StreamIt

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StreamIt

Overview

- Motivation

- Stream Programming

- StreamIt Language

- StreamIt Parallelization

- Conclusions

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Motivation

- As argued by "StreamIt Cookbook":

- Why von Neumann languages (C/C++/Java) successful?

    - Abstract out differences of von Neumann machines

    - Efficient mapping to von Neumann machine

- "Today von Neumann languages are a curse!"

    - Efficient mapping to parallel architectures difficult

    - Force programmer to take target architecture into account

    - Force programmer to *explicitly* parallelize: deal with threads, communication and synchronization
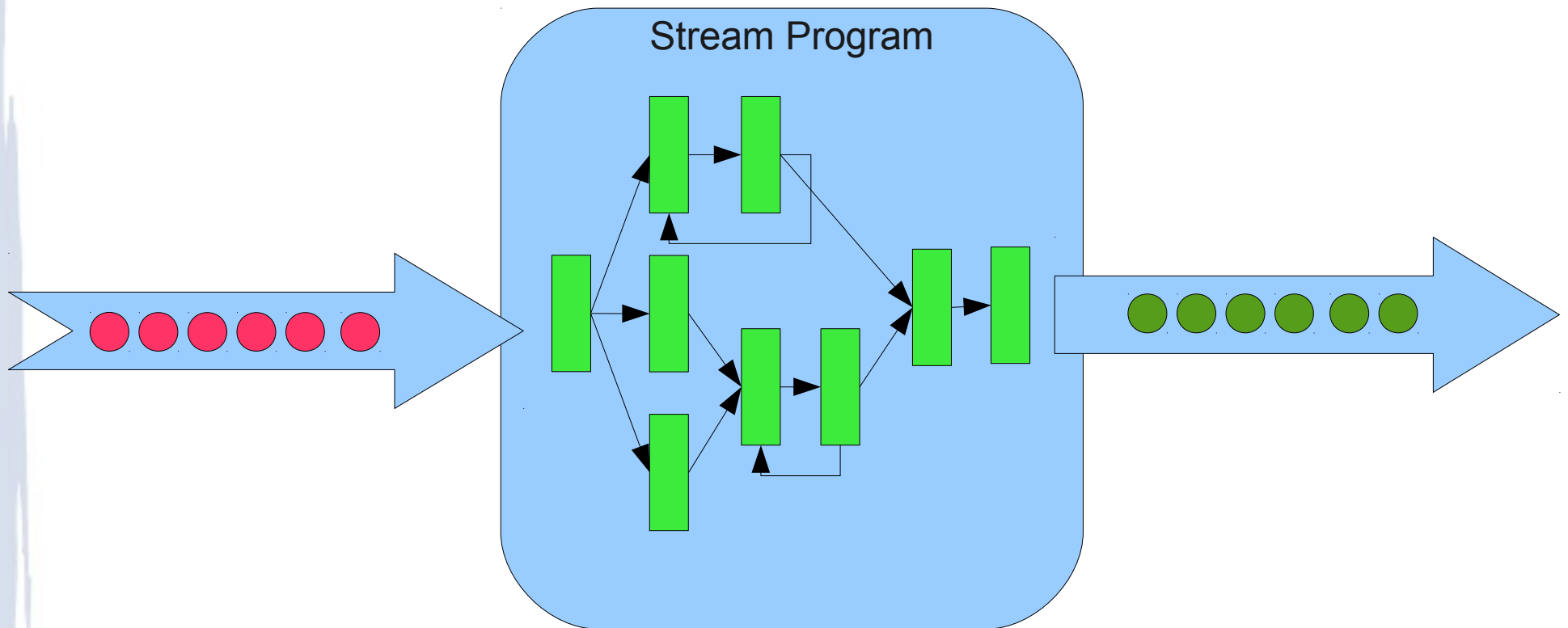
Motivation

