

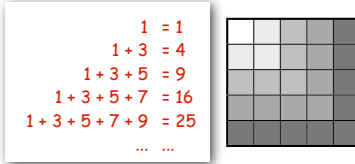


## Mathematical Induction

**Mathematical induction.** Prove a statement involving an integer  $N$  by

- **base case:** Prove it for some specific  $N$  (usually 0 or 1).
- **induction step:** Assume it to be true for all positive integers less than  $N$ , use that fact to prove it for  $N$ .

Ex. Sum of the first  $N$  odd integers is  $N^2$ .



Base case: True for  $N = 1$ .

Induction step:

- ▣ Let  $T(N)$  be the sum of the first  $N$  odd integers:  $1 + 3 + 5 + \dots + (2N - 1)$ .
- ▣ Assume that  $T(N-1) = (N-1)^2$ .
- ▣  $T(N) = T(N-1) + (2N - 1)$   
 $= (N-1)^2 + (2N - 1)$   
 $= N^2 - 2N + 1 + (2N - 1)$   
 $= N^2$

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## Recursive Program

**Recursive Program.** Implement a function having integer arguments by

- **base case:** Implementing it for some specific values of the arguments.
- **reduction step:** Assume the function works for smaller values of its arguments and use it to implement it for the given values.

Ex.  $\text{gcd}(p, q)$ .

Base case:  $\text{gcd}(p, 0) = p$ .

Reduction step:  $\text{gcd}(p, q) = \text{gcd}(p, p-q)$  if  $p > q > 0$   
 $= \text{gcd}(p, p-2q)$  if  $p-2q > 0$   
 $\dots$   
 $= \text{gcd}(p, p \% q)$

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## Greatest Common Divisor

**GCD.** Find largest integer that evenly divides into  $p$  and  $q$ .

**Euclid's algorithm.** [Euclid 300 BCE]

$$\text{gcd}(p, q) = \begin{cases} p & \text{if } q = 0 \\ \text{gcd}(q, p \% q) & \text{otherwise} \end{cases}$$

← base case  
 ← reduction step, converges to base case

$$\begin{aligned} \text{gcd}(4032, 1272) &= \text{gcd}(1272, 216) \\ &= \text{gcd}(216, 192) \\ &= \text{gcd}(192, 24) \\ &= \text{gcd}(24, 0) \\ &= 24. \end{aligned}$$

$4032 = 3 \times 1272 + 216$

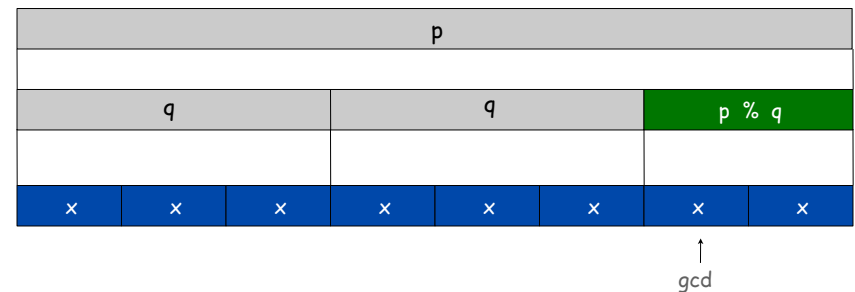
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## Euclid's Algorithm

**GCD.** Find largest integer  $d$  that evenly divides into  $p$  and  $q$ .

$$\text{gcd}(p, q) = \begin{cases} p & \text{if } q = 0 \\ \text{gcd}(q, p \% q) & \text{otherwise} \end{cases}$$

