

**COS 126**

**Spring 2016**

**Written Exam 2 solutions**

**Q11. Searching and Sorting (7 points).** To the right of each option, fill in the circle corresponding to the one-word characterization that best describes the order of growth of the *worst-case* running time. *Correct answers are worth 1 point; incorrect answers are worth -1 point; blank answers are worth 0 points.*

logarithmic ( $\log n$ )    linear ( $n$ )    linearithmic ( $n \log n$ )    quadratic ( $n^2$ )

A. mergesort                                                                                       

~~→~~ B. merge *negating arrays*                                                                                          
*2 (1) 1000000000*

C. binary search                                                                                       

D. BST search *1000*                                                                                       

E. insertion sort                                                                                       

F. sequential search *1000*                                                                                       

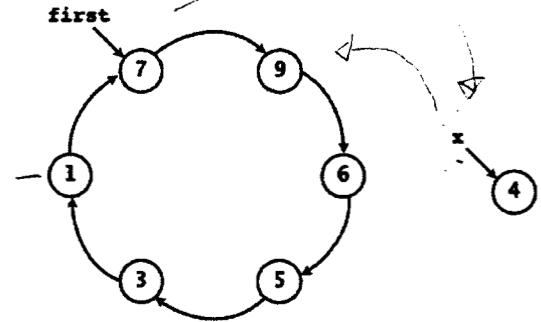
G. bubble sort

Q12. Linked Structures (7 points). Suppose that the Node data type is defined as:

```
private class Node {
    private int item;
    private Node next;
}
```

and that first is a variable of type Node that refers to one node in a circular linked list. Let x be a variable that refers to a newly created Node.

```
Node x = new Node();
x.item = 4;
```



For each of the independent code fragments below, fill in the circle that describes its effect. Fill in "neither" if the code invalidates the property that first refers to a node in a circular linked list. Correct answers are worth 1 point; incorrect answers are worth -1 point; blank answers are worth 0 points.

inserts x after first    inserts x before first    neither

```
x.next = first.next;
first.next = x;
```

```
Node z = first;
while(z.next != first) z = z.next;
z.next = x;
x.next = first;
```

```
first.next = x;
x.next = first.next;
```

```
x = first;
x.next = first;
```

*x refers to same Node as first  
so setting next to itself*

```
first.next = x;
```

*breaks circle*

```
Node z = first;
x.next = first.next;
z.next = x;
```

```
x.item = first.item;
first.item = x.item;
first.next = x;
```

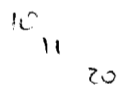
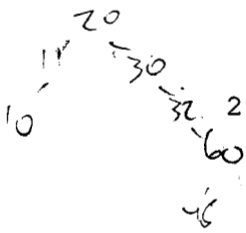
*first.item = 7  
first.next = x*

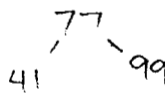
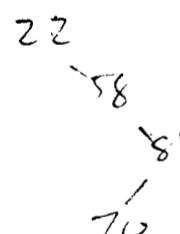


**Q13. Binary Search Trees (7 points).** Correct answers are worth 1 point; incorrect answers are worth -1 point; blank answers are worth 0 points.

A. Suppose that we create a BST by inserting integers into an initially empty tree. For each of the following insertion orders, fill in the bubble corresponding to the height of the binary tree that is produced (max distance from the root to any node) when keys are inserted in that order into an initially empty tree. The first answer is provided for you.

