Algorithms

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2.2 MERGESORT

- *mergesort*
- *bottom-up mergesort*
- *sorting complexity*
- *divide-and-conquer*

Two classic sorting algorithms: mergesort and quicksort

Critical components in the world's computational infrastructure.

- Full scientific understanding of their properties has enabled us to develop them into practical system sorts.
- Quicksort honored as one of top 10 algorithms of 20th century in science and engineering.

Mergesort. [this lecture]

JS ...

Quicksort. [next lecture]

JS ...

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2.2 MERGESORT

- *mergesort*
- *bottom-up mergesort*
- *sorting complexity*
- *divide-and-conquer*

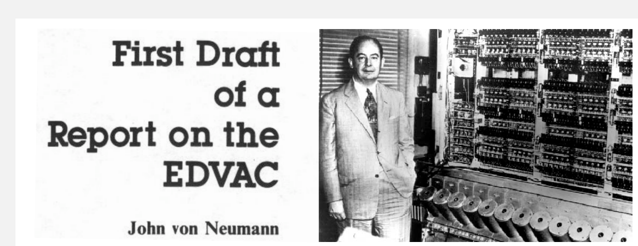
Mergesort

Basic plan.

- Divide array into two halves.
- Recursively sort each half.
- Merge two halves.

input	M	E	R	G	E	S	O	R	T	E	X	A	M	P	L	E
sort left half	E	E	G	M	O	R	R	S	T	E	X	A	M	P	L	E
sort right half	E	E	G	M	O	R	R	S	A	E	E	L	M	P	T	X
merge results	A	E	E	E	E	G	L	M	M	O	P	R	R	S	T	X

Mergesort overview



Abstract in-place merge demo

Goal. Given two sorted subarrays `a[lo]` to `a[mid]` and `a[mid+1]` to `a[hi]`, replace with sorted subarray `a[lo]` to `a[hi]`.

	lo				mid	mid+1				hi
a[]	E	E	G	M	R	A	C	E	R	T

sorted sorted

Abstract in-place merge demo

Goal. Given two sorted subarrays `a[lo]` to `a[mid]` and `a[mid+1]` to `a[hi]`, replace with sorted subarray `a[lo]` to `a[hi]`.

	lo									hi
a[]	A	C	E	E	E	G	M	R	R	T

sorted

Merging: Java implementation

```
private static void merge(Comparable[] a, Comparable[] aux, int lo, int mid, int hi)
{
   for (int k = lo; k <= hi; k++)                                   copy
      aux[k] = a[k];

   int i = lo, j = mid+1;
   for (int k = lo; k <= hi; k++)
   {
      if      (i > mid)              a[k] = aux[j++];               merge
      else if (j > hi)               a[k] = aux[i++];
      else if (less(aux[j], aux[i])) a[k] = aux[j++];
      else                           a[k] = aux[i++];
   }
}
```

	lo			i	mid			j		hi
aux[]	A	G	L	O	R	H	I	M	S	T

						k				
a[]	A	G	H	I	L	M				

Mergesort quiz

How many calls to `less()` does `merge()` make in the worst case to merge two subarrays of length $N/2$ into a single array of length N.

A. $N/2$

B. $N/2+1$

C. $N-1$

D. N

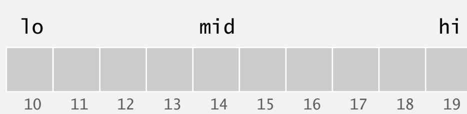
E. *I don't know.*

Mergesort: Java implementation

```
public class Merge
{
   private static void merge(...)
   {  /* as before */  }

   private static void sort(Comparable[] a, Comparable[] aux, int lo, int hi)
   {
      if (hi <= lo) return;
      int mid = lo + (hi - lo) / 2;
      sort(a, aux, lo, mid);
      sort(a, aux, mid+1, hi);
      merge(a, aux, lo, mid, hi);
   }

   public static void sort(Comparable[] a)
   {
      Comparable[] aux = new Comparable[a.length];
      sort(a, aux, 0, a.length - 1);
   }
}
```



