

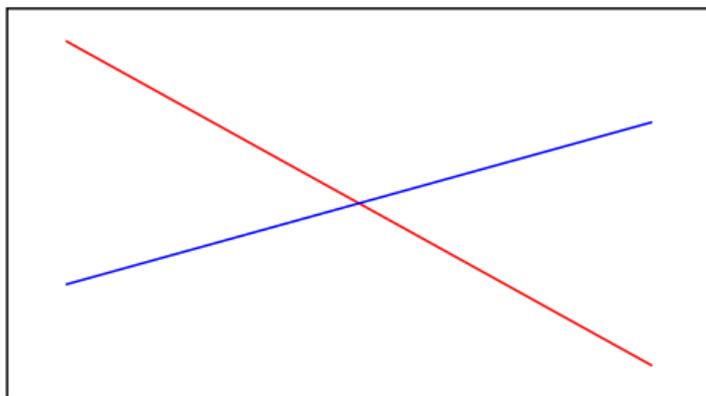
# Linear Systems

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COS 302, Fall 2020



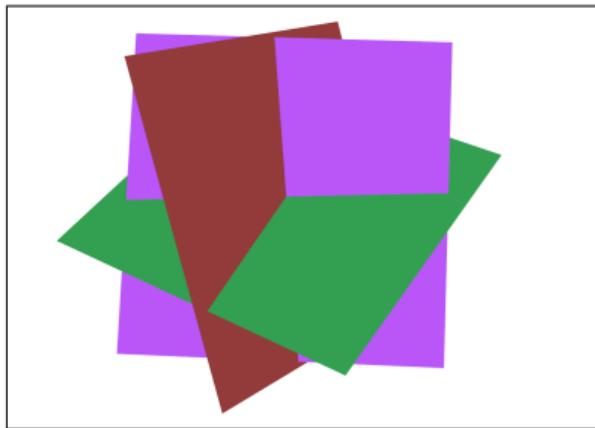
# Linear Systems

- Simultaneously satisfy a set of *linear* equations
- In 2D, a linear equation in 2 variables defines a line
  - 2 equations *might* intersect in 1 point, giving a unique solution



# Linear Systems

- Simultaneously satisfy a set of *linear* equations
- In 3D, a linear equation in 3 variables defines a plane
  - 3 equations *might* intersect in 1 point, giving a unique solution



# Applications of Linear Systems

- Regression (“fitting a model to data”)
- Simulation (e.g., mass-spring systems)
- Analysis (e.g., how much stress is there in a beam in a building or bridge)
- Subroutine in other numerical algorithms (e.g., Newton’s method for optimization or implicit Euler method for differential equations)

# Linear Systems

$$a_{11}x_1 + a_{12}x_2 + a_{13}x_3 + \cdots = b_1$$

$$a_{21}x_1 + a_{22}x_2 + a_{23}x_3 + \cdots = b_2$$

$$a_{31}x_1 + a_{32}x_2 + a_{33}x_3 + \cdots = b_3$$

$$\vdots$$

$$\begin{bmatrix} a_{11} & a_{12} & a_{13} & \cdots \\ a_{21} & a_{22} & a_{23} & \cdots \\ a_{31} & a_{32} & a_{33} & \cdots \\ \vdots & \vdots & \vdots & \ddots \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \\ \vdots \end{bmatrix} = \begin{bmatrix} b_1 \\ b_2 \\ b_3 \\ \vdots \end{bmatrix}$$

# Linear Systems

- Solve  $A\mathbf{x} = \mathbf{b}$ , where  $A$  is an  $n \times n$  matrix and  $\mathbf{b}$  is an  $n \times 1$  column vector
- Can also talk about non-square systems where  $A$  is  $m \times n$ ,  $\mathbf{b}$  is  $m \times 1$ , and  $\mathbf{x}$  is  $n \times 1$ 
  - Usually *overdetermined* if  $m > n$ : “more constraints than unknowns”  
(Can look for “best” solution using *least squares*.)
  - *Underdetermined* if  $n > m$ : “more unknowns than constraints”  
(Can compute all solutions, as in textbook, but it is more common to look for “best” solution using *regularization*.)

# Singular Systems

- A square matrix  $A$  is *singular* if some row is linear combination of other rows
- Singular systems might have infinitely many solutions:

$$2x_1 + 3x_2 = 5$$

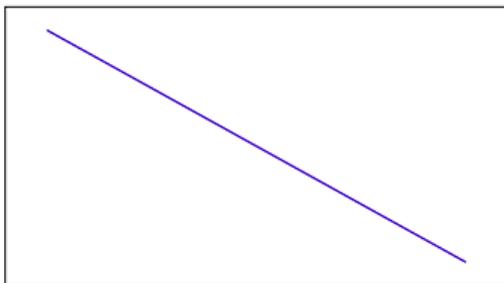
$$4x_1 + 6x_2 = 10$$

or no solutions:

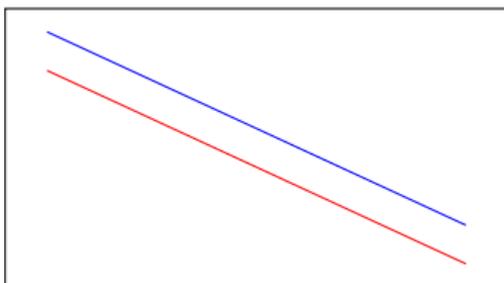
$$2x_1 + 3x_2 = 5$$

$$4x_1 + 6x_2 = 11$$

# Singular Systems

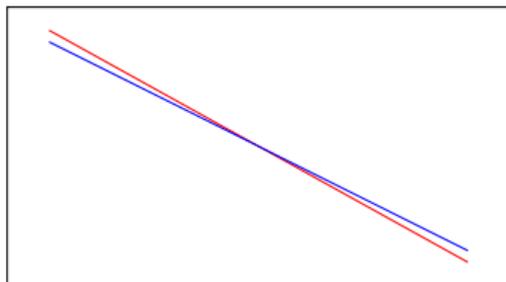


Singular with infinite solutions



Singular with no solutions

# Near-Singular Systems



Near-singular or ill-conditioned:  
noise in inputs, or roundoff error  
in computation, may result in large  
changes to solution

# Solving Linear Systems

- Seemingly the most direct way to solve a well-determined, square system is to use the matrix *inverse*:

$$Ax = b$$

$$A^{-1}Ax = A^{-1}b$$

$$x = A^{-1}b$$

Notes:

- The inverse of a square matrix need not exist.  
But if it does, it is unique and has the property  $AA^{-1} = A^{-1}A = I$ .
- Matrix multiplication is associative, so  $A^{-1}(Ax) = (A^{-1}A)x = Ix = x$ .
- Matrix multiplication is **not** commutative, so we were careful to multiply both  $Ax$  and  $b$  on the *left* by  $A^{-1}$ .

# Inverses and Linear Systems

- In fact, using  $\mathbf{x} = \mathbf{A}^{-1}\mathbf{b}$ , and computing the inverse, is usually a bad idea:
  - Inefficient
  - Prone to roundoff error
- Linear solver algorithms
  - *Direct*: nested loops over matrix, get solution at end
  - *Iterative*: get approximate answer, then each iteration improves it
- In fact compute inverse using linear solver
  - Solve  $\mathbf{A}\mathbf{x}_i = \mathbf{b}_i$ , where  $\mathbf{b}_i$  are columns of identity,  $\mathbf{x}_i$  are columns of inverse
  - Many solvers can solve several Right Hand Sides (RHS) at once

# Gauss-Jordan or Gaussian Elimination

- Simple-to-understand *direct* solver (though not used in practice)
- Transforms matrix to *reduced row-echelon form*<sup>\*</sup>, while simultaneously manipulating one or more right-hand side(s)
  - \* First nonzero entry in each row is 1, it's to the right of the 1 in the row above, and it's the only nonzero entry in that column.
- Fundamental operations:
  1. Replace one row with linear combination of it and other rows
  2. Interchange two rows
  3. Re-label two variables (interchange two columns)
- Simplest variant uses only #1 operations but numerical stability improved by adding #2 (partial pivoting) or both #2 and #3 (full pivoting).

# Gaussian Elimination

- Solve:

$$2x_1 + 3x_2 = 7$$

$$4x_1 + 5x_2 = 13$$

- Only care about numbers – form “tableau” or “augmented matrix”:

$$\left[ \begin{array}{cc|c} 2 & 3 & 7 \\ 4 & 5 & 13 \end{array} \right]$$

- Could have multiple right-hand sides (e.g. for computing an inverse)

# Gaussian Elimination

- Given:

$$\left[ \begin{array}{cc|c} 2 & 3 & 7 \\ 4 & 5 & 13 \end{array} \right]$$

- Goal: reduce this to the following form:

$$\left[ \begin{array}{cc|c} 1 & 0 & ? \\ 0 & 1 & ? \end{array} \right]$$

and read off answer from right column

# Gaussian Elimination

$$\left[ \begin{array}{cc|c} 2 & 3 & 7 \\ 4 & 5 & 13 \end{array} \right]$$

- Basic operation: replace any row by linear combination with any other row
- Here, replace first row with  $\frac{1}{2}$  times first row plus 0 times second row:

$$\left[ \begin{array}{cc|c} 1 & 3/2 & 7/2 \\ 4 & 5 & 13 \end{array} \right]$$

# Gaussian Elimination

$$\left[ \begin{array}{cc|c} 1 & 3/2 & 7/2 \\ 4 & 5 & 13 \end{array} \right]$$

- Replace second row with -4 times first row plus 1 times second row:

$$\left[ \begin{array}{cc|c} 1 & 3/2 & 7/2 \\ 0 & -1 & -1 \end{array} \right]$$

- Negate second row:

$$\left[ \begin{array}{cc|c} 1 & 3/2 & 7/2 \\ 0 & 1 & 1 \end{array} \right]$$

# Gaussian Elimination

$$\left[ \begin{array}{cc|c} 1 & 3/2 & 7/2 \\ 0 & 1 & 1 \end{array} \right]$$

- Add  $-3/2$  times second row to first row:

$$\left[ \begin{array}{cc|c} 1 & 0 & 2 \\ 0 & 1 & 1 \end{array} \right]$$

- ...aaaaand we're done. This is in reduced row-echelon form!