← base case  
 ← reduction step, converges to base case



$$\begin{aligned} p &= 8x \\ q &= 3x \end{aligned}$$

$$\text{gcd}(p, q) = \text{gcd}(3x, 2x) = x$$

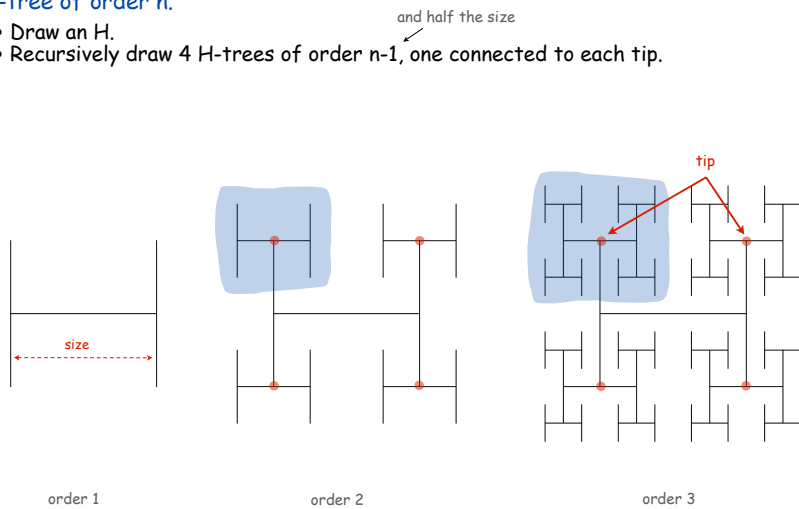
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## Htree

### H-tree of order n.

- Draw an H.
- Recursively draw 4 H-trees of order n-1, one connected to each tip.



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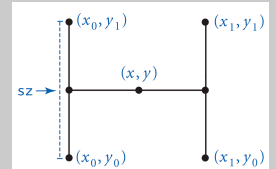
## Htree in Java

```
public class Htree
{
    public static void draw(int n, double sz, double x, double y)
    {
        if (n == 0) return;
        double x0 = x - sz/2, x1 = x + sz/2;
        double y0 = y - sz/2, y1 = y + sz/2;

        StdDraw.line(x0, y, x1, y);
        StdDraw.line(x0, y0, x0, y1); ← draw the H, centered on (x, y)
        StdDraw.line(x1, y0, x1, y1);

        draw(n-1, sz/2, x0, y0);
        draw(n-1, sz/2, x0, y1); ← recursively draw 4 half-size Hs
        draw(n-1, sz/2, x1, y0);
        draw(n-1, sz/2, x1, y1);
    }

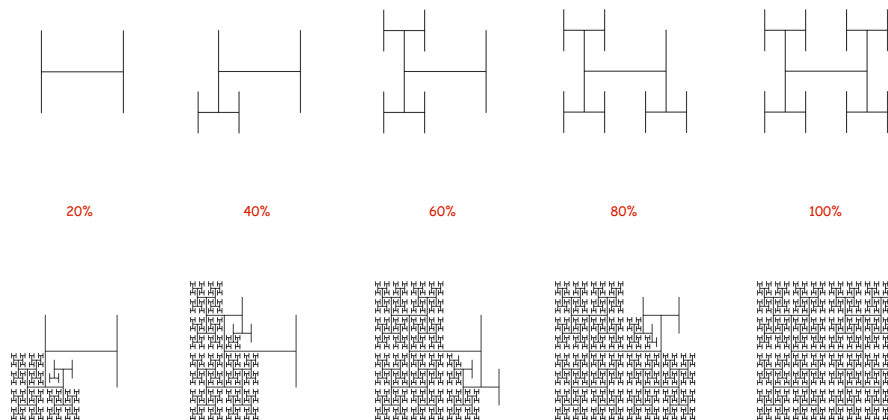
    public static void main(String[] args)
    {
        int n = Integer.parseInt(args[0]);
        draw(n, .5, .5, .5);
    }
}
```



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## Animated H-tree

Animated H-tree. Pause for 1 second after drawing each H.



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## Divide-and-Conquer

### Divide-and-conquer paradigm.

- Break up problem into smaller subproblems of same structure.
- Solve subproblems recursively using same method.
- Combine results to produce solution to original problem.

Divide et impera. Veni, vidi, vici. - Julius Caesar

### Many important problems succumb to divide-and-conquer.

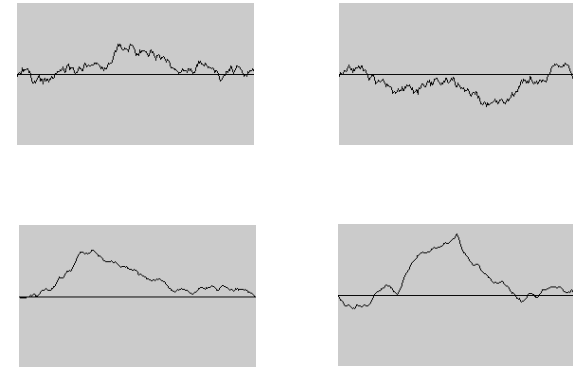
- FFT for signal processing.
- Parsers for programming languages.
- Multigrid methods for solving PDEs.
- **Quicksort and mergesort for sorting.**
- Hilbert curve for domain decomposition.
- Quad-tree for efficient N-body simulation.
- **Midpoint displacement method for fractional Brownian motion.**

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# Application: Fractional Brownian Motion

Physical process which models many natural and artificial phenomenon.

