



COS 318: Operating Systems

Message Passing

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(<http://www.cs.princeton.edu/courses/cos318/>)

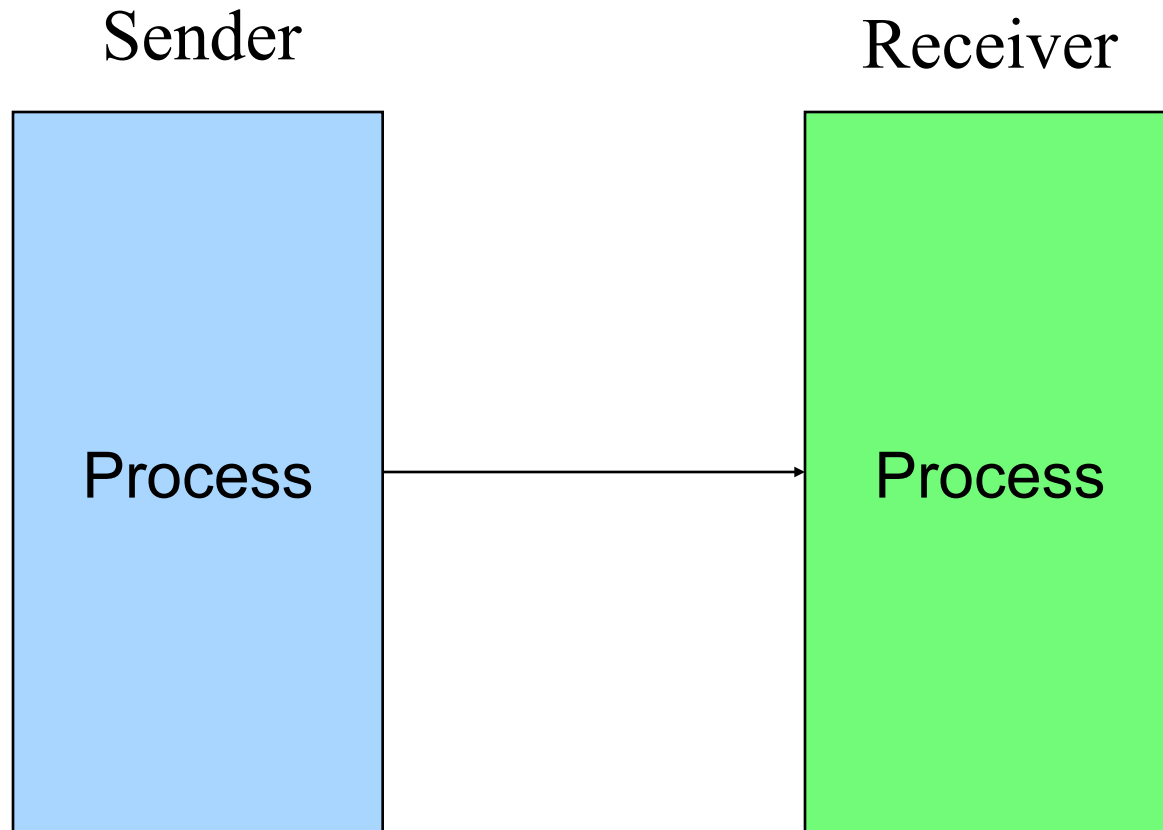


