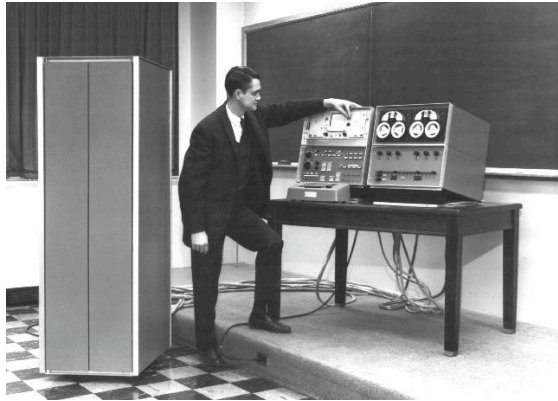


Lecture 9: TOY II



LINC

COS 126: General Computer Science · <http://www.Princeton.EDU/~cos126>

What We've Learned About TOY

Data representation.

- Binary and hex.

TOY: what's in it, how to use it.

- Box with switches and lights.
- $4,328 \text{ bits} = (255 \times 16) + (15 \times 16) + (8)$. 541 bytes!
- von Neumann architecture.

TOY instruction set architecture.

- 16 instruction types.

Sample TOY machine language programs.

- Arithmetic.
- Loops.

2

What We Do Today

Binary add, subtract.

Standard input, standard output.

Manipulate addresses.

- References (pointers).
- Arrays.

TOY simulator in Java.

3

How to add and subtract binary numbers

Binary addition facts:

- $0 + 0 = 0$
- $0 + 1 = 1 + 0 = 1$
- $1 + 1 = 10$
- $1 + 1 + 1 = 11$ (needed for carries)

Bigger numbers example:

$$\begin{array}{r} 1 \leftarrow \text{carries} \rightarrow 1 \ 1 \\ 013 \\ + 092 \\ \hline 105 \end{array} \qquad \begin{array}{r} 00001101 \\ + 01011100 \\ \hline 01101001 \end{array}$$

OK, but: subtract?

- Subtract by adding a negative integer (e.g., $6 - 4 = 6 + (-4)$)
- OK, but: negative integers?

4

How to Represent Negative Integers

TOY words are 16 bits each.

- We could use 16 bits to represent 0 to $2^{16} - 1$.
- But we want negative integers too.
- Reserving half the possible bit-patterns for negative seems fair.

Highly desirable property:

- If X is a positive integer, then the representation of $-X$, when added to X, had better yield zero.

$$\begin{array}{r}
 X \quad \quad 00110100 \\
 +(-X) \quad +?????? \\
 \hline
 0 \quad \quad 00000000
 \end{array}$$

$$\begin{array}{r}
 X \quad \quad 00110100 \\
 +(-X) \quad +11001011 \\
 \hline
 0 \quad \quad 00000000
 \end{array}$$

(Note: In the second diagram, red arrows point from the 1s in the second row to the 0s in the first row, illustrating bit flipping.)

5

"Two's Complement" Integers

Properties:

- Leading bit (bit 15) signifies sign.
- Negative integer $-N$ represented by $2^{16} - N$.
- Trick to compute $-N$:

1. Start with N.

	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
+4	0	0	0	0	0	0	0	0	0	0	0	0	0	1	0	0

2. Flip bits.

	1	1	1	1	1	1	1	1	1	1	1	1	1	0	1	1
--	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---

3. Add 1.

-4	1	1	1	1	1	1	1	1	1	1	1	1	1	1	0	0
----	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---

6

Two's Complement Integers

		15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
Dec	Hex	Binary															
+32767	7FFF	0	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1

...

+4	0004	0	0	0	0	0	0	0	0	0	0	0	0	0	1	0	0
+3	0003	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	1
+2	0002	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	0
+1	0001	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1
+0	0000	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
-1	FFFF	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1
-2	FFFE	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	0
-3	FFFD	1	1	1	1	1	1	1	1	1	1	1	1	1	1	0	1
-4	FFFC	1	1	1	1	1	1	1	1	1	1	1	1	1	1	0	0

...

