

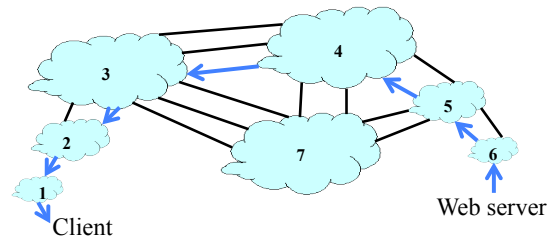
Interdomain Routing Security

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COS 461: Computer Networks

<http://www.cs.princeton.edu/courses/archive/spr14/cos461/>

Interdomain Routing

- **AS-level topology**
 - Nodes are Autonomous Systems (ASes)
 - Edges are links and business relationships



Border Gateway Protocol (BGP)

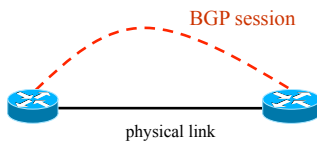
- **ASes exchange reachability information**
 - Destination: block of addresses (an "IP prefix")
 - AS path: sequence of ASes along the path
- **Policies configured by network operators**
 - Path selection: which of the paths to use?
 - Path export: which neighbors to tell?



BGP Session Security

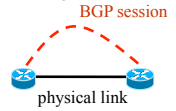
TCP Connection Underlying BGP Session

- **BGP session runs over TCP**
 - TCP connection between neighboring routers
 - BGP messages sent over TCP connection
 - Makes BGP vulnerable to attacks on TCP



Attacks on Session Security

- **Confidentiality**
 - Eavesdropping by tapping the link
 - Inferring routing policies and stability
- **Integrity**
 - Tampering by dropping, modifying, adding packets
 - Changing, filtering, or replaying BGP routes
- **Availability**
 - Resetting the session or congesting the link
 - Disrupting communication and overloading routers



Defending Session Security is Easy

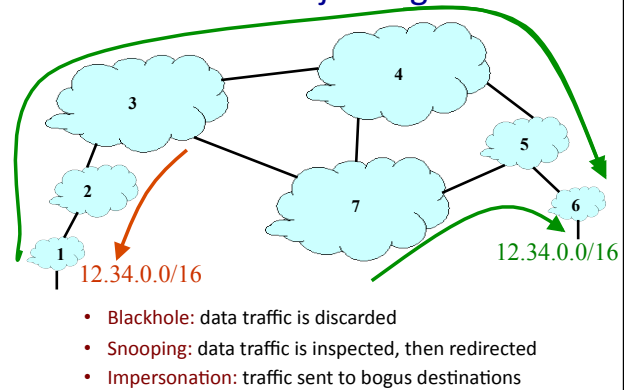
- **BGP routing information is propagated widely**
 - Confidentiality isn't all that important
- **Two end-points have a business relationship**
 - Use known IP addresses and ports to communicate
 - Can agree to sign and encrypt messages
- **Limited physical access to the path**
 - Direct physical link, often in same building
- **Low volume of special traffic**
 - Filter packets from unexpected senders
 - Can give BGP packets higher priority

Validity of the routing information:
Origin authentication

IP Address Ownership and Hijacking

- **IP address block assignment**
 - ICANN -> Regional Internet Registries -> ISPs
- **Proper origination of a prefix into BGP**
 - By the AS who owns the prefix
 - ... or, by its upstream provider(s) in its behalf
- **However, what's to stop someone else?**
 - Prefix hijacking: another AS originates the prefix
 - BGP does not verify that the AS is authorized
 - Registries of prefix ownership are inaccurate