Today's Topics

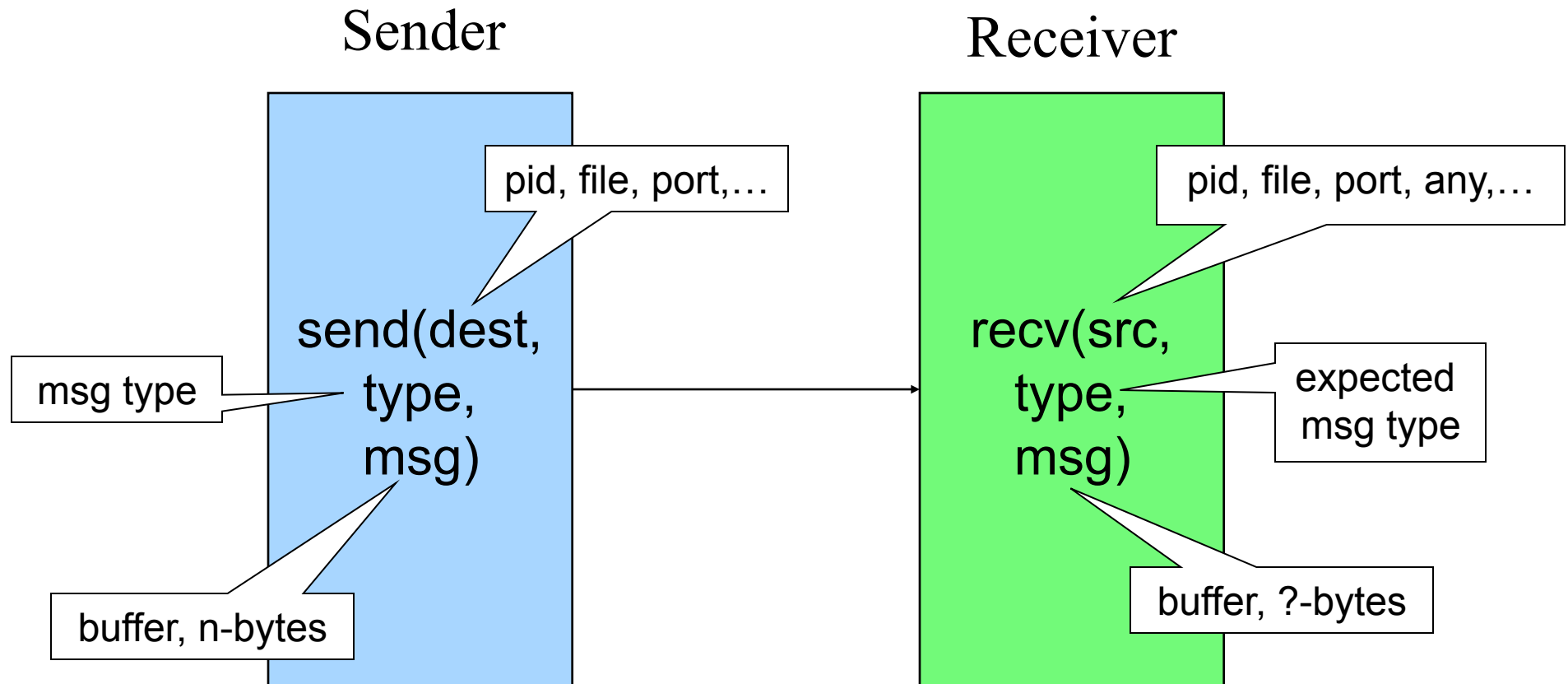
- ◆ Message passing
 - Semantics
 - How to use
- ◆ Implementation issues
 - Synchronous vs. asynchronous
 - Buffering
 - Indirection
 - Exceptions



