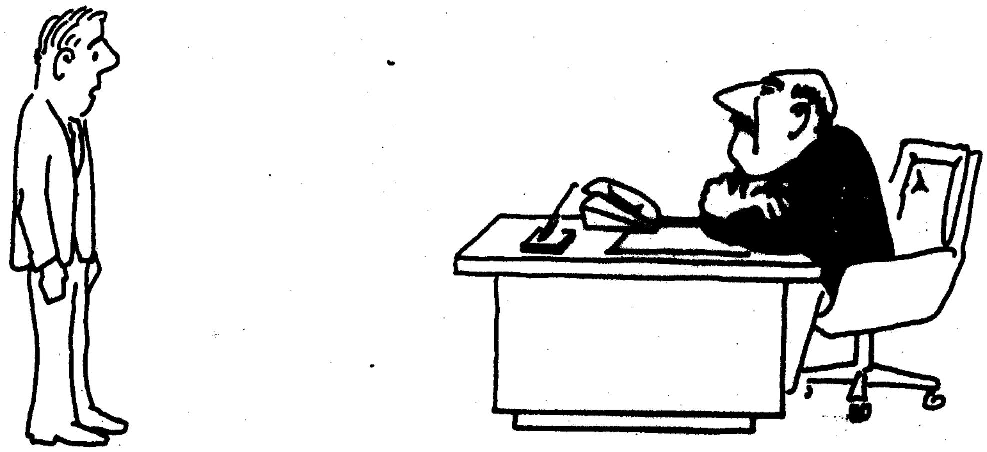
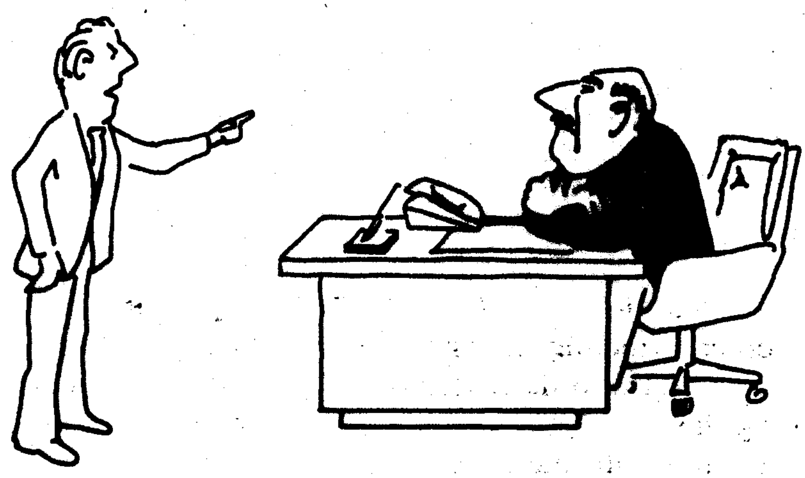
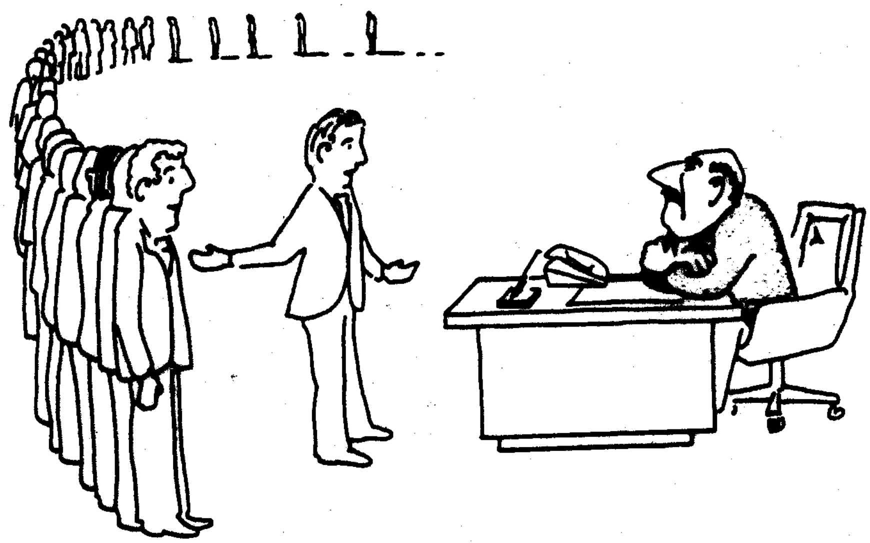


