

## Network Security Protocols

Mike Freedman  
COS 461: Computer Networks

<http://www.cs.princeton.edu/courses/archive/spr14/cos461/>

## Network Security

- **Application layer**
  - E-mail: PGP, using a web-of-trust
  - Web: HTTP-S, using a certificate hierarchy
- **Transport layer**
  - Transport Layer Security/ Secure Socket Layer
- **Network layer**
  - IP Sec
- **Network infrastructure**
  - DNS-Sec and BGP-Sec

## Basic Security Properties

- **Confidentiality:**
- **Authenticity:**
- **Integrity:**
- **Availability:**
- **Non-repudiation:**
- **Access control:**

## Basic Security Properties

- **Confidentiality:** Concealment of information or resources
- **Authenticity:** Identification and assurance of origin of info
- **Integrity:** Trustworthiness of data or resources in terms of preventing improper and unauthorized changes
- **Availability:** Ability to use desired information or resource
- **Non-repudiation:** Offer of evidence that a party indeed is sender or a receiver of certain information
- **Access control:** Facilities to determine and enforce who is allowed access to what resources (host, software, network, ...)

## Encryption and MAC/Signatures

### Confidentiality (Encryption)

Sender:

- Compute  $C = \text{Enc}_k(M)$
- Send  $C$

Receiver:

- Recover  $M = \text{Dec}_k(C)$

### Auth/Integrity (MAC / Signature)

Sender:

- Compute  $s = \text{Sig}_k(\text{Hash}(M))$
- Send  $\langle M, s \rangle$

Receiver:

- Compute  $s' = \text{Ver}_k(\text{Hash}(M))$
- Check  $s' == s$

These are simplified forms of the actual algorithms

## Email Security: Pretty Good Privacy (PGP)

## E-Mail Security

- **Security goals**
  - Confidentiality: only intended recipient sees data
  - Integrity: data cannot be modified en route
  - Authenticity: sender and recipient are who they say
- **Security non-goals**
  - Timely or successful message delivery
  - Avoiding duplicate (replayed) message
  - (Since e-mail doesn't provide this anyway!)

## Sender and Receiver Keys

- If the sender knows the receiver's public key
  - Confidentiality
  - Receiver authentication
- If the receiver knows the sender's public key
  - Sender authentication
  - Sender non-repudiation



## Sending an E-Mail Securely

- **Sender digitally signs the message**
  - Using the sender's private key
- **Sender encrypts the data**
  - Using a one-time session key
  - Sending the session key, encrypted with the receiver's public key
- **Sender converts to an ASCII format**
  - Converting the message to base64 encoding
  - (Email messages must be sent in ASCII)

## Public Key Certificate

- **Binding between identity and a public key**
  - "Identity" is, for example, an e-mail address
  - "Binding" ensured using a digital signature
- **Contents of a certificate**
  - Identity of the entity being certified
  - Public key of the entity being certified
  - Identity of the signer
  - Digital signature
  - Digital signature algorithm id



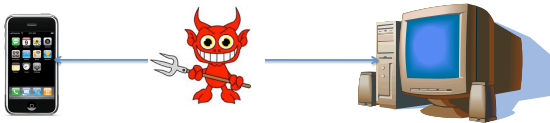
## Web of Trust for PGP

- **Decentralized solution**
  - Protection against government intrusion
  - No central certificate authorities
- **Customized solution**
  - Individual decides whom to trust, and how much
  - Multiple certificates with different confidence levels
- **Key-signing parties!**
  - Collect and provide public keys in person
  - Sign other's keys, and get your key signed by others

## HTTP Security

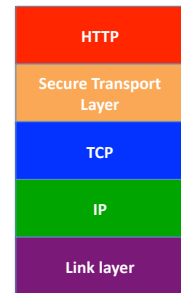
## HTTP Threat Model

- **Eavesdropper**
  - Listening on conversation (confidentiality)
- **Man-in-the-middle**
  - Modifying content (integrity)
- **Impersonation**
  - Bogus website (authentication, confidentiality)



## HTTP-S: Securing HTTP

- HTTP sits on top of secure channel (SSL/TLS)
  - https:// vs. http://
  - TCP port 443 vs. 80
- All (HTTP) bytes encrypted and authenticated
  - No change to HTTP itself!
- Where to get the key???



## Learning a Valid Public Key



- **What is that lock?**
  - Securely binds domain name to public key (PK)
    - If PK is authenticated, then any message signed by that PK cannot be forged by non-authorized party
  - Believable only if you trust the attesting body
    - Bootstrapping problem: Who to trust, and how to tell if this message is actually from them?

