

Light at the
end of the tunnel



Course Overview

Mike Freedman
COS 461: Computer Networks

<http://www.cs.princeton.edu/courses/archive/spr14/cos461/>

Course Logistics

- **Last assignment**
 - Due on Dean's Date (11:59pm Tuesday May 13)
- **Final exam**
 - Cumulative, emphasis on second half of the class
 - Wednesday May 21th at 9am-11am
- **Questions?**
 - Ask on piazza
 - Office hours will be posted to piazza

Key Concepts in Networking

(Exam preparation idea: look for other examples)

Some Key Concepts

- **Course was organized around protocols**
 - But a small set of concepts recur in many protocols
- **General CS concepts**
 - Hierarchy, indirection, caching, randomization
- **Networking-specific concepts**
 - Soft state, layering, (de)multiplexing
 - End-to-end argument

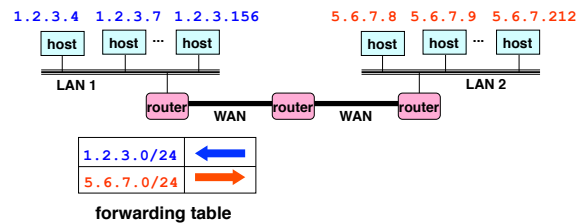
Hierarchy

- Scalability of large systems
 - Cannot store all information everywhere
 - Cannot centrally coordinate everything
- Hierarchy to manage scale
 - Divide system into smaller pieces
- Hierarchy to divide control
 - Decentralized management
- Examples in the Internet
 - IP addresses, routing protocols, DNS, peer-to-peer



Hierarchy: IP Address Blocks

- Number related hosts from a common subnet
 - 1.2.3.0/24 on the left LAN
 - 5.6.7.0/24 on the right LAN

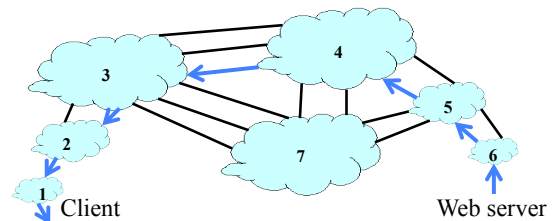


Hierarchy: IP Address Blocks

- Separation of control
 - Prefix: assigned to an institution
 - Addresses: assigned by institution to its nodes
- Who assigns prefixes?
 - Internet Corporation for Assigned Names & Numbers
 - Regional Internet Registries (RIRs)
 - Internet Service Providers (ISPs)
 - Stub networks
 - Regions within an enterprise

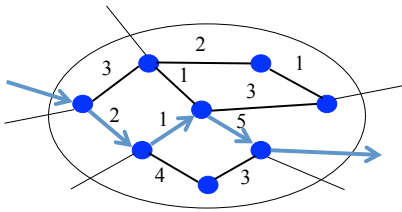
Hierarchy: Routing Protocols

- AS-level topology
 - Nodes are Autonomous Systems (ASes)
 - Edges are links and business relationships
 - Hides the detail within each AS's network



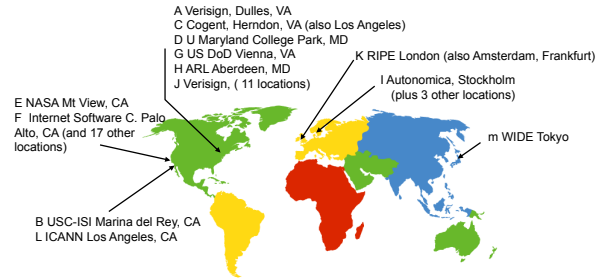
Hierarchy: Routing Protocols

- Interdomain routing ignores details in an AS
 - Routers flood information to learn the topology
 - Routers determine “next hop” to other routers...
 - By computing shortest paths based on link weights