-32768	8000	1	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
--------	------	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---

7

Properties of Two's Complement Integers

Nice properties:

- 0000000000000000 represents 0.
- -0 and $+0$ are the same.
- Addition is easy (see next slide).
- Checking for arithmetic overflow is easy.

$$-N = \sim N + 1$$

Not-so-nice properties.

- Can represent one more negative integer than positive integer. ($-32,768 = -2^{15}$ but not $32,768 = 2^{15}$).

8

Two's Complement Arithmetic

Addition is carried out as if all integers were positive.

- It usually works.

-3	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	0	1	
+																			
4	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	0	0
=																			
1	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1

Two's Complement Arithmetic

Addition is carried out as if all integers were positive.

- It usually works.
- But overflow can occur:
 - carry into sign (left most) bit with no carry out

+32,767	0	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1
+																			
2	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1
=																			
-32,767	1	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1

Standard Output

Standard output.

- Writing to memory location `FF` sends one word to TOY stdout.
- `9AFF` writes the integer in register `A` to stdout.

```

00: 0000  0
01: 0001  1

10: 8A00  RA <- mem[00]      a = 0
11: 8B01  RB <- mem[01]      b = 1
                                while(a > 0) {
12: 9AFF  print RA          print a
13: 1AAB  RA <- RA + RB      a = a + b
14: 2BAB  RB <- RA - RB      b = a - b
15: DA12  if (RA > 0) goto 12  }
16: 0000  halt
    
```

Standard Output

Standard output.

- Writing to memory location `FF` sends one word to TOY stdout.
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```

00: 0000  0
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                                while(a > 0) {
12: 9AFF  print RA          print a
13: 1AAB  RA <- RA + RB      a = a + b
14: 2BAB  RB <- RA - RB      b = a - b
15: DA12  if (RA > 0) goto 12  }
16: 0000  halt
    
```

- 0000
- 0001
- 0001
- 0002
- 0003
- 0005
- 0008
- 000D
- 0015
- 0022
- 0037
- 0059
- 0090
- 00E9
- 0179
- 0262
- 03DB
- 063D
- 0A18
- 1055
- 1A6D
- 2AC2
- 452F
- 6FF1

Standard Input

Standard input.

- Loading from memory address `FF` loads one word from TOY stdin.
- `8AFF` reads in an integer from stdin and stores it in register `A`.

Ex: read in a sequence of integers and print their sum.

- In Java, stop reading when EOF.
- In TOY, stop reading when user enters `0000`.

```
while(!StdIn.isEmpty()) {
    a = StdIn.readInt();
    sum = sum + a;
}
System.out.println(sum);
```

```
00: 0000 0
10: 8C00 RC ← mem[00]
11: 8AFF read RA
12: CA15 if (RA == 0) pc ← 15
13: 1CCA RC ← RC + RA
14: C011 pc ← 11
15: 9CFF write RC
16: 0000 halt
```

```
00AE
0046
0003
0000
00F7
```

14

Standard Input and Output: Implications

Standard input and output enable you to:

- Put information from real world into machine.
- Get information out of machine.
- Process more information than fits in memory.
- Interact with the computer while it is running.

Information can be instructions!

- Booting a computer.

15

Load Address (a.k.a. Load Constant)

Load address. (opcode 7)

- Loads an 8-bit integer into a register.
- `7A30` means load the value `30` into register `A`.

Applications.

- Load a small constant into a register.
- Load a 8-bit memory address into a register.
 - register stores "pointer" to a memory cell

```
a = 30;
```

Java code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	1	1	1	0	1	0	0	0	1	1	0	0	0	0
7 ₁₆				A ₁₆				3 ₁₆				0 ₁₆			
opcode				dest d				addr							

16

Arrays in TOY

TOY main memory is a giant array.

- Can access memory cell `30` using load and store.
- `8C30` means load `mem[30]` into register `C`.
- Goal: access memory cell `i` where `i` is a variable.

Load indirect. (opcode B) a variable index

- `BC06` means load `mem[R6]` into register `C`.

Store indirect. (opcode A) a variable index

- `AC06` means store contents of register `C` into `mem[R6]`.