## Hierarchical Public Key Infrastructure

- **Public key certificate**
  - Binding between identity and a public key
  - “Identity” is, for example, a domain name
  - Digital signature to ensure integrity
- **Certificate authority**
  - Issues public key certificates and verifies identities
  - Trusted parties (e.g., VeriSign, GoDaddy, Comodo)
  - Preconfigured certificates in Web browsers

## Public Key Certificate

**Public Key Certificate**

This certificate has been verified for the following uses:  
SSL Server Certificate

**Issued To**  
Common Name (CN) www.wellsfargo.com  
Organization (O) Wells Fargo and Company  
Organizational Unit (OU) ISG  
Serial Number 41:C5:CD:90:95:3C:A1:48:C1:8A...

**Issued By**  
Common Name (CN) <Not Part Of Certificate>  
Organization (O) VeriSign Trust Network  
Organizational Unit (OU) VeriSign, Inc.

**Validity**  
Issued On 5/12/10  
Expires On 5/13/11

**Fingerprints**  
SHA1 Fingerprint C5:EC:18:24:50:9D:90:93:96:69:5...  
MD5 Fingerprint 1C:51:99:C9:EA:78:FB:64:3F:92:F...

**Certificate Hierarchy**  
Builtin Object Token-VeriSign Class 3 Public Primary Certifi...  
VeriSign, Inc.  
www.wellsfargo.com

**Certificate Fields**  
Not After  
Subject  
Subject Public Key Info  
Subject Public Key Algorithm  
Subject's Public Key

**Extensions**  
Certificate Basic Constraints  
Certificate Key Usage  
CRL Distribution Points

**Field Value**  
Modulus (1024 bits):  
c9 b3 f9 c0 4a 42 be 1a c4 0a a0 b5 e0 9c 79 89  
52 82 b1 99 b3 92 de 2d 03 2b 1e 77 c3 4c 7d 97  
37 62 c6 7b 31 b5 6b 25 d3 9e 7e 7e 07 95 7e f6  
ab 6a 5e 88 ec 27 9d 72 3e a0 80 0c a5 ea d4 ff

## Transport Layer Security (TLS)

Based on the earlier Secure Socket Layer (SSL) originally developed by Netscape

## TLS Handshake Protocol

- Send new random value, list of supported ciphers
- Send pre-secret, encrypted under PK
- Create shared secret key from pre-secret and random
- Switch to new symmetric-key cipher using shared key

- Send new random value, digital certificate with PK
- Create shared secret key from pre-secret and random
- Switch to new symmetric-key cipher using shared key

## TLS Record Protocol

- Messages from application layer are:
  - Fragmented or coalesced into blocks
  - Optionally compressed
  - Integrity-protected using an HMAC
  - Encrypted using symmetric-key cipher
  - Passed to the transport layer (usually TCP)
- Sequence #s on record-protocol messages
  - Prevents replays and reorderings of messages

## Comments on HTTPS

21

- **HTTPS authenticates server, not content**
  - If CDN (Akamai) serves content over HTTPS, customer must trust Akamai not to change content
- **Symmetric-key crypto after public-key ops**
  - Handshake protocol using public key crypto
  - Symmetric-key crypto much faster (100-1000x)
- **HTTPS on top of TCP, so reliable byte stream**
  - Can leverage fact that transmission is reliable to ensure: each data segment received exactly once
  - Adversary can't successfully drop or replay packets

## IP Security

22

## IP Security

23

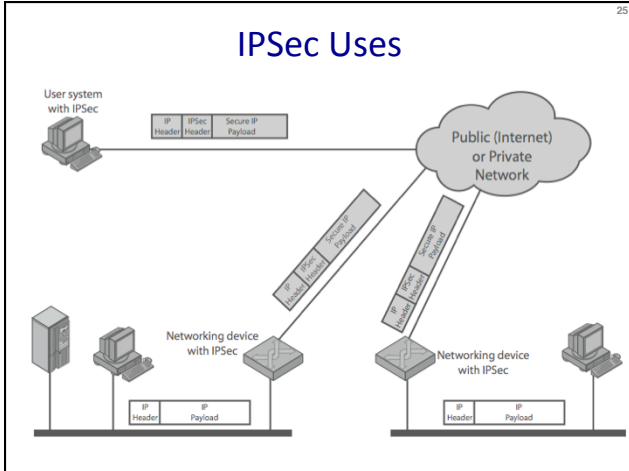
- There are range of app-specific security mechanisms
  - eg. TLS/HTTPS, S/MIME, PGP, Kerberos, ...
- But security concerns that cut across protocol layers
- Implement by the network for all applications?

