

# **Point Set Alignment**

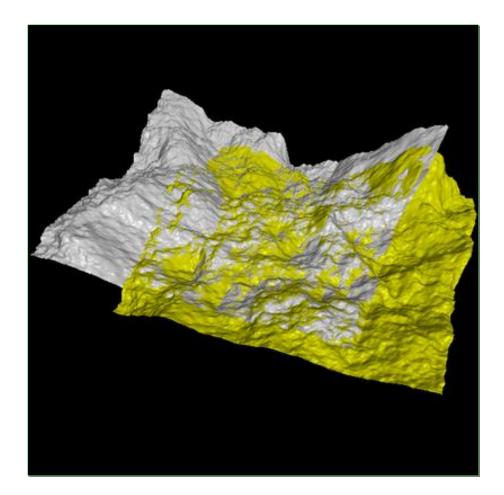
Thomas Funkhouser COS 526, Fall 2016

## Motivation



#### Point sets to be aligned

- ≻Range scans
- Image features
- Molecules
- etc.



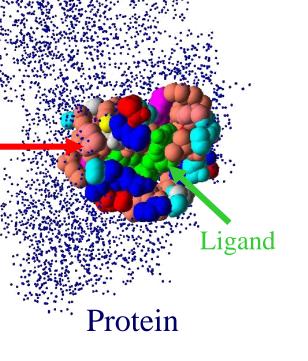
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- Range scans
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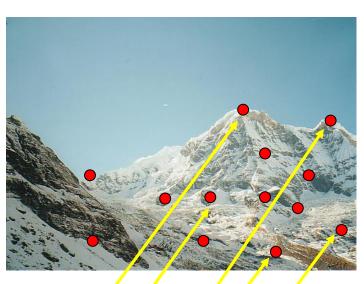


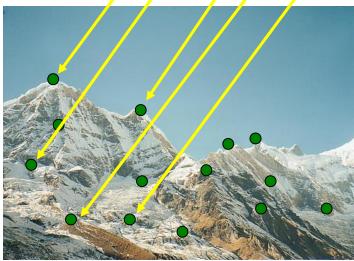
## Motivation



#### Point sets to be aligned

- Range scans
- Image features
- Molecules
- etc.

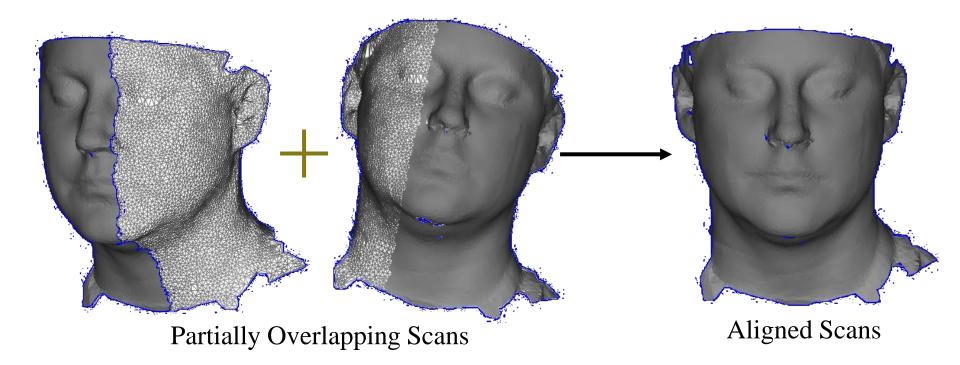








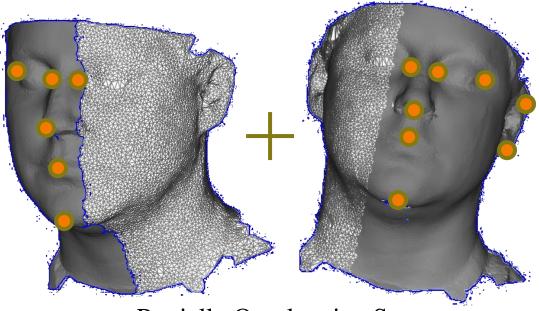
# Given two partially overlapping point sets, compute the transformation that merges the two



## **General Approach**



1. Find feature points

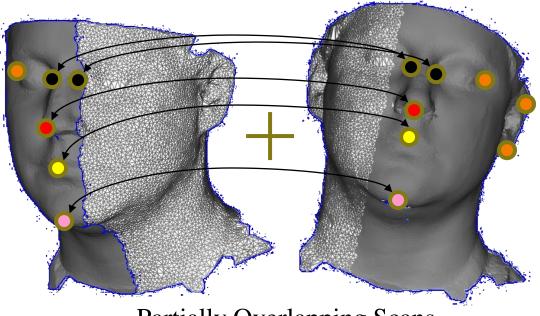


#### Partially Overlapping Scans

## **General Approach**



- 1. Find feature points
- 2. Establish correspondences

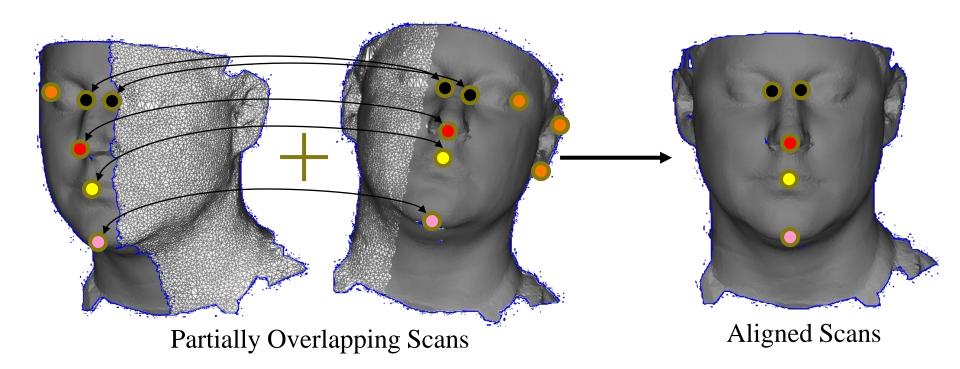


#### Partially Overlapping Scans

## **General Approach**



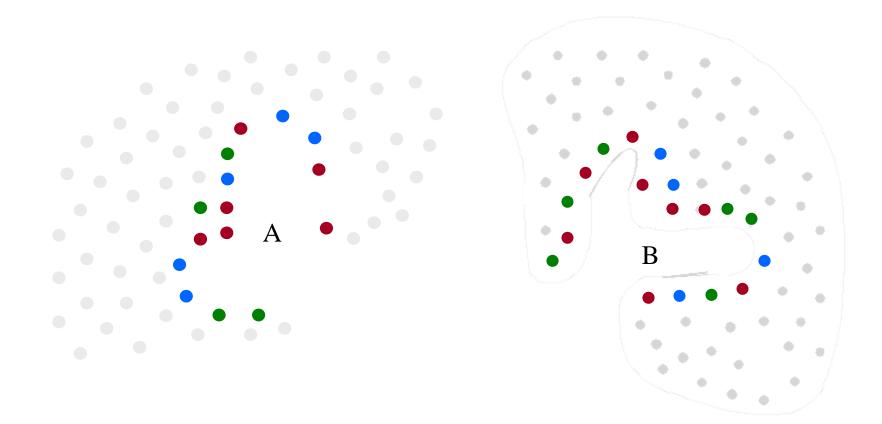
- 1. Find feature points
- 2. Establish correspondences
- 3. Compute the aligning transformation