# Gaussian Elimination Analysis

- For each row  $i$ :
  - Multiply row  $i$  by  $1/a_{ii}$
  - For each other row  $j$ :
    - Add  $-a_{ji}$  times row  $i$  to row  $j$
- Innermost loop executed  $n(n - 1)$  times, and requires  $n + 1$  additions and multiplications
  - Asymptotic behavior: when  $n$  is large, that's about  $2n$  arithmetic operations in inner loop, or about  $2n^3$  total
- Can solve any number of RHS at once (but must be known ahead of time)

# Pivoting

- Consider this system:

$$\left[ \begin{array}{cc|c} 0 & 1 & 2 \\ 2 & 3 & 8 \end{array} \right]$$

- Immediately run into problem: algorithm wants us to divide by zero!
- More subtle version:

$$\left[ \begin{array}{cc|c} 0.001 & 1 & 2 \\ 2 & 3 & 8 \end{array} \right]$$

- Small diagonal (“pivot”) elements bad!
  - Swap in larger element from somewhere else...

# Partial Pivoting

$$\left[ \begin{array}{cc|c} 0 & 1 & 2 \\ 2 & 3 & 8 \end{array} \right]$$

- Swap rows 1 and 2:

$$\left[ \begin{array}{cc|c} 2 & 3 & 8 \\ 0 & 1 & 2 \end{array} \right]$$

- Now continue:

$$\left[ \begin{array}{cc|c} 1 & 3/2 & 4 \\ 0 & 1 & 2 \end{array} \right] \quad \left[ \begin{array}{cc|c} 1 & 0 & 1 \\ 0 & 1 & 2 \end{array} \right]$$

# Real-World Linear Solvers

- In practice, partial pivoting widely implemented
- To speed things up, don't go all the way to reduced row-echelon form
  - Also, it would be nice to be able to specify a new  $\mathbf{b}$  *after* the expensive computation on  $\mathbf{A}$  has been done...

# Triangular Systems

- Special case: *lower-triangular system*

$$\left[ \begin{array}{cccc|c} a_{11} & 0 & 0 & \cdots & b_1 \\ a_{21} & a_{22} & 0 & \cdots & b_2 \\ a_{31} & a_{32} & a_{33} & \cdots & b_3 \\ \vdots & \vdots & \vdots & \ddots & \vdots \end{array} \right]$$

# Triangular Systems

- Solve by forward substitution

$$\left[ \begin{array}{cccc|c} a_{11} & 0 & 0 & \cdots & b_1 \\ a_{21} & a_{22} & 0 & \cdots & b_2 \\ a_{31} & a_{32} & a_{33} & \cdots & b_3 \\ \vdots & \vdots & \vdots & \ddots & \vdots \end{array} \right]$$

$$x_1 = \frac{b_1}{a_{11}} \quad x_2 = \frac{b_2 - a_{21}x_1}{a_{22}} \quad x_3 = \frac{b_3 - a_{31}x_1 - a_{32}x_2}{a_{33}} \quad \dots$$

# Triangular Systems

- Similarly, *upper triangular* systems solved by back-substitution

$$\left[ \begin{array}{cccc|c} a_{11} & a_{12} & a_{13} & a_{14} & b_1 \\ 0 & a_{22} & a_{23} & a_{24} & b_2 \\ 0 & 0 & a_{33} & a_{34} & b_3 \\ 0 & 0 & 0 & a_{44} & b_4 \end{array} \right]$$

$$x_4 = \frac{b_4}{a_{44}} \quad x_3 = \frac{b_3 - a_{34}x_4}{a_{33}} \quad x_2 = \frac{b_2 - a_{24}x_4 - a_{23}x_3}{a_{22}} \quad \dots$$

# LU Decomposition

- Both special cases can be solved in time  $\sim n^2$
- This motivates a factorization approach:
  - Find a way of writing  $A$  as  $LU$ , where  $L$  is lower-triangular and  $U$  is upper-triangular
  - $Ax = b \Rightarrow LUx = b \Rightarrow Ly = b \Rightarrow Ux = y$
  - Time to factor matrix dominates computation
  - Turns out to be faster than Gaussian elimination (but still cubic in  $n$ )
  - Can solve for new  $b$  at any time after factorization
- Real-world, general-purpose linear solvers (such as `numpy.linalg.solve`) use LU Decomposition with partial pivoting
  - ...but for big ( $n \sim$  thousands or millions) problems, approximate iterative methods are common