		1	2	3	4	5	6	7
	10, 11, 20, 30, 32, 48, 60	<input type="radio"/>	<input type="radio"/>	<input type="radio"/>	<input type="radio"/>	<input type="radio"/>	<input type="radio"/>	<input checked="" type="radio"/>
	10, 20, 32, 11, 30, 48, 60	<input type="radio"/>	<input type="radio"/>	<input type="radio"/>	<input type="radio"/>	<input checked="" type="radio"/>	<input type="radio"/>	<input type="radio"/>
	30, 10, 20, 11, 60, 48, 32	<input type="radio"/>	<input type="radio"/>	<input type="radio"/>	<input checked="" type="radio"/>	<input type="radio"/>	<input type="radio"/>	<input type="radio"/>
	20, 11, 10, 30, 32, 60, 48	<input type="radio"/>	<input type="radio"/>	<input type="radio"/>	<input type="radio"/>	<input checked="" type="radio"/>	<input type="radio"/>	<input type="radio"/>

B. Suppose that you are searching for the key 70 in a binary search tree. In the following list, indicate for each sequence whether it could be the sequence of keys examined, by filling in the corresponding bubble.

		Possible	Not possible
	77, 41, 99, 20, 85, 70	<input type="radio"/>	<input checked="" type="radio"/>
	99, 10, 80, 20, 60, 70	<input checked="" type="radio"/>	<input type="radio"/>
	5, 10, 80, 40, 32, 50, 70	<input type="radio"/>	<input checked="" type="radio"/>
	22, 58, 81, 70	<input checked="" type="radio"/>	<input type="radio"/>
			

**Q14. REs (7 points).** Let  $L = \{aaaba, aabaa, abbba, ababa, aaaaa\}$ . For each of the regular expressions below fill in the only answer that applies. *Correct answers are worth 1 point; incorrect answers are worth -1 point; blank answers are worth 0 points.*

The possible options (and their shortnames) are:

- [NONE] Matches no strings in  $L$ .
- [SOME] Matches only some strings in  $L$  and some other strings.
- [MORE] Matches all strings in  $L$  and some other strings.
- [EXACT] Matches all strings in  $L$  and no other strings.

	NONE	SOME	MORE	EXACT
$abbba$ A. $a(a b)^*abb(a b)^*$	<input checked="" type="radio"/>	<input type="radio"/>	<input type="radio"/>	<input type="radio"/>
B. $a(a b)^*a$	<input type="radio"/>	<input type="radio"/>	<input checked="" type="radio"/>	<input type="radio"/>
$aabaa$ C. $a^*b^*aba$	<input type="radio"/>	<input checked="" type="radio"/>	<input type="radio"/>	<input type="radio"/>
$ababa$ D. $a((a^* b^*) (b^*aba^*))a$	<input type="radio"/>	<input checked="" type="radio"/>	<input type="radio"/>	<input type="radio"/>
E. $a^*b^*a^*b^*a^*b^*a^*b^*a^*b^*$	<input type="radio"/>	<input checked="" type="radio"/>	<input type="radio"/>	<input type="radio"/>
$aaaaa$ F. $(a b)(a b)(a b)(a b)a$	<input type="radio"/>	<input type="radio"/>	<input checked="" type="radio"/>	<input type="radio"/>
G. $(a aa aaa)(ba aa bbb)a$	<input type="radio"/>	<input checked="" type="radio"/>	<input type="radio"/>	<input type="radio"/>

