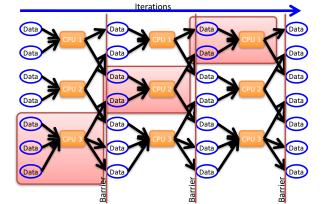
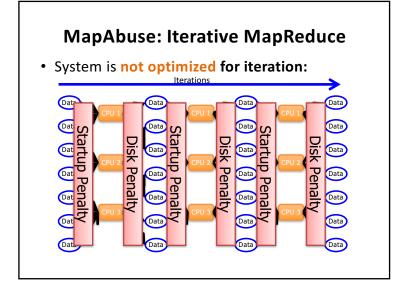
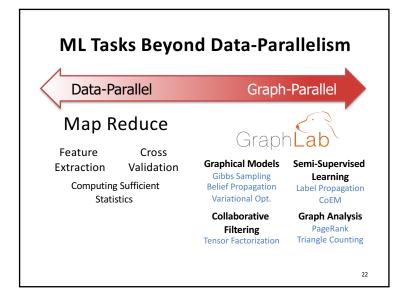


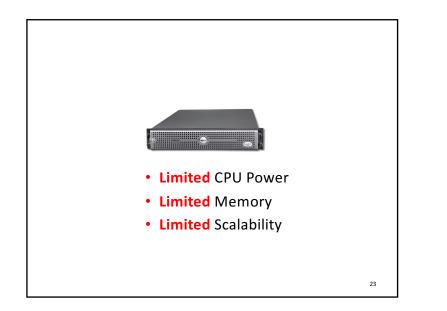
MapAbuse: Iterative MapReduce

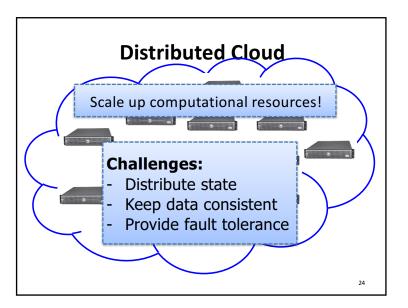
• Only a subset of data needs computation:

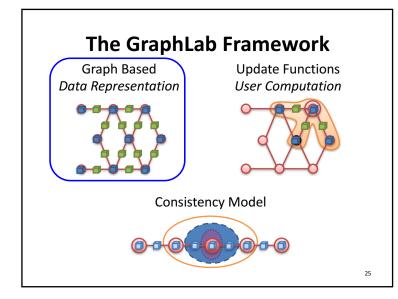


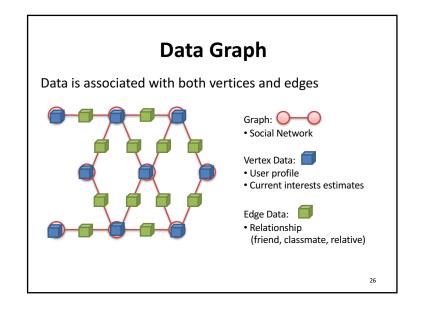


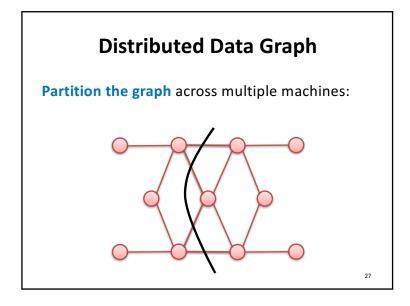


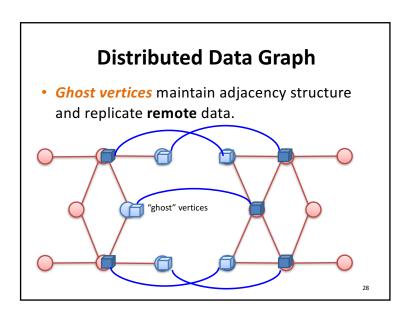


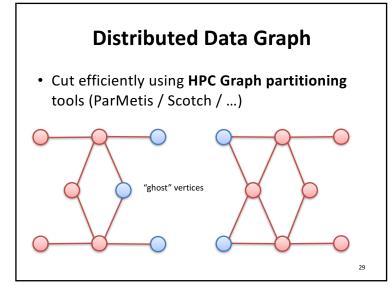


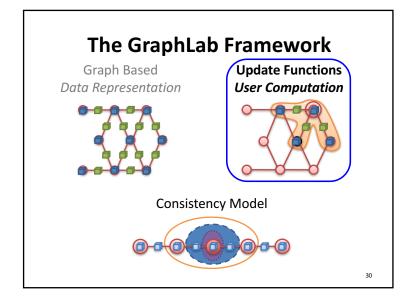


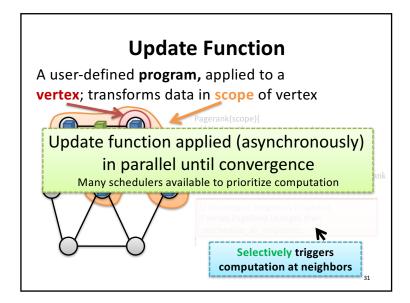


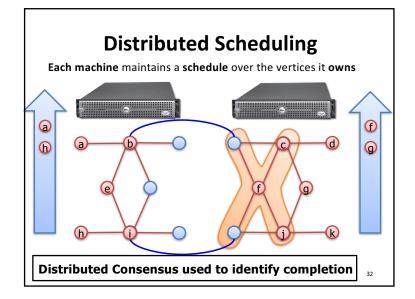


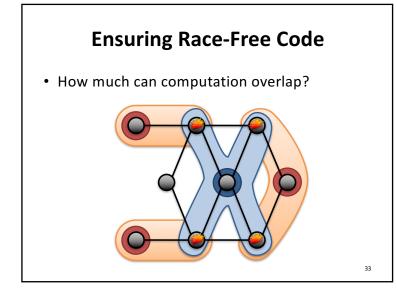


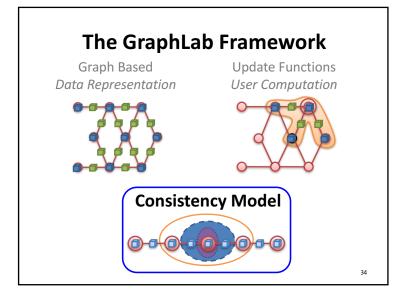


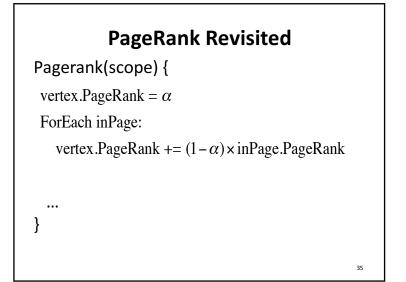


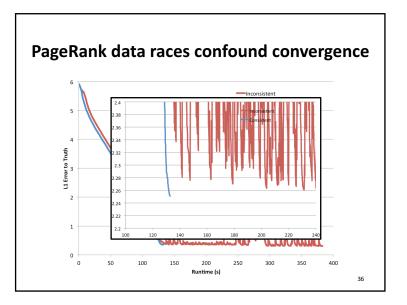












Racing PageRank: Bug

Pagerank(scope) {

vertex.PageRank = α

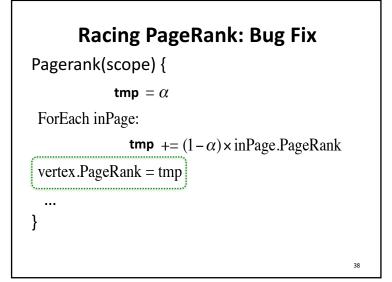
ForEach inPage:

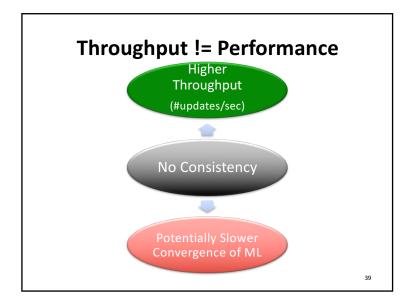
• • •

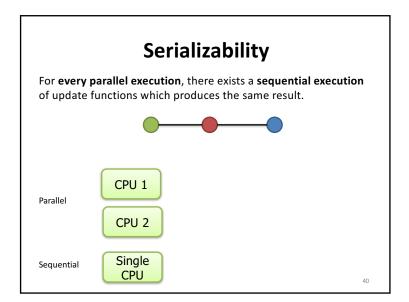
}

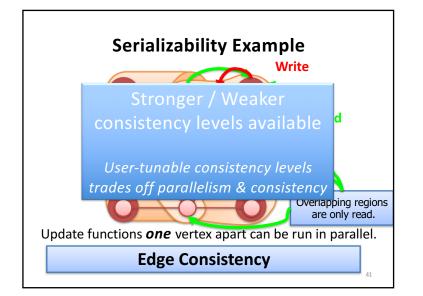
vertex.PageRank += $(1 - \alpha) \times inPage.PageRank$