- Price of stocks.
- Dispersion of ink flowing in water.
- Rugged shapes of mountains and clouds.
- Fractal landscapes and textures for computer graphics.

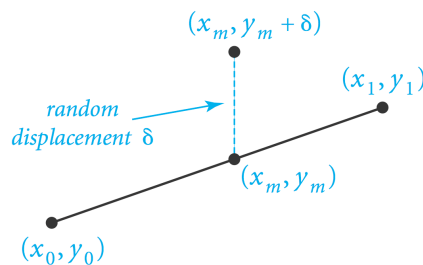


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## Simulating Brownian Motion

### Midpoint displacement method.

- Maintain an interval with endpoints  $(x_0, y_0)$  and  $(x_1, y_1)$ .
- Divide the interval in half.
- Choose  $\delta$  at random from Gaussian distribution.
- Set  $x_m = (x_0 + x_1)/2$  and  $y_m = (y_0 + y_1)/2 + \delta$ .
- Recur on the left and right intervals.



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## Simulating Brownian Motion: Java Implementation

### Midpoint displacement method.

- Maintain an interval with endpoints  $(x_0, y_0)$  and  $(x_1, y_1)$ .
- Choose  $\delta$  at random from Gaussian distribution.
- Divide the interval in half: Set  $x_m = (x_0 + x_1)/2$  and  $y_m = (y_0 + y_1)/2 + \delta$ .
- Recur on the left and right intervals.

```
public static void curve(double x0, double y0,
                        double x1, double y1, double var)
{
    if (x1 - x0 < 0.01)
    {
        StdDraw.line(x0, y0, x1, y1);
        return;
    }
    double xm = (x0 + x1) / 2;
    double ym = (y0 + y1) / 2;
    ym += StdRandom.gaussian(0, Math.sqrt(var));
    curve(x0, y0, xm, ym, var/2);
    curve(xm, ym, x1, y1, var/2);
}
```

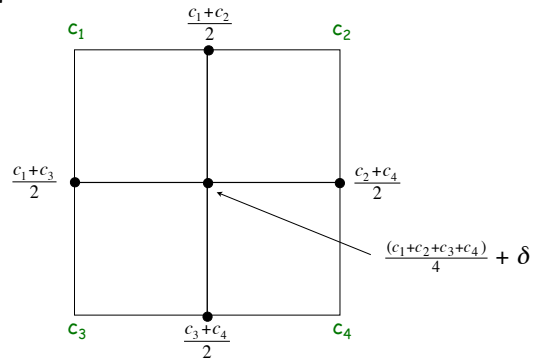
variance halves at each level;  
change factor to get different shapes

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## Plasma Cloud

Plasma cloud centered at  $(x, y)$  of size  $s$ .

- Each corner labeled with some grayscale value.
- Divide square into four quadrants.
- The grayscale of each new corner is the average of others.
  - center: average of the four corners + random displacement
  - others: average of two original corners
- Recur on the four quadrants.



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## Plasma Cloud



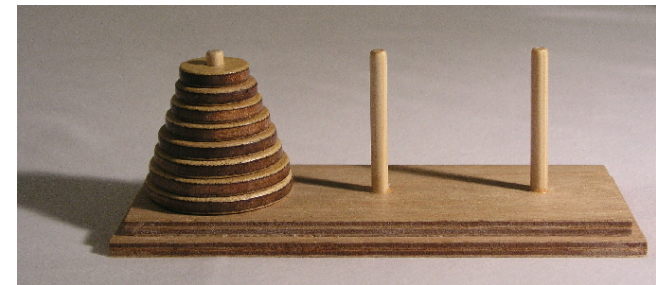
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## Brownian Landscape



Reference: [http://www.geocities.com/aaron\\_torpy/gallery.htm](http://www.geocities.com/aaron_torpy/gallery.htm)

## Towers of Hanoi



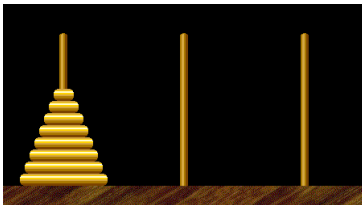
<http://en.wikipedia.org/wiki/Image:Hanoiklein.jpg>

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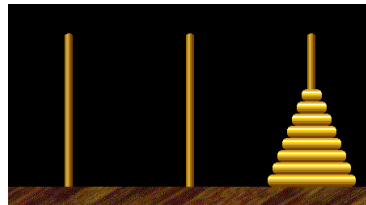
## Towers of Hanoi

Move all the discs from the leftmost peg to the rightmost one.