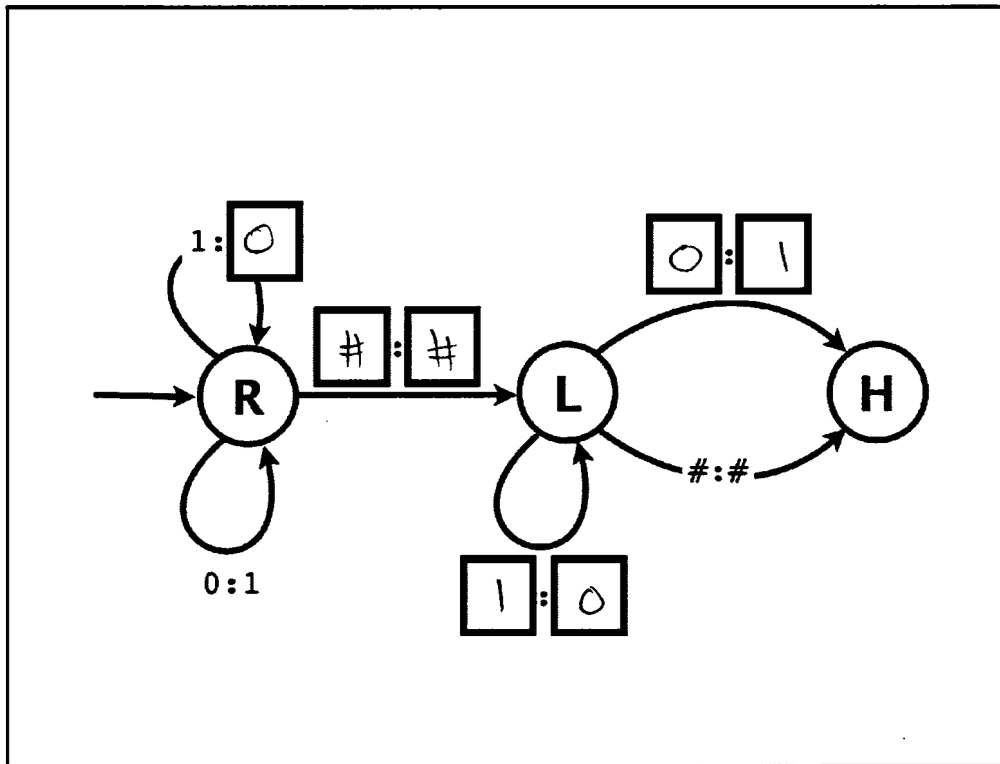
$aaaba$   
 $aabaa$   
 $abbba$   
 $ababa$   
 $aaaaa$



**Q15. Turing Machines (7 points).** Below is an incomplete diagram of a Turing Machine. Complete the diagram so that it satisfies the following specification:

- Assume the initial tape consists of a binary string with infinitely many # symbols on both sides.
- Assume the initial head location is the leftmost bit of the binary string.
- Assume the initial state is the state labelled R on the left, and recall that the Turing Machine's first step is reading/writing, not moving the head.
- Interpret the input as a two's-complement binary integer. After the Turing machine halts, the binary integer that remains on the tape should be the negative of the original input. (Recall that computing the negative of a number involves flipping all of the bits and then adding one.)

Fill in exactly one symbol in each of the seven empty square boxes below so as to satisfy this specification. Do not add new states or new transitions, and do not use any tape symbols other than #, 0, or 1.



flip  
 /  
 right of tape

# 10100 #  
 # 01011 #  
 # 01100 #

01011  
 01100

**Q16. Intractability (7 points).** For each of the computational problems below, indicate its difficulty by filling in the most appropriate choice among *True*, *False* or *Nobody Knows*. Correct answers are worth 1 point; incorrect answers are worth -1 point; blank answers are worth 0 points.

- |   | True                             | False                            | Nobody Knows          |
|---|----------------------------------|----------------------------------|-----------------------|
| <p><i>then about decidible problem</i></p> <p>A. If <u>P is not equal to NP</u> there is no polynomial-time algorithm for integer linear programming (ILP). <i>is it recursive?</i></p> | <input checked="" type="radio"/> | <input type="radio"/>            | <input type="radio"/> |
| <p>B. If <math>P = NP</math>, every problem in P is NP-complete. <i>can all problem in P be NP-complete?</i></p>  | <input checked="" type="radio"/> | <input type="radio"/>            | <input type="radio"/> |
| <p>C. If <math>P = NP</math> there is a polynomial-time algorithm for factoring. <i>because of NP</i></p>   | <input checked="" type="radio"/> | <input type="radio"/>            | <input type="radio"/> |
| <p>D. No problem is in both P and NP. <i>all P are in NP</i></p>  | <input type="radio"/>            | <input checked="" type="radio"/> | <input type="radio"/> |
| <p>E. There exists a deterministic Turing machine that can solve every problem in NP. <i>but it's false</i></p>   | <input checked="" type="radio"/> | <input type="radio"/>            | <input type="radio"/> |
| <p>F. A Universal Turing Machine can simulate the operation of any Turing machine, including itself, <u>in polynomial time</u>.</p>   | <input type="radio"/>            | <input checked="" type="radio"/> | <input type="radio"/> |
| <p>G. No polynomial-time algorithm can solve the Halting Problem. <i>nothing can</i></p>  | <input checked="" type="radio"/> | <input type="radio"/>            | <input type="radio"/> |