Mergesort: trace

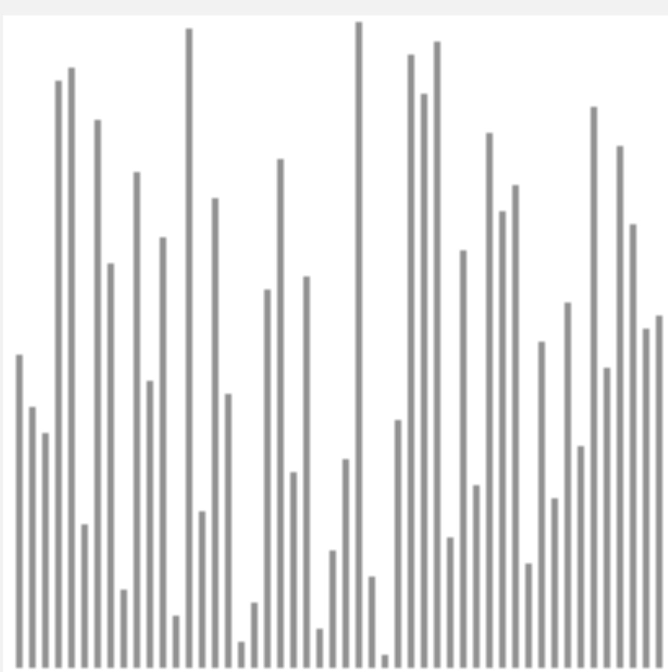
a[]

	lo		hi	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
				M	E	R	G	E	S	O	R	T	E	X	A	M	P	L	E
merge(a, aux,	0,	0,	1)	E	M	R	G	E	S	O	R	T	E	X	A	M	P	L	E
merge(a, aux,	2,	2,	3)	E	M	G	R	E	S	O	R	T	E	X	A	M	P	L	E
merge(a, aux,	0,	1,	3)	E	G	M	R	E	S	O	R	T	E	X	A	M	P	L	E
merge(a, aux,	4,	4,	5)	E	G	M	R	E	S	O	R	T	E	X	A	M	P	L	E
merge(a, aux,	6,	6,	7)	E	G	M	R	E	S	O	R	T	E	X	A	M	P	L	E
merge(a, aux,	4,	5,	7)	E	G	M	R	E	O	R	S	T	E	X	A	M	P	L	E
merge(a, aux,	0,	3,	7)	E	E	G	M	O	R	R	S	T	E	X	A	M	P	L	E
merge(a, aux,	8,	8,	9)	E	E	G	M	O	R	R	S	E	T	X	A	M	P	L	E
merge(a, aux,	10,	10,	11)	E	E	G	M	O	R	R	S	E	T	A	X	M	P	L	E
merge(a, aux,	8,	9,	11)	E	E	G	M	O	R	R	S	A	E	T	X	M	P	L	E
merge(a, aux,	12,	12,	13)	E	E	G	M	O	R	R	S	A	E	T	X	M	P	L	E
merge(a, aux,	14,	14,	15)	E	E	G	M	O	R	R	S	A	E	T	X	M	P	E	L
merge(a, aux,	12,	13,	15)	E	E	G	M	O	R	R	S	A	E	T	X	E	L	M	P
merge(a, aux,	8,	11,	15)	E	E	G	M	O	R	R	S	A	E	E	L	M	P	T	X
merge(a, aux,	0,	7,	15)	A	E	E	E	E	G	L	M	M	O	P	R	R	S	T	X

result after recursive call

Mergesort: animation

50 random items

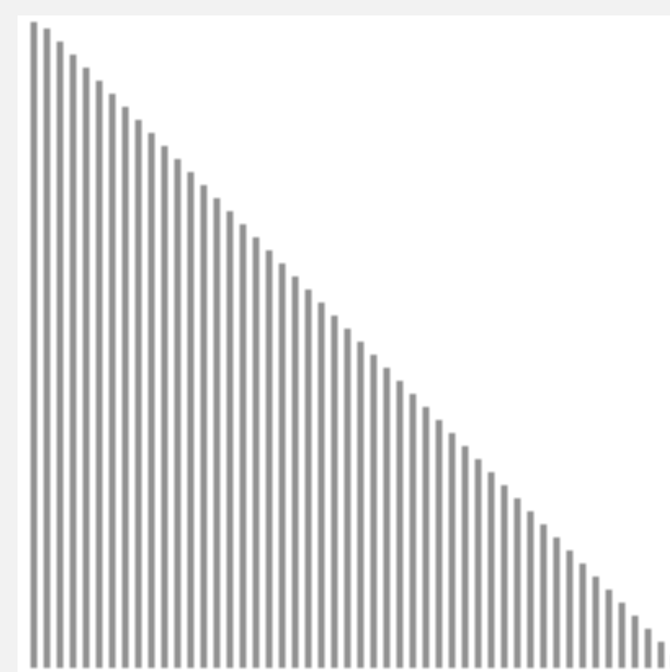


▲ algorithm position
in order
current subarray
not in order

http://www.sorting-algorithms.com/merge-sort

Mergesort: animation

50 reverse-sorted items



▲ algorithm position
in order
current subarray
not in order

http://www.sorting-algorithms.com/merge-sort

Mergesort: empirical analysis

Running time estimates:

- Laptop executes 10^8 compares/second.
- Supercomputer executes 10^{12} compares/second.

	insertion sort (N^2)			mergesort (N log N)		
computer	thousand	million	billion	thousand	million	billion
home	instant	2.8 hours	317 years	instant	1 second	18 min
super	instant	1 second	1 week	instant	instant	instant

Bottom line. Good algorithms are better than supercomputers.

Mergesort: number of compares

Proposition. Mergesort uses $\leq N \lg N$ compares to sort an array of length N.

Pf sketch. The number of compares $C(N)$ to mergesort an array of length N satisfies the recurrence:

$$C(N) \leq \underbrace{C(\lceil N/2 \rceil)}_{\text{left half}} + \underbrace{C(\lfloor N/2 \rfloor)}_{\text{right half}} + \underbrace{N-1}_{\text{merge}} \quad \text{for } N > 1, \text{ with } C(1) = 0.$$

We solve this simpler recurrence, and assume N is a power of 2:

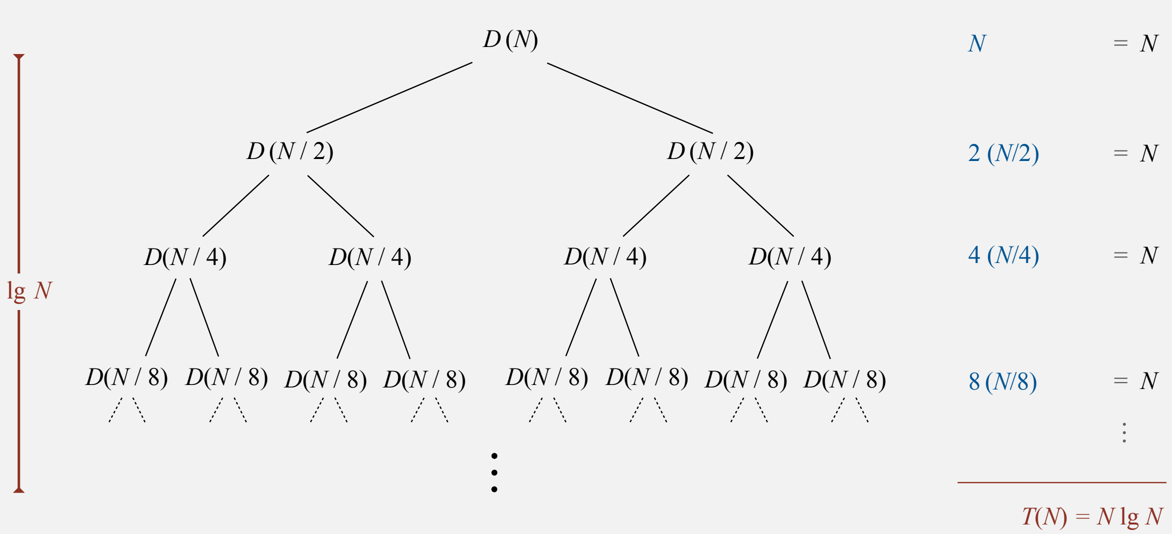
result holds for all N (analysis cleaner in this case)

$$D(N) = 2\,D(N/2) + N, \text{ for } N > 1, \text{ with } D(1) = 0.$$

Divide-and-conquer recurrence

Proposition. If $D(N)$ satisfies $D(N) = 2\,D(N/2) + N$ for $N > 1$, with $D(1) = 0$, then $D(N) = N \lg N$.

Pf by picture. [assuming N is a power of 2]



Mergesort: number of array accesses

Proposition. Mergesort uses $\leq 6 N \lg N$ array accesses to sort an array of length N.