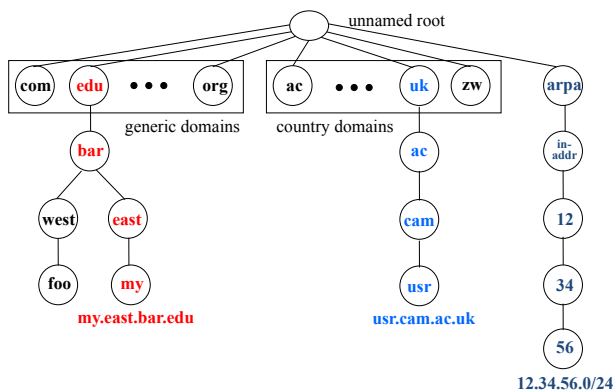


Hierarchy: Domain Name System

- 13 root servers (see <http://www.root-servers.org/>)
- Labeled A through M

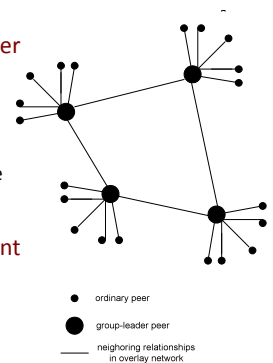


Hierarchy: Domain Name System



Hierarchy: Super Peers in KaZaA

- Each peer is either group leader or assigned to group leader
 - TCP connection between peer and its group leader
 - TCP connections between some pairs of group leaders
- Group leader tracks the content in all its children



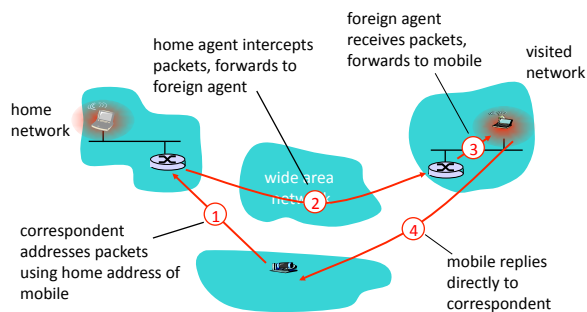
Indirection

- **Referencing by name**
 - Rather than the value itself
 - E.g., manipulating a variable through a pointer
- **Benefits of indirection**
 - Human convenience
 - Reducing overhead when things change
- **Examples of indirection in the Internet**
 - Names vs. addresses
 - Mobile IP

Indirection: Names vs. Addresses

- **Host name to IP address**
 - Mnemonic names to location-dependent addresses
 - E.g., from www.cnn.com to 64.236.16.20
 - Using the Domain Name System (DNS)
- **From IP address to MAC address**
 - From hierarchical global address to interface card
 - E.g., from 64.236.16.20 to 00-15-C5-49-04-A9
 - Using the Address Resolution Protocol (ARP)

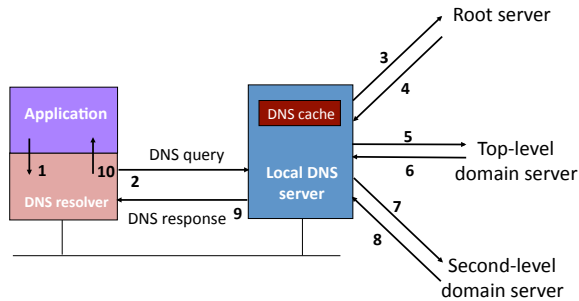
Indirection: Mobile IP



Caching

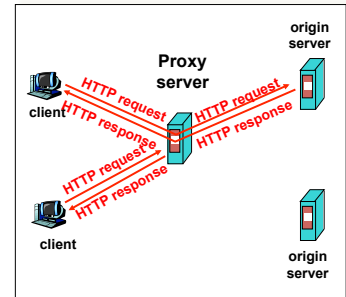
- **Duplicating data stored elsewhere**
 - To reduce latency for accessing the data
 - To reduce resources consumed
- **Caching is often quite effective**
 - Speed difference between cache and primary copy
 - Locality of reference, and small set of popular data
- **Examples from the Internet**
 - DNS caching, Web caching

Caching: DNS Caching



Caching: Web Caching

- **Caching location**
 - Proxy cache
 - Browser cache
- **Better performance**
 - Lower RTT
 - Existing connection
 - Less network load