- Want *implicitly* parallel abstraction that
  - Abstracts out differences in parallel architectures (number of cores, communication methods, synchronization methods, etc.)
  - Allows efficient mapping to parallel architecture
  - Directly exposes tasks that can run in parallel
- Stream Programming Abstraction
  - Trade off generality for performance and ease of programming
  - Many applications naturally fit paradigm
  - Implicitly parallel
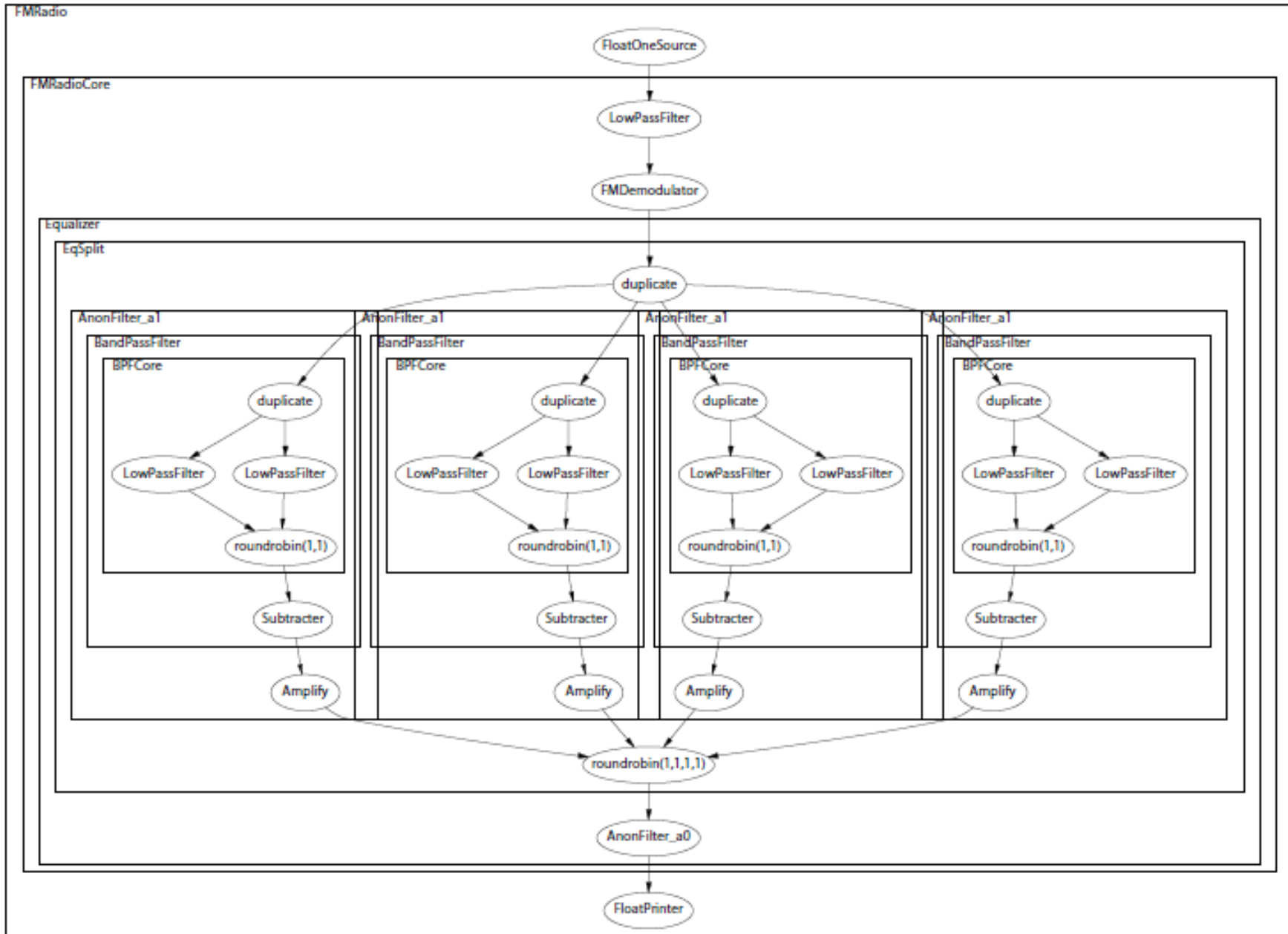 - Allows efficient mapping by compiler

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Stream Programming

Stream Programming: Example

Stream Programming: Applications

- Mobile
    - Compression (LZW)
    - Encryption (DES)

- Desktop
    - Streaming audio / video (MPEG-2)
    - Graphics (Depth of Field)

- Servers
    - Software routers
    - Modulation / Demodulation (Cell phone)

Stream Programming: Properties

- Large/Unbounded amount of data
    - Short lifetime per data item
    - Minimal processing per data item

- Regular, repeating computation
    - Static structured graph of filters
    - Independent actors
    - Explicit communication

