sublist that ends at  $x[i]$

Suppose  $max$  is the maximum sum in  $x[0..i-1]$ ; extend that solution to  $x[i]$

31 -41 **59 26** -53 58 97 -93 -23 84  
                   85 32

31 -41 **59 26 -53 58** 97 -93 -23 84  
                                   90

```
void sublist(int x[], int n) {
    int i, l, r, max = 0, maxi = 0;

    for (i = 0; i < n; i++) {
        if (maxi + x[i] > 0)
            maxi += x[i];
        } else
            maxi = 0;
        if (maxi > max)
            max = maxi;
    }
    printf("x[%d..%d] = %d\n", l, r, max);
}
```

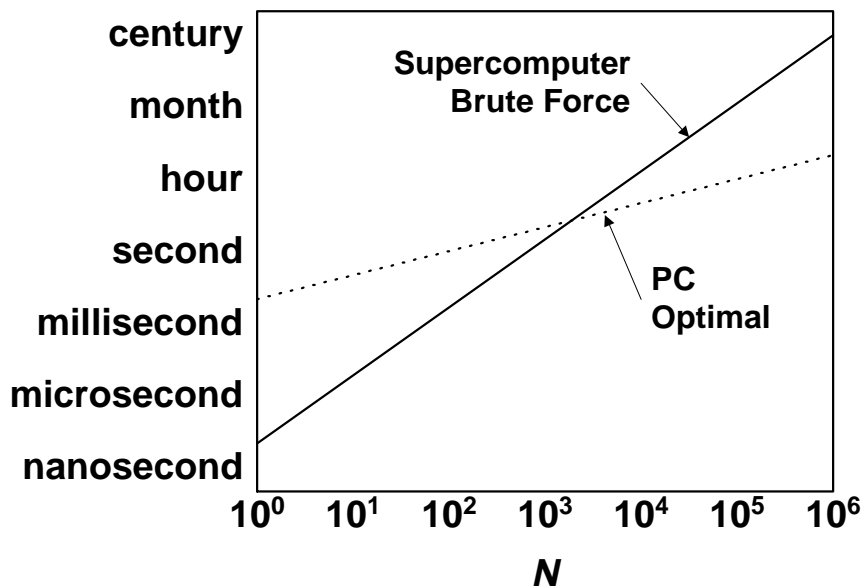
31 -41 **59 26 -53 58 97** -93 -23 84  
 31 0 59 85 32 90 **187** 94 71 155  
 31 31 59 85 85 90 187 187 187 187

- Execution time  $\approx N$ , because there's just one loop;  $N = 1,000,000$  takes  $\approx 1$  second
- See `sublistn.c` for details of computing  $l$  and  $r$

## Summary

- A good algorithm can be more powerful than a supercomputer

		Thousand	Million
Brute Force	$N^3$	17 min	300 centuries!
Better	$N^2$	1 sec	11 days
Divide and Conquer	$N \lg N$	0.01 sec	20 sec
Optimal	$N$	0.001 sec	1 sec



- For more, see J. Bentley, *Programming Pearls*, Addison Wesley, 1986