



Binary Numbers

COS 217



Goals of Today's Lecture

- **Binary numbers**
 - Why binary?
 - Converting base 10 to base 2
 - Octal and hexadecimal
- **Integers**
 - Unsigned integers
 - Integer addition
 - Signed integers
- **C bit operators**
 - And, or, not, and xor
 - Shift-left and shift-right
 - Function for counting the number of 1 bits
 - Function for XOR encryption of a message



Why Bits (Binary Digits)?

- Computers are built using digital circuits
 - Inputs and outputs can have only two values
 - True (high voltage) or false (low voltage)
 - Represented as 1 and 0
- Can represent many kinds of information
 - Boolean (true or false)
 - Numbers (23, 79, ...)
 - Characters ('a', 'z', ...)
 - Pixels
 - Sound
- Can manipulate in many ways
 - Read and write
 - Logical operations
 - Arithmetic
 - ...



Base 10 and Base 2

- Base 10

- Each digit represents a power of 10
- $4173 = 4 \times 10^3 + 1 \times 10^2 + 7 \times 10^1 + 3 \times 10^0$

- Base 2

- Each bit represents a power of 2
- $10110 = 1 \times 2^4 + 0 \times 2^3 + 1 \times 2^2 + 1 \times 2^1 + 0 \times 2^0 = 22$

Divide repeatedly by 2 and keep remainders

$$12/2 = 6 \quad R = 0$$

$$6/2 = 3 \quad R = 0$$

$$3/2 = 1 \quad R = 1$$

$$1/2 = 0 \quad R = 1$$

$$\text{Result} = 1100$$

Writing Bits is Tedious for People



- Octal (base 8)
 - Digits 0, 1, ..., 7
 - In C: 00, 01, ..., 07
- Hexadecimal (base 16)
 - Digits 0, 1, ..., 9, A, B, C, D, E, F
 - In C: 0x0, 0x1, ..., 0xf

0000 = 0	1000 = 8
0001 = 1	1001 = 9
0010 = 2	1010 = A
0011 = 3	1011 = B
0100 = 4	1100 = C
0101 = 5	1101 = D
0110 = 6	1110 = E
0111 = 7	1111 = F

Thus the 16-bit binary number

1011 0010 1010 1001

converted to hex is

B2A9

Representing Colors: RGB



- Three primary colors
 - Red
 - Green
 - Blue
- Strength
 - 8-bit number for each color (e.g., two hex digits)
 - So, 24 bits to specify a color
- In HTML, on the course Web page
 - Red: `<i>Symbol Table Assignment Due</i>`
 - Blue: `<i>Fall Recess</i>`
- Same thing in digital cameras
 - Each pixel is a mixture of red, green, and blue

Storing Integers on the Computer



- Fixed number of bits in memory

- Short: usually 16 bits
- Int: 16 or 32 bits
- Long: 32 bits

- Unsigned integer

- No sign bit
- Always positive or 0
- All arithmetic is modulo 2^n

- Example of unsigned int

- 00000001 → 1
- 00001111 → 15
- 00010000 → 16
- 00100001 → 33
- 11111111 → 255

Adding Two Integers: Base 10



- From right to left, we add each pair of digits
- We write the sum, and add the carry to the next column

$$\begin{array}{r} \\ + \\ \hline \text{Sum} \\ \text{Carry} \end{array}$$

1 9 8
+ 2 6 4

Sum 4 6 2
Carry 0 1 1

$$\begin{array}{r} \\ + \\ \hline \text{Sum} \\ \text{Carry} \end{array}$$

0 1 1
+ 0 0 1

Sum 1 0 0
Carry 0 1 1



Binary Sums and Carries

a	b	Sum
0	0	0
0	1	1
1	0	1
1	1	0

XOR

a	b	Carry
0	0	0
0	1	0
1	0	0
1	1	1

AND

$$\begin{array}{r} 0100 \ 0101 \longleftarrow 69 \\ + 0110 \ 0111 \longleftarrow 103 \\ \hline 1010 \ 1100 \longleftarrow 172 \end{array}$$