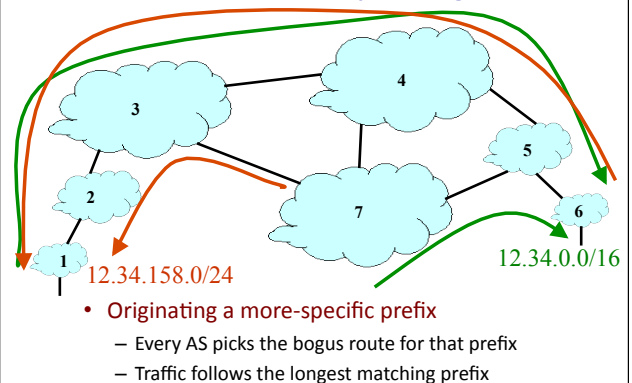
Prefix Hijacking



Hijacking is Hard to Debug

- **The victim AS doesn't see the problem**
 - Picks its own route, might not learn the bogus route
- **May not cause loss of connectivity**
 - Snooping, with minor performance degradation
- **Or, loss of connectivity is isolated**
 - E.g., only for sources in parts of the Internet
- **Diagnosing prefix hijacking**
 - Analyzing updates from many vantage points
 - Launching traceroute from many vantage points

Sub-Prefix Hijacking



How to Hijack a Prefix

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- **The hijacking AS has**
 - Router with BGP session(s)
 - Configured to originate the prefix
- **Getting access to the router**
 - Network operator makes configuration mistake
 - Disgruntled operator launches an attack
 - Outsider breaks in to the router and reconfigures
- **Getting other ASes to believe bogus route**
 - Neighbor ASes do not discard the bogus route
 - E.g., not doing protective filtering

YouTube Outage on Feb 24, 2008

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- **YouTube (AS 36561): 208.65.152.0/22**
- **Pakistan Telecom (AS 17557)**
 - Government order to block access to YouTube
 - Announces 208.65.153.0/24 to PCCW (AS 3491)
 - All packets to YouTube get dropped on the floor
- **Mistakes were made**
 - AS 17557: announce to everyone, not just customers
 - AS 3491: not filtering routes announced by AS 17557
- **Lasted 100 minutes for some, 2 hours for others**

Timeline (UTC Time)

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- 18:47:45: First evidence of hijacked /24 route in Asia
- 18:48:00: Several big trans-Pacific providers carrying route
- 18:49:30: Bogus route fully propagated
- 20:07:25: YouTube advertising /24 to attract traffic back
- 20:08:30: Many (but not all) providers are using valid route
- 20:18:43: YouTube announces two more-specific /25 routes
- 20:19:37: Some more providers start using the /25 routes
- 20:50:59: AS 17557 starts prepending (“3491 17557 17557”)
- 20:59:39: AS 3491 disconnects AS 17557
- 21:00:00: Videos of cats flushing toilets are available again!

Another Example: Spammers

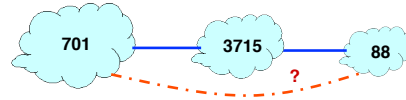
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- **Spammers sending spam**
 - Form a (bidirectional) TCP connection to mail server
 - Send a bunch of spam e-mail, then disconnect
- **But, best not to use your real IP address**
 - Relatively easy to trace back to you
- **Could hijack someone’s address space**
 - But you might not receive all the (TCP) return traffic
- **How to evade detection**
 - Hijack unused (i.e., unallocated) address block
 - Temporarily use the IP addresses to send your spam

BGP AS Path

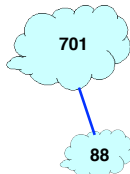
Bogus AS Paths

- Remove ASes from the AS path
 - E.g., turn “701 3715 88” into “701 88”
- Motivations
 - Attract sources that normally try to avoid AS 3715
 - Help AS 88 look like it is closer to the Internet’s core
- Who can tell that this AS path is a lie?
 - Maybe AS 88 *does* connect to AS 701 directly



Bogus AS Paths

- Add ASes to the path
 - E.g., turn “701 88” into “701 3715 88”
- Motivations
 - Trigger loop detection in AS 3715
 - Denial-of-service attack on AS 3715
 - Or, blocking unwanted traffic coming from AS 3715!
 - Make your AS look like it has richer connectivity
- Who can tell the AS path is a lie?
 - AS 3715 could, if it could see the route
 - AS 88 could, but would it really care?