- Only one disc may be moved at a time.
- A disc can be placed either on empty peg or on top of a larger disc.



start



finish



Towers of Hanoi demo



Edouard Lucas (1883)

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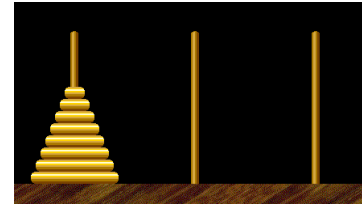
## Towers of Hanoi Legend

Q. Is world going to end (according to legend)?

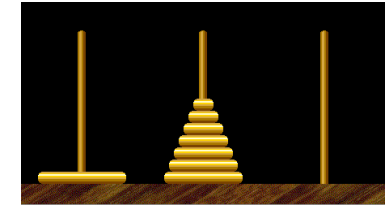
- 64 golden discs on 3 diamond pegs.
- World ends when certain group of monks accomplish task.

Q. Will computer algorithms help?

## Towers of Hanoi: Recursive Solution

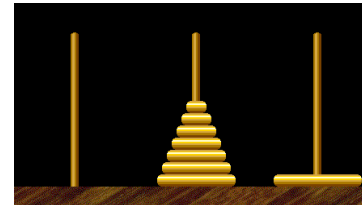


Move n-1 smallest discs right.

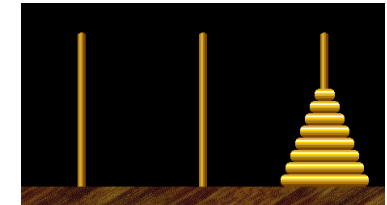


Move largest disc left.

cyclic wrap-around



Move n-1 smallest discs right.



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## Towers of Hanoi: Recursive Solution

```
public class TowersOfHanoi
{
    public static void moves(int n, boolean left)
    {
        if (n == 0) return;
        moves(n-1, !left);
        if (left) System.out.println(n + " left");
        else     System.out.println(n + " right");
        moves(n-1, !left);
    }

    public static void main(String[] args)
    {
        int N = Integer.parseInt(args[0]);
        moves(N, true);
    }
}
```

moves(n, true) : move discs 1 to n one pole to the left  
 moves(n, false) : move discs 1 to n one pole to the right

← smallest disc

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## Towers of Hanoi: Recursive Solution

```

% java TowersOfHanoi 3
1 left
2 right
1 left
3 left
1 left
2 right
1 left
        
```

```

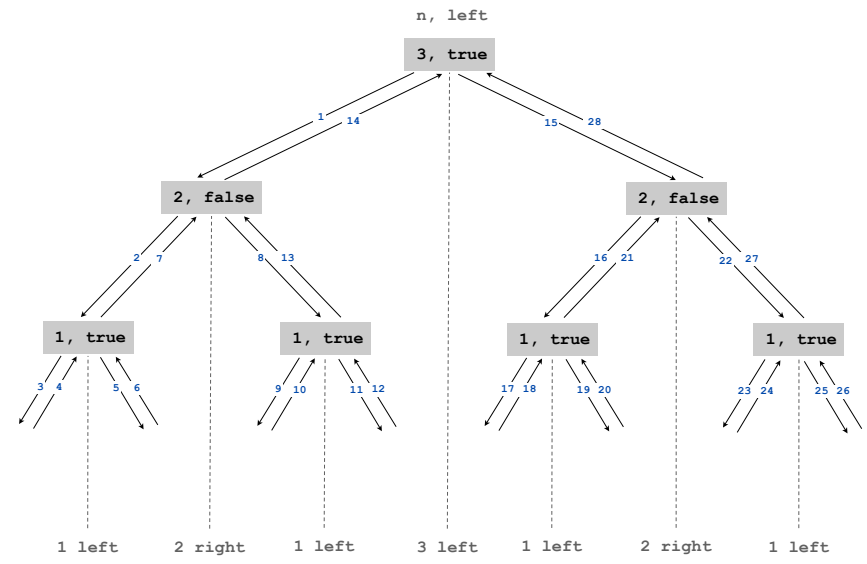
% java TowersOfHanoi 4
1 right
2 left
1 right
3 right
1 right
2 left
1 right
4 left
1 right
2 left
1 right
3 right
1 right
2 left
1 right
        
```

every other move is smallest disc

subdivisions of ruler

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## Towers of Hanoi: Recursion Tree



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## Towers of Hanoi: Properties of Solution

### Remarkable properties of recursive solution.

- Takes  $2^n - 1$  moves to solve n disc problem.
- Sequence of discs is same as subdivisions of ruler.
- Every other move involves smallest disc.