#### Problem



#### Most problems require aligning a subset of features

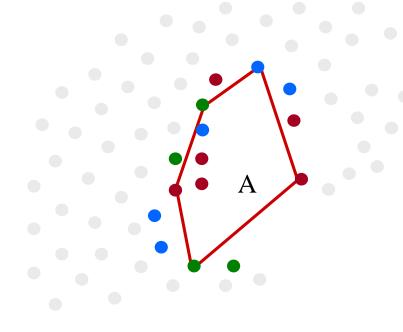


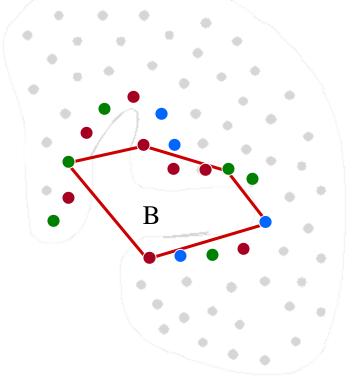
## Problem



Most problems require aligning a subset of features

- Find the maximal subsets of points that align with error E
- Find the minimum misalignment for any subset of a size S

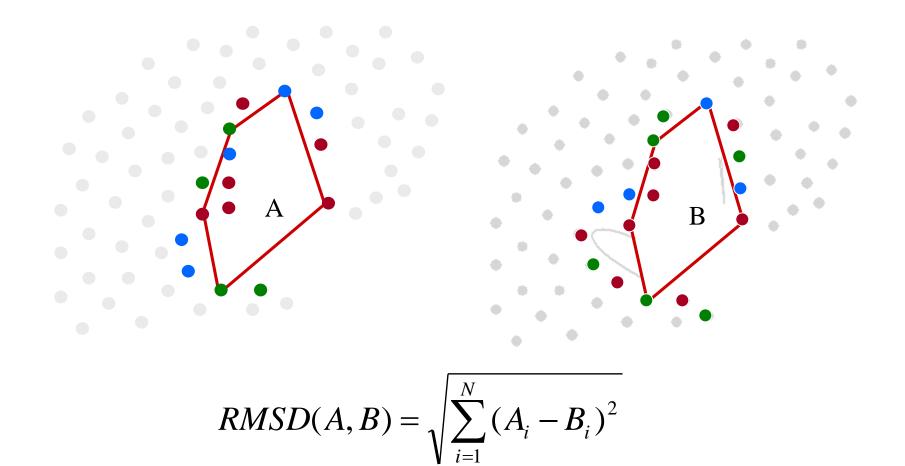




## **Observation I**



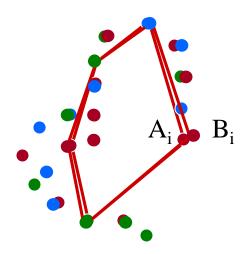
Calculating the aligning transformation is usually easy if correspondences are known (proposed)



## **Observation II**



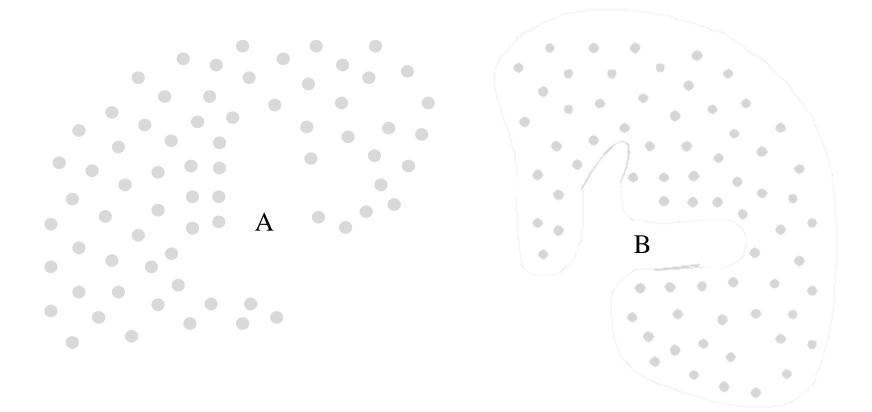
Calculating the correspondences is usually easy if the aligning transformation is known (proposed)



## Challenge



The challenge is to discover the correspondences and aligning transformation together



## Outline

#### Introduction

#### Point set matching

- Brute force search
- RANSAC
- Geometric hashing
- Assocation graphs
- Generalized Hough transform
- Iterative closest point

#### Methods used for RGB-D scanning

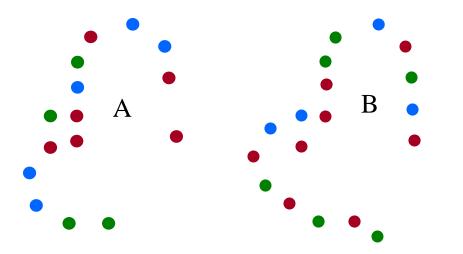
Discussion

## **Brute Force Search**



Simple method:

- Try all possible sets of point correspondences
- Score the alignment for each one



Problem:

 O(n<sup>m</sup>) possible sets of m correspondences among n points

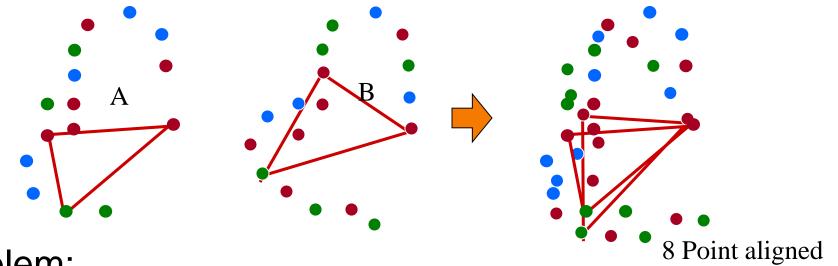
## **Brute Force Search**



RMSD = 3.1

Simple method:

- Try all possible sets of point correspondences
- Score the alignment for each one