```
for (int i = 0; i < N; i++)
    a[i] = StdIn.readInt();

for (int i = 0; i < N; i++)
    System.out.println(a[N-i-1]);
```

Reverse.java

17

TOY Implementation of Reverse

TOY implementation of reverse.

- Read in a sequence of integers and store in memory 30, 31, 32, ...
 - stop reading if 0000
- Print sequence in reverse order.

```

10: 7101 R1 ← 0001      constant 1
11: 7A30 RA ← 0030      a[]
12: 7B00 RB ← 0000      n

13: 8CFF read RC          while(true) {
14: CC19 if (RC == 0) goto 19   c = StdIn.readInt();
15: 16AB R6 ← RA + RB         if (c == 0) break;
16: BC06 mem[R6] ← RC         address of a[n]
17: 1BB1 RB ← RB + R1         a[n] = c;
18: C013 goto 13             n++;
                                }
                                read in the data

```

18

TOY Implementation of Reverse

TOY implementation of reverse.

- Read in a sequence of integers and store in memory 30, 31, 32, ...
 - stop reading if 0000
- Print sequence in reverse order.

```

19: CB20 if (RB == 0) goto 20   while (n > 0) {
1A: 16AB R6 ← RA + RB         address of a[n]
1B: 2661 R6 ← R6 - R1         address of a[n-1]
1C: AC06 RC ← mem[R6]         c = a[n-1];
1D: 9CFF write RC             System.out.println(c);
1E: 2BB1 RB ← RB - R1         n--;
1F: C019 goto 19              }
20: 0000 halt

```

19

Unsafe Code at any Speed

What happens if we make array start at 00 instead of 30?

- Self modifying program.
- Exploit buffer overflow and run arbitrary code!

```

10: 7101 R1 ← 0001      constant 1
11: 7A00 RA ← 0000      a[]
12: 7B00 RB ← 0000      n

13: 8CFF read RC          while(true) {
14: CC19 if (RC == 0) goto 19   c = StdIn.readInt();
15: 16AB R6 ← RA + RB         if (c == 0) break;
16: BC06 mem[R6] ← RC         address of a[n]
17: 1BB1 RB ← RB + R1         a[n] = c;
18: C013 goto 13             n++;
                                }

```

Crazy 8s Input

```

1 1 1 1 1 1 1 1 1
1 1 1 1 1 1 1 1 1
8888 8810
98FF C011

```

20

Dumping

Dumping.

- Work all day to develop operating system.
- How do you save it for tomorrow?
 - leave computer on?
 - write short program dump.toy
 - run dump.toy to dump contents of memory onto tape

```

00: 7001 R1 ← 0001
01: 7210 R2 ← 0010      i = 10
02: 73FF R3 ← 00FF

03: AA02 RA ← mem[R2]
04: 9AFF write RA        do {
05: 1221 R2 ← R2 + R1    a = mem[i]
06: 2432 R4 ← R3 - R2    print a
07: D403 if (R4 > 0) goto 03   i++
08: 0000 halt           } while (i < 255)

```

dump.toy

21

Booting



Booting.

- How do you get it back?
 - turn on computer, old memory values gone
 - write short program `boot.toy`
 - read contents of memory from tape by running `boot.toy`
 - use original program

```

00: 7001  R1 ← 0001
01: 7210  R2 ← 0010          i = 10
02: 73FF  R3 ← 00FF

03: 8AFF  read RA          do {
04: BA02  mem[R2] ← RA      read a
                                mem[i] = a
05: 1221  R2 ← R2 + R1      i++
06: 2432  R4 ← R3 - R2
07: D403  if (R4 > 0) goto 03 } while (i < 255)
08: 0000  halt
    
```

`boot.toy`

22

TOY Simulator

Write a program to "simulate" the behavior of the TOY machine.

- TOY simulator in Java.
- TOY simulator in TOY!

