



COS 126 Exam Review

- Exams overview
- Example programming exam
- Example written exam questions (part 1)

Exams tab on the booksite

See Exams tab for full details and old exams.

- Read carefully *before* each exam.
- Policies are the contract between us and you.

Watch this space for details

Policies (written exam).

- Closed book/notes/computer.
- 1 page (one side) cheatsheet.
- [two sides for Exam 2.]

The screenshot shows the 'EXAMS' page on the COS126 website. The page includes a navigation bar with links for Syllabus, Meetings, Lectures, Precepts, Assignments, Exams, and Help!. Below the navigation bar, there is a section titled 'EXAMS' with introductory text and a table of exam dates and links for previous semesters.

	PROGRAMMING EXAM 1	WRITTEN EXAM 1	PROGRAMMING EXAM 2	WRITTEN EXAM 2
FALL 2018	In class on Oct. 11th.	In class on Oct. 18th.	In class on Dec. 6th.	In class on Dec. 13th.
SPRING 2018	Programming Exam 1 Files Rainfall.java , Precipitation.java	Written Exam 1 Solutions	EXAM AND SOLUTION POSTED HERE SOON	EXAM AND SOLUTION POSTED HERE SOON
FALL 2017	Programming Exam 1 Files Submit! Prices.java , MovingAverage3.java	Written Exam 1 Solutions	Programming Exam 2 Files Submit! Path.java	Written Exam 2 Solutions

Things to remember about inclass exams

We know that you don't have much time.

- Exams are 50 minutes.
- "One page" programming exams.
- Five-minute questions on written exams.

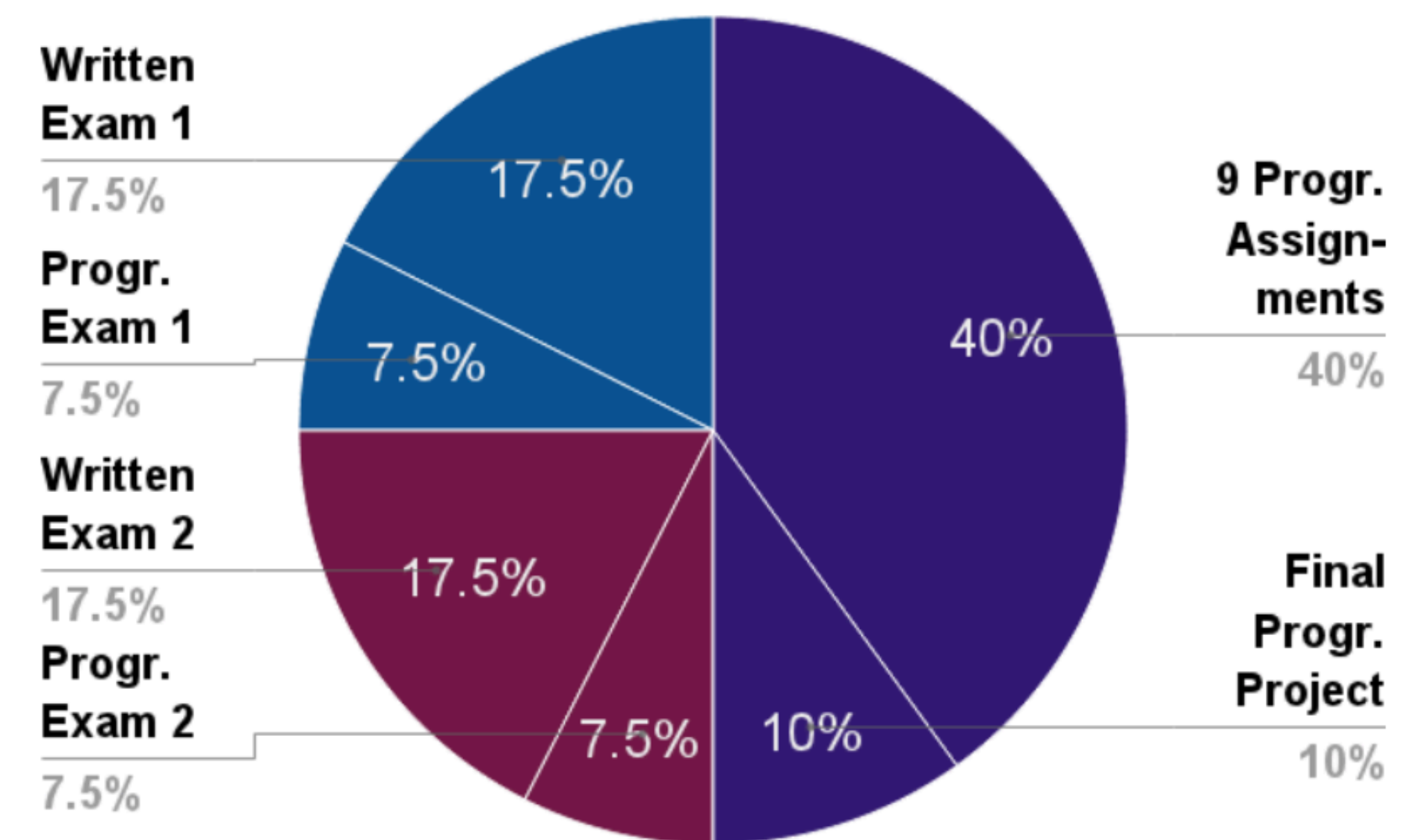
We have to grade the exams.