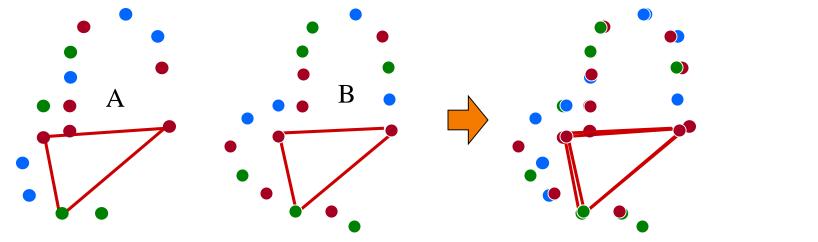
Problem:

 O(n<sup>m</sup>) possible sets of m correspondences among n points



Simple method:

- Try all possible sets of point correspondences
- Score the alignment for each one (e.g., RMSD)



Problem:

All points aligned RMSD = 0.2

 O(n<sup>m</sup>) possible sets of m correspondences among n points

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#### Point set matching

- Brute force search
- ►RANSAC
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- Generalized Hough transform
- Iterative closest point

#### Methods used for RGB-D scanning

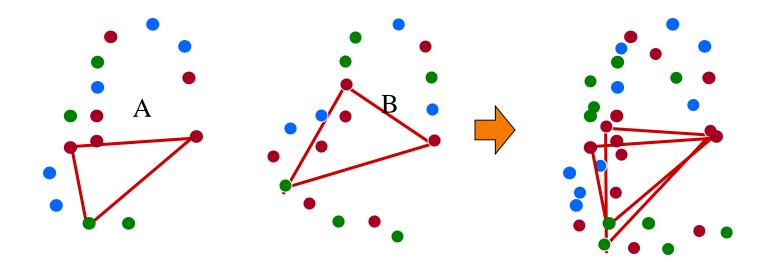
Discussion

## RANSAC



Randomly sample set of possible correspondences

- Randomly generate a small set of point correspondences
- Compute the aligning transformation for correspondences Score how well other points align after that transformation
- Remember the best transformation



## RANSAC



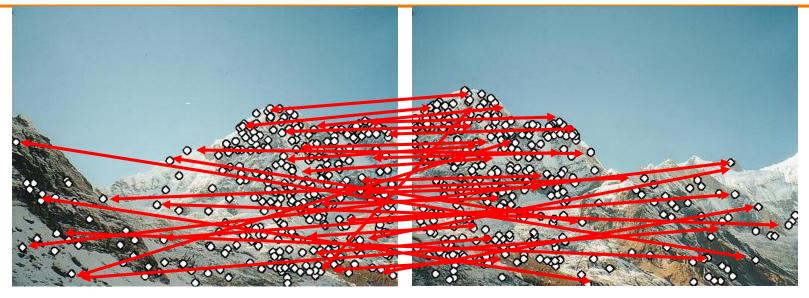
RANSAC loop:

- 1. Select k matches (at random)
- 2. Compute transformation T aligning those matches
- 3. Find *inlier matches* where  $d(p_i, Tp_i) < \varepsilon$
- 4. Re-compute T to align on all of its inliers
- 5. Re-find *inlier matches* where  $d(p_i, Tp_i) < \varepsilon$
- 6. T\*=T if has T largest set of inliers seen so far

Warp image by T\*

Composite images



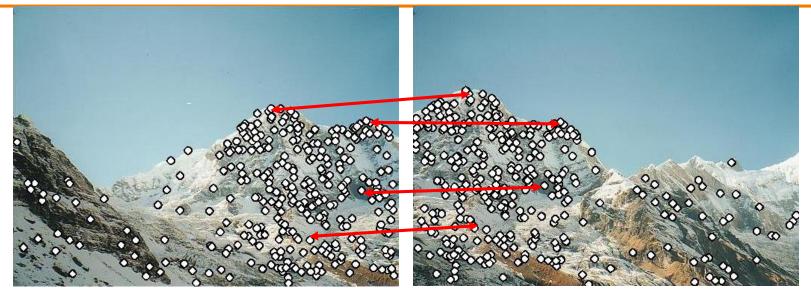


RANSAC loop:

- 1. Select four matches (at random)
- 2. Compute homography H aligning those matches
- 3. Find *inlier matches* where  $d(p_i, Hp_i) < \varepsilon$
- 4. Re-compute H to align on all of its inliers (least squares)
- 5. Re-find *inlier matches* where  $d(p_i, Hp_i) < \varepsilon$
- 6. H\*=H if has H largest set of inliers seen so far

Warp image by H\* and composite images



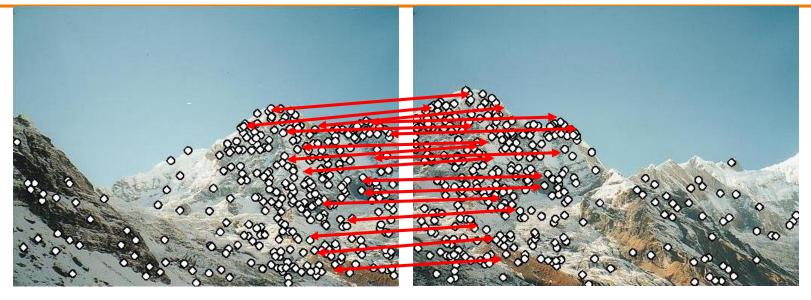


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Warp image by H\* and composite images



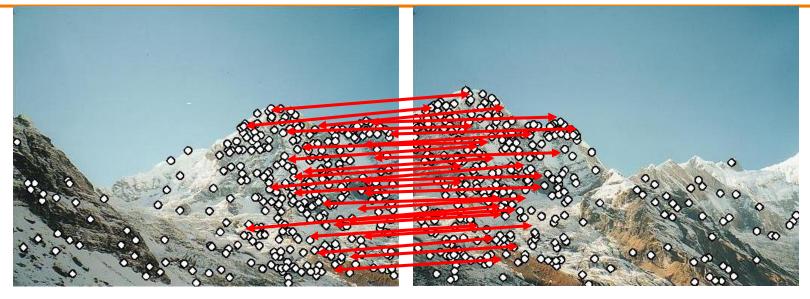


RANSAC loop:

- 1. Select four matches (at random)
- 2. Compute homography H aligning those matches
- 3. Find *inlier matches* where  $d(p_i, Hp_i) \le \varepsilon$
- 4. Re-compute H to align on all of its inliers (least squares)
- 5. Re-find *inlier matches* where  $d(p_i) < \varepsilon$
- 6. H\*=H if has H largest set of inliers seen so far

Warp image by H\* and composite images





RANSAC loop:

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Warp image by H\* and composite images





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Warp image by H\* and composite images

