Algorithms

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Flipped learning

Metacognition

- Thinking about how you think!
- Blackboard and PollEverywhere questions test only basic comprehension.





pollEv.com/jhug	text to 37607
n=floor(lg N), what positio	ns are valid for the key 2?
N. 1,, N-2	[668544]
3. h,, N	[688545]
. floor(N/2),, N	[688546]
). ceil(N/2),, N	[688547]
Q: Given a heap with N eleme	ents, how many valid orderings are there?







Seeing the forest from the heap

Work in groups of 2 or 3 (no more!).

- 1. What is the difference between a priority queue and a heap?
- 2. Give a specific example (real world or fantastical) of a situation where a heap would not be the best way to implement a priority queue. How would you implement the PQ?
- 3. What tasks utilize the sink method?
- 4. What tasks utilize the swim method?
- 5. How would you implement a MaxPQ that also has a constant time min() method?
- 6. Bonus question: If you used stacks instead of a heap, what is the minimum number of stacks you'd need to implement a priority queue? What are the run times of your methods? Would having more stacks improve run time?

Debriefing

15

What is the difference between a priority queue and a heap?

• A heap as an efficient implementation of a priority queue.

16

- Priority queue is an abstract data type
- Heap is a data structure

In 5 minutes we will 'debrief'.

Debriefing

Give a specific example (real world or fantastical) of a situation where a heap would not be the best way to implement a priority queue?

- a. 3 way sort (or more generally to support a different PQ-sort)
- b. List of items is known to be provided in sorted order
- Fantastical!
- c. Only two distinct key values (more generally, only k distinct keys)
- d. When you're worried about cache performnace
- e. If you want to track only say top 10 elements
 - But still ok to use heap (though an array is just fine for N=10)
- Find average item (median)
 - Heap is still the way to go (see end of slides)
- If you want to use a PQ for sorting AND want stability
- Canned answer we had in mind #1: If almost everything was insert, and almost never ask for the max: Maintained PQ as unsorted array
- Canned answer #2r: If almost every operation was get max, very few inserts: Maintain as sorted array the whole time

Debriefing

What tasks utilize the sink method?

- deleteMax
 - Heapsort
- If you change a key value (not allowed by our API, but if keys mutable)
- Top down heapification (swimming every item starting from the leftmost item -- bad! N lg N!)

What tasks utilize the swim method?

- Inserting an element
- Heapification (bottom up heap construction) (sinking every item starting from the right most item -- good! N)

Debriefing

How would you implement a MaxPQ that also has a constant time min() method?

- One way: maintain total order at all times (slower insert, but constant min() as required by problem)
- To maintain logarithmic insert(), keep an instance variable that tracks **min**
 - When you delete or insert have to make sure min is correct
 - Follow up question after class: How do you actually do that?
 - insert(): Check and see if the insert item is less than the stored minimum. If so, replace it.
 - delete():Exercise for the reader (hint: it's trivially easy!)

Debriefing

17

19

If you used stacks instead of a heap, what is the minimum number of stacks you'd need to implement a priority queue?

- 2 stacks
- Insert: Add to one of your stacks
- deleteMax: pop everything off, and track the biggest thing you see -push to the other stack as you go (find the diamond in stack of pancakes)

What are the run times of your methods?

- delete: linear
- insert: constant

Would having more stacks improve run time?

- Even more griddles
- API can only see one pancake at a time
- Arvind's crazy bonus answer: N stacks, put in heap (but we said stacks instead of a heap so Arvind is breaking the rules)



Big picture mini-lecture

Heapsort

- Given PQ, you can trivially sort N items: Insert them all, then delete them.
- With a max heap, you can sort in place!

Basic Idea

- Given arbitrary array (i.e. not a heap):
- Max-heapify the array (using sink and/or swim).
- Delete max items one by one (thus moving max to end of array).
- Items take a round trip (but across a logarithmic space).

Heapsort has sometimes been described as the " \oint " algorithm, because of the motion of l and r. The upper triangle represents the heap creation phase, when r = N and l decreases to 1; and the lower triangle represents the selection phase, when l = 1 and r decreases to 1.

Donald Knuth - The Art of Computer Programming Volume 3

22

Modern Heapsort

• Invented by Bob "W." Floyd (was best buddies with Donald Knuth).











Heap Construction

Quicksort

- Partition once on every item (pivot)
 - Choice 1: How do you select order of pivots?
 - Choice 2: What partitioning algorithm will you use?

Heap Construction (Heapification)

- Sink or swim items in array
 - Choice 1: How do you select order of items?
 - Choice 2: How do you decide when to swim or sink?

Bottom-up Heapification

- Choice 1: Always rightmost element
- Choice 2: Always swim
- Fewer than 2N compares and N exchanges!











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Heapsort demo

















Sink-based (bottom up) heapification

Observation

- Given two heaps of height 1.
- A heap of height 2 results by:
 - Pointing the root of each heap at a new item.
 - Sinking that new item.
- Cost: 4 compares (2 * height of new tree).