Pf sketch. The number of array accesses $A(N)$ satisfies the recurrence:

$$A(N) \leq A(\lceil N/2 \rceil) + A(\lfloor N/2 \rfloor) + 6N \quad \text{for } N > 1, \text{ with } A(1) = 0.$$

Key point. Any algorithm with the following structure takes $N \log N$ time:

```
public static void f(int N)
{
    if (N == 0) return;
    f(N/2);          // solve two problems
    f(N/2);          // of half the size
    linear(N);       // do a linear amount of work
}
```

Notable examples. FFT, hidden-line removal, Kendall-tau distance, …

Mergesort analysis: memory

Proposition. Mergesort uses extra space proportional to N.
Pf. The array `aux[]` needs to be of length N for the last merge.

two sorted subarrays

```
A C D G H I M N U V   B E F J O P Q R S T

A B C D E F G H I J   M N O P Q R S T U V
```

merged result

Def. A sorting algorithm is in-place if it uses $\leq c \log N$ extra memory.
Ex. Insertion sort, selection sort, shellsort.

Challenge 1 (not hard). Use `aux[]` array of length $\sim \frac{1}{2} N$ instead of N.
Challenge 2 (very hard). In-place merge. [Kronrod 1969]

Mergesort quiz 2

Is our implementation of mergesort stable?

A. Yes.

B. No, but it can be modified to be stable.

C. No, mergesort is inherently unstable.

D. *I don't remember what stability means.*

E. *I don't know.*

a sorting algorithm is stable if it preserves the relative order of equal keys

input	C	A_1	B	A_2	A_3
sorted	A_3	A_1	A_2	B	C

not stable

Stability: mergesort

Proposition. Mergesort is stable.

```
public class Merge
{
   private static void merge(...)
   {  /* as before */  }

   private static void sort(Comparable[] a, Comparable[] aux, int lo, int hi)
   {
      if (hi <= lo) return;
      int mid = lo + (hi - lo) / 2;
      sort(a, aux, lo, mid);
      sort(a, aux, mid+1, hi);
      merge(a, aux, lo, mid, hi);
   }

   public static void sort(Comparable[] a)
   {  /* as before */  }
}
```

Pf. Suffices to verify that merge operation is stable.

Stability: mergesort

Proposition. Merge operation is stable.

```
private static void merge(...)
{
   for (int k = lo; k <= hi; k++)
      aux[k] = a[k];

   int i = lo, j = mid+1;
   for (int k = lo; k <= hi; k++)
   {
      if      (i > mid)              a[k] = aux[j++];
      else if (j > hi)               a[k] = aux[i++];
      else if (less(aux[j], aux[i])) a[k] = aux[j++];
      else                           a[k] = aux[i++];
   }
}
```

0	1	2	3	4	5	6	7	8	9	10
A_1	A_2	A_3	B	D	A_4	A_5	C	E	F	G

Pf. Takes from left subarray if equal keys.

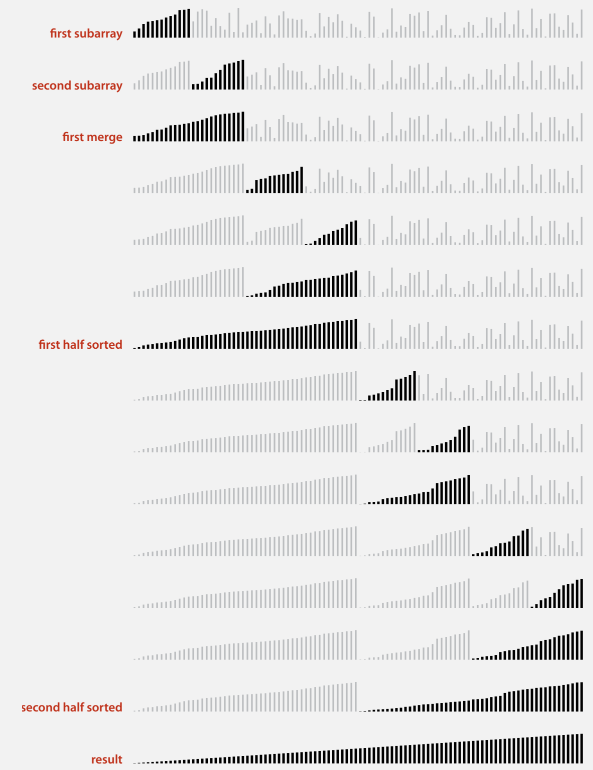
Mergesort: practical improvements

Use insertion sort for small subarrays.

- Mergesort has too much overhead for tiny subarrays.
- Cutoff to insertion sort for ≈ 10 items.

```
private static void sort(Comparable[] a, Comparable[] aux, int lo, int hi)
{
   if (hi <= lo + CUTOFF - 1)
   {
      Insertion.sort(a, lo, hi);
      return;
   }
   int mid = lo + (hi - lo) / 2;
   sort (a, aux, lo, mid);
   sort (a, aux, mid+1, hi);
   merge(a, aux, lo, mid, hi);
}
```


Mergesort with cutoff to insertion sort: visualization



Mergesort: practical improvements

Stop if already sorted.