"I can't find an efficient algorithm, I guess I'm just too dumb."



"I can't find an efficient algorithm, because no such algorithm is possible!"



"I can't find an efficient algorithm, but neither can all these famous people."

P = problems solvable in p-time

NP = yes-no problems s.t. if answer is "yes", can be verified in p-time given a (p-length) "proof" (hint).

p-time on a Turing machine
or random-access machine

Sat $\in NP$: proof = satisfying assignment

Graph coloring $\in NP$: proof = coloring

k-clique, k-vertex cover, k-independent set $\in NP$

Tautology $\in$ co-NP ("no" instances have a proof)

NP complete:

1) in NP

2) <u>every</u> problem in NP reducible to it.

Transitivity of p-time reduction implies

NP complete iff

1) in NP

2') <u>some</u> NP-complete problem reducible to it.

We need one NP-complete problem to get started.

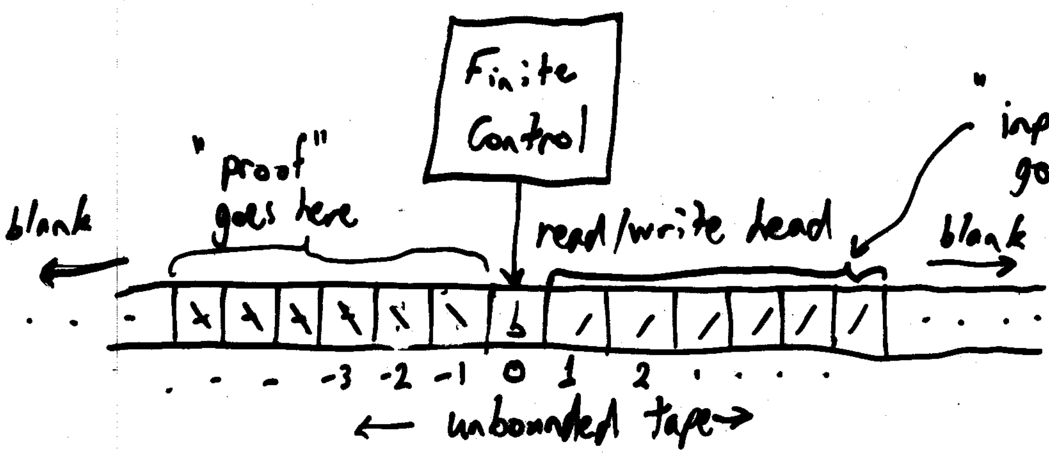
Cook-Levin Theorem: Sat is NP-complete.

Given any p-time verifier, construct (in p-time)

an instance of Sat s.t. verifier answers "yes"

iff formula is satisfiable.

Verifier: Turing Machine



In one step, machine can write a symbol, move head one position, change state.

What to do is based on state, symbol read.

Fixed # of states, fixed # of tape symbols, including blank; start state, "yes" state, ("no" state)

Explicitly given polynomial time bound $p(n)$.

Input (of size n) is a "yes" instance iff

for some "proof" and given input, the machine

reaches "yes" state within $p(n)$ steps from

start state.

Must construct a formula that is satisfiable

iff this happens.

Note: input is specified, proof is not (nondeterministic

part)

Proof can't exceed length $p(n)$: machine can't get

farther in $p(n)$ steps.