Stream Programming: Properties

- Filter is autonomous unit of computation

- Each filter
    - Own PC
    - Own Address space

- Filters
    - Unaware of execution order
    - Communicate explicitly

- Stream program consists of a static structured graph of filters

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StreamIt Language

- Program consists of filters connected by structured graph

- Filter

    - Autonomous unit of computation

    - Single input, single output

    - Stateful / Stateless

    - Composable

- Pipeline

    - Filters connected by FIFO queues

- Structured Stream

    - Pipeline
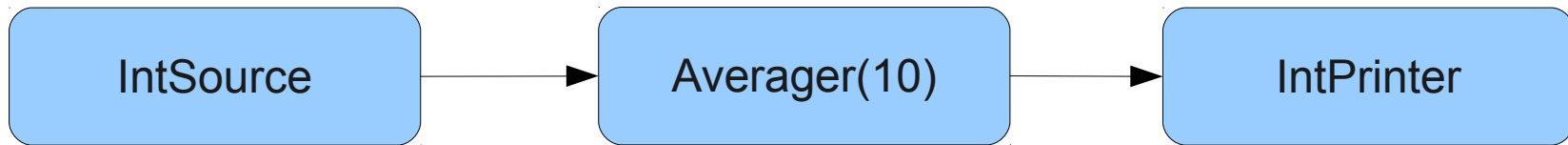
    - Splitjoin and Feedback Loop

StreamIt Language: Filter


```
int->int filter Averager(int n) {
  work pop 1 push 1 peek n {
    int sum = 0;
    for (int i = 0; i < n; i++)
      sum += peek(i);
    push(sum/n);
    pop();
  }
}
```

StreamIt Language: Filter

- Declaration

  - Defines input / output data type (int, float, bit, complex, struct)

  - Parameters for filter instantiation (parameters constant in init/work)

- Init Function

  - Called once to set up state

- Work Function

  - Called forever

  - Defines number of items popped from input stream, pushed to output stream, peeked at from input stream

  - Push/pop rates do not have to match

StreamIt Language: Pipeline


```
void->void pipeline MovingAverage {
    add IntSource();
    add Averager(10);
    add IntPrinter();
}
```

StreamIt Language: SplitJoin


```
float->float splitjoin
dualAverager(int n, int m) {
  split duplicate;
  add Averager(n);
  add Averager(m);
  join roundrobin;
}
```

```
float->int filter Trender{
  work pop 2 push 1 {
    float a = pop();
    float b = pop();
    if (a > b) { push(1); }
    else { push(0); }
  }
}
```

StreamIt Language: SplitJoin

- Split: divide stream into multiple streams
  - Duplicate
  - Round robin
- Join: combine streams into single stream
  - Round robin
- Can specify flow rate from each input/output filter

StreamIt Language: FeedbackLoop

```
int->int feedbackloop
EdgeDetecter {
    join roundrobin(1, 1);
    body int->int filter {
        work pop 2 push 2 peek 2 {
            push(peek(0)-peek(1));
            push(peek(0));
            pop();
            pop();
        }
    }
    loop Identity<int>;
    split roundrobin(1, 1);
    enqueue(0);
}
```

StreamIt Language: FeedbackLoop

- Join: combine input and loop feedback

- Body: filter for forward operation

- Loop: filter for reverse operation

- Split: split forward operation output into output and loop feedback input

- Enqueue: initial values on joiner feedback path

- Deadlock if loop filter output becomes empty

- Prequeue values with enqueue to ensure joiner always has input at startup

- Immediate data dependency with slow body can slow execution of feedback block

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StreamIt Parallelization

- Structured Graph Exposes Parallelism
    - Task parallelism (example Threads)
    - Data parallelism (example Do All)
    - Pipeline parallelism (example ILP)

- Target architectures vary
    - Granularity
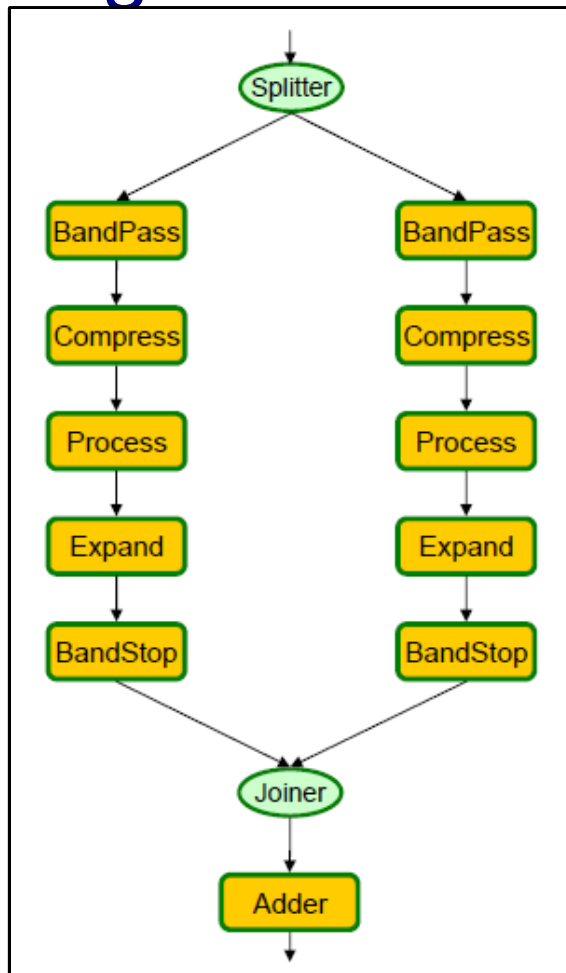   - Topology
    - Communication
    - Memory