Enter IPSec!

## IPSec

24

- **General IP Security framework**
- **Allows one to provide**
  - Access control, integrity, authentication, originality, and confidentiality
- **Applicable to different settings**
  - Narrow streams: Specific TCP connections
  - Wide streams: All packets between two gateways



- Benefits of IPSec** 26
- **If in a firewall/router:**
    - Strong security to all traffic crossing perimeter
    - Resistant to bypass
  - **Below transport layer**
    - Transparent to applications
    - Can be transparent to end users
  - **Can provide security for individual users**

- IP Security Architecture** 27
- **Specification quite complex**
    - Mandatory in IPv6, optional in IPv4
  - **Two security header extensions:**
    - **Authentication Header (AH)**
      - Connectionless integrity, origin authentication
        - MAC over most header fields and packet body
      - Anti-replay protection
    - **Encapsulating Security Payload (ESP)**
      - These properties, plus confidentiality

- Encapsulating Security Payload (ESP)** 28
- **Transport mode: Data encrypted, but not header**
    - After all, network headers needed for routing!
    - Can still do traffic analysis, but is efficient
    - Good for host-to-host traffic
  - **Tunnel mode: Encrypts entire IP packet**
    - Add new header for next hop
    - Good for VPNs, gateway-to-gateway security

## Replay Protection is Hard

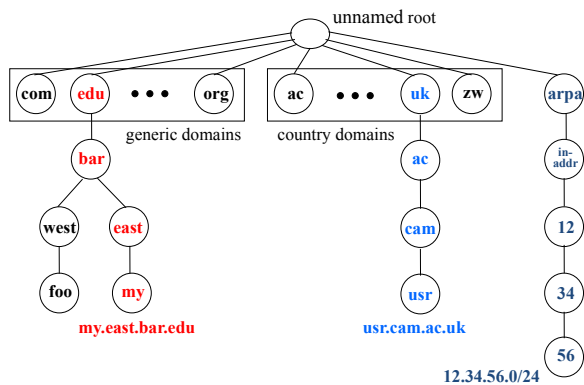
- **Goal: Eavesdropper can't capture encrypted packet and duplicate later**
  - Easy with TLS/HTTP on TCP: Reliable byte stream
  - But IP Sec at packet layer; transport may not be reliable
- **IP Sec solution: Sliding window on sequence #'s**
  - All IPsec packets have a 64-bit monotonic sequence number
  - Receiver keeps track of which seqno's seen before
    - [latest – windowsize + 1, latest]; windowsize typically 64 packets
  - Accept packet if
    - seqno > latest (and update latest)
    - Within window but has not been seen before
  - If reliable, could just remember last, and accept iff last + 1

29

## DNS Security

30

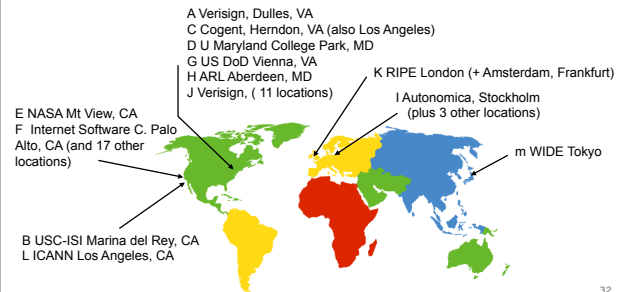
## Hierarchical Naming in DNS



31

## DNS Root Servers

- 13 root servers (see <http://www.root-servers.org/>)
- Labeled A through M



32



## DoS attacks on DNS Availability

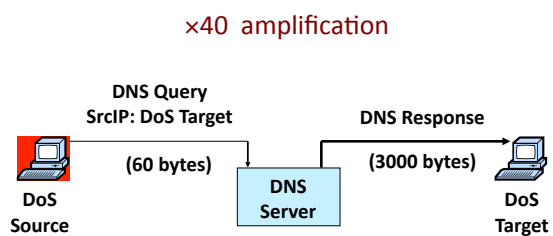
- Feb. 6, 2007
  - Botnet attack on the 13 Internet DNS root servers
  - Lasted 2.5 hours
  - None crashed, but two performed badly:
    - g-root (DoD), I-root (ICANN)
    - Most other root servers use anycast