- Is largest item in first half ≤ smallest item in second half?
- Helps for partially-ordered arrays.

A B C D E F G H I J M N O P Q R S T U V

A B C D E F G H I J M N O P Q R S T U V

```
private static void sort(Comparable[] a, Comparable[] aux, int lo, int hi)
{
   if (hi <= lo) return;
   int mid = lo + (hi - lo) / 2;
   sort (a, aux, lo, mid);
   sort (a, aux, mid+1, hi);
   if (!less(a[mid+1], a[mid])) return;
   merge(a, aux, lo, mid, hi);
}
```


Mergesort: practical improvements

Eliminate the copy to the auxiliary array. Save time (but not space) by switching the role of the input and auxiliary array in each recursive call.

```
private static void merge(Comparable[] a, Comparable[] aux, int lo, int mid, int hi)
{
   int i = lo, j = mid+1;
   for (int k = lo; k <= hi; k++)
   {
      if      (i > mid)           aux[k] = a[j++];
      else if (j > hi)            aux[k] = a[i++];
      else if (less(a[j], a[i])) aux[k] = a[j++];
      else                        aux[k] = a[i++];
   }
}

private static void sort(Comparable[] a, Comparable[] aux, int lo, int hi)
{
   if (hi <= lo) return;
   int mid = lo + (hi - lo) / 2;
   sort (aux, a, lo, mid);
   sort (aux, a, mid+1, hi);
   merge(a, aux, lo, mid, hi);
}
```

merge from a[] to aux[]

assumes aux[] is initialize to a[] once, before recursive calls

switch roles of aux[] and a[]

Java 6 system sort

Basic algorithm for sorting objects = mergesort.

- Cutoff to insertion sort = 7.
- Stop-if-already-sorted test.
- Eliminate-the-copy-to-the-auxiliary-array trick.

Arrays.sort(a)

http://hg.openjdk.java.net/jdk6/jdk6/jdk/file/tip/src/share/classes/java/util/Arrays.java

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2.2 Mergesort

- mergesort
- **bottom-up mergesort**
- sorting complexity
- divide-and-conquer

Bottom-up mergesort

Basic plan.

- Pass through array, merging subarrays of size 1.
- Repeat for subarrays of size 2, 4, 8, ...

	a[i] 0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
sz = 1	M	E	R	G	E	S	O	R	T	E	X	A	M	P	L	E
merge(a, aux, 0, 0, 1)	E	M	R	G	E	S	O	R	T	E	X	A	M	P	L	E
merge(a, aux, 2, 2, 3)	E	M	G	R	E	S	O	R	T	E	X	A	M	P	L	E
merge(a, aux, 4, 4, 5)	E	M	G	R	E	S	O	R	T	E	X	A	M	P	L	E
merge(a, aux, 6, 6, 7)	E	M	G	R	E	S	O	R	T	E	X	A	M	P	L	E
merge(a, aux, 8, 8, 9)	E	M	G	R	E	S	O	R	E	T	X	A	M	P	L	E
merge(a, aux, 10, 10, 11)	E	M	G	R	E	S	O	R	E	T	A	X	M	P	L	E
merge(a, aux, 12, 12, 13)	E	M	G	R	E	S	O	R	E	T	A	X	M	P	L	E
merge(a, aux, 14, 14, 15)	E	M	G	R	E	S	O	R	E	T	A	X	M	P	E	L
sz = 2																
merge(a, aux, 0, 1, 3)	E	G	M	R	E	S	O	R	E	T	A	X	M	P	E	L
merge(a, aux, 4, 5, 7)	E	G	M	R	E	O	R	S	E	T	A	X	M	P	E	L
merge(a, aux, 8, 9, 11)	E	G	M	R	E	O	R	S	A	E	T	X	M	P	E	L
merge(a, aux, 12, 13, 15)	E	G	M	R	E	O	R	S	A	E	T	X	E	L	M	P
sz = 4																
merge(a, aux, 0, 3, 7)	E	E	G	M	O	R	R	S	A	E	T	X	E	L	M	P
merge(a, aux, 8, 11, 15)	E	E	G	M	O	R	R	S	A	E	E	L	M	P	T	X
sz = 8																
merge(a, aux, 0, 7, 15)	A	E	E	E	E	G	L	M	M	O	P	R	R	S	T	X