- 400+ exams.
- No open-ended questions.
- Fully prepared rubrics.

Old exams are not completely reliable.

- Course offerings differ slightly.
- We have made mistakes in the past.

Exams are only part of the story.



Written Exam Logistics

The first exam is on Thursday Oct. 18.

You don't all fit in this room.

- Pay attention and know where to go.
- Arrive early.
- No calculator/phone/computer/headphones

Advice.

- Review lectures/reading.
- Try an old exam (untimed).
- Try another one (timed).
- Review a few more.



Example question: Input and output

Q. Do you understand basic ways of communicating with your programs ?

Ex. (S2011 Q4) Give the results of invoking this program with the given commands.

```
public class Q4
{
    public static void main(String[] args)
    {
        int curr = StdIn.readInt();
        StdOut.print(curr + " ");
        int prev = curr;
        while (!StdIn.isEmpty())
        {
            curr = StdIn.readInt();
            StdOut.print((prev + curr) / 2 + " ");
            prev = curr;
        }
        StdOut.println();
    }
}
```

```
% more input.txt
```

```
2 4 6 8 10 12 8 2
```

```
% java Q4 < input.txt
```

```
2 3 5 7 9 11 10 5
```

```
% java Q4 < input.txt | java Q4
```

```
2 2 4 6 8 10 10 7
```

Note: It prints the first number, then the average of each number and its predecessor.

Example question: Functions

Q. Do you understand basic mechanisms for invoking functions ?

Ex. (S2018 Q7) Give the contents of the array a[] after executing the given code.

```
public static int halve1(int x)
{
    x = x / 2;
    return x;
}

public static void halve2(int[] a)
{
    for (int i = 0; i < a.length; i++)
    {
        halve1(a[i]);
        a[i] = halve1(a[i]);
    }
}
```

```
int[] a = { 16, 32, 48, 64 };
halve2(a);
```

8 16 24 32 NOT 4 8 12 16

```
int[] a = { 16, 32, 48, 64 };
halve2(a);
halve2(a);
```

4 8 12 16

Example question: Functions

Ex. (S2018 Q7) Give the contents of the array `a[]` after executing the given code.

```
public static void halve3(int[] a)
{
    int n = a.length;
    int[] b = new int[n/2];
    for (int i = 0; i < n/2; i++)
        b[i] = a[i];
    a = b;
}
```

```
int[] a = { 16, 32, 48, 64 };
halve3(a);
halve3(a);
```

16 32 48 64

Example question: Recursion

Q. Can you figure out the effect of a simple recursive program (or two)?

Ex. (Fall 2017 Q5) Fill in the values returned by these mutually recursive functions:

```
public static int mystery1(int n)
{
    if (n == 0) return 0;
    else return mystery2(n - 1);
}
public static int mystery2(int n)
{
    if (n == 0) return 1;
    else return mystery1(n - 1);
}
```

n	mystery1(n)	mystery2(n)
0	0	1
1	1	0
2	0	1
3	1	0
4	0	1
5	1	0

Write *one line of code* that could replace the body of `mystery1`.

return n % 2;

Example question: Binary operations

Q. Why is ~ 0 equal to -1 and not 1 ? (Fall 2014 Q1B)

A (wrong).