```

public class TOY {
    public static void main(String[] args) {
        int pc = 0x10; // program counter
        int[] R = new int[16]; // registers
        int[] mem = new int[256]; // main memory

        // READ IN .toy FILE

        while(true) {
            // FETCH INSTRUCTION and DECODE
            ...
            // EXECUTE
            ...
        }
    }
}
    
```

```

% java TOY add-stdin.toy
A012
002B
A03D
    
```

23

TOY Simulator: Fetch

Extract destination register of `1CAB` by shifting and masking.

0	0	0	1	1	1	0	0	1	0	1	0	1	0	1	1	<code>inst</code>
1_{16}				C_{16}				A_{16}				B_{16}				
0	0	0	0	0	0	0	0	0	0	0	1	1	1	0	0	<code>inst >> 8</code>
0_{16}				0_{16}				1				C_{16}				
0	0	0	0	0	0	0	0	0	0	0	1	1	1	1	1	15
0_{16}				0_{16}				0_{16}				F_{16}				
0	0	0	0	0	0	0	0	0	0	0	1	1	0	0	0	<code>(inst >> 8) & 15</code>
0_{16}				0_{16}				0				C_{16}				

```

int inst = mem[pc++]; // fetch and increment
int op = (inst >> 12) & 15; // opcode (bits 12-15)
int d = (inst >> 8) & 15; // dest d (bits 08-11)
int s = (inst >> 4) & 15; // source s (bits 04-07)
int t = (inst >> 0) & 15; // source t (bits 00-03)
int addr = (inst >> 0) & 255; // addr (bits 00-07)
    
```

24

TOY Simulator: Execute

```

if (op == 0) break; // halt

switch (op) {
    case 1: R[d] = R[s] + R[t]; break;
    case 2: R[d] = R[s] - R[t]; break;
    case 3: R[d] = R[s] & R[t]; break;
    case 4: R[d] = R[s] ^ R[t]; break;
    case 5: R[d] = R[s] << R[t]; break;
    case 6: R[d] = R[s] >> R[t]; break;
    case 7: R[d] = addr; break;
    case 8: R[d] = mem[addr]; break;
    case 9: mem[addr] = R[d]; break;
    case 10: R[d] = mem[R[t]]; break;
    case 11: mem[R[t]] = R[d]; break;
    case 12: if (R[d] == 0) pc = addr; break;
    case 13: if (R[d] > 0) pc = addr; break;
    case 14: pc = R[d]; break;
    case 15: R[d] = pc; pc = addr; break;
}
    
```

25

TOY Simulator: Missing Details

Omitted details.

- Register 0 is always 0.
 - reset to 0000 after each fetch-execute step
- Standard input and output.
 - if `addr` is `FF` and opcode is load (indirect) then read in data
 - if `addr` is `FF` and opcode is store (indirect) then write out data
- TOY registers are 16-bit integers; program counter is 8-bit.
 - Java `int` is 32 bits

See `TOY.java` for full details.

26

Simulation

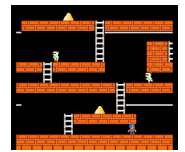
Consequences of simulation.

- Test out new machine or microprocessor using simulator.
 - cheaper and faster than building actual machine
- Easy to add new functionality to simulator.
 - trace, single-step, breakpoint debugging
 - simulator more useful than TOY itself
- Reuse software from old machines.



Ancient programs still running on modern computers.

- Ticketron.
- Lode Runner on Apple IIe.



27

Announcements

Not-exactly Midterm Exam

- Not, repeat not next week!
- Wed March 23, 7:30 PM, right here
- Closed book, but
- You can bring one *cheatsheet*
 - one side of one (8.5 by 11) sheet, handwritten by you
- P.S. No calculators, laptops, Palm Pilots, talking watches, etc.

Helpful review session

- Tuesday March 22, 7:30 PM, COS 105
- Not a canned presentation
- Driven by your questions

28