# 4.2 Sorting and Searching



I DRAG THIS PILE OF UNSORTED PAPERWORK AROUND THE SUPERMARKET AS A SIGN OF MY SHAME

## Scan through array, looking for key.

- search hit: return array index
- search miss: return -1

```
public static int search(String key, String[] a)
{
  for (int i = 0; i < a.length; i++)
    if ( a[i].compareTo(key) == 0 ) return i;
  return -1;
}</pre>
```

Search Client: Exception Filter

Exception filter. Read a list of strings from a whitelist file, then print out all strings from standard input not in the whitelist.

| <pre>public static void main(String[] args) {</pre> |
|---|
| <pre>In in = new In(args[0]);</pre>                 |
| <pre>String s = in.readAll();</pre>                 |
| <pre>String[] words = s.split("\\s+");</pre>        |
| <pre>while (!StdIn.isEmpty())</pre>                 |
| {   |
| <pre>String key = StdIn.readString();</pre>         |
| if (search(key, words) == -1)                       |
| <pre>StdOut.println(key);</pre>                     |
| }   |
| }   |

| more test.txt<br>bob@office<br>carl@beach<br>marvin@spam<br>bob@office<br>bob@office | % more whitelist.txt<br>alice@home<br>bob@office<br>carl@beach<br>dave@boat |
|--|---|
| mallory@spam   | % java BinarySearch whitelist.txt < test.txt                                |
| dave@boat  | marvin@spam   |
| eve@airport  | mallory@spam  |
| alice@home   | eve@airport   |

# TEQ on Searching 1

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A credit card company needs to whitelist 10 million customer accounts, processing 1000 transactions per second. Using sequential search, what kind of computer is needed?

- A. Toaster
- B. Cellphone
- C. Your laptop
- D. Supercomputer
- E. Google server farm

TEQ on Searching 1

A credit card company needs to whitelist 10 million customer accounts, processing 1000 transactions per second. Using sequential search, what kind of computer is needed?

#### A. Toaster

- B. Cellphone
- C. Your laptop
- D. Supercomputer
- E. Google server farm

#### D. or E.

need enough memory for 10M accounts

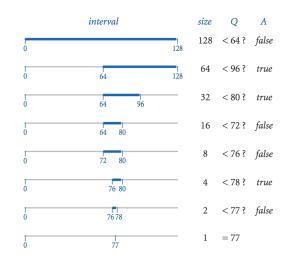
• BOE rule of thumb for any computer:

N bytes in memory, ~N memory accesses per second.

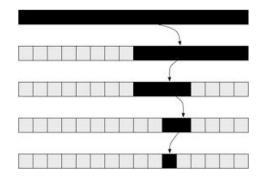
- sequential search touches about half the memory
- 2 transactions per second, 500 seconds for 1000 transactions
- fix 1: Increase memory (and speed) by factor of 1000 (supercomputer)
- fix 2: Increase number of processors by factor of 1000 (server farm)
- fix 3: Use a better algorithm (stay tuned)

#### Twenty Questions

Intuition. Find a hidden integer.



# **Binary Search**



**Binary Search** 

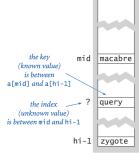
# Idea:

• Sort the array (stay tuned)

- Play "20 questions" to determine the index associated with a given key.
- Ex. Dictionary, phone book, book index, credit card numbers, ...

#### Binary search.

- Examine the middle key.
- If it matches, return its index.
- Otherwise, search either the left or right half.



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lo aback

Binary search in a sorted array (one step)

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Binary Search: Java Implementation

Invariant. Algorithm maintains  $a[lo] \leq key \leq a[hi-1]$ .

```
public static int search(String key, String[] a)
{
    return search(key, a, 0, a.length);
}
public static int search(String key, String[] a, int lo, int hi)
{
    if (hi <= lo) return -1;
    int mid = lo + (hi - lo) / 2;
    int cmp = a[mid].compareTo(key);
    if (cmp > 0) return search(key, a, lo, mid);
    else if (cmp < 0) return search(key, a, mid+1, hi);
    else return mid;
}</pre>
```

Java library implementation: Arrays.binarySearch()

Binary Search: Mathematical Analysis

Analysis. To binary search in an array of size N: do one comparison, then binary search in an array of size N / 2.