Image Source:http://groups.csail.mit.edu/cag/streamit/talks/pact03tutorial/index.html

StreamIt Parallelization: Compiler

- Partitioning/Placement
    - Coarsen Granularity
    - Data Parallelize
    - Software Pipeline
- Scheduling
- Code Generation (output C/Java code)

StreamIt Parallelization: Compiler

- Naïve Partitioning / Scheduling

- Fine-grained data parallelism

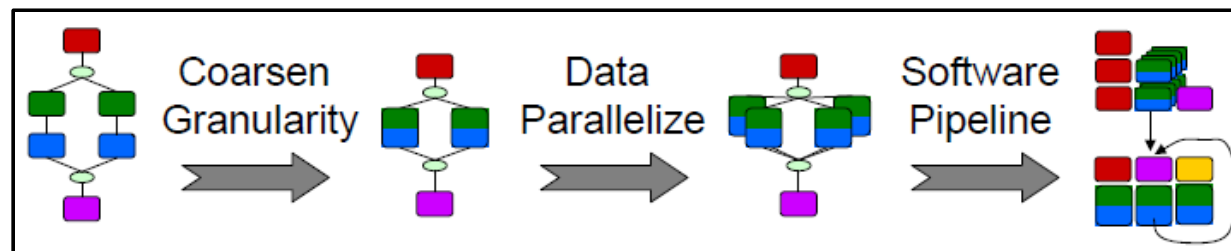# StreamIt Parallelization: Compiler

- Large synchronization overhead

Image Source:http://groups.csail.mit.edu/cag/streamit/talks/pact03tutorial/index.html

StreamIt Parallelization: Compiler

- ## Ideal Partitioning

  - Each filter has dedicated tile
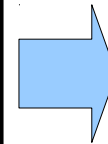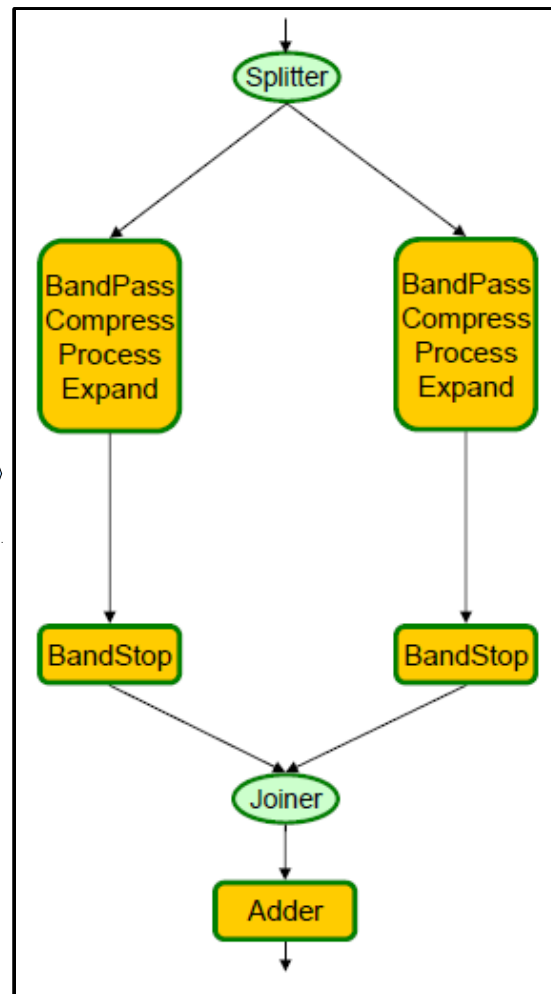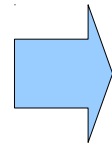
  - Each filter performs same amount of work