Big Picture



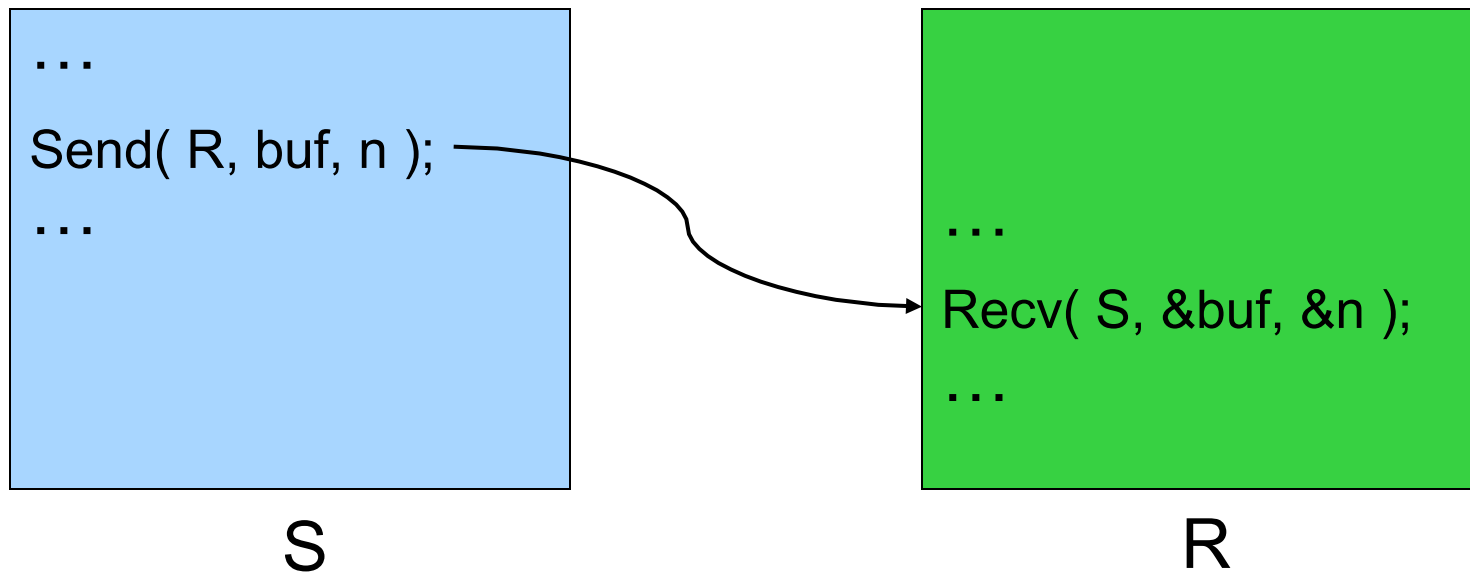
Send and Receive Primitives



Many ways to design the message passing API

Synchronous Message Passing

- ◆ Move data between processes
 - Sender: when data is ready, send it to the receiver process
 - Receiver: when the data has arrived and when the receive process is ready to take the data, move the data
- ◆ Synchronization
 - Sender: signal the receiver process that a particular event happens
 - Receiver: block until the event has happened



Example: Producer-Consumer

```
Producer () {  
    ...  
    while (1) {  
        produce item;  
        rcv(Consumer, &credit);  
        send(Consumer, item);  
    }  
}
```

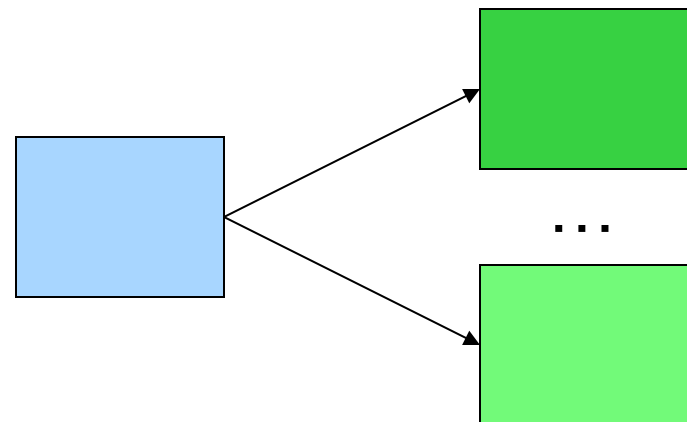
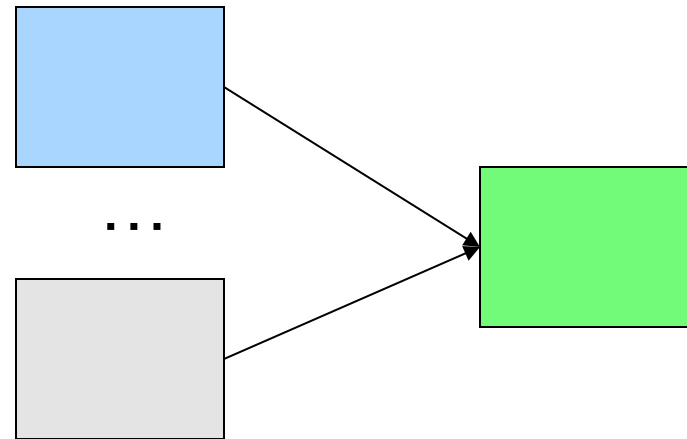
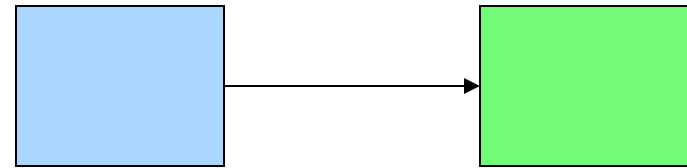
```
Consumer () {  
    ...  
    for (i=0; i<N; i++)  
        send(Producer, credit);  
    while (1) {  
        rcv(Producer, &item);  
        send(Producer, credit);  
        consume item;  
    }  
}
```