 $N \rightarrow N/2 \rightarrow N/4 \rightarrow N/8 \rightarrow \dots \rightarrow 1$ 

Q. How many times can you divide a number by 2 until you reach 1? A.  $\log_2 N$ .

```
\begin{array}{c} 1\\ 2 \rightarrow 1\\ 4 \rightarrow 2 \rightarrow 1\\ 8 \rightarrow 4 \rightarrow 2 \rightarrow 1\\ 16 \rightarrow 8 \rightarrow 4 \rightarrow 2 \rightarrow 1\\ 32 \rightarrow 16 \rightarrow 8 \rightarrow 4 \rightarrow 2 \rightarrow 1\\ 32 \rightarrow 16 \rightarrow 8 \rightarrow 4 \rightarrow 2 \rightarrow 1\\ 128 \rightarrow 64 \rightarrow 32 \rightarrow 16 \rightarrow 8 \rightarrow 4 \rightarrow 2 \rightarrow 1\\ 128 \rightarrow 64 \rightarrow 32 \rightarrow 16 \rightarrow 8 \rightarrow 4 \rightarrow 2 \rightarrow 1\\ 256 \rightarrow 128 \rightarrow 64 \rightarrow 32 \rightarrow 16 \rightarrow 8 \rightarrow 4 \rightarrow 2 \rightarrow 1\\ 512 \rightarrow 256 \rightarrow 128 \rightarrow 64 \rightarrow 32 \rightarrow 16 \rightarrow 8 \rightarrow 4 \rightarrow 2 \rightarrow 1\\ 1024 \rightarrow 512 \rightarrow 256 \rightarrow 128 \rightarrow 64 \rightarrow 32 \rightarrow 16 \rightarrow 8 \rightarrow 4 \rightarrow 2 \rightarrow 1\end{array}
```

TEQ on Searching 2

A credit card company needs to whitelist 10 million customer accounts, processing 1 thousand transactions per second. Using binary search, what kind of computer is needed?

- A. Toaster
- B. Cellphone
- C. Your laptop
- D. Supercomputer
- E. Google server farm

# Sorting

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# Sorting

## Q. What's the fastest way to sort 1 million 32-bit integers?



Sorting problem. Rearrange N items in ascending order.

Applications. Binary search, statistics, databases, data compression, bioinformatics, computer graphics, scientific computing, (too numerous to list) ...

| Hauser  | Hanley  |
|---------|---------|
| Hong    | Haskell |
| Hsu     | Hauser  |
| Hayes   | Hayes   |
| Haskell | Hong    |
| Hanley  | Hornet  |
| Hornet  | Hsu     |

Insertion Sort

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#### Insertion sort.

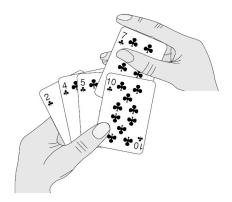
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- Brute-force sorting solution.
- Move left-to-right through array.
- Exchange next element with larger elements to its left, one-by-one.

|   |   |     |     |     | а   |     |     |     |     |
|---|---|-----|-----|-----|-----|-----|-----|-----|-----|
| 1 | J | 0   | 1   | 2   | 3   | 4   | 5   | 6   | 7   |
| 6 | 6 | and | had | him | his | was | you | the | but |
| 6 | 5 | and | had | him | his | was | the | you | but |
| 6 | 4 | and | had | him | his | the | was | you | but |
|   |   | and | had | him | his | the | was | you | but |
|   |   |     |     |     |     |     |     |     | 1.0 |

Inserting a[6] into position by exchanging with larger entries to its left

# **Insertion Sort**



#### Insertion sort.

- Brute-force sorting solution.
- Move left-to-right through array.
- Exchange next element with larger elements to its left, one-by-one.