Bogus AS Paths

- Adds AS hop(s) at the end of the path
 - E.g., turns “701 88” into “701 88 3”
- Motivations
 - Evade detection for a bogus route
 - E.g., by adding the legitimate AS to the end
- Hard to tell that the AS path is bogus...
 - Even if other ASes filter based on prefix ownership



Invalid Paths

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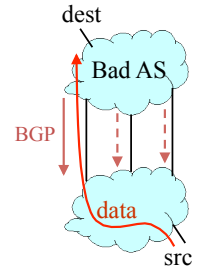
- **AS exports a route it shouldn't**
 - AS path is a valid sequence, but violated policy
- **Example: customer misconfiguration**
 - Exports routes from one provider to another
- **Interacts with provider policy**
 - Provider prefers customer routes
 - Directing all traffic through customer
- **Main defense**
 - Filtering routes based on prefixes and AS path



Missing/Inconsistent Routes

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- **Peers require consistent export**
 - Prefix advertised at all peering points
 - Prefix advertised with same AS path length
- **Reasons for violating the policy**
 - Trick neighbor into “cold potato”
 - Configuration mistake
- **Main defense**
 - Analyzing BGP updates or traffic for signs of inconsistency



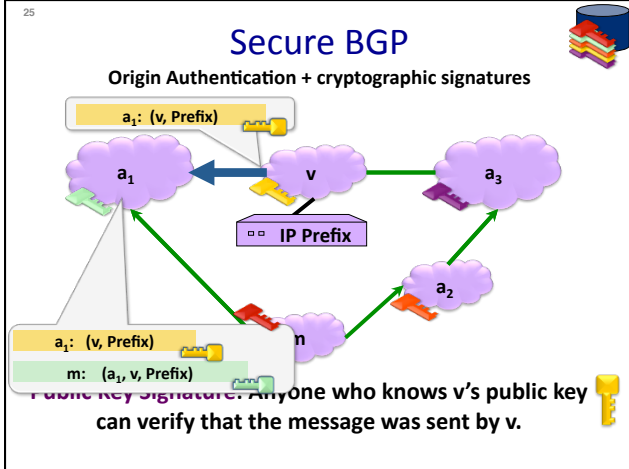
BGP Security Today

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- **Applying “best common practices”**
 - Securing the session (authentication, encryption)
 - Filtering routes by prefix and AS path
 - Packet filters to block unexpected control traffic
- **This is not good enough**
 - Depends on vigilant application of practices
 - Doesn't address fundamental problems
 - Can't tell who owns the IP address block
 - Can't tell if the AS path is bogus or invalid
 - Can't be sure the data packets follow the chosen route

Proposed Enhancements to BGP

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- ## “Secure BGP”
- **Route attestations**
 - Distributed as an attribute in BGP update message
 - Signed by each AS as route traverses the network
 - **Address attestations**
 - Claim the right to originate a prefix
 - Signed and distributed out-of-band
 - Checked through delegation chain from ICANN
 - **S-BGP can validate**
 - AS path indicates the order ASes were traversed
 - No intermediate ASes were added or removed
 - Proper ASes originate prefixes

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- ## S-BGP Deployment Challenges
- **Complete, accurate registries of prefix “owner”**
 - **Public Key Infrastructure**
 - To know the public key for any given AS
 - **Cryptographic operations**
 - E.g., digital signatures on BGP messages
 - **Need to perform operations quickly**
 - To avoid delaying response to routing changes
 - **Difficulty of incremental deployment**
 - Hard to have a “flag day” to deploy S-BGP

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- ## Incrementally Deployable Solutions?
- **Backwards compatible**
 - No changes to router hardware or software
 - No cooperation from other ASes
 - **Incentives for early adopters**
 - Security benefits for ASes that deploy the solution
 - ... and further incentives for others to deploy
 - **What kind of solutions are possible?**
 - Detecting suspicious routes
 - ... and then filtering or depreferencing them

Detecting Suspicious Routes

- **Monitoring BGP update messages**
 - Use past history as an implicit registry
- **E.g., AS that announces each address block**
 - Prefix 18.0.0.0/8 usually originated by AS 3
- **E.g., AS-level edges and paths**
 - Never seen the subpath “7018 88 1785”
- **Out-of-band detection mechanism**
 - Generate reports and alerts
 - Internet Alert Registry: <http://iar.cs.unm.edu/>
 - Prefix Hijack Alert System: <http://phas.netsec.colostate.edu/>