Bottom-up mergesort: Java implementation

```
public class MergeBU
{
   private static void merge(...)
   {  /* as before */  }

   public static void sort(Comparable[] a)
   {
      int N = a.length;
      Comparable[] aux = new Comparable[N];
      for (int sz = 1; sz < N; sz = sz+sz)
         for (int lo = 0; lo < N-sz; lo += sz+sz)
            merge(a, aux, lo, lo+sz-1, Math.min(lo+sz+sz-1, N-1));
   }
}
```

but about 10% slower than recursive, top-down mergesort on typical systems

Bottom line. Simple and non-recursive version of mergesort.

Mergesort quiz 3

Which is faster in practice: top-down mergesort or bottom-up mergesort?

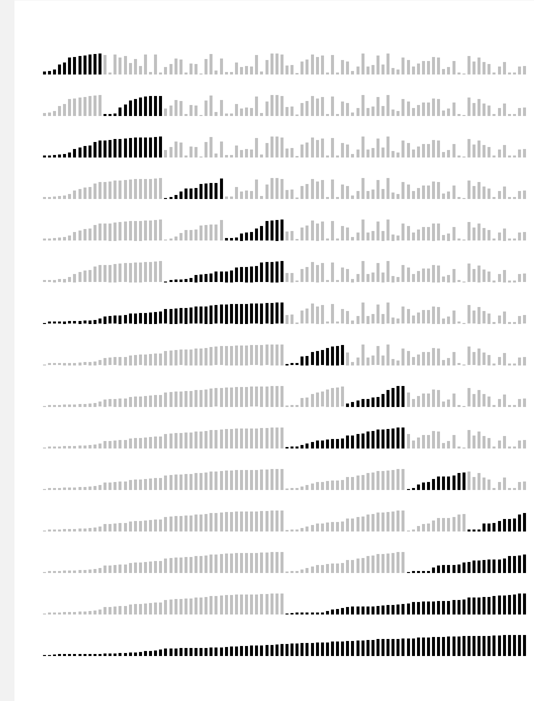
A. Top-down (recursive) mergesort.

B. Bottom-up (nonrecursive) mergesort.

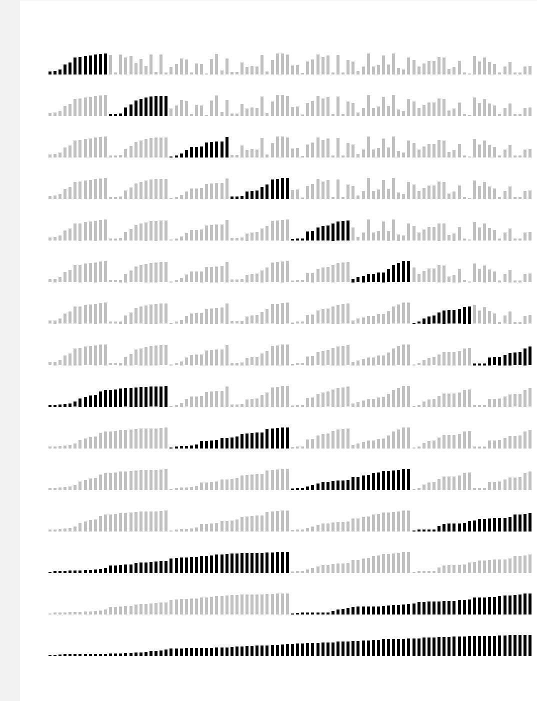
C. A tie.

D. *I don't know.*

Mergesort: visualizations



top-down mergesort (cutoff = 12)



bottom-up mergesort (cutoff = 12)

Natural mergesort

Idea. Exploit pre-existing order by identifying naturally-occurring runs.

input

1	5	10	16	3	4	23	9	13	2	7	8	12	14

first run

1	5	10	16	3	4	23	9	13	2	7	8	12	14

second run

1	5	10	16	3	4	23	9	13	2	7	8	12	14

merge two runs

1	3	4	5	10	16	23	9	13	2	7	8	12	14

Tradeoff. Fewer passes vs. extra compares per pass to identify runs.

Timsort

- Natural mergesort.
- Use binary insertion sort to make initial runs (if needed).
- A few more clever optimizations.