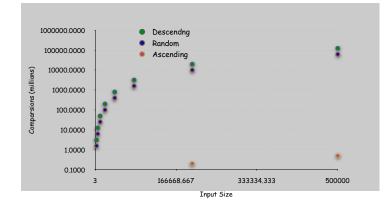
| i | 4           |                 |        |              | i        | a        |         |         |     |
|---|-------------|-----------------|--------|--------------|----------|----------|---------|---------|-----|
|   | j           | 0               | 1      | 2            | 3        | 4        | 5       | 6       | 7   |
|   |             | was             | had    | him          | and      | you      | his     | the     | but |
| 1 | 0           | had             | was    | him          | and      | you      | his     | the     | but |
| 2 | 1           | had             | him    | was          | and      | you      | his     | the     | but |
| 3 | 0           | and             | had    | him          | was      | you      | his     | the     | but |
| 4 | 4           | and             | had    | him          | was      | you      | his     | the     | but |
| 5 | 3           | and             | had    | him          | his      | was      | you     | the     | but |
| 6 | 4           | and             | had    | him          | his      | the      | was     | you     | but |
| 7 | 1           | and             | but    | had          | him      | his      | the     | was     | you |
|   |             | and             | but    | had          | him      | his      | the     | was     | you |
|   | Inserting : | а Г17 <i>th</i> | ough a | <b>FN-11</b> | into pos | ition (i | nsertio | n sort) |     |

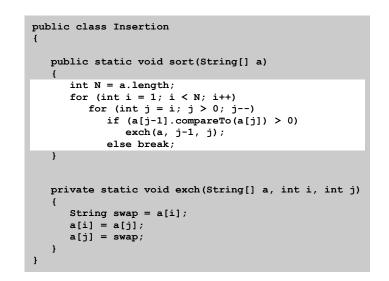
Inserting a[1] through a[N-1] into position (insertion sort)

# Insertion Sort: Empirical Analysis

# Observation. Number of comparisons depends on input family.

- Descending: ~  $N^2/2$ .
- Random: ~  $N^2/4$ .
- Ascending:  $\sim N$ .





# Insertion Sort: Mathematical Analysis

#### Worst case. [descending]

- Iteration *i* requires *i* comparisons.
- Total =  $(0 + 1 + 2 + ... + N-1) \sim N^2/2$  compares.



#### Average case. [random]

- Iteration i requires i / 2 comparisons on average.
- Total =  $(0 + 1 + 2 + ... + N-1)/2 \sim N^2/4$  compares



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## Hypothesis: Running time is ~ a N <sup>b</sup> seconds

#### Initial experiments:

| N      | Comparisons  | Time         | Ratio |
|--------|--------------|--------------|-------|
| 5,000  | 6.2 million  | 0.13 seconds |       |
| 10,000 | 25 million   | 0.43 seconds | 3.3   |
| 20,000 | 99 million   | 1.5 seconds  | 3.5   |
| 40,000 | 400 million  | 5.6 seconds  | 3.7   |
| 80,000 | 1600 million | 23 seconds   | 4.1   |

• Data source: N random numbers between 0 and 1. • Machine: Apple G5 1.8GHz with 1.5GB

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• Timing: Skagen wristwatch.

#### Doubling hypothesis:

- b = lg 4 = 2, so running time is  $\sim a N^2$
- checks with math analysis
- $a \approx 23 / 80000^2 = 3.5 \times 10^{-9}$

## Refined hypothesis: Running time is $\approx 3.5 \times 10^{-9}$ N <sup>2</sup> seconds

TEQ on Sorting 1

A credit card company uses insertion sort to sort 10 million customer account numbers, for use in whitelisting with binary search. What kind of computer is needed?

- A. Toaster
- B. Cellphone
- C. Your laptop
- D. Supercomputer
- E. Google server farm

Insertion Sort: Scientific Analysis (continued)

# Refined hypothesis: Running time is $\approx 3.5 \times 10^{-9}$ N<sup>2</sup> seconds

Prediction: Running time for N = 200,000 should be  $3.5 \times 10^{-9} \times 4 \times 10^{10} \approx 140$  seconds

#### Observation:

| N       | Time        |
|---------|-------------|
| 200,000 | 145 seconds |

Observation matches prediction and validates refined hypothesis.