## Outline



#### Introduction

#### Point set matching

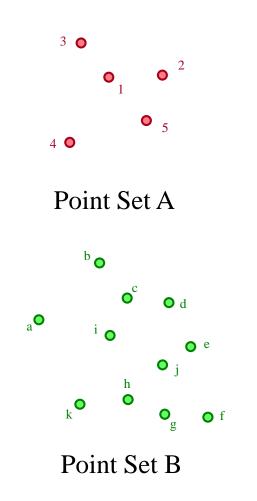
- Brute force search
- RANSAC
- ➤Geometric hashing
- Assocation graphs
- Generalized Hough transform
- Iterative closest point

#### Methods used for RGB-D scanning

Discussion

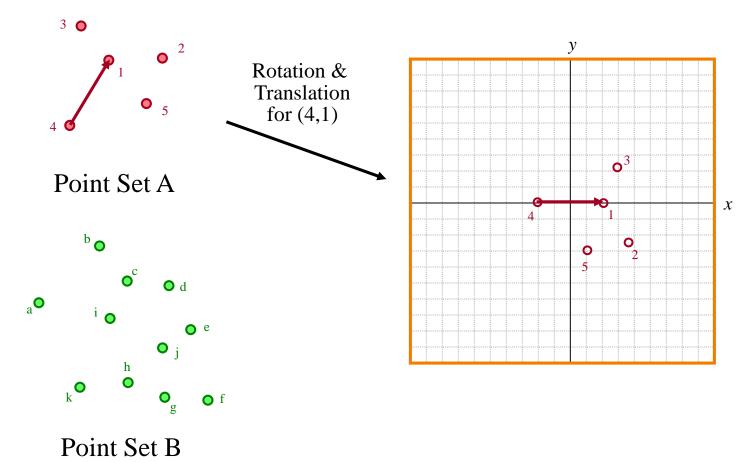


Discretize transformations and scoring



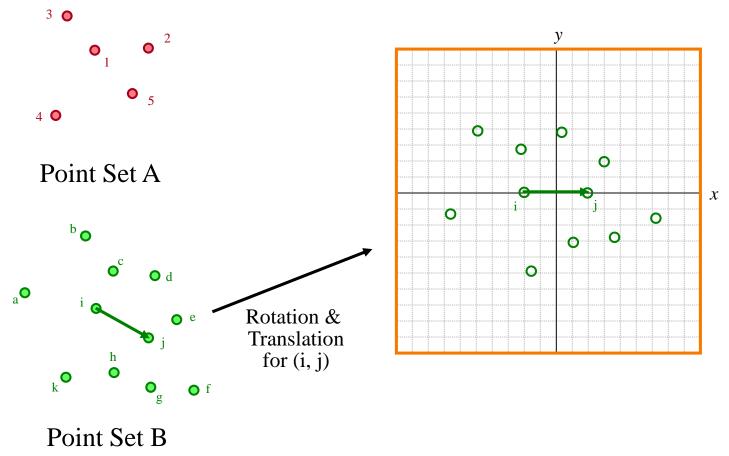


#### Discretize transformations and scoring



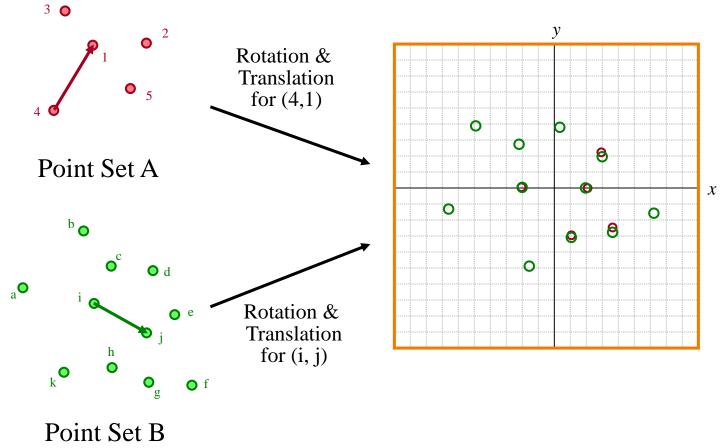


Discretize transformations and scoring



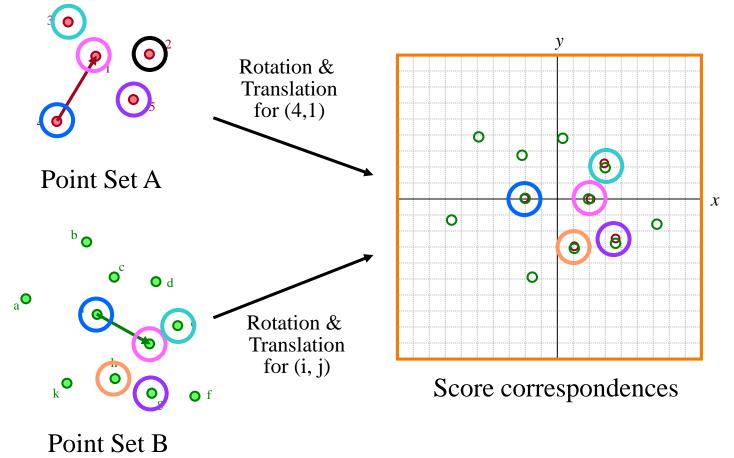


#### Discretize transformations and scoring

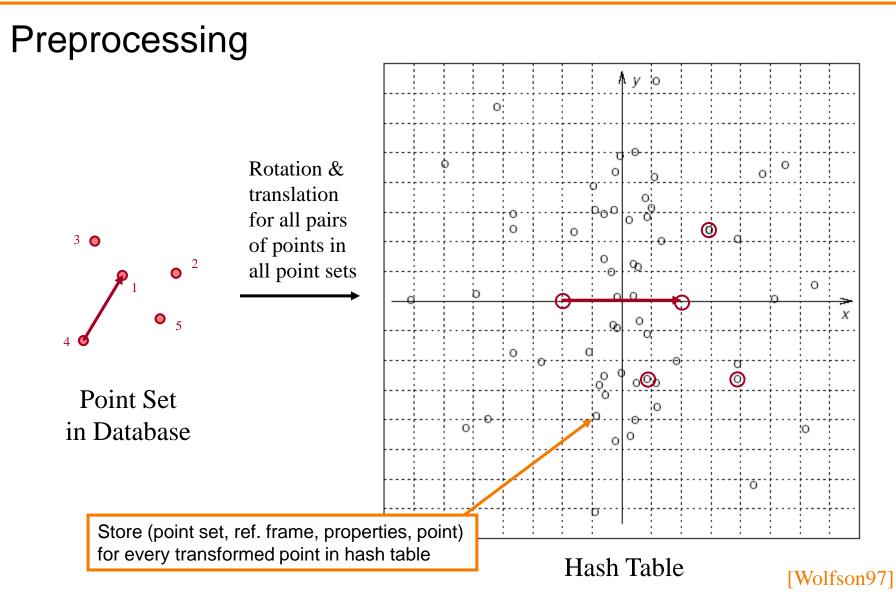




#### Discretize transformations and scoring

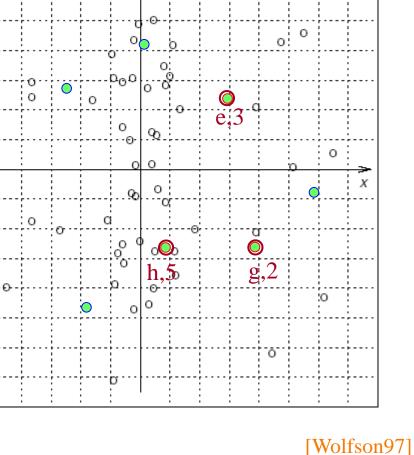








#### Query processing łу 0 Ġ۵, Ο Rotation & $\bigcirc$ <sup>b</sup> **o** 000 translation 0 $\bigcirc$ 0 0 for one pair o **O** d 0 of points a · 0· b o 0 90 h 0 0 k k 0 α **O** f Ο 000 g Point Set Ð 0 C 0 Query