- ## Compiler Algorithm

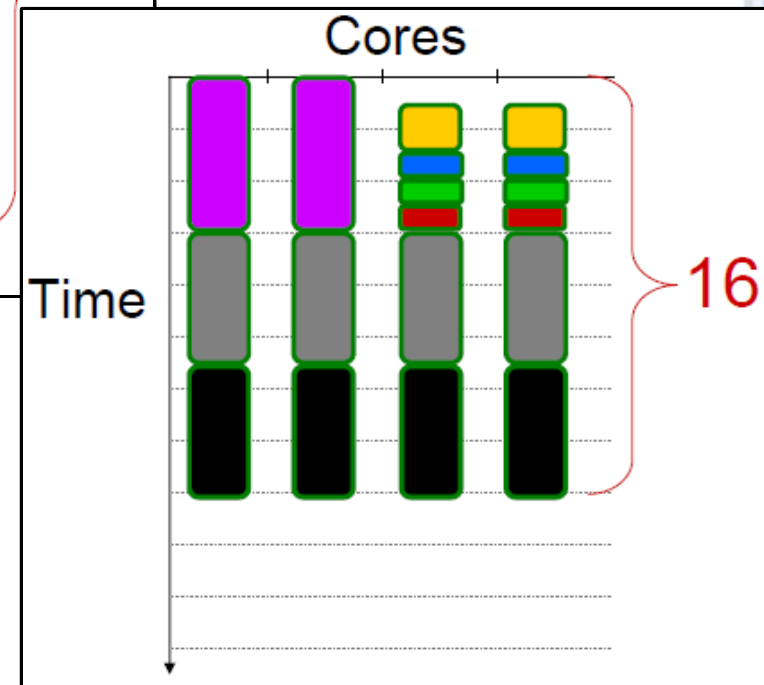Image Source:http://groups.csail.mit.edu/cag/streamit/talks/pact03tutorial/index.html