Can assume machine loops in "yes" state: if

ever in "yes", will be in "yes" at step $p(n)$.

States: $1, \ldots, y$ $\quad$ 1 = start, y = yes

Symbols: $1, \ldots, z$ $\quad$ 1 = blank

Tape cells, $-p(n), \ldots, 0, \ldots, p(n)$

Time: $0, 1, \ldots, p(n)$

Variables for formula:

h_{it}: true if head on tape cell i at time t

$$-p(n) \leq i \leq p(n), \quad 0 \leq t \leq p(n)$$

s_{jt}: true if state j at time t

$$1 \leq j \leq y, \quad 0 \leq t \leq p(n)$$

c_{ikt}: true if tape cell i holds symbol k at time t

$$-p(n) \leq i \leq p(n), \; 1 \leq k \leq z, \; 0 \leq t \leq p(n)$$

What does the formula need to say?

At most one state, head position, and symbol per cell at each time:

$(\bar{h}_{it} \vee \bar{h}_{i't})$ $i \neq i'$, all t

$(\bar{s}_{jt} \vee \bar{s}_{j't})$ $j \neq j'$, all t

$(\bar{c}_{ikt} \vee \bar{c}_{ik't})$ $k \neq k'$, all i, all t

Correct initial state, head position, and tape contents:

$$h_{00} \wedge s_{10} \wedge c_{010} \wedge c_{1k_1 0} \wedge c_{2k_2 0} \wedge \cdots \wedge c_{nk_n 0}$$

$$\wedge\, c_{(n+1)10} \wedge \cdots \wedge c_{p(n)10}$$

Input is $k_1 k_2 \cdots k_n$, rest of right side of tape is blank

Correct final state: $s_{y\,p(n)}$

Correct transitions:

E.g. if machine in state j reads k, it then writes k', moves head right, and changes to state j':

$$s_{jt} \wedge h_{it} \wedge c_{ikt} \supset s_{j't+1} \wedge h_{i+1\,t+1} \wedge c_{ik't+1}$$

($\supset$ = "implies") (for each i,t)

$$h_{it} \wedge c_{i'kt} \supset c_{i'kt+1} \text{ (for } i \neq i', \text{ each } k, t)$$

(unread tape cells are unaffected)

CNF?

$$(x \wedge y \wedge z) \supset (a \wedge b \wedge c)$$

$$\Rightarrow$$

$$((x \wedge y \wedge z) \supset a)$$
$$\wedge$$
$$((x \wedge y \wedge z) \supset b)$$
$$\wedge$$
$$((x \wedge y \wedge z) \supset c)$$