Modulo Arithmetic

- Consider only numbers in a range
 - E.g., five-digit car odometer: 0, 1, ..., 99999
 - E.g., eight-bit numbers 0, 1, ..., 255
- Roll-over when you run out of space
 - E.g., car odometer goes from 99999 to 0, 1, ...
 - E.g., eight-bit number goes from 255 to 0, 1, ...
- Adding 2^n doesn't change the answer
 - For eight-bit number, $n=8$ and $2^n=256$
 - E.g., $(37 + 256) \bmod 256$ is simply 37
- This can help us do subtraction...
 - Suppose you want to compute $a - b$
 - Note that this equals $a + (256 - 1 - b) + 1$

One's and Two's Complement



- One's complement: flip every bit
 - E.g., b is 01000101 (i.e., 69 in base 10)
 - One's complement is 10111010
 - That's simply $255 - 69$
- Subtracting from 11111111 is easy (no carry needed!)

$$\begin{array}{r} 1111 \ 1111 \\ - 0100 \ 0101 \longleftarrow b \\ \hline 1011 \ 1010 \longleftarrow \text{one's complement} \end{array}$$

- Two's complement
 - Add 1 to the one's complement
 - E.g., $(255 - 69) + 1 \rightarrow 1011 \ 1011$



Putting it All Together

- Computing “ $a - b$ ” for unsigned integers

- Same as “ $a + 256 - b$ ”
- Same as “ $a + (255 - b) + 1$ ”
- Same as “ $a + \text{oncomplement}(b) + 1$ ”
- Same as “ $a + \text{twocomplement}(b)$ ”

- Example: $172 - 69$

- The original number 69: 0100 0101
- One’s complement of 69: 1011 1010
- Two’s complement of 69: 1011 1011
- Add to the number 172: 1010 1100
- The sum comes to: 0110 0111
- Equals: 103 in base 10

$$\begin{array}{r} 1010\ 1100 \\ +\ 1011\ 1011 \\ \hline 1\ 0110\ 0111 \end{array}$$



Signed Integers

- Sign-magnitude representation

- Use one bit to store the sign
 - Zero for positive number
 - One for negative number
- Examples
 - E.g., 0010 1100 → 44
 - E.g., 1010 1100 → -44
- Hard to do arithmetic this way, so it is rarely used

- Complement representation

- One's complement
 - Flip every bit
 - E.g., 1101 0011 → -44
- Two's complement
 - Flip every bit, then add 1
 - E.g., 1101 0100 → -44

Overflow: Running Out of Room



- Adding two large integers together
 - Sum might be too large to store in the number of bits allowed
 - What happens?
- Unsigned numbers
 - All arithmetic is “modulo” arithmetic
 - Sum would just wrap around
- Signed integers
 - Can get nonsense values
 - Example with 16-bit integers
 - Sum: $10000+20000+30000$
 - Result: -5536
 - In this case, fixable by using “long”...



Bitwise Operators: AND and OR

- Bitwise AND (&)

&	0	1
0	0	0
1	0	1

- Bitwise OR (|)

	0	1
0	0	1
1	1	1

- o Mod on the cheap!
 - E.g., $h = 53 \& 15$;

53

0	0	1	1	0	1	0	1
---	---	---	---	---	---	---	---

& 15

0	0	0	0	1	1	1	1
---	---	---	---	---	---	---	---

5

0	0	0	0	0	1	0	1
---	---	---	---	---	---	---	---

Bitwise Operators: Not and XOR



- One's complement (\sim)
 - Turns 0 to 1, and 1 to 0
 - E.g., set last three bits to 0
– $x = x \& \sim 7$;
- XOR (\wedge)
 - 0 if both bits are the same
 - 1 if the two bits are different

\wedge	0	1
0	0	1
1	1	0

Bitwise Operators: Shift Left/Right



- Shift left (<<): Multiply by powers of 2

- Shift some # of bits to the left, filling the blanks with 0

53 0 0 1 1 0 1 0 0

53<<2 1 1 0 1 0 0 0 0

- Shift right (>>): Divide by powers of 2

- Shift some # of bits to the right
 - For unsigned integer, fill in blanks with 0
 - What about signed integers? Varies across machines...
 - Can vary from one machine to another!

