# Logic: From Greeks to philosophers to circuits. 

COS 116: 3/15/2005
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## Question from Last Lecture

Boolean Expression $E=S$ AND $\bar{D}$
Can we deduce an expression for D? Let's look at truth table...

| D | S | E |
| :---: | :---: | :---: |
| 0 | 0 | 0 |
| 0 | 1 | 1 |
| 1 | 0 | 0 |
| 1 | 1 | 0 |$\rightarrow$| E | S | D |
| :---: | :---: | :---: |
| 0 | 0 | $?$ |
| 0 | 1 | $?$ |
| 1 | 0 | $?$ |
| 1 | 1 | $?$ |

## Three Equivalent Representations

Boolean Expression $E=S A N D \bar{D}$

Boolean Circuit


Truth table:
Value of E for every possible D, S.
TRUE=1; FALSE= 0.

| $D$ | $S$ | $E$ |
| :---: | :---: | :---: |
| 0 | 0 | 0 |
| 0 | 1 | 1 |
| 1 | 0 | 0 |
| 1 | 1 | 0 |

## Boolean "algebra"

A AND B written as A.B
A OR B written as A + B

$$
\begin{array}{ll}
0 \cdot 0=0 & 0+0=0 \\
0 \cdot 1=0 & 1+0=1 \\
1 \cdot 1=1 & 1+1=1
\end{array}
$$

Funny arithmetic

Will provide readings on this...

## Boolean gates

## Shannon (1939)

High voltage = 1
Low voltage $=0$


Output voltage is high
if both of the input voltages are high; otherwise output voltage low.

Output voltage is high
if either of the input voltages are high; otherwise output voltage low.

Output voltage is high
if the input voltage is low; otherwise output voltage high.
(implicit extra wires for power)

## Claude Shannon (1916-2001)

Founder of many fields
(circuits, information theory, artificial intelligence...)


## Combinational circuit

- Boolean gates connected by wires


Wires: transmit voltage (and hence value)

■ Important: no cycles allowed

## Examples

4-way AND

(Sometimes we use this for shorthand)


More complicated example
 are sometimes drawn like this.

## Combinational circuits and control

- "If data has arrived and packet has not been sent, send a signal"



## Circuits compute functions

- Every combinational circuit computes a Boolean function of its inputs


Inputs


## Ben Revisited

Ben only rides to class if he overslept, but even then if it is raining he'll walk and show up late (he hates to bike in the rain). But if there's an exam that day he'll bike if he overslept, even in the rain.

B: Ben Bikes
$\mathbf{R}$ : It is raining
$E$ : There is an exam today
O: Ben overslept

Ben's truth table

| O | R | E | B |
| :---: | :---: | :---: | :---: |
| 0 | 0 | 0 | 0 |
| 0 | 0 | 1 | 0 |
| 0 | 1 | 0 | 0 |
| 0 | 1 | 1 | 0 |
| 1 | 0 | 0 | 1 |
| 1 | 0 | 1 | 1 |
| 1 | 1 | 0 | 0 |
| 1 | 1 | 1 | 1 |

## Truth table $\rightarrow$ Boolean expression

Use OR of all input combinations that lead to TRUE

| O | R | E | B |
| :---: | :---: | :---: | :---: |
| 0 | 0 | 0 | 0 |
| 0 | 0 | 1 | 0 |
| 0 | 1 | 0 | 0 |
| 0 | 1 | 1 | 0 |
| 1 | 0 | 0 | 1 |
| 1 | 0 | 1 | 1 |
| 1 | 1 | 0 | 0 |
| 1 | 1 | 1 | 1 |

$B=O \cdot \bar{R} \cdot \bar{E}+O \cdot \bar{R} \cdot E+O \cdot R \cdot E$

Note:
AND, OR, and NOT gates suffice to implement every Boolean function!

## Sizes of representations

■ For $k$ variables:

| $k$ | 10 | 20 | 30 |
| :--- | ---: | ---: | ---: |
| $2^{k}$ | 1024 | 1048576 | 1073741824 |

For an arbitrary function, expect roughly half of X's to be 1 (for 30 inputs roughly $1 / 2$ billion!)
$2^{k}\left\{\begin{array}{c|c|c|c|}\hline \mathrm{A} & \mathrm{B} & \ldots & \mathrm{X} \\ \hline 0 & 0 & \ldots & 0 \\ \hline 0 & 0 & \ldots & 0 \\ \hline 0 & 1 & \ldots & 0 \\ \hline 0 & 1 & \ldots & 1 \\ \hline \ldots & \ldots & \ldots & \ldots \\ \hline \ldots & \ldots & \ldots & \ldots \\ \hline 1 & 1 & \ldots & 1 \\ \hline\end{array}\right.$


Tools for reducing size:
(a) circuit optimization (b) modular design

## Expression simplification

- Some simple rules:

$$
\begin{aligned}
& x+\bar{x}=1 \\
& x \cdot 1=x \\
& x \cdot 0=0
\end{aligned}
$$

$$
\begin{array}{rl}
x & y+x \cdot \bar{y} \\
& =x \cdot(y+\bar{y}) \\
& =x \cdot 1 \\
& =x
\end{array}
$$

$$
x+0=x
$$

$$
x+1=1
$$

De Morgan's Laws:

$$
x+x=x \cdot x=x
$$

$\overline{x \cdot y}=\bar{x}+\bar{y}$

$$
x \cdot(y+z)=x \cdot y+x \cdot z
$$

$\overline{x+y}=\bar{x} \cdot \bar{y}$
$x+(y \cdot z)=(x+y) \cdot(x+z)$

## Simplifying Ben's circuit

$$
\begin{aligned}
\square & =O \cdot \bar{R} \cdot \bar{E}+O \cdot \bar{R} \cdot E+O \cdot R \cdot E \\
& =O \cdot(\bar{R} \cdot \bar{E}+\bar{R} \cdot E+R \cdot E) \\
& =O \cdot(\bar{R} \cdot(\bar{E}+E)+R \cdot E) \\
& =O \cdot(\bar{R}+R \cdot E) \\
& \ldots \\
& =O \cdot(\bar{R}+E)
\end{aligned}
$$

## Something to think about: How hard is Circuit Verification?



■ Given a circuit, decide if it is "trivial" (no matter the input, it either always outputs 1 or always outputs 0 )


- Alternative statement: Decide if there is any setting of the inputs that makes the circuit evaluate to 1.

Time required?

## Boole's reworking of Clarke's "proof" of existence of God (see handout)



- General idea: Try to prove that Boolean expressions $\mathrm{E}_{1}, \mathrm{E}_{2}, \ldots, \mathrm{E}_{\mathrm{k}}$ cannot simultaneously be true
- Method: Show $E_{1} \cdot E_{2} \cdot \ldots \cdot E_{k}=0$
- Discussion for next time: What exactly does Clarke's "proof" prove? How convincing is such a proof to you?

Also: Do Google search for "Proof of God's Existence."

## Beyond combinational circuits ...

- Need 2-way communication (must allow cycles!)
- Need memory (scratchpad)



# Circuit for Binary Addition? 

$25=11001$<br>$29=11101$<br>110110

Read handout: Will discuss next time.