lo



Preprocessing complexity

- O(n<sup>4</sup>) for n points per binding site
  - O(n<sup>3</sup>) possible triples \* O(n) transformations per triple

#### Query complexity

- O(m) \* binsize for m points in query binding site
  - 1 triple \* O(m) transformations per triple \* binsize hash processing per transformation



## Outline



#### Introduction

#### Point set matching

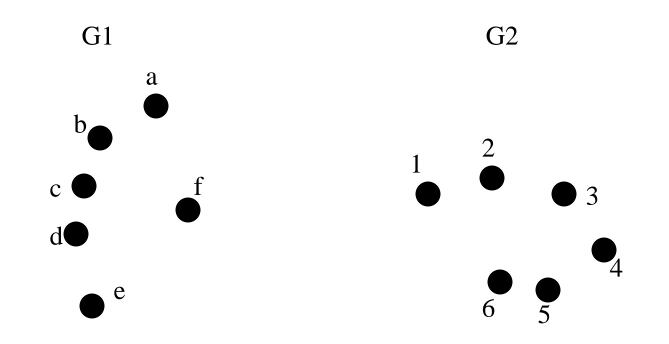
- Brute force search
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- Geometric hashing
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- Generalized Hough transform
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Discussion

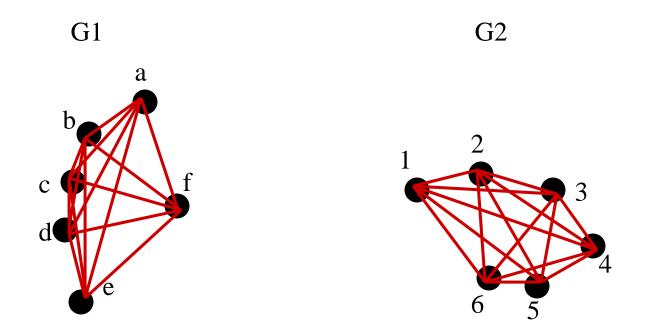
#### **Association Graphs**





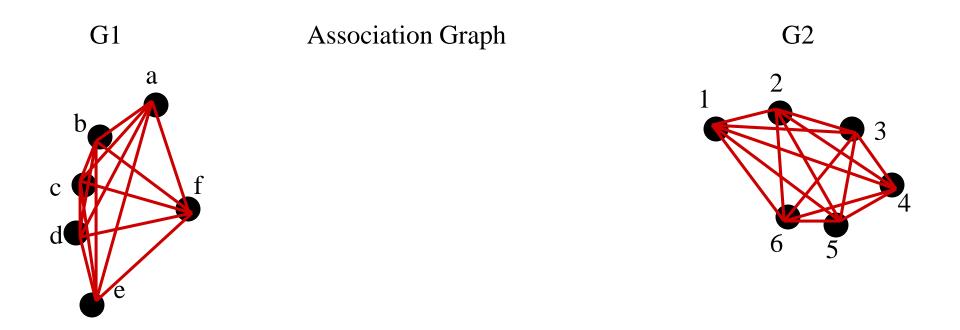
[Schmitt02, Brown82]





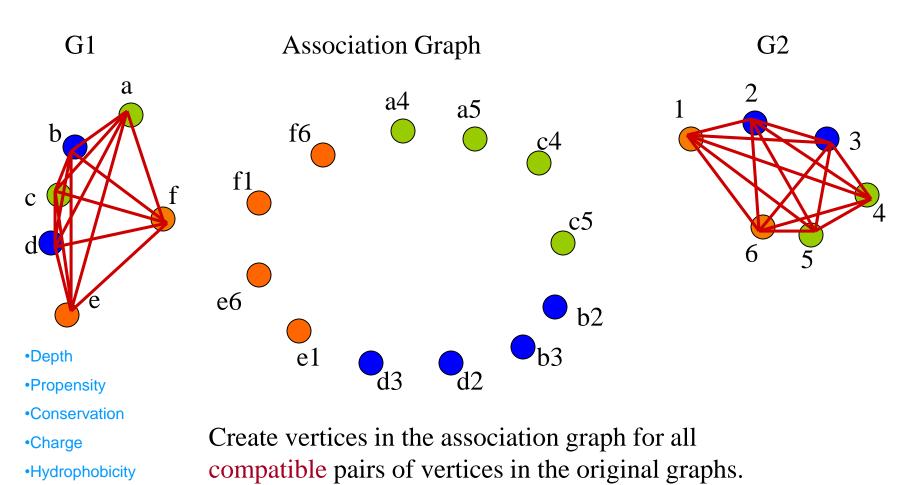
Represent both points sets as complete graphs (G1 and G2). (edges connect all pairs of vertices within each point set)





Create vertices in the association graph for all compatible pairs of vertices in the original graphs. This can lead to a large number of vertices.

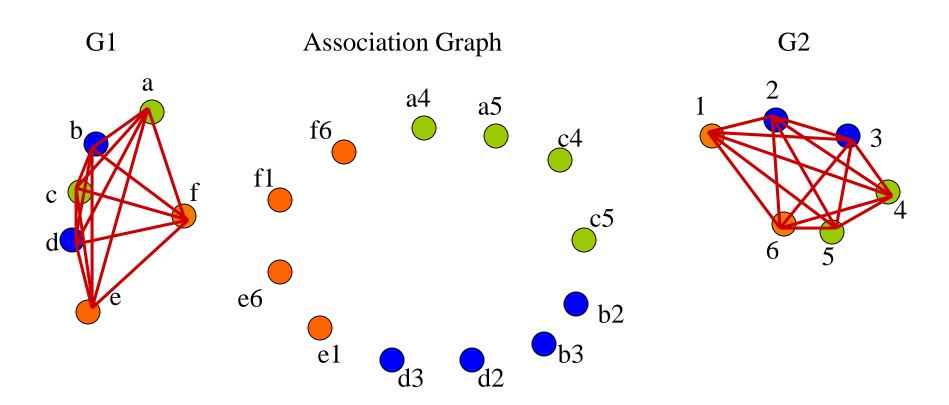




•Secondary structure type Compatibility could refer to chemical properties.