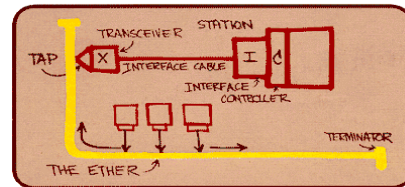


Randomization

- **Distributed adaptive algorithms**
 - Multiple distributed parties
 - Adapting independently
- **Risk of synchronization**
 - Many parties reacting at the same time
 - Leading to bad aggregate behavior
- **Randomization can desynchronize**
 - Ethernet back-off, Random Early Detection
- **Rather than imposing centralized control**

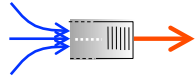
Randomization: Ethernet Back-off

- **Random access: exponential back-off**
 - After collision, wait random time before retrying
 - After m^{th} , choose K randomly from $\{0, \dots, 2^m - 1\}$
 - Wait for $K * 512$ bit times before trying again



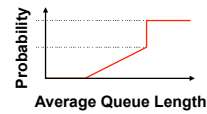
Randomization: Dropping Packets Early

- **Congestion on a link**
 - Eventually the queue becomes full
 - And new packets must be dropped
- **Drop-tail queuing leads to bursty loss**
 - Many packets encounter a full queue
 - Many TCP senders reduce their sending rates



Randomization: Dropping Packets Early

- **Better to give early feedback**
 - Get a few connections to slow down
 - ... before it is too late
- **Random Early Detection (RED)**
 - Randomly drop packets when queue (near) full
 - Drop rate increases as function of queue length



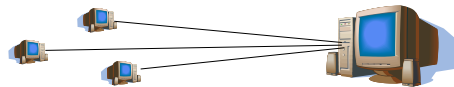
Soft State

- **State: stored in nodes by network protocols**
 - Installed by receiver of a set-up message
 - Updated when conditions change
- **Hard state: valid unless told otherwise**
 - Removed by receiver of tear-down message
 - Requires error handling to deal with sender failure
- **Soft state: invalid if not told to refresh**
 - Periodically refreshed, removed by timeout
- **Soft state reduces complexity**
 - DNS caching, DHCP leases



Soft State: DNS Caching

- **Cache consistency is a hard problem**
 - Ensuring the cached copy is not out of date
- **Strawman: explicit revocation or updates**
 - Keep track of everyone who has cached information
 - If name-to-host mapping changes, update caches
- **Soft state solution**
 - DNS responses include a “time to live” (TTL) field
 - Cached entry is deleted after TTL expires

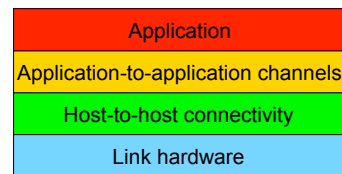


Soft State: DHCP Leases

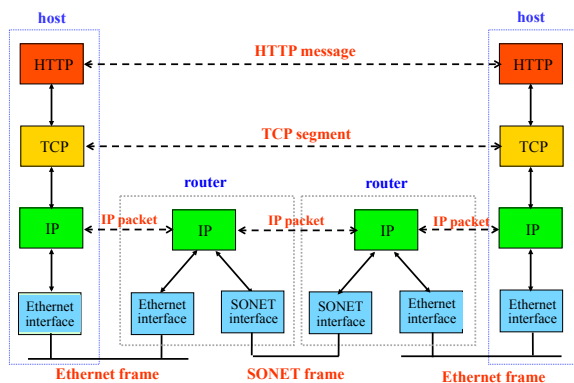
- DHCP “offer message” from the server
 - Configuration parameters (proposed IP address, mask, gateway router, DNS server, ...)
 - Lease time (the time information remains valid)
- Why is a lease time necessary?
 - Client can release address (DHCP RELEASE)
 - E.g., “ipconfig /release” or clean shutdown of computer
 - But, the host might not release the address
 - E.g., the host crashes or buggy client software
 - You don’t want address to be allocated forever

Layering: A Modular Approach

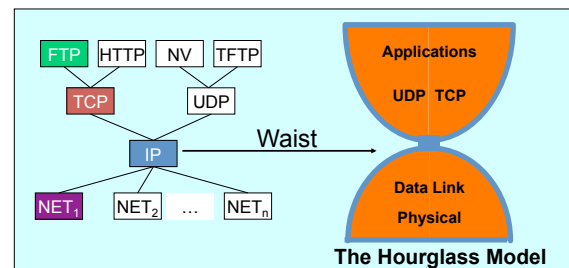
- Sub-divide the problem
 - Each layer relies on services from layer below
 - Each layer exports services to layer above
- Interface between layers defines interaction
 - Hides implementation details
 - Layers can change without disturbing other layers