Tim Peters

```
Intro
-----
This describes an adaptive, stable, natural mergesort, modestly called
timsort (hey, I earned it <wink>).  It has supernatural performance on many
kinds of partially ordered arrays (less than lg(N!) comparisons needed, and
as few as N-1), yet as fast as Python's previous highly tuned samplesort
hybrid on random arrays.

In a nutshell, the main routine marches over the array once, left to right,
alternately identifying the next run, then merging it into the previous
runs "intelligently".  Everything else is complication for speed, and some
hard-won measure of memory efficiency.
...
```

Consequence. Linear time on many arrays with pre-existing order.

Now widely used. Python, Java 7, GNU Octave, Android, ….

http://hg.openjdk.java.net/jdk7/jdk7/jdk/file/tip/src/share/classes/java/util/Arrays.java

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2.2 MERGESORT

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- sorting complexity
- divide-and-conquer

Complexity of sorting

Computational complexity. Framework to study efficiency of algorithms for solving a particular problem X.

Model of computation. Allowable operations.
Cost model. Operation count(s).
Upper bound. Cost guarantee provided by some algorithm for X.
Lower bound. Proven limit on cost guarantee of all algorithms for X.
Optimal algorithm. Algorithm with best possible cost guarantee for X.

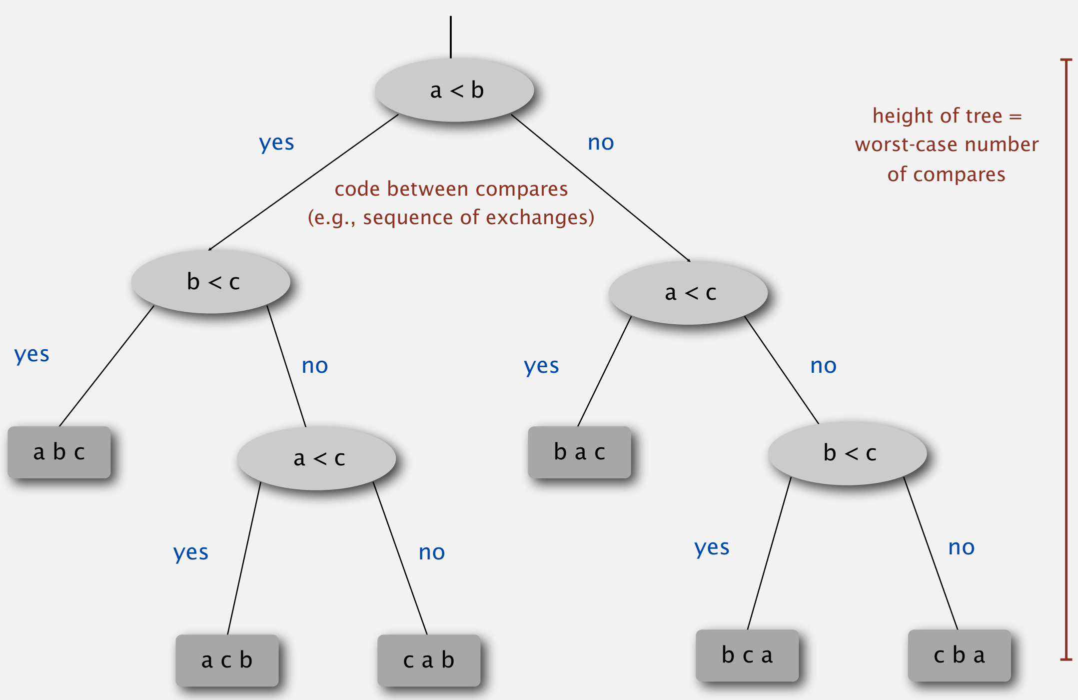
lower bound ~ upper bound

Example: sorting.

- Model of computation: decision tree.
- Cost model: # compares.
- Upper bound: $\sim N \lg N$ from mergesort.
- Lower bound:
- Optimal algorithm:

can access information only through compares (e.g., Java Comparable framework)

Decision tree (for 3 distinct keys a, b, and c)