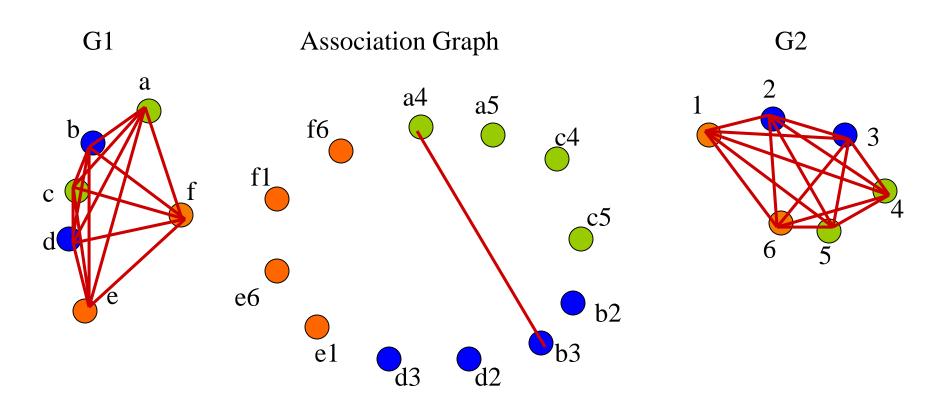
Destabilization





Create edges between (uv) and (wx) if the edges between (u) and (w) as well as between (v) and (x) match.

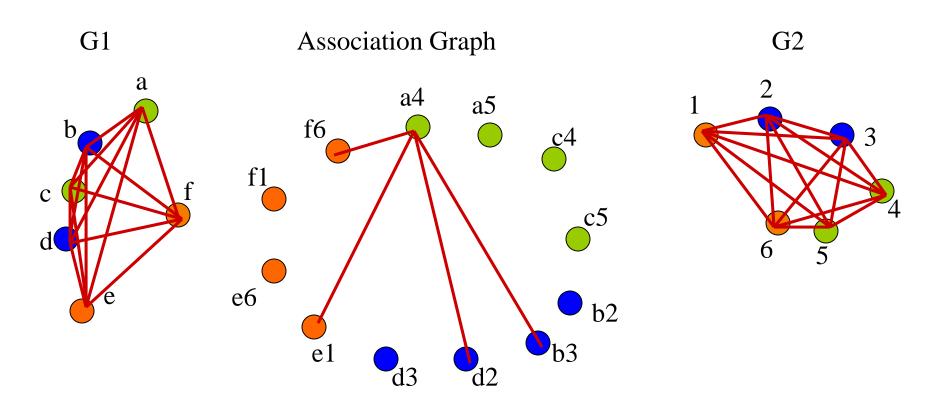




Create edges between (uv) and (wx) if the edges between (u) and (w) as well as between (v) and (x) match.

For this example, edge length is the only consideration

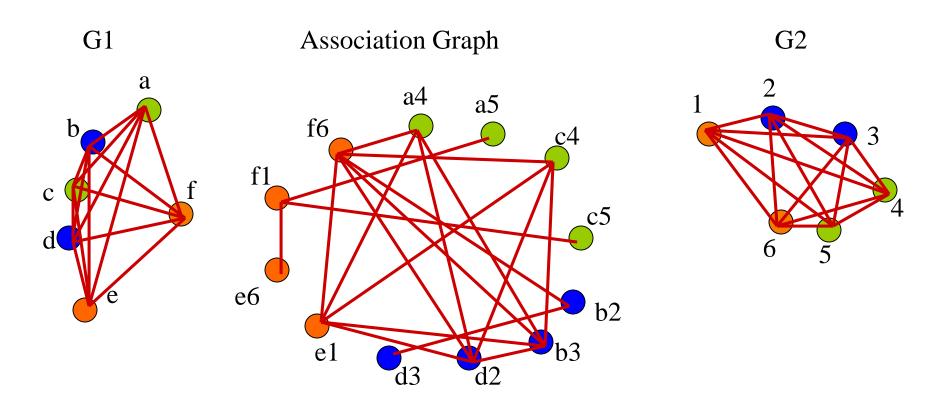




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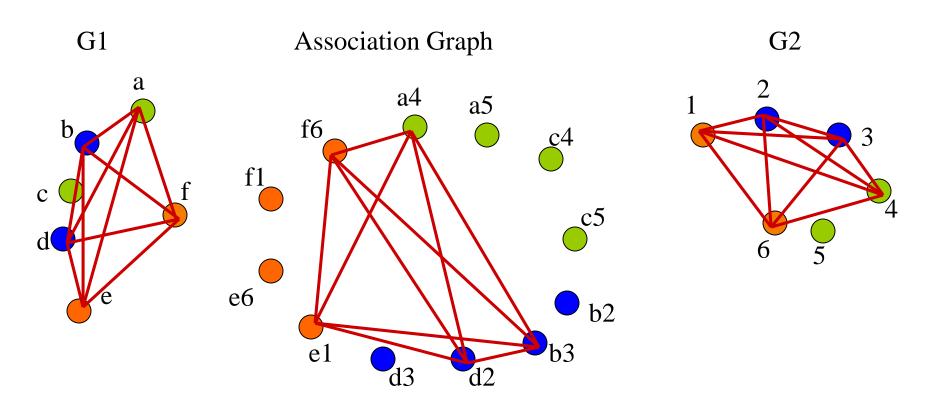




Create edges between (uv) and (wx) if the edges between (u) and (w) as well as between (v) and (x) match.

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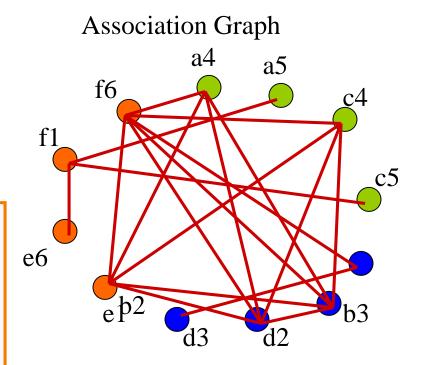


**Finding correspondences:** The the largest set of corresponding nodes in the same configuration is the maximal clique in the association graph



### Computational complexity:

- O(2<sup>n</sup>) for n points
- NP-complete
- Branch and bound algorithms



# Outline



### Introduction

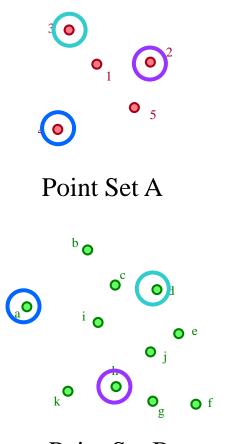
### Point set matching

- Brute force search
- RANSAC
- Geometric hashing
- Association graphs
- Generalized Hough transform
- Iterative closest point