## Defense: Replication and Caching

Letter	Old name	Operator	Location
A	ns.internic.net	VeriSign	Dulles, Virginia, USA
B	ns1.isi.edu	ISI	Marina Del Rey, California, USA
C	c.psi.net	Cogent Communications	distributed using anycast
D	terp.umd.edu	University of Maryland	College Park, Maryland, USA
E	ns.nasa.gov	NASA	Mountain View, California, USA
F	ns.isc.org	ISC	distributed using anycast
G	ns.nic.ddn.mil	U.S. DoD NIC	Columbus, Ohio, USA
H	aos.arl.army.mil	U.S. Army Research Lab	Aberdeen Proving Ground, Maryland, USA
I	nic.nordu.net	Autonomica	distributed using anycast
J		VeriSign	distributed using anycast
K		RIPE NCC	distributed using anycast
L		ICANN	Los Angeles, California, USA
M		WIDE Project	distributed using anycast

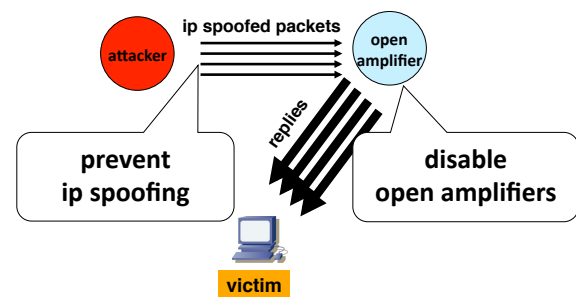
source: wikipedia

## Denial-of-Service Attacks on Hosts



580,000 open resolvers on Internet (Kaminsky-Shiffman'06)

## Preventing Amplification Attacks



## DNS Integrity and the TLD Operators

- If domain name doesn't exist, DNS should return NXDOMAIN (non-existent domain) msg
- Verisign instead creates wildcard records for all [.com](#) and [.net](#) names not yet registered
  - September 15 – October 4, 2003
- Redirection for these domain names to Verisign web portal: "to help you search"
  - And serve you ads...and get "sponsored" search
  - Verisign and online advertising companies make \$\$

## DNS Integrity: Cache Poisoning

- Was answer from an authoritative server?
  - Or from somebody else?
- DNS cache poisoning
  - Client asks for www.evil.com
  - Nameserver authoritative for www.evil.com returns additional section for (www.cnn.com, 1.2.3.4, A)
  - Thanks! I won't bother check what I asked for

## DNS Integrity: DNS Hijacking

- To prevent cache poisoning, client remembers:
  - The domain name in the request
  - A 16-bit request ID (used to demux UDP response)
- DNS hijacking
  - 16 bits: 65K possible IDs
  - What rate to enumerate all in 1 sec? 64B/packet
  - $64 * 65536 * 8 / 1024 / 1024 = 32$  Mbps
- Prevention: also randomize DNS source port
  - Kaminsky attack: this source port... wasn't random

<http://unixwiz.net/techtips/iguide-kaminsky-dns-vuln.html>

## Let's strongly believe the answer! Enter DNSSEC

- DNSSEC protects against data spoofing and corruption
- DNSSEC also provides mechanisms to authenticate servers and requests
- DNSSEC provides mechanisms to establish authenticity and integrity

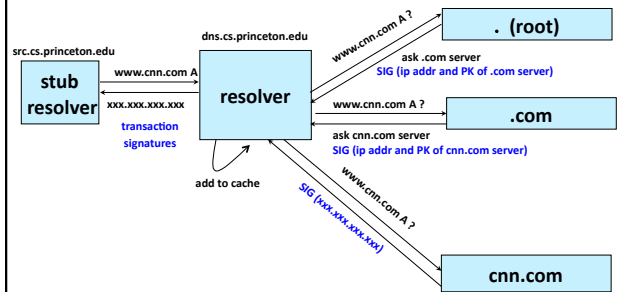
## PK-DNSSEC (Public Key)

- The DNS servers sign the hash of resource record set with its private (signature) keys
  - Public keys can be used to verify the SIGs
- Leverages hierarchy:
  - Authenticity of name server's public keys is established by a signature over the keys by the parent's private key
  - In ideal case, only roots' public keys need to be distributed out-of-band

41

## Verifying the Tree

Question: **www.cnn.com** ?



42

## Conclusions

- Security at many layers
  - Application, transport, and network layers
  - Customized to the properties and requirements
- Exchanging keys
  - Public key certificates
  - Certificate authorities vs. Web of trust
- Next time
  - Interdomain routing security
- Learn more: take COS 432 in the fall!

43