53 0 0 1 1 0 1 0 0

53>>2 0 0 0 0 1 1 0 1

Count Number of 1s in an Integer



- Function `bitcount(unsigned x)`
 - Input: unsigned integer
 - Output: number of bits set to 1 in the binary representation of `x`
- Main idea
 - Isolate the last bit and see if it is equal to 1
 - Shift to the right by one bit, and repeat

```
int bitcount(unsigned x) {  
    int b;  
    for (b = 0; x != 0; x >>= 1)  
        if (x & 1)  
            b++;  
    return b;  
}
```



XOR Encryption

- Program to encrypt text with a key
 - Input: original text in stdin
 - Output: encrypted text in stdout
- Use the same program to decrypt text with a key
 - Input: encrypted text in stdin
 - Output: original text in stdout
- Basic idea
 - Start with a key, some 8-bit number (e.g., 0110 0111)
 - Do an operation that can be inverted
 - E.g., XOR each character with the 8-bit number

$$\begin{array}{r} 0100 \ 0101 \\ \wedge \ 0110 \ 0111 \\ \hline 0010 \ 0010 \end{array}$$

$$\begin{array}{r} 0010 \ 0010 \\ \wedge \ 0110 \ 0111 \\ \hline 0100 \ 0101 \end{array}$$

XOR Encryption, Continued



- But, we have a problem
 - Some characters are control characters
 - These characters don't print
- So, let's play it safe
 - If the encrypted character would be a control character
 - ... just print the original, unencrypted character
 - Note: the same thing will happen when decrypting, so we're okay
- C function `iscntrl()`
 - Returns true if the character is a control character

XOR Encryption, C Code



```
#define KEY '&'
int main() {
    int orig_char, new_char;

    while ((orig_char = getchar()) != EOF) {
        new_char = orig_char ^ KEY;
        if (iscntrl(new_char))
            putchar(orig_char);
        else
            putchar(new_char);
    }
    return 0;
}
```

Next Week



- Wednesday lecture time
 - Midterm exam
 - Open book and open notes
 - Practice exams online

Stupid Programmer Tricks



- Where do I use bitwise & most?
 - Bit vectors
- What's a bit vector?
 - Lots of booleans packed into an int/long
 - Often used to indicate some condition(s)
 - Less storage space than lots of fields
 - More explicit storage than compiled-defined bit fields

- Your compiler can do this?

```
typedef struct blah {  
    int b_onoff:1;  
    int b_temperature:7;  
    char b_someChar;  
}
```



Example From Real Code

- `#define DONTCACHE_REQNOSTORE` 0x000001
- `#define DONTCACHE_AUTHORIZED` 0x000002
- `#define DONTCACHE_MISSINGVARIANTHDR` 0x000004
- `#define DONTCACHE_USERORPASS` 0x000008
- `#define DONTCACHE_BYPASSFILTER` 0x000010
- `#define DONTCACHE_NONCACHEMETHOD` 0x000020
- `#define DONTCACHE_CTLPRIVATE` 0x000040
- `#define DONTCACHE_CTLNOSTORE` 0x000080
- `#define DONTCACHE_ISQUERY` 0x000100
- `#define DONTCACHE_EARLYEXPIRE` 0x000200
- `#define DONTCACHE_NOLASTMOD` 0x000400
- `#define DONTCACHE_NONEGCACHING` 0x000800
- `#define DONTCACHE_INSTANTEXPIRE` 0x001000
- `#define DONTCACHE_FILETOOBIG` 0x002000
- `#define DONTCACHE_FILEGREWTOOBIG` 0x004000
- `#define DONTCACHE_ICPPROXYONLY` 0x008000
- `#define DONTCACHE_LARGEFILEBLAST` 0x010000
- `#define DONTCACHE_PERSISTLOGLOADING` 0x020000
- `#define DONTCACHE_NEWERCOPYEXISTS` 0x040000
- `#define DONTCACHE_BADVARYFIELDS` 0x080000
- `#define DONTCACHE_SETCOOKIE` 0x100000
- `#define DONTCACHE_HTTPSTATUSCODE` 0x200000
- `#define DONTCACHE_OBJECTINCOMPLETE` 0x400000

Conclusions



- Computer represents everything in binary
 - Integers, floating-point numbers, characters, addresses, ...
 - Pixels, sounds, colors, etc.
- Binary arithmetic through logic operations
 - Sum (XOR) and Carry (AND)
 - Two's complement for subtraction
- Binary operations in C
 - AND, OR, NOT, and XOR
 - Shift left and shift right
 - Useful for efficient and concise code, though sometimes cryptic