Q17. TOY (7 points). Fill in the blanks in the following table.

Hint:  $7 \times 16^3 = 28,672$ .

Note: The notation here and in the text and lectures is slightly different than the notation used in precepts and online. For example,  $R[2] = R[3] + R[4]$  means precisely the same thing as  $R[2] \leftarrow R[3] + R[4]$ .

hex	decimal integer	16-bit two's complement	TOY instruction pseudo-code
FFFE	-2	1111111111111110	$R[F] = PC; PC = FE$
1234	4,660	0001001000110100	$R[2] = R[3] + R[4]$
1101	4,353	0001000100000001	$R[1] = R[0] + R[1]$
77FF	30,719	0111011111111111	$R[7] = 00FF$
FF01	-255	1111111100000001	$R[F] = PC; PC = 01$
7A00	31232	0111101000000000	$R[A] = 0000$

$$\rightarrow 16^0 + 16^1 + 16^2 + 16^3$$

2860

28672

$$10 \times 16^3 = ?$$

$$\begin{array}{r} 31232 \\ - 28672 \\ \hline 2560 \end{array}$$

$$\begin{array}{r} 2560 \\ + 28672 \\ \hline 31232 \end{array}$$

$$\begin{array}{r} 0000000011111110 \\ + 1 \\ \hline 0000000011111111 \end{array}$$

$$\begin{array}{r} 128 \\ 64 \\ 32 \\ 16 \\ 8 \\ 4 \\ 2 \\ 1 \end{array}$$



TYPO

Q18. TOY Programming (7 points). Consider the following TOY program:

```

20: 81FF R[0] = stdin
21:      SEE BELOW
22: 1211 R[2] = R[1] + R[1] R[1] replaced with 1111
23: 0000 halt

```

To the right of each of the possibilities below for the instruction at M[21], write the contents of R[1] when the machine halts after being started at 20 with 1111 on standard input. Your answers must each be four hex digits.

Note: The notation used here and in the text and lectures is slightly different than the notation used in precepts and online. For example, R[1] = M[21] means precisely the same thing as R[1] ← mem[21].

Possibilities for M[21]

Contents of R[1] when halting

```

21: 1111 R[1] = R[1] + R[1]

```

2222

2222  
line 22 does nothing

```

21: 0000 halt

```

1111

```

21: 1211 R[2] = R[1] + R[1]

```

1111

does nothing  
0000

```

21: 7100 R[1] = 0000

```

0000

```

21: C023 PC = 23

```

1111

```

21: 8121 R[1] = M[21]

```

8121

```

21: 9122 M[22] = R[1]

```

2222

1111



**Q19. Combinational Circuits (7 points).** The *IS-XOR?* function of 3 boolean variables  $x$ ,  $y$  and  $z$  is 1 if and only if  $(x \text{ XOR } y) == z$ .