Layering: Standing on Shoulders



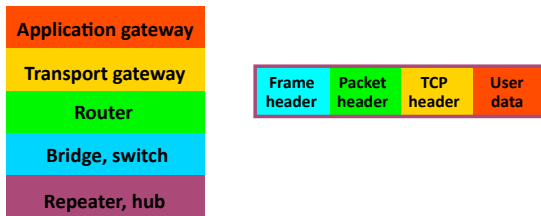
Layering: Internet Protocol Suite



The waist facilitates interoperability

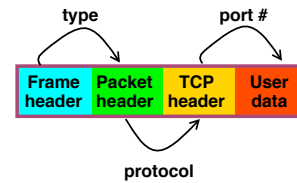
Layering: Encapsulation of Data

- Different devices switch different things
 - Physical layer: electrical signals (repeaters and hubs)
 - Link layer: frames (bridges and switches)
 - Network layer: packets (routers)

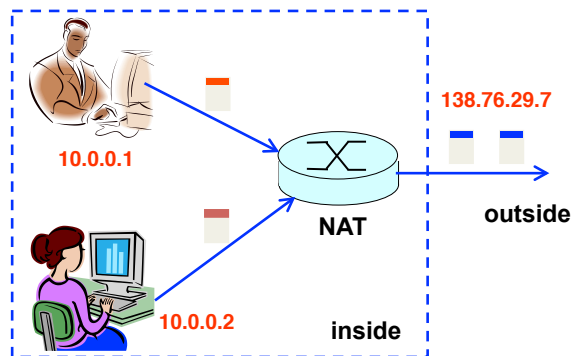


Demultiplexing

- Separating multiple streams out of one
 - Recognizing the separate streams
 - Treating the separate streams accordingly
- Examples in the Internet



(De)multiplexing: With a NAT



Power at the End Host

End-to-End Principle

Whenever possible, communications protocol operations should be defined to occur at the **end-points** of a communications system.

Programmability

With programmable end hosts, new network services can be added at **any time, by anyone**.

Why No Math in This Course?

- **Hypothesis #1: theory not relevant to Internet**
 - Body of math created for telephone networks
 - Many of these models don't work in data networks
- **Hypothesis #2: too many kinds of theory**
 - Queuing: statistical multiplexing works
 - Control: TCP congestion control works
 - Optimization: TCP maximizes aggregate utility
 - Game: reasoning about competing ASes

What Will Happen to the Internet

No Strict Notions of Identity



- **Leads to**
 - Spam
 - Spoofing
 - Denial-of-service
 - Route hijacking

Protocols Designed Based on Trust

- **That you don't spoof your addresses**
 - MAC spoofing, IP address spoofing, spam, ...
- **That port numbers correspond to applications**
 - Rather than being arbitrary, meaningless numbers
- **That you adhere to the protocol**
 - Ethernet exponential back-off after a collision
 - TCP additive increase, multiplicative decrease
- **That protocol specifications are public**
 - So others can build interoperable implementations

Nobody in Charge

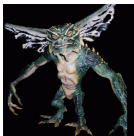
- Traffic traverses many Autonomous Systems
 - Who's fault is it when things go wrong?
 - How do you upgrade functionality?
- Implicit trust in the end host
 - What if some hosts violate congestion control?
- Anyone can add any application
 - Whether or not it is legal, moral, good, etc.
- Spans many countries
 - So no one government can be in charge

Challenging New Requirements

- Disseminating data
- Mobile, multi-homed hosts
- Sometimes-connected hosts
- Large number of hosts
- Real-time applications

The Internet of the Future

- Can we fix what ails the Internet
 - Security, performance, reliability
 - Upgradability, managability
 - <Your favorite gripe here>
- Without throwing out baby with bathwater
 - Ease of adding new hosts
 - Ease of adding new services
 - Ease of adding new link technologies
- An open technical and policy question...



Thank You!