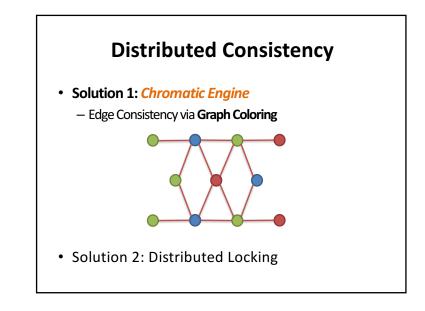
37

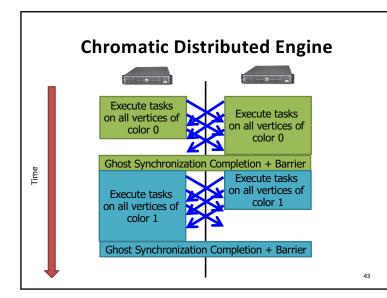


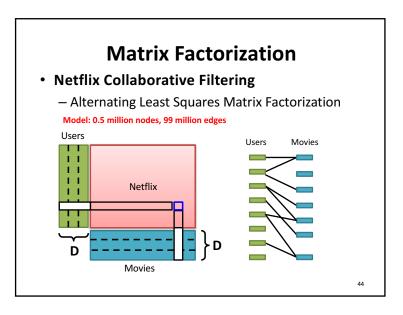


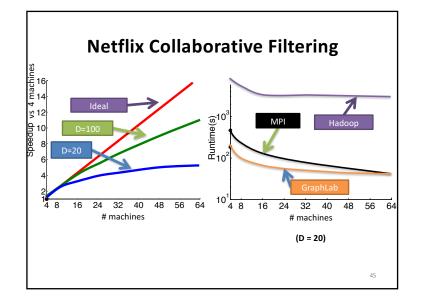


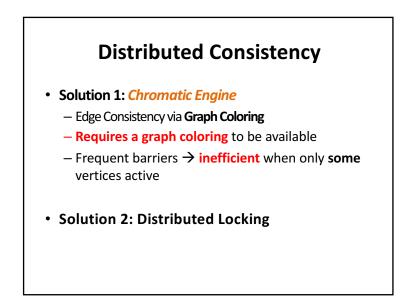


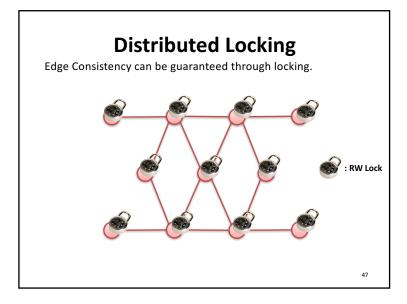


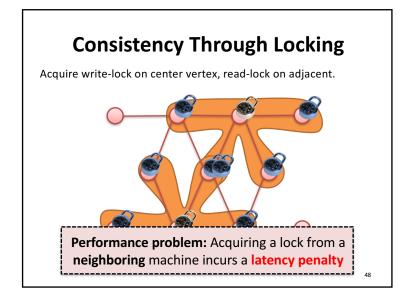


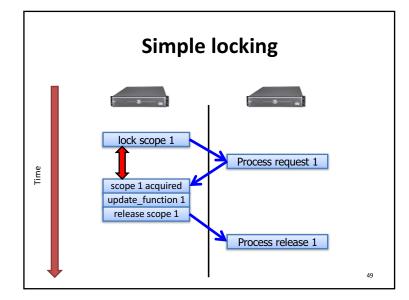






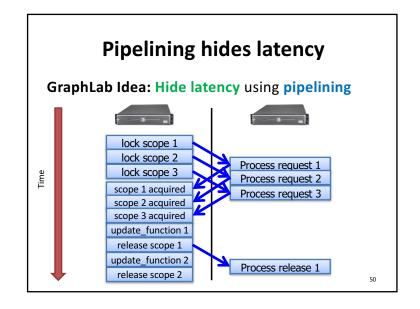






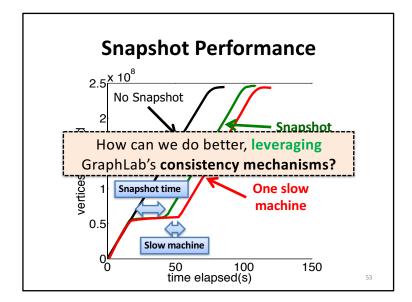
Distributed Consistency

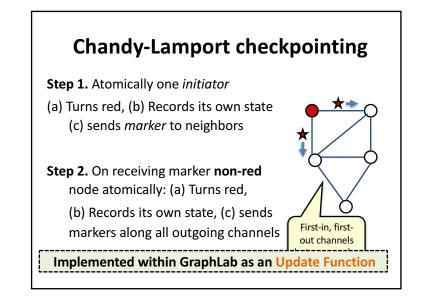
- Solution 1: Chromatic Engine
 - Edge Consistency via Graph Coloring
 - Requires a graph coloring to be available
 - − Frequent barriers → inefficient when only some vertices active
- Solution 2: Distributed Locking
 - Residual BP on 190K-vertex/560K-edge graph, 4 machines
 - No pipelining: 472 sec; with pipelining: 10 sec

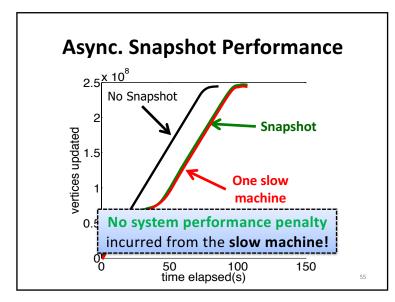


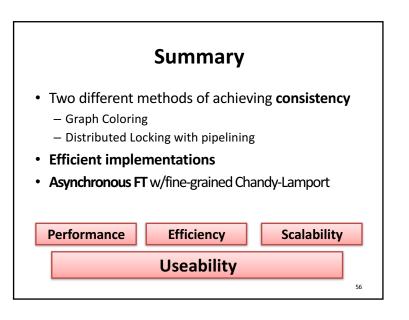
How to handle machine failure?

- What when machines fail? How do we provide fault tolerance?
- Strawman scheme: Synchronous snapshot checkpointing
 - 1. Stop the world
 - 2. Write each machines' state to disk









Friday Precept: Roofnet performance More Graph Processing

Monday topic: Streaming Data Processing