$$\Rightarrow$$

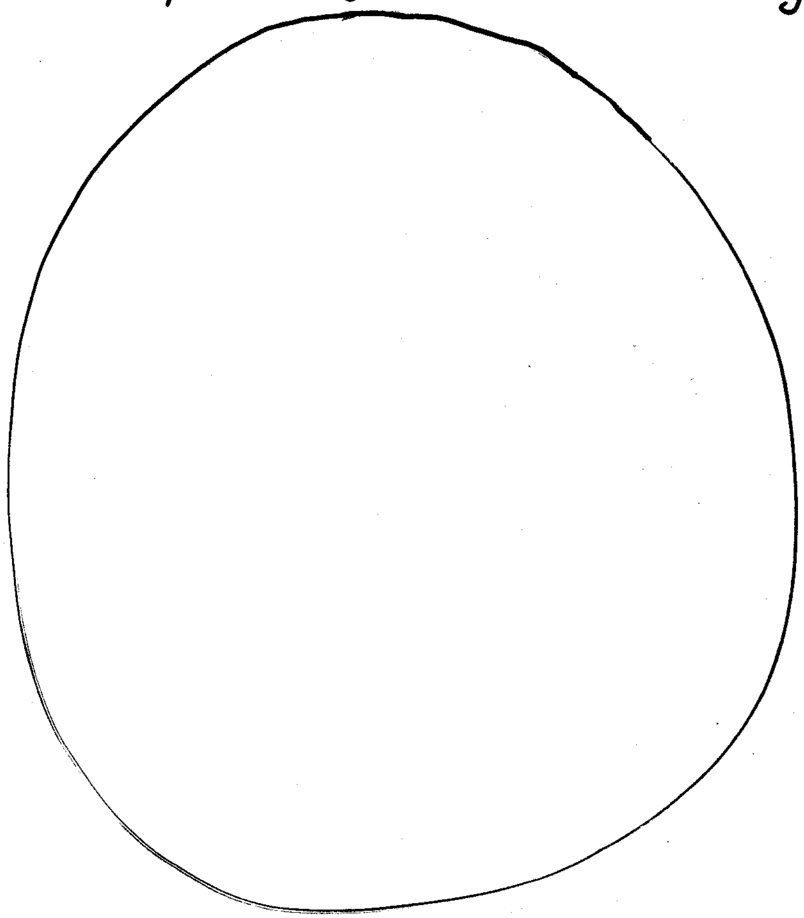
$$(\bar{x} \vee \bar{y} \vee \bar{z} \vee a)$$
$$\wedge$$
$$(\bar{x} \vee \bar{y} \vee \bar{z} \vee b)$$
$$\wedge$$
$$(\bar{x} \vee \bar{y} \vee \bar{z} \vee c)$$

Any proof that gives a "yes" execution gives a satisfying assignment, and vice-versa.

Conclusion: SAT is NP-complete

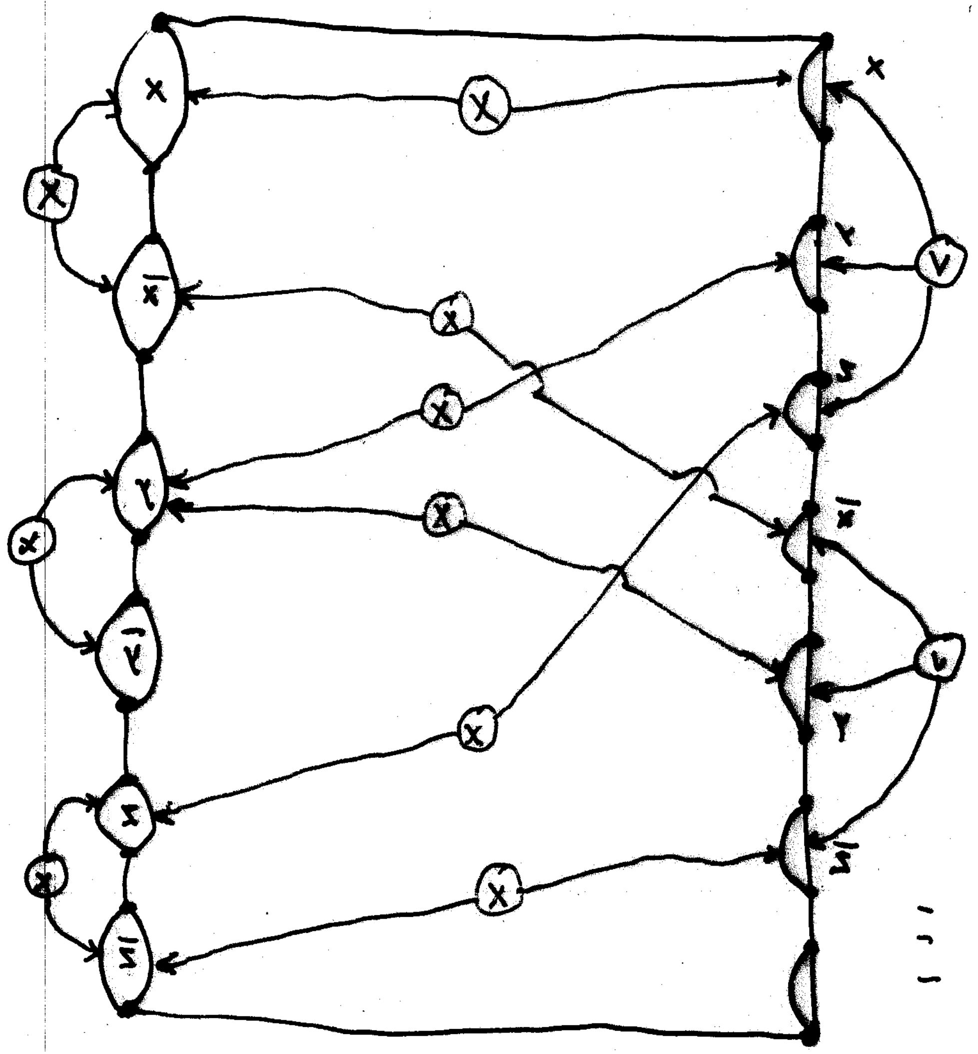
(and k-coloring, k-clique, k-independent set, k-vertex cover).