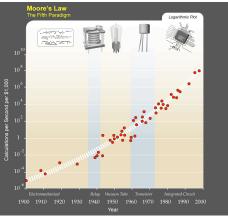
Insertion Sort: Lesson

Lesson. Supercomputer can't rescue a bad algorithm.

| Computer | Comparisons<br>Per Second | Thousand | Million  | Billion     |
|----------|---------------------------|----------|----------|-------------|
| laptop   | 107                       | instant  | 1 day    | 3 centuries |
| super    | 10 <sup>12</sup>          | instant  | 1 second | 2 weeks     |

# Moore's law. Transistor density on a chip doubles every 2 years.

Variants. Memory, disk space, bandwidth, computing power per \$.



http://en.wikipedia.org/wiki/Moore's\_law

#### Quadratic algorithms do not scale with technology.

- New computer may be 10x as fast.
- But, has 10x as much memory so problem may be 10x bigger.
- With quadratic algorithm, takes 10x as long!

"Software inefficiency can always outpace Moore's Law. Moore's Law isn't a match for our bad coding." – Jaron Lanier



Lesson. Need linear (or linearithmic) algorithm to keep pace with Moore's law.

Mergesort

#### Mergesort.

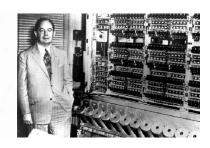
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- Divide array into two halves.
- Recursively sort each half.
- Merge two halves to make sorted whole.

First Draft of a Report on the EDVAC

John von Neumann



#### input

was had him and you his the but

sort left

and had him was you his the but

#### sort right

and had him was but his the you

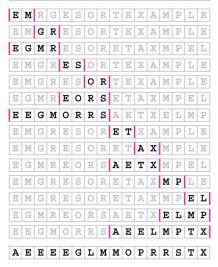
#### merge

and but had him his the was you

# Mergesort

#### Merging

MERGESORTEXAMPLE



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#### Merging

Merging. Combine two pre-sorted lists into a sorted whole.

How to merge efficiently? Use an auxiliary array.

```
String[] aux = new String[N];
// Merge into auxiliary array.
int i = lo, j = mid;
for (int k = 0; k < N; k++)
  if
           (i == mid) aux[k] = a[j++];
  else if (j == hi) aux[k] = a[i++];
   else if (a[j].compareTo(a[i]) < 0) aux[k] = a[j++];
                                      aux[k] = a[i++];
   else
}
// Copy back.
```

for (int k = 0; k < N; k++) a[lo + k] = aux[k];

Merging. Combine two pre-sorted lists into a sorted whole.

How to merge efficiently? Use an auxiliary array.

| 2 |   | Ŀ     | aux Ek T    |           |          |          | a        | ı      |          |         |     |
|---|---|-------|-------------|-----------|----------|----------|----------|--------|----------|---------|-----|
| i | j | k     | aux[k]      | 0         | 1        | 2        | 3        | 4      | 5        | 6       | 7   |
|   |   |       |             | and       | had      | him      | was      | but    | his      | the     | you |
| 0 | 4 | 0     | and         | and       | had      | him      | was      | but    | his      | the     | you |
| 1 | 4 | 1     | but         | and       | had      | him      | was      | but    | his      | the     | you |
| 1 | 5 | 2     | had         | and       | had      | him      | was      | but    | his      | the     | you |
| 2 | 5 | 3     | him         | and       | had      | him      | was      | but    | his      | the     | you |
| 3 | 5 | 4     | his         | and       | had      | him      | was      | but    | his      | the     | you |
| 3 | 6 | 5     | the         | and       | had      | him      | was      | but    | his      | the     | you |
| 3 | 6 | 6     | was         | and       | had      | him      | was      | but    | his      | the     | you |
| 4 | 7 | 7     | you         | and       | had      | him      | was      | but    | his      | the     | you |
|   |   | Trace | e of the me | rge of ti | he sorte | d left h | alf with | the so | rted rig | ht half |     |

Mergesort: Java Implementation

```
public class Merge
   public static void sort(String[] a)
   { sort(a, 0, a.length); }
   // Sort a[lo, hi).
   public static void sort(String[] a, int lo, int hi)
      int N = hi - lo;
      if (N <= 1) return;
      // Recursively sort left and right halves.
      int mid = lo + N/2;
      sort(a, lo, mid);
      sort(a, mid, hi);
      // Merge sorted halves (see previous slide).
   }
                  lo
                                          hi
                             mid
```