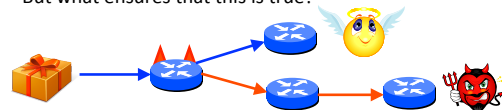
Avoiding Suspicious Routes

- **Soft response to suspicious routes**
 - Prefer routes that agree with the past
 - Delay adoption of unfamiliar routes when possible
- **Why is this good enough?**
 - Some attacks will go away on their own
 - Let someone else be the victim instead of you
 - Give network operators time to investigate
- **How well would it work?**
 - If top ~40 largest ASes applied the technique
 - ... most other ASes are protected, too

What About Packet Forwarding?

Control Plane vs. Data Plane

- **Control plane**
 - BGP security concerns validity of routing messages
 - I.e., did the BGP message follow the sequence of ASes listed in the AS-path attribute
- **Data plane**
 - Routers forward data packets
 - Supposedly along path chosen in the control plane
 - But what ensures that this is true?



Data-Plane Attacks, Part 1

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- **Drop packets in the data plane**
 - While still sending the routing announcements
- **Easier to evade detection**
 - Especially if you only drop some packets
 - Like, oh, say, BitTorrent or Skype traffic
- **Even easier if you just slow down some traffic**
 - How different are normal congestion and an attack?
 - Especially if you let traceroute packets through?

Data-Plane Attacks, Part 2

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- **Send packets in a different direction**
 - Disagreeing with the routing announcements
- **Direct packets to a different destination**
 - E.g., one the adversary controls
- **What to do at that bogus destination?**
 - Impersonate the legitimate destination
 - Snoop on traffic and forward along to real destination
- **How to detect?**
 - Traceroute? Longer than usual delays?
 - End-to-end checks, like site certificate or encryption?

Data-Plane Attacks are Harder

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- **Adversary must control a router along the path**
 - So that the traffic flows through him
- **How to get control a router**
 - Buy access to a compromised router online
 - Guess the password, exploit router vulnerabilities
 - Insider attack (disgruntled network operator)
- **Malice vs. greed**
 - Malice: gain control of someone else's router
 - Greed: say, Verizon DSL blocks Skype to encourage me to use (Verizon) landline phone

What's the Internet to Do?

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BGP is So Vulnerable

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- **Several high-profile outages**
 - <http://merit.edu/mail.archives/nanog/1997-04/msg00380.html>
 - http://www.renesys.com/blog/2005/12/internetwide_nearcatastrophela.shtml
 - http://www.renesys.com/blog/2006/01/coned_steals_the_net.shtml
 - http://www.renesys.com/blog/2008/02/pakistan_hijacks_youtube_1.shtml
 - http://www.theregister.co.uk/2010/04/09/china_bgp_interweb_snafu/
- **Many smaller examples**
 - Blackholing a single destination prefix
 - Hijacking unallocated addresses to send spam
- **Why isn't it an even bigger deal?**
 - Really, most big outages are configuration errors
 - Most bad guys want the Internet to stay up

BGP is So Hard to Fix

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- **Complex system**
 - Large, with around 40,000 ASes
 - Decentralized control among competitive Ases
- **Hard to reach agreement on the right solution**
 - S-BGP with PKI, registries, and crypto?
 - Who should be in charge of running PKI & registries?
 - Worry about data-plane attacks or just control plane?
- **Hard to deploy the solution once you pick it**
 - Hard enough to get ASes to apply route filters
 - Now you want them to upgrade to a new protocol

Conclusions

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- **Internet protocols designed based on trust**
 - Insiders are good guys, bad guys on the outside
- **Border Gateway Protocol is very vulnerable**
 - Glue that holds the Internet together
 - Hard for an AS to locally identify bogus routes
 - Attacks can have very serious global consequences
- **Proposed solutions/approaches**
 - Secure variants of the Border Gateway Protocol
 - Anomaly detection, with automated response
 - Broader focus on data-plane availability