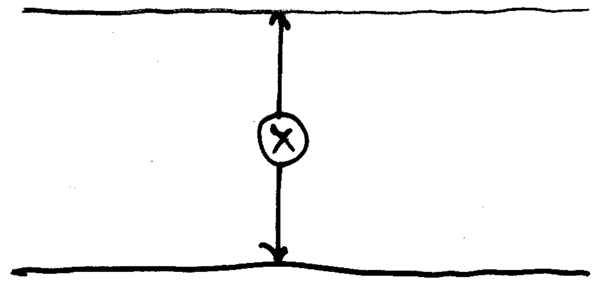
Hamiltonian Cycle: find a simple cycle through all vertices of a graph.

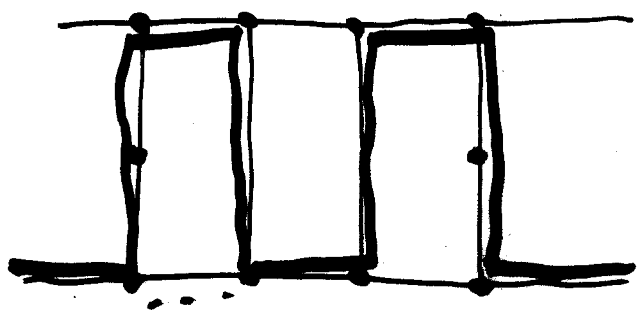


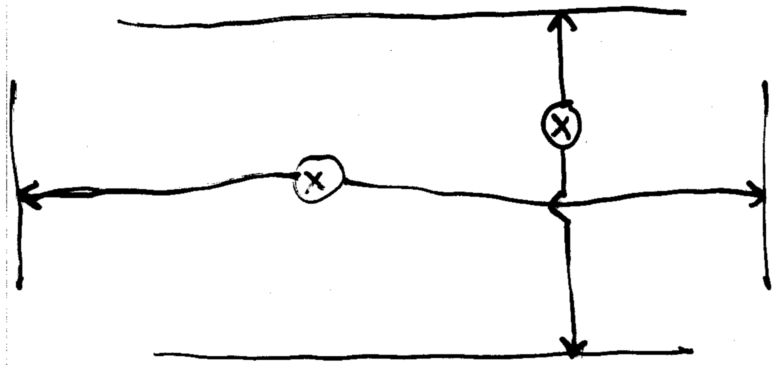
NP-complete: reduction from 3-CNF sat via "gadgets"