◆ Questions

- Does this work?
- Would it work with multiple producers and 1 consumer?
- Would it work with 1 producer and multiple consumers?
- What about multiple producers and multiple consumers?

Implementation Issues

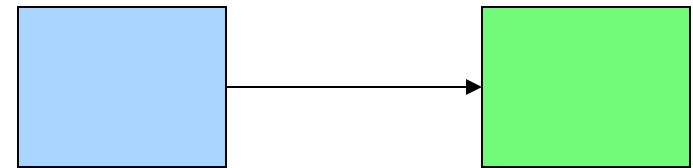
- ◆ Buffering messages
- ◆ Direct vs. indirect
- ◆ Unidirectional vs. bidirectional
- ◆ Asynchronous vs. synchronous
- ◆ Event handler vs. receive
- ◆ How to handle exceptions?



Buffering Messages

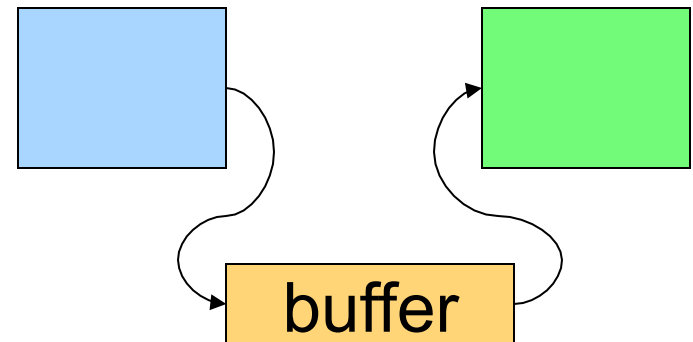
◆ No buffering

- Sender must wait until the receiver receives the message
- Rendezvous on each message



◆ Bounded buffer

- Finite size
- Sender blocks on buffer full
- Use mesa-monitor to solve the problem



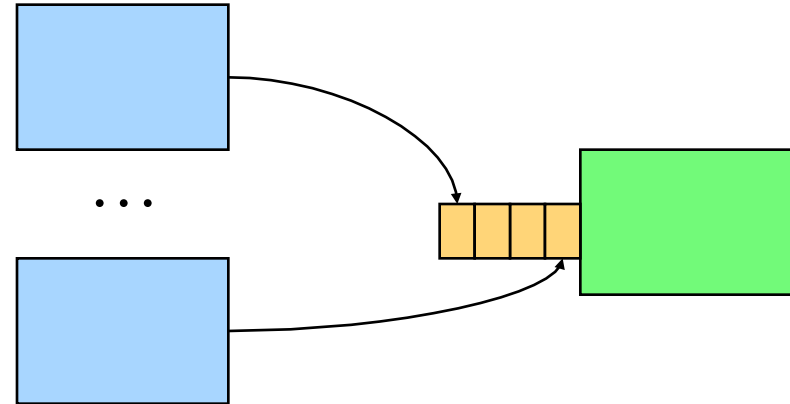
◆ Unbounded buffer

- “Infinite” size
- Sender never blocks

Direct Communication