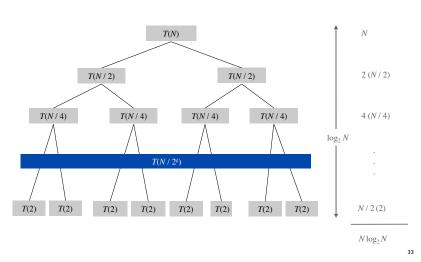


}

# Mergesort: Mathematical Analysis

# Analysis. To mergesort array of size N, mergesort two subarrays of size N/2, and merge them together using $\leq N$ comparisons.

we assume N is a power of 2



# Mergesort: Scientific Analysis

#### Hypothesis. Running time is a N lg N seconds

| Initial experiments:   | N         | Time     |
|--|-----------|----------|
|  | 4 million | 3.13 sec |
| • a ≈ 3.2 / (4 × 10 <sup>6</sup> × 32) = 2.5 × 10 <sup>-8</sup>  | 4 million | 3.25 sec |
| <sup>3</sup> <sup>4</sup> <sup>3</sup> <sup>1</sup> | 4 million | 3.22 sec |

# Refined hypothesis. Running time is $2.5 \times 10^{-7}$ N lg N seconds.

Prediction: Running time for N = 20,000,000

```
should be about 2.5 \times 10^{-8} \times 2 \times 10^7 \times 35 \approx 17.5 seconds
```

| Observation: | Ν          | Time     |
|--------------|------------|----------|
|              | 20 million | 17.5 sec |

Observation matches prediction and validates refined hypothesis.

#### Mathematical analysis.

| analysis | comparisons      |
|----------|------------------|
| worst    | $N \log_2 N$     |
| average  | $N \log_2 N$     |
| best     | $1/2 N \log_2 N$ |

#### Validation. Theory agrees with observations.

| Ν          | actual        | predicted     |
|------------|---------------|---------------|
| 10,000     | 120 thousand  | 133 thousand  |
| 20 million | 460 million   | 485 million   |
| 50 million | 1,216 million | 1,279 million |

#### TEQ on Sorting 2

A credit card company uses mergesort to sort 10 million customer account numbers, for use in whitelisting with binary search. What kind of computer is needed?

- A. Toaster
- B. Cellphone
- C. Your laptop
- D. Supercomputer
- E. Google server farm

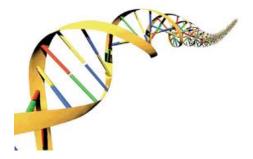
Mergesort: Lesson

#### Lesson. Great algorithms can be more powerful than supercomputers.

| Computer | Comparisons<br>Per Second | Insertion   | Mergesort |  |  |
|----------|---------------------------|-------------|-----------|--|--|
| laptop   | 107                       | 3 centuries | 3 hours   |  |  |
| super    | 10 <sup>12</sup>          | 2 weeks     | instant   |  |  |

N = 1 billion

# Longest Repeated Substring



Redundancy Detector

Longest repeated substring. Given a string, find the longest substring that appears at least twice.

a a c a a g t t t a c a a g c

#### Brute force.

- Try all indices *i* and *j* for start of possible match.
- Compute longest common prefix for each pair (quadratic+).



Applications. Bioinformatics, cryptography, ...

LRS application: patterns in music

#### Music is characterized by its repetitive structure



Source: http://www.bewitched.com/match/

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LRS applications: patterns in sequences

#### Repeated sequences in real-world data are causal.

#### Ex 1. Digits of pi

- Q. are they "random"?
- A. No, but we can't tell the difference
- Ex. Length of LRS in first 10 million digits is 14

#### Ex 2. Cryptography

- Find LRS
- Check for "known" message header identifying place, date, person, etc.
- Break code

## Ex 3. DNA

- Find LRS
- Look somewhere else for causal mechanisms
- Ex. Chromosome 11 has 7.1 million nucleotides

Longest repeated substring. Given a string, find the longest substring that

Brute-force solution

a a c a a g t t t a c a a g c

#### Brute force.

appears at least twice.