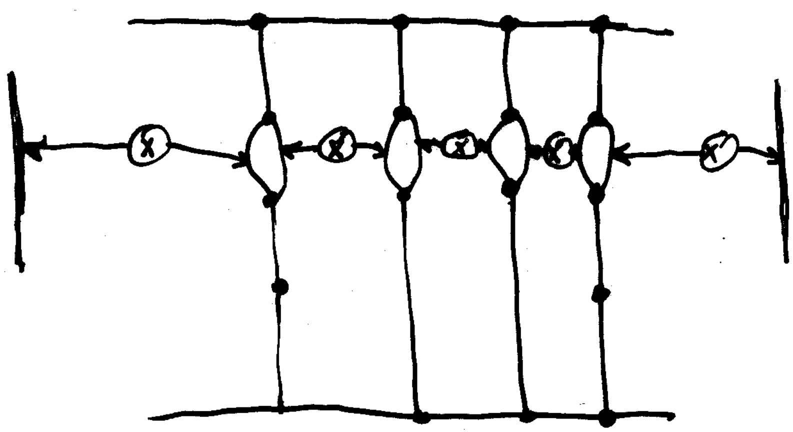
(even for planar graphs, all vertices of degree 3, all faces of 5 or more sides)

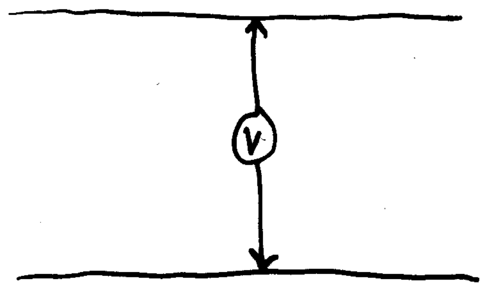


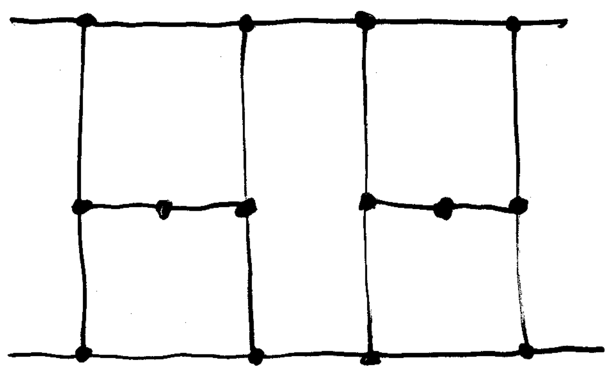


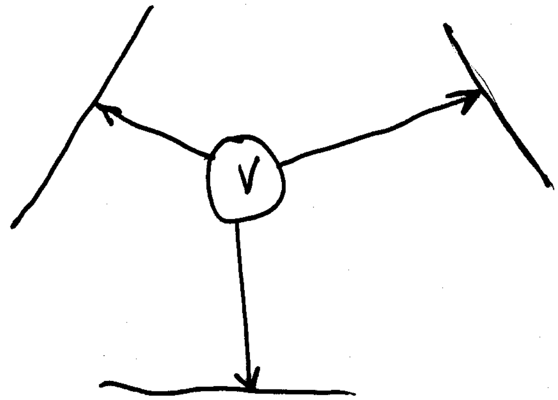


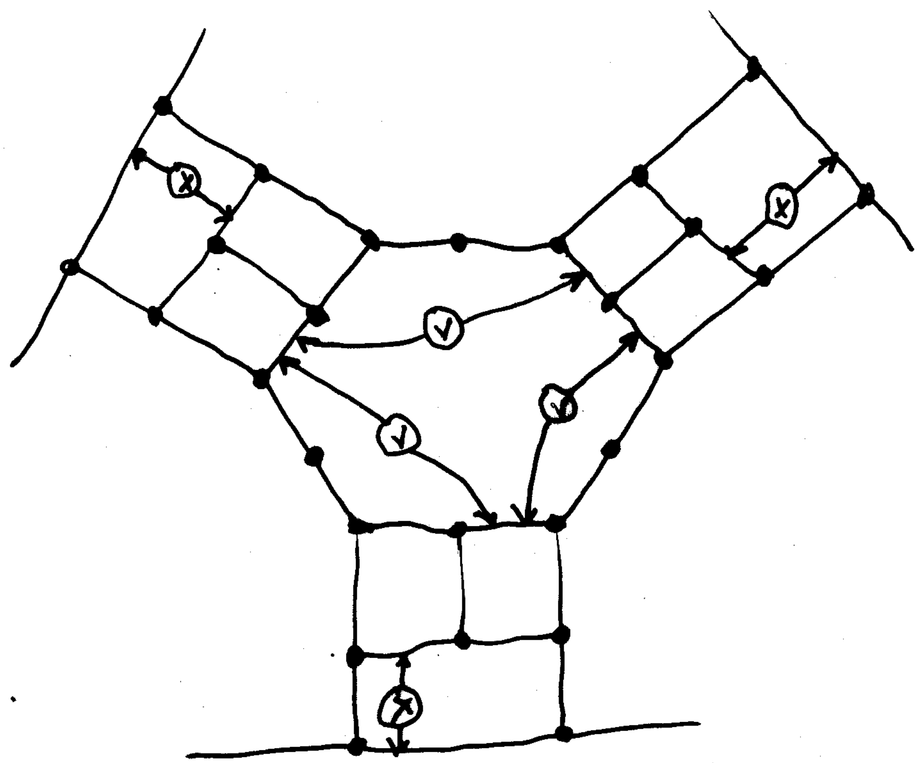

x
x

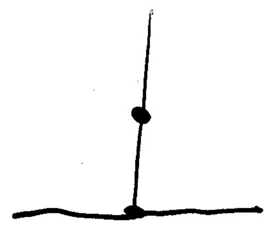


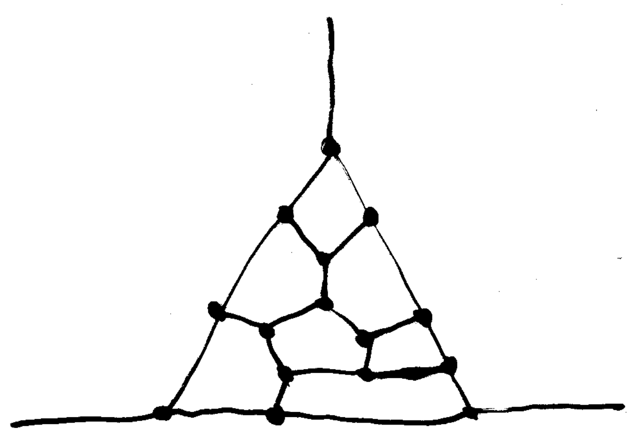

V

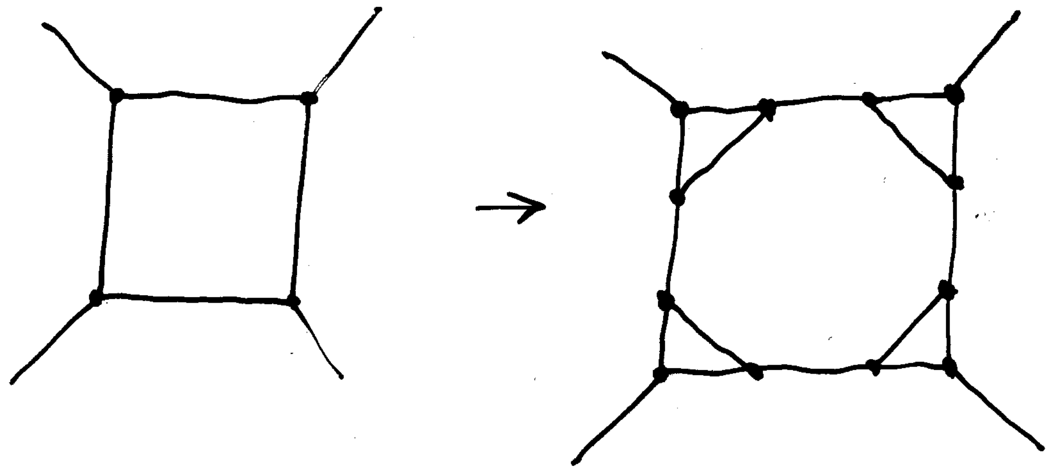


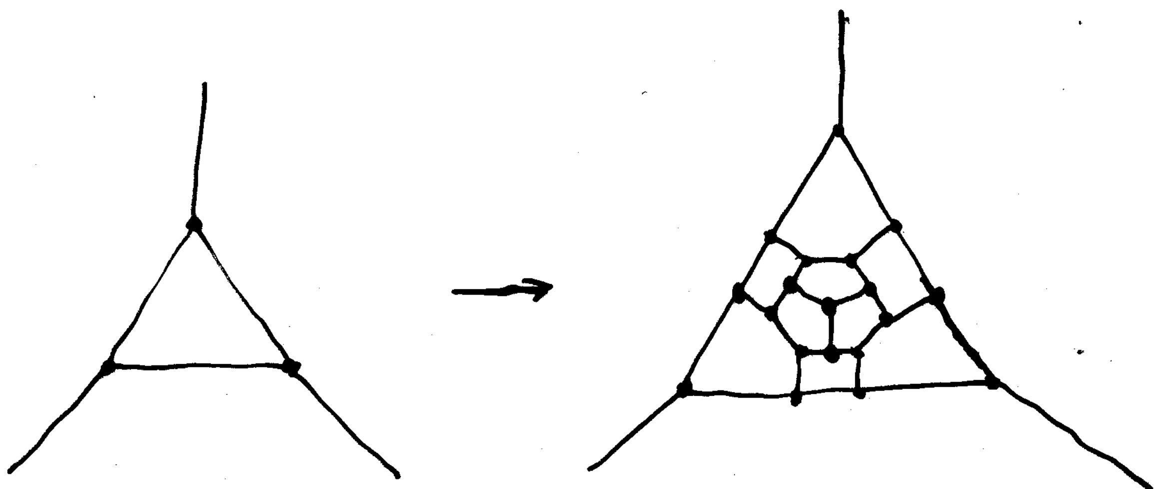