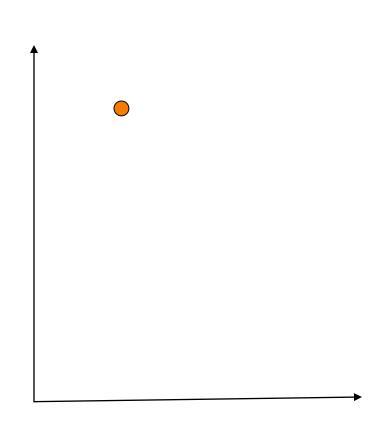
### Methods used for RGB-D scanning

Discussion



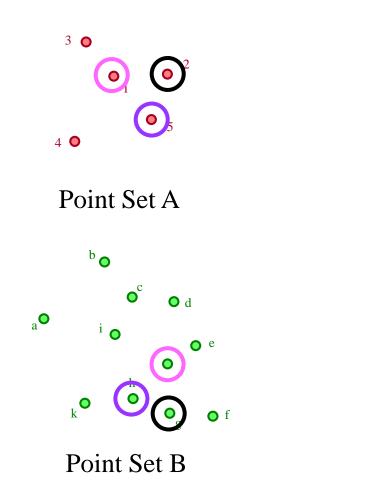


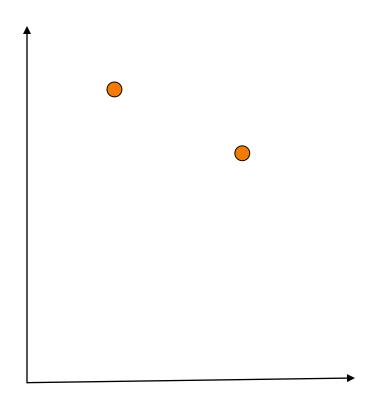




#### Hough Transformation Space

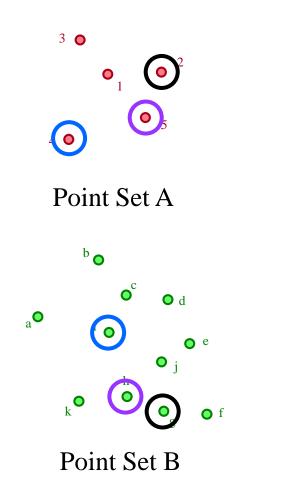


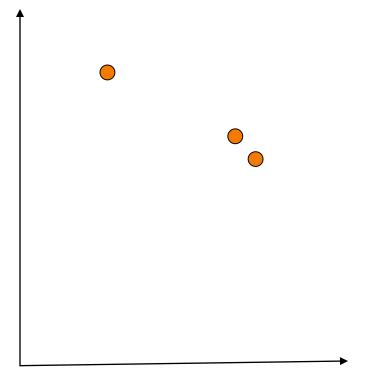




#### Hough Transformation Space

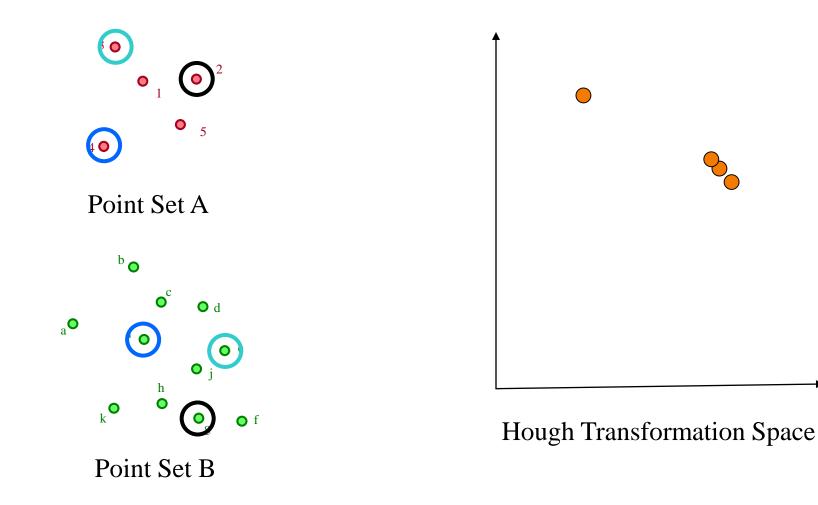
Vote for transformations



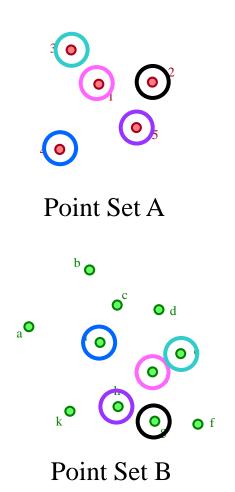


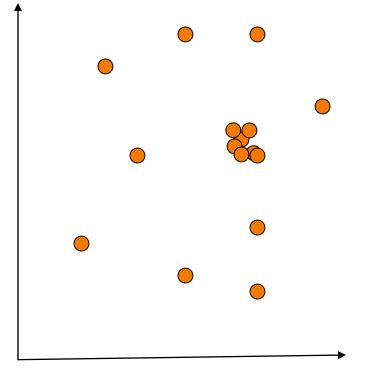
#### Hough Transformation Space

Vote for transformations



Vote for transformations





Hough Transformation Space





Simple to implement

- Can use grid to represent transformation space
- Expensive for high-dimensional transformations
  - Storage and number of samples is exponential in dimensionality of transformation space
    - Translation (3D)
    - Rotation (3D)
    - Translation & rotation (6D)
    - Translation & rotation & scale (7D)

# Outline



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### Point set matching

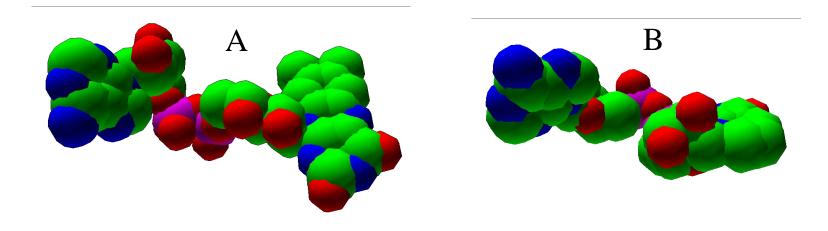
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### Methods used for RGB-D scanning