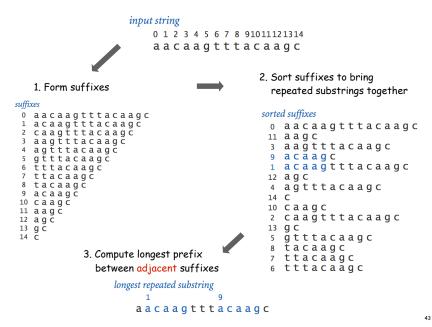
- Try all indices i and j for start of possible match.
- Compute longest common prefix (LCP) for each pair

a a c a a g t t t a c a a g c

#### Analysis.

- all pairs: 1 + 2 + ... + N ~ N<sup>2</sup>/2 calls on LCP
- too slow for long strings

Longest Repeated Substring: A Sorting Solution



Longest Repeated Substring: Java Implementation

#### Suffix sorting implementation.

```
int N = s.length();
String[] suffixes = new String[N];
for (int i = 0; i < N; i++)
    suffixes[i] = s.substring(i, N);
Arrays.sort(suffixes);
```

#### Longest common prefix: lcp(s, t).

- longest string that is a prefix of both s and t
- Ex: lcp("acaagtttac", "acaagc") = "acaag".
- easy to implement (you could write this one).

Longest repeated substring. Search only adjacent suffixes.

```
String lrs = "";
for (int i = 0; i < N-1; i++)
{
    String x = lcp(suffixes[i], suffixes[i+1]);
    if (x.length() > lrs.length()) lrs = x;
}
```

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#### Memory representation of strings.

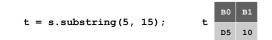
#### s = "aacaagtttacaagc";

| D0 | D1 | D2 | D3 | D4 | D5 | D6 | D7 | D8 | D9 | DA | DB | DC | DD | DE | - | A0 | A1 |
|----|----|----|----|----|----|----|----|----|----|----|----|----|----|----|---|----|----|
| a  | a  | с  | a  | a  | g  | t  | t  | t  | a  | с  | a  | a  | g  | с  | 3 | D0 | 15 |
|    |    |    |    | -  |    |    |    |    |    |    |    |    |    |    |   | 1  | 1  |

• A string is an address and a length.

• Characters can be shared among strings.

• substring() computes address, length (instead of copying chars).



#### Consequences.

- substring() is a constant-time operation (instead of linear).
- Creating suffixes takes linear space (instead of quadratic).
- Running time of LRS is dominated by the string sort.

#### Longest Repeated Substring: Empirical Analysis

| Input File     | Characters  | Brute                 | Suffix Sort | Length |
|----------------|-------------|-----------------------|-------------|--------|
| LRS.java       | 2,162       | 0.6 sec               | 0.14 sec    | 73     |
| Amendments     | 18,369      | 37 sec                | 0.25 sec    | 216    |
| Aesop's Fables | 191,945     | 3958 sec              | 1.0 sec     | 58     |
| Moby Dick      | 1.2 million | 43 hours <sup>†</sup> | 7.6 sec     | 79     |
| Bible          | 4.0 million | 20 days †             | 34 sec      | 11     |
| Chromosome 11  | 7.1 million | 2 months <sup>†</sup> | 61 sec      | 12,567 |
| Pi             | 10 million  | 4 months <sup>†</sup> | 84 sec      | 14     |

† estimated

Lesson. Sorting to the rescue; enables new research.

Many, many, many other things enabled by fast sort and search!

# TEQ on Sorting 3

# Q. Four researchers A, B, C and D are looking for long repeated subsequences in a genome with over 1 billion characters.

- A. has a grad student do it.
- B. uses brute force (check all pairs) solution.
- C. uses sorting solution with insertion sort.
- D. uses sorting solution with mergesort.

Which one is more likely to find a cancer cure?

#### Summary

Binary search. Efficient algorithm to search a sorted array.

Mergesort. Efficient algorithm to sort an array.

Applications. Many, many, many things are enabled by fast sort and search.

address

length

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