V


x
x
V
V
V
x









Subset Sum is NP-complete

Given m integers, and a target k, is there a subset that sums to exactly k?

$\{2, 5, 6, 8, 9, 12\}$ $k = 31$

yes: 5, 6, 8, 12 $n = \#$ bits

(no for $k = 30$)

In NP: subset is proof (verifiable in p-time)

Some NPC problem reducible to subset sum

reduce 3-CNF sat to subset sum

Write numbers base 10

$$(x \vee \bar{y} \vee \bar{z}) \wedge (\bar{x} \vee y \vee z) \wedge (y \vee \bar{z})$$

C_1 C_2 C_3

	x	y	z	C_1	C_2	C_3	
x	1	0	0	1	0	0	
$\bar{x}$	1	0	0	0	1	0	
y	0	1	0	0	1	1	← y makes C_2, C_3 true
$\bar{y}$	0	1	0	1	0	0	
z	0	0	1	0	1	0	
$\bar{z}$	0	0	1	1	0	1	← $\bar{z}$ makes C_1, C_2 true
Dummies to get clause columns to sum to 4	0	0	0	1	0	0	
	0	0	0	2	0	0	
	0	0	0	0	1	0	
	0	0	0	0	2	0	
	0	0	0	0	0	1	
	0	0	0	0	0	2	
	1	1	1	4	4	4	= k Required sum

Interpret each row as a base 10 number.

Subset sum has a solution iff formula is satisfiable.