Discussion



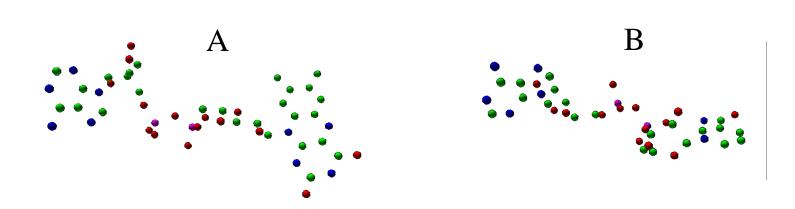
Given two point sets





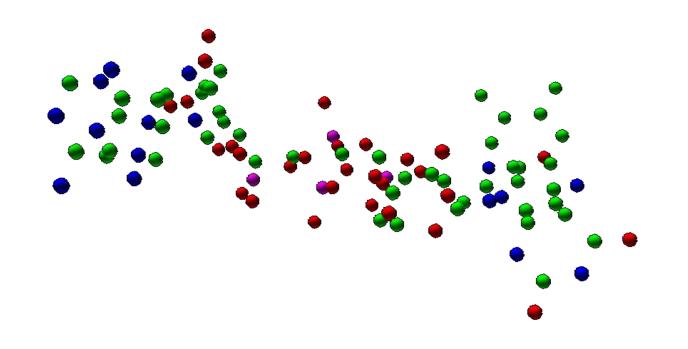


Given two point sets



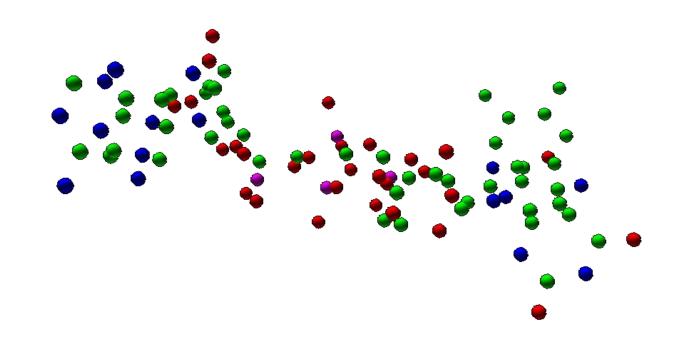


Given two point sets and an initial guess for the transformation that aligns them



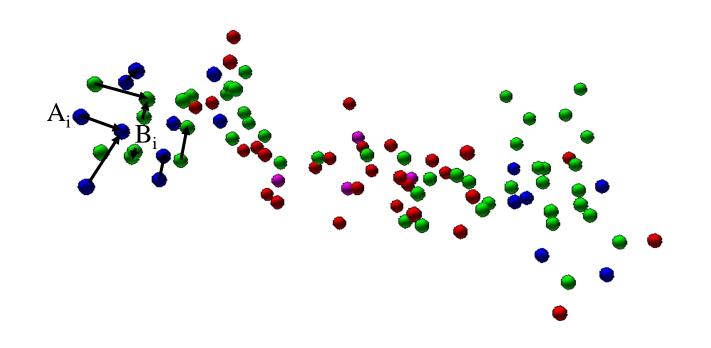


Assume closest points correspond



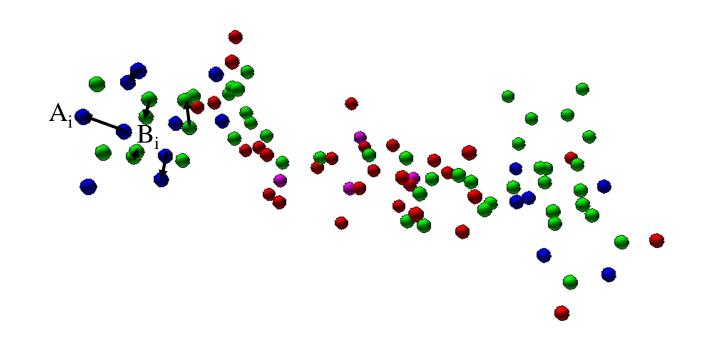


Assume closest points correspond:  $A \rightarrow B$ 



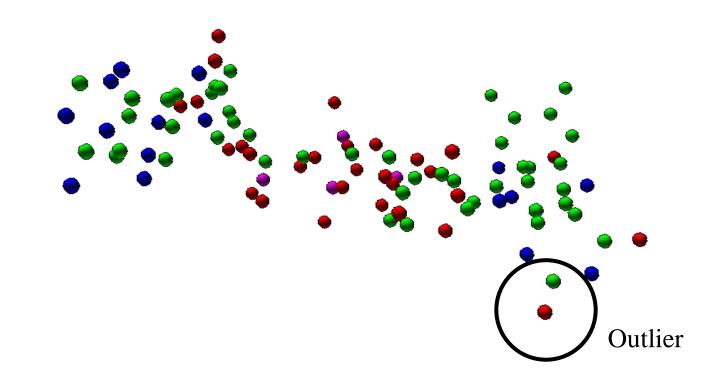


Assume closest points correspond:  $A \rightarrow B$  and  $B \rightarrow A$ 



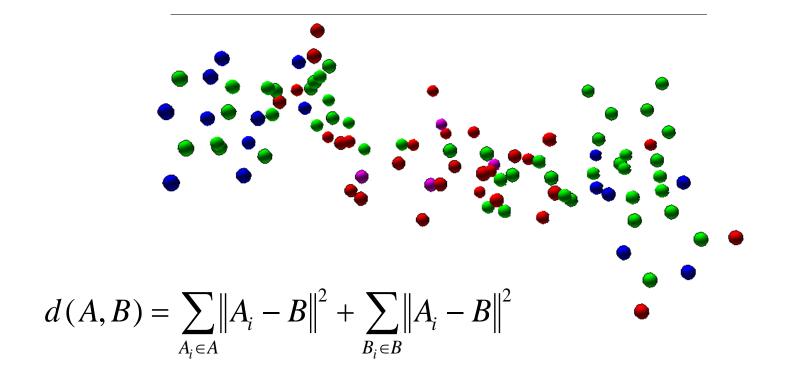


**Rejecting outliers** 





Find the transformation that optimally aligns proposed correspondences (superposition)





Iterate until convergence

- 1. Select source points (from one or both point sets)
- 2. Match to points in the other point set
- 3. Weight the correspondences
- 4. Reject outlier point pairs
- 5. Compute an error metric for the current transform
- 6. Minimize the error metric w.r.t. transformation

### **Computational complexity**

- O(k \* nlogn) for n points per binding site and k iterations
  - k iterations \* O(n) points \* O(logn) to find closest point

# Summary

#### Brute force

• Accurate, slow

### RANSAC

Approximate

### Geometric hashing

- Fast query, after slow preprocessing
- Distance threshold implicit in hash bucket sizes

#### Association graphs

- Expensive for large point sets
- Distance threshold for "associations"

#### Generalized Hough transform

• Requires lots of space/samples for high dimensional transformations

#### Iterative closest points

- Fast, in practice
- Requires good initial guess