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Sink-based (bottom up) heapification

Observation

- Given two heaps of height h-1.
- A heap of height h results by
 - Pointing the root of each heap at a new item.
 - Sinking that new item.
- Cost to sink: At most 2h compares.
- Total heap construction cost: 4*2 + 2*4 + 6 = 22 compares



Sink-based (bottom up) heapification

Total Heap Construction Cost

- For h=1: C₁ = 2
- For h=2: C₂ = 2C₁ + 2*2
- For h: $C_h = 2C_{h-1} + 2h$
- Total cost: Doubles with h (plus a small constant factor): Exponential in h
- Total cost: Linear in N



Heapsort

Order of growth of running time

- Heap construction: N
- N calls to delete max: N lg N

What does your computer look like inside?

Total Extra Space

• Constant (in-place)



Heapsort summary

The good news:

• Heap sort: In place and theoretically fast (not in place)



The bad news:

- (Almost) nobody uses Heapsort in the real world. Why?
- Like Miss Manners, Heapsort is very well-behaved, but is unable to handle the stresses of the real world
- In particular, performance on real computers is heavily impacted by really messy factors like cache performance

Play with it!

51



52

Levels of caches

	Processor package	e
CPU	L2 cache	Memory
L1 cache 11010011 11011101	address value FFFF 11010011 C032 11011101 0167 10010110 0000 11100001	address value FFFF 11010111 0002 10110100 0001 11100001 0000 01001101

We'll assume there's just one cache, to keep things simple

That's bad enough...

Key idea behind caching



When fetching one memory address, fetch everything nearby

Because memory access patterns of most programs/algorithms are highly localized!

Which of these is faster?

A. sum=0 for (i = 0 to size) for (j = 0 to size) sum += array[i][j]

B. sum=0 for (i = 0 to size) for (j = 0 to size) sum += array[j][i]

Answer: A is faster, sometimes by an order of magnitude or more.

Cache and memory latencies: an analogy

Cache Get up and get something from the kitchen



RAM Walk down the block to borrow from neighbor

Hard drive Drive around the world...

...twice

53

55





Sort algorithms and cache performanceMergesort: sort subarrays firstQuicksort: partition into subarraysImage: the the placeImage: the the place



Another real world issue

Work in groups of 3.

• What is the primary implementation issue that would affect the real world usability of the MaxPQ class?

```
private class MaxPQ<Key extends<Comparable<Key>> {
   public MaxPQ(int maxN)
   public boolean isEmpty()
   public int size()
   public void insert(Key v)
   public key delMax()
}
```

- Give two distinct solutions for handling this issue in plain English.
 - What is the WORST case run time of each solution?
 - What is the AMORTIZED run time of each solution?

Debriefing.

What is the primary real world issue that would affect the real world usability of the MaxPQ class?

60

• Size is fixed!!

Another real world issue

Give two distinct solutions for handling this issue in plain English.

- Resizing array
- Triply linked list: Each node has a parent and two chidren

What is the WORST case run time of each solution?

- Insert for resizing array: Linear [copy the whole array]
- Insert for Triple: Logarithmic

What is the AMORTIZED run time of each solution?

- Insert (amortized): AMORTIZED LG!!
- Insert (amortized TLL): Logarithmic [no amortization about it]

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64

rule identifying your sort(s). Don't share yet.

A foray into crowd sourcing

Go to the web address: www.reddit.com/r/226

- Create an account (if you don't have one) takes 30 seconds (no email verification!)
- Post your sort identification heuristics.
- Upvote your favorites.
- Downvote your least favorites (or incorrect ones).

This has (almost certainly) never been tried.

• Hopefully it works!

Implementing new abstract data types

Randomized Priority Queue.

 Describe how you would implement the sample() and delRandom() methods in the following class.

```
private class MinPQ<Key extends<Comparable<Key>> {
   public MinPQ()
   public void insert(Key key)
   Key min()
   Key delMin()
   Key sample() // return random item, constant time
   Key delRandom() // del random item, logarithmic time
}
```

In Groups:

- Design a sample() and delRandom() method.
- Between groups:
 - Compare and critique designs.

Repurposing existing data structures

Work in groups of 3.

· Consider the class below which tracks the median of inserted items.



• Solo (1 minute):

- How would you trivially implement each method if given linear time for each operation?
- In constant for insertItem(), but average linear for the other two operations?
- Group (5 minutes): Compare solo answers. Devise an algorithm for solving the problem.
 - · Hint: How would you track the 2nd largest?
- Between groups: Compare solutions.

The order of the day

Summary

65

67

- The priority queue is a very powerful abstract datatype
 - Can process items in order without storing in sorted order (and without even storing everything at once! [online algorithms])
 - Can be used as a sorting algorithm
- · The heap is the best data structure for almost any PQ
 - Heap-based PQ leads to heapsort
- Heapsort
 - Theoretically important: Optimal bounds, in-place, but non-stable
 - Infrequently used in practice (particularly due to caching)
- · Real world heaps
 - Resizing array: Amortized logarithmic time
 - Linked list: Prevents sampling, uses more memory
- Reddit sorting experiment
 - Patterns invoked by various sorts
- · Randomized PQ