### Recursive algorithm yields non-recursive solution!

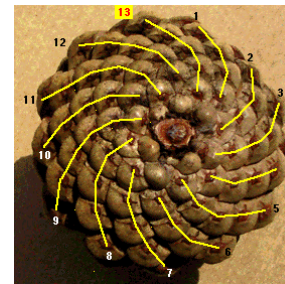
- Alternate between two moves:
    - move smallest disc to right if n is even
    - make only legal move not involving smallest disc
- ↙ to left if n is odd

### Recursive algorithm may reveal fate of world.

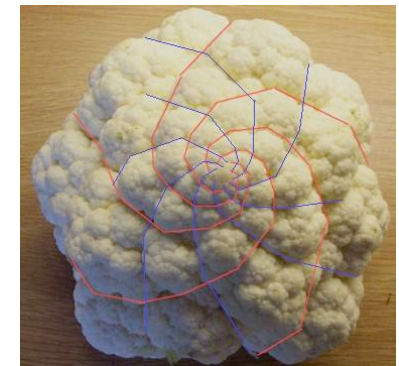
- Takes 585 billion years for  $n = 64$  (at rate of 1 disc per second).
- Reassuring fact: any solution takes at least this long!

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## Fibonacci Numbers



pinecone



cauliflower

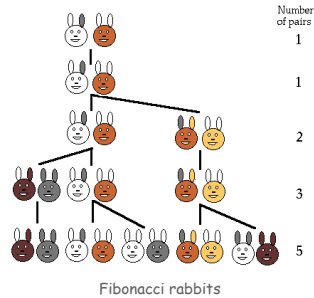
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## Fibonacci Numbers

Fibonacci numbers. 0, 1, 1, 2, 3, 5, 8, 13, 21, 34, ...

$$F_n = \begin{cases} 0 & \text{if } n = 0 \\ 1 & \text{if } n = 1 \\ F_{n-1} + F_{n-2} & \text{otherwise} \end{cases}$$



L. P. Fibonacci  
(1170 - 1250)

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TEQ on Recursion 1.1 (difficult but important)

Is this an efficient way to compute F(50)?

```
public static long F(int n)
{
    if (n == 0) return 0;
    if (n == 1) return 1;
    return F(n-1) + F(n-2);
}
```

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## A Possible Pitfall With Recursion

Fibonacci numbers. 0, 1, 1, 2, 3, 5, 8, 13, 21, 34, ...

$$F_n = \begin{cases} 0 & \text{if } n = 0 \\ 1 & \text{if } n = 1 \\ F_{n-1} + F_{n-2} & \text{otherwise} \end{cases}$$

FYI (classical math):

$$F(n) = \frac{\phi^n - (1-\phi)^n}{\sqrt{5}} = \lfloor \phi^n / \sqrt{5} \rfloor$$

$\phi = \text{golden ratio} \approx 1.618$

Ex:  $F(50) \approx 1.2 \times 10^{10}$

A natural for recursion?

```
public static long F(int n)
{
    if (n == 0) return 0;
    if (n == 1) return 1;
    return F(n-1) + F(n-2);
}
```

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TEQ on Recursion 1.2 (easy and also important)

Is this an efficient way to compute F(50)?

```
long[] F = new long[51];
F[0] = 0; F[1] = 1;
if (n == 1) return 1;
for (int i = 2; i <= 50; i++)
    F[i] = F[i-1] + F[i-2];
```

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## Summary

### How to write simple recursive programs?

- Base case, reduction step.
- Trace the execution of a recursive program.
- Use pictures.

### Why learn recursion?

Towers of Hanoi by W. A. Schloss.

- New mode of thinking.
- Powerful programming tool.

**Divide-and-conquer.** Elegant solution to many important problems.

### Exponential time.

- Easy to specify recursive program that takes exponential time.
- Don't do it unless you plan to (and are working on a small problem).