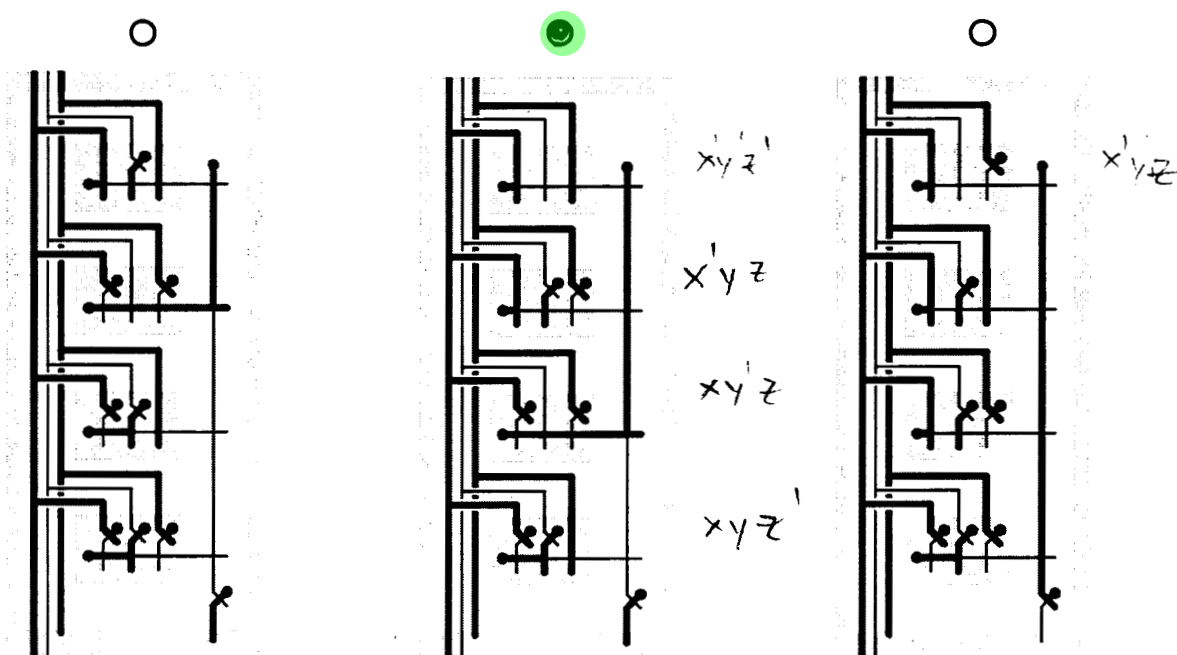
A. (2 points) Fill in the missing entries in this truth table for the 3-variable *IS-XOR?* function.

$x$	$y$	$z$	<i>IS-XOR?</i>
0	0	0	1
0	0	1	0
0	1	0	0
0	1	1	1
1	0	0	0
1	0	1	1
1	1	0	1
1	1	1	0

B. (3 points) In the box, write out the sum-of-products form of the 3-variable *IS-XOR?* function.

$x'y'z' + x'yz + xy'z + xyz'$

C. (2 points) Which of the circuits below is computing *IS-XOR?* for 3 variables with the inputs 1 0 1? In each circuit, assume that the inputs  $x$   $y$   $z$  are provided in that order to the three lines at the upper left and the output is the line at the bottom right. Fill in the circle above the correct circuit. *Correct answers are worth 1 point; incorrect answers are worth -1 point; blank answers are worth 0 points.*



**Q20. Central Processing Unit (7 points).** Consider the following list of CPU components. In the box to the right of each option, write the letter corresponding to the description at right that best matches. Each letter should be used **at most once**. *Correct answers are worth 1 point.*

Bus connection

I

ALU

G

MUX

?

F

IR

H

Control line

C

PC

D

Clock

J

- A. Sequential circuit that carries information.
- B. Combinational circuit that holds data.
- C. Component input that determines behavior.
- D. Holds address of current instruction.
- E. Allows switching among component inputs.
- F. Allows switching among component outputs.
- G. Computes boolean function values.
- H. Holds instruction being executed.
- I. Path carrying data between components.
- J. Sequential circuit that produces periodic pulses.
- K. Memory unit extension.