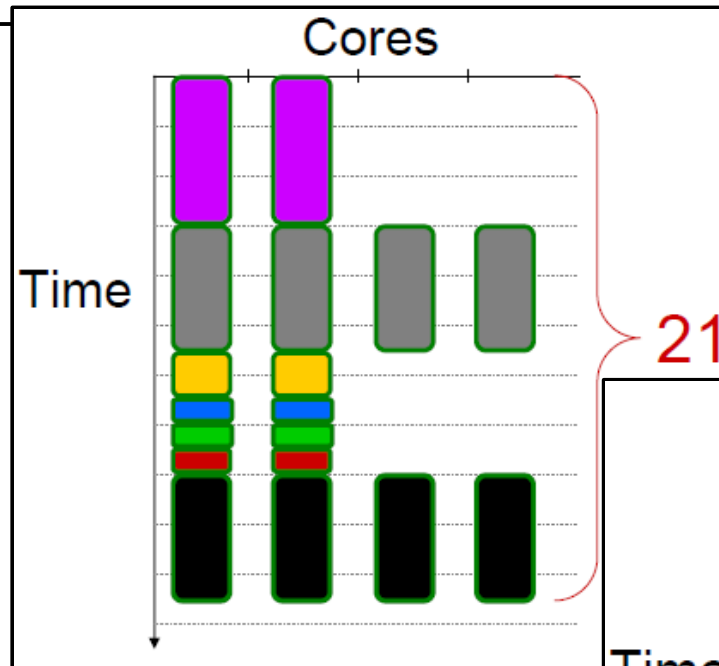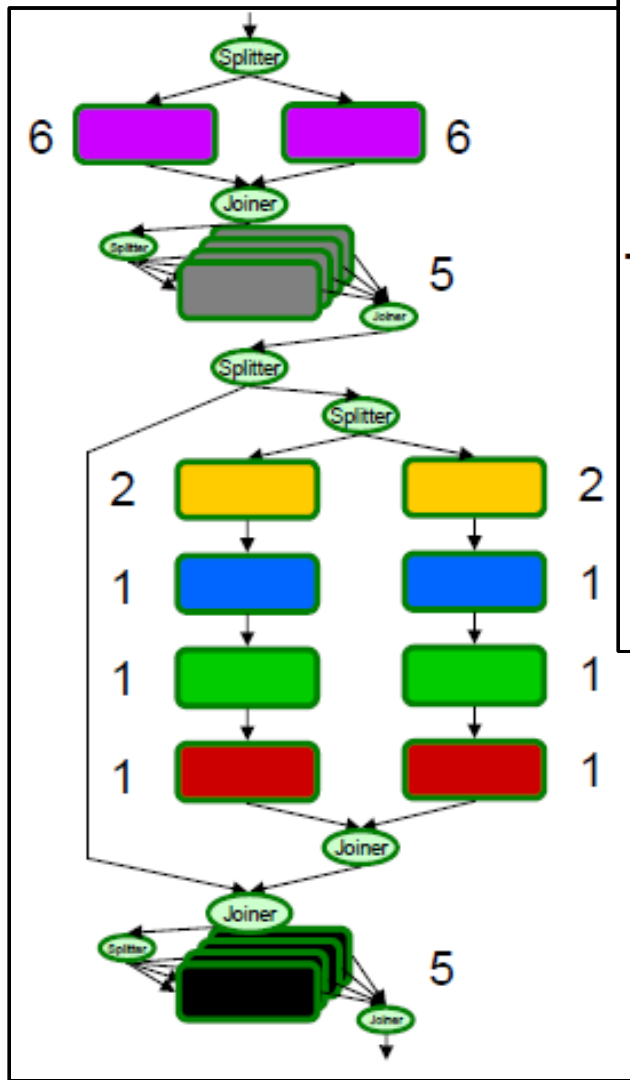
StreamIt Parallelization: Compiler
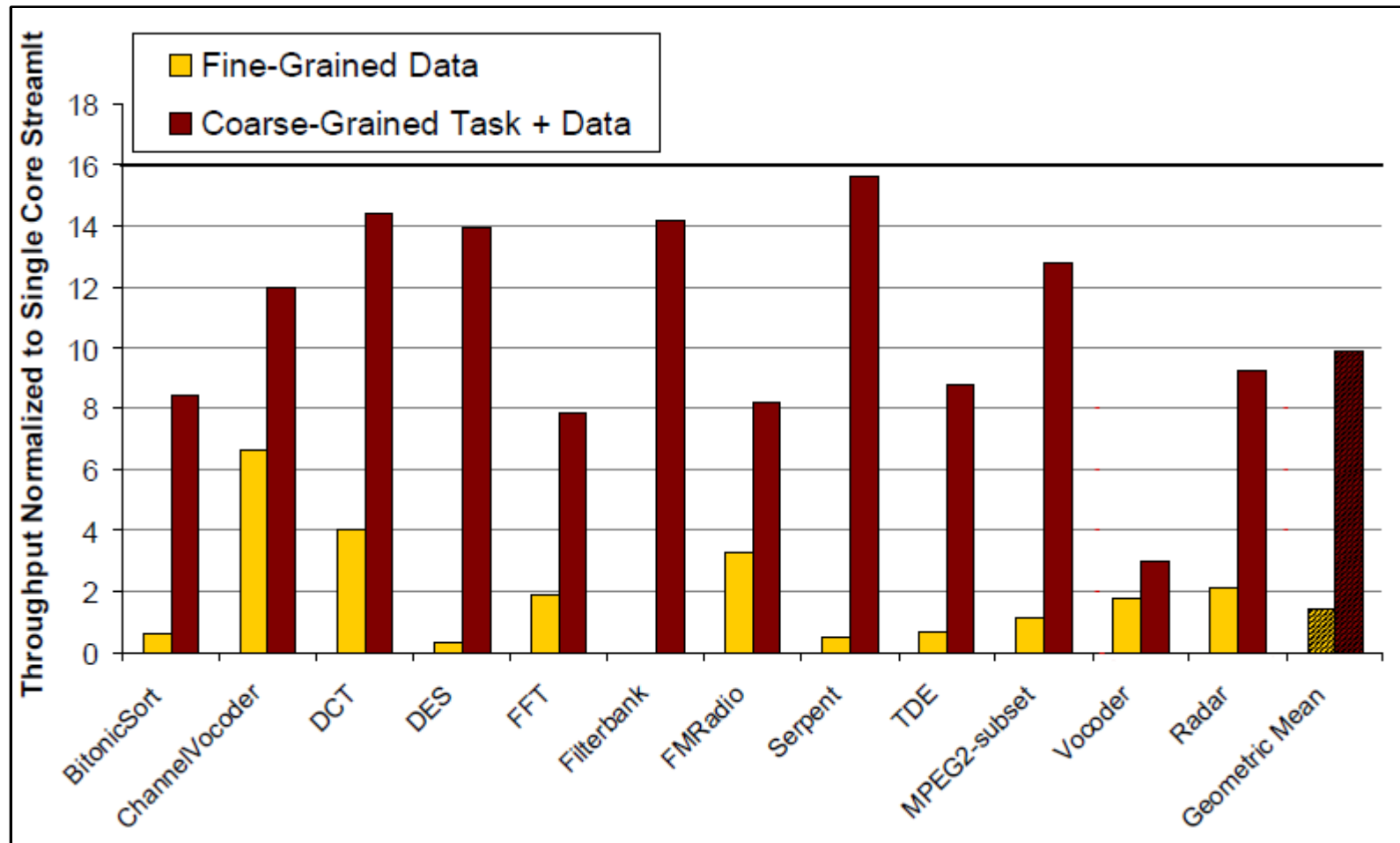
- Coarsen Granularity, Data Parallelization