\sim is "not"

0 is "false"

"not false" is "true"

"true" is 1

A (correct).

\sim is **BITWISE** "not"

0 is 000000000000000000000000000000000000

~ 0 is 111111111111111111111111111111111111

111111111111111111111111111111111111 is -1 (2s complement)

Example question: TOY/number representation

Q. (Fall 2013 Q8) Consider this sequence of TOY instructions:

7 1 0 1

$R[1] = 0001$

2 2 0 1

$R[2] = R[0] - R[1]$ sets R[2] to all 1s

4 7 7 2

$R[7] = R[7] \wedge R[2]$ sets R[7] to bitwise XOR of R[7] with all 1s

1 7 7 1

$R[7] = R[7] + R[1]$ adds 1 to R[7]

Q. What is the value of R[7] after this sequence if it was initially 0025 ?

000000000100101
111111111011010
111111111011011

FFDB

Q. In English, what does sequence do to R[7] ?

Negates it.

to negate a 2s complement number: *flip its bits and add 1*

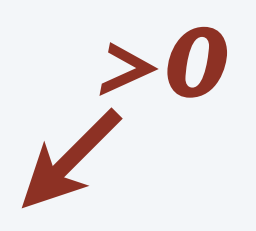
Example question: TOY

Q. Can you simulate the effect of a very simple TOY program?

Ex. (Fall 2016 Q7) Suppose that you load the following into memory locations 10-17 of TOY, set the PC to 10, and press RUN. Give the result in 01 when 00 is initially **0001**.

```
10: 8A00 R[A] <- M[00]
11: 7101 R[1] <- 1
12: 221A R[2] <- R[1] - R[A]
13: D216 if (R[2] > 0) PC <- 16
14: 1111 R[1] <- R[1] + R[1]
15: C012 PC <- 12
16: 9101 M[01] <- R[1]
17: 0000 halt
```

PC	R[A]	R[1]	R[2]
10	0001		
11	0001	0001	
12	0001	0001	0000
13	0001	0001	0000
14	0001	0002	0000
12	0001	0002	0001
13	0001	0002	0001
16	0001	0002	0001
17	0001	0002	0001



Example question: TOY

Q. Can you simulate the effect of a simple TOY program?

Ex. (Fall 2016 Q7) Suppose that you load the following into memory locations 10-17 of TOY, set the PC to 10, and press RUN. Give the result in 01 when 00 is initially **0006**.

```
10: 8A00 R[A] <- M[00]
11: 7101 R[1] <- 1
12: 221A R[2] <- R[1] - R[A]
13: D216 if (R[2] > 0) PC <- 16
14: 1111 R[1] <- R[1] + R[1]
15: C012 PC <- 12
16: 9101 M[01] <- R[1]
17: 0000 halt
```

PC	R[1]	R[2]	PC	R[1]	R[2]
10			12	0004	FFFC
11	0001		13	0004	FFFC
12	0001	FFFA	14	0008	FFFC
13	0001	FFFA	12	0008	FFF9
14	0002	FFFA	13	0008	FFF9
12	0002	FFFE	14	0010	FFF9
13	0002	FFFE	12	0010	0004
14	0004	FFFE	13	0010	0004
			16	0010	0004

Annotations:

- Red arrow from **0001-0006** to R[2] at PC 11.
- Red arrow from **<0** to R[2] at PC 12.
- Red arrow from **NOT 0016** to R[1] at PC 13.
- Red arrow from **>0** to R[2] at PC 14.

Example question: TOY

Q. Can you **reason about** the effect of a simple TOY program?

Ex. (Fall 2016 Q7) Suppose that you load the following into memory locations 10-17 of TOY, set the PC to 10, and press RUN. Give the result in M[01] when M[00] is initially **1EAF**.

```
10: 8A00 R[A] <- M[00]
11: 7101 R[1] <- 1
12: 221A R[2] <- R[1] - R[A]
13: D216 if (R[2] > 0) PC <- 16
14: 1111 R[1] <- R[1] + R[1]
15: C012 PC <- 12
16: 9101 M[01] <- R[1]
17: 0000 halt
```

```
load limit from M[00]
x = 1
while (x <= limit)
{
    x = 2*x
}
store x to M[01]
```

```
1
2
4
8
10
20
40
80
100
...
2000
```

Good luck!