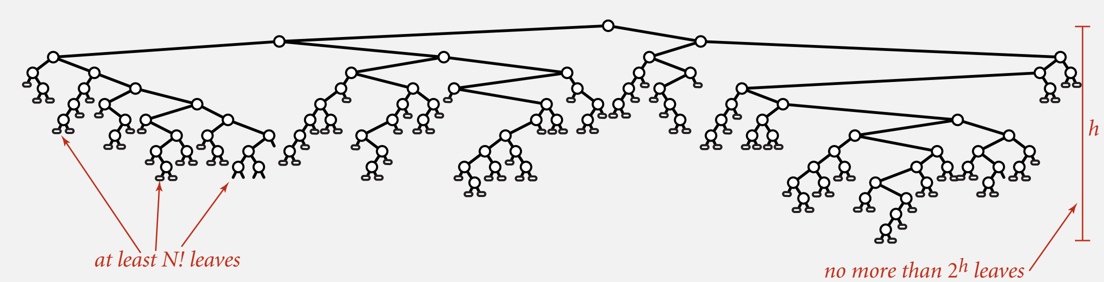
each leaf corresponds to one (and only one) ordering;
(at least) one leaf for each possible ordering

Compare-based lower bound for sorting

Proposition. Any compare-based sorting algorithm must use at least $\lg(N!) \sim N \lg N$ compares in the worst-case.

Pf.

- Assume array consists of N distinct values a_1 through a_N.
- Worst case dictated by height h of decision tree.
- Binary tree of height h has at most 2^h leaves.
- N ! different orderings $\Rightarrow$ at least N ! leaves.



Compare-based lower bound for sorting

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Pf.

- Assume array consists of N distinct values a_1 through a_N.
- Worst case dictated by height h of decision tree.
- Binary tree of height h has at most 2^h leaves.
- $N!$ different orderings $\Rightarrow$ at least $N!$ leaves.

$$2^h \geq \#\text{ leaves} \geq N!$$
$$\Rightarrow \; h \geq \lg(N!) \sim N \lg N$$

Stirling's formula

Complexity of sorting

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Upper bound. Cost guarantee provided by some algorithm for X.
Lower bound. Proven limit on cost guarantee of all algorithms for X.
Optimal algorithm. Algorithm with best possible cost guarantee for X.

Example: sorting.

- Model of computation: decision tree.
- Cost model: # compares.
- Upper bound: $\sim N \lg N$ from mergesort.
- Lower bound: $\sim N \lg N$.
- Optimal algorithm = mergesort.

First goal of algorithm design: optimal algorithms.

Complexity results in context

Compares? Mergesort is optimal with respect to number compares.
Space? Mergesort is not optimal with respect to space usage.

Lessons. Use theory as a guide.
Ex. Design sorting algorithm that guarantees $\frac{1}{2} N \lg N$ compares?
Ex. Design sorting algorithm that is both time- and space-optimal?

Complexity results in context (continued)

Lower bound may not hold if the algorithm can take advantage of:

- The initial order of the input.
  Ex: insert sort requires only a linear number of compares on partially-sorted arrays.
- The distribution of key values.
  Ex: 3-way quicksort requires only a linear number of compares on arrays with a constant number of distinct keys. [stay tuned]
- The representation of the keys.
  Ex: radix sort requires no key compares — it accesses the data via character/digit compares.

Sorting summary

	inplace?	stable?	best	average	worst	remarks
selection	✔		$\frac{1}{2} N^2$	$\frac{1}{2} N^2$	$\frac{1}{2} N^2$	N exchanges
insertion	✔	✔	N	$\frac{1}{4} N^2$	$\frac{1}{2} N^2$	use for small N or partially ordered
shell	✔		$N \log_3 N$	?	$c N^{3/2}$	tight code; subquadratic
merge		✔	$\frac{1}{2} N \lg N$	$N \lg N$	$N \lg N$	$N \log N$ guarantee; stable
timsort		✔	N	$N \lg N$	$N \lg N$	improves mergesort when preexisting order
?	✔	✔	N	$N \lg N$	$N \lg N$	holy sorting grail

http://algs4.cs.princeton.edu

2.2 Mergesort

- mergesort
- bottom-up mergesort
- sorting complexity
- divide-and-conquer

Interview question: shuffle a linked list

Problem. Given a singly-linked list, rearrange its nodes uniformly at random.

Assumption. Access to a perfect random-number generator.

all N! permutations equally likely

Version 1. Linear time, linear extra space.

Version 2. Linearithmic time, logarithmic or constant extra space.

before: first → 2♣ → 3♣ → 4♣ → 5♣ → 6♣ → 7♣ → *null*

after: first → 5♣ → 6♣ → 2♣ → 7♣ → 3♣ → 4♣ → *null*

Interview question: counting inversions

Problem. Given a permutation of length N, count the number of inversions.

Version 1. N^2 time.

Version 2. $N \log N$ time.

0	2	3	1	4	5	7	6

3 inversions: 2-1, 3-1, 7-6