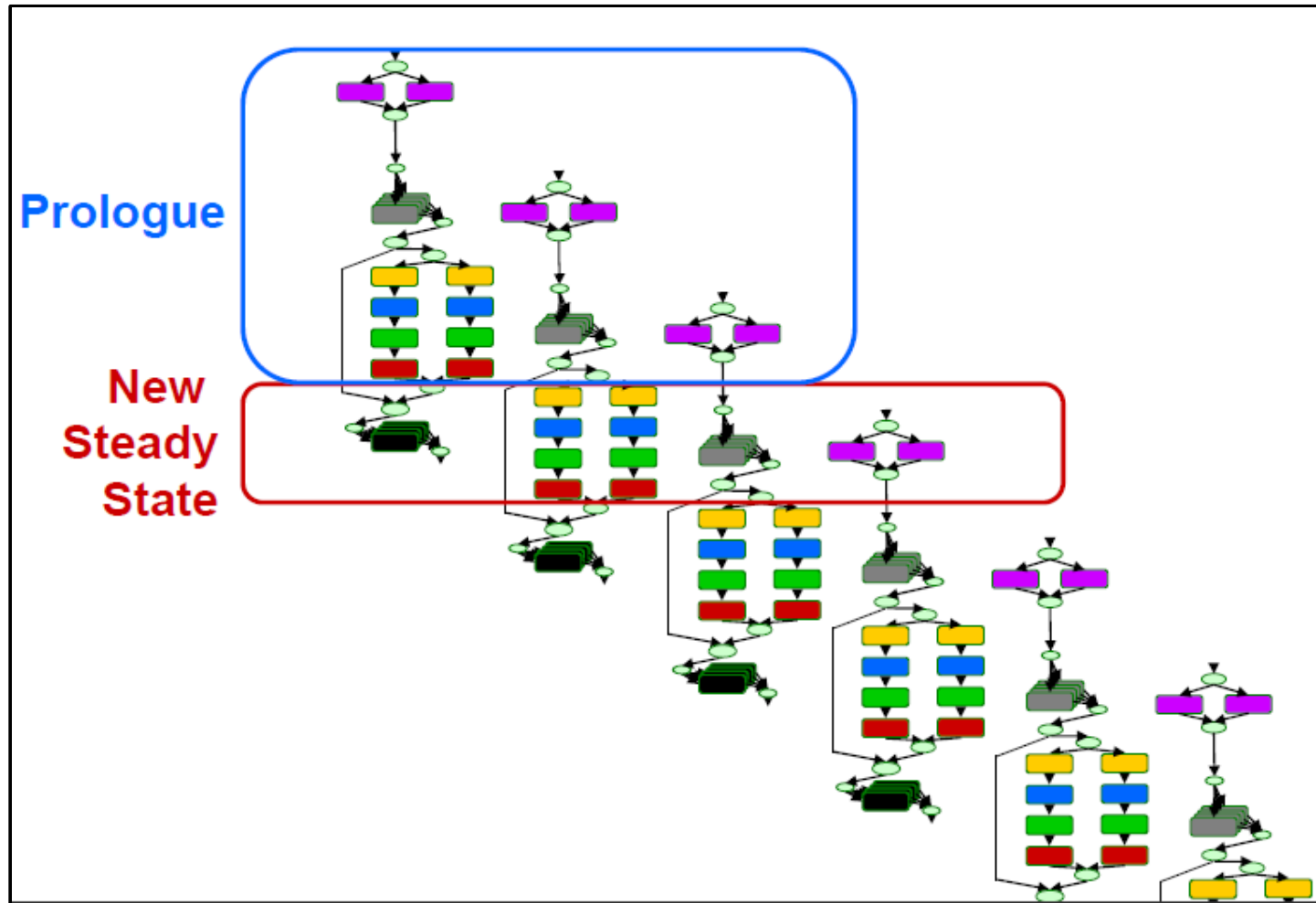# StreamIt Parallelization: Compiler

- Task Parallelization

StreamIt Parallelization: Compiler

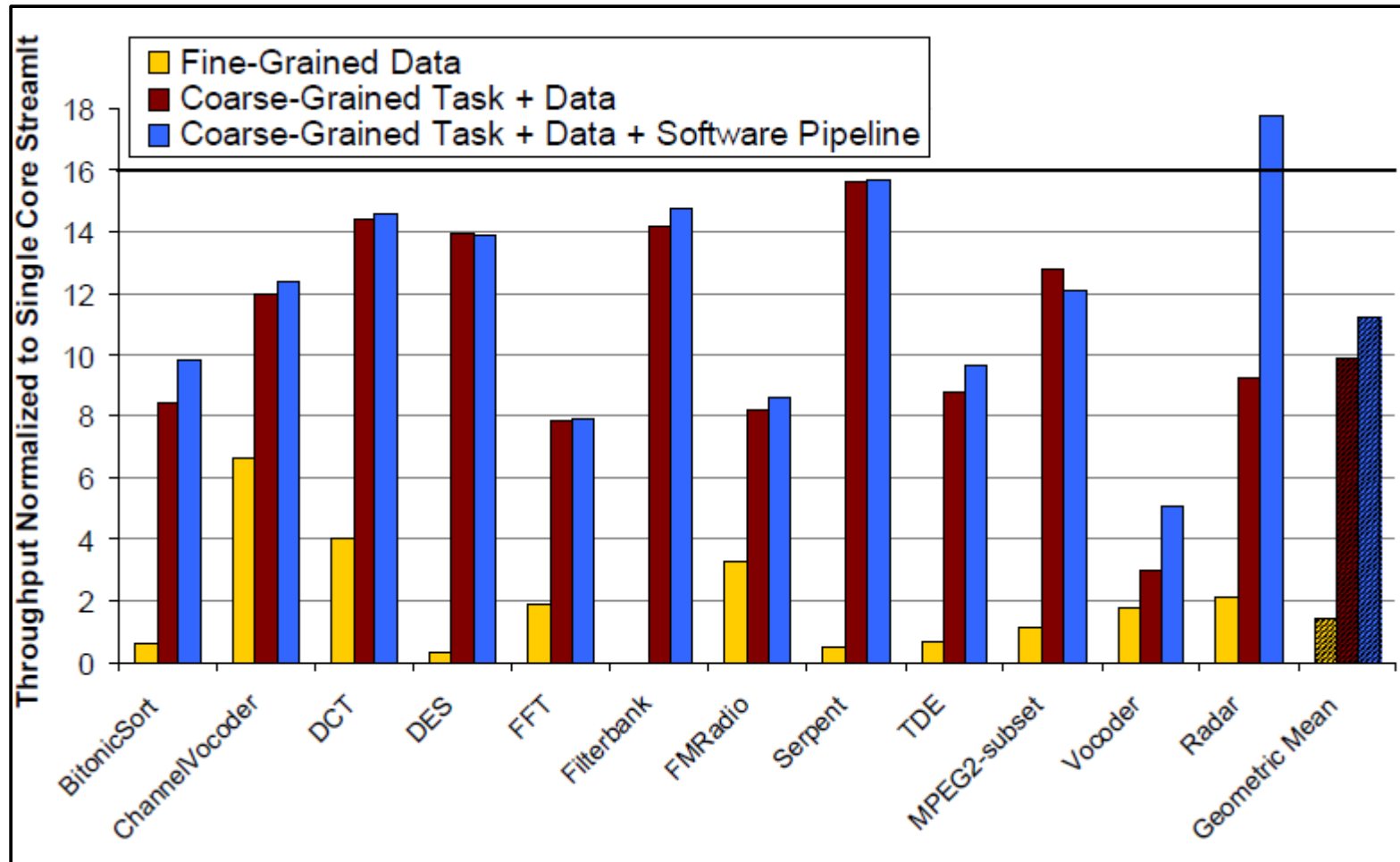- Lower synchronization overhead

StreamIt Parallelization: Compiler

- Software Pipeline

StreamIt Parallelization: Compiler

Streamit Parallelization: Scheduler

- Steady state remains same for each cycle

    – Find steady state schedule at compile time

    – All data rates known at compile time

    – Solve system of linear equations to find steady state schedule

- Find prologue schedule

- Low scheduling overhead at runtime

Conclusions

- Stream Programming
  - Programmer does not need to focus on concurrency
  - Programmer does need to be aware of which techniques run more efficiently in parallel
  - Exposes task, data and pipeline parallelism
- Compiler can manage parallelism
  - Choose granularity
  - Perform load balancing
  - Take care of concurrency issues
  - Optimize for given architecture

References

- Language Reference
  - http://groups.csail.mit.edu/cag/streamit/papers/streamit-lang-spec.pdf

- StreamIt Cookbook
  - http://groups.csail.mit.edu/cag/streamit/papers/streamit-cookbook.pdf

- Presentations
  - http://people.csail.mit.edu/mgordon/mgordon-phd-defense.pdf
  - http://groups.csail.mit.edu/cag/streamit/talks/StreamIt-IBM-PL-Day-04.pdf
  - http://research.microsoft.com/en-us/um/people/thies/thesis-defense.pdf
  - http://groups.csail.mit.edu/cag/streamit/talks/pact03tutorial/index.html
  - http://groups.csail.mit.edu/cag/streamit/talks/StreamIt-NEPLS-8-02.ppt

- Images and line of thought on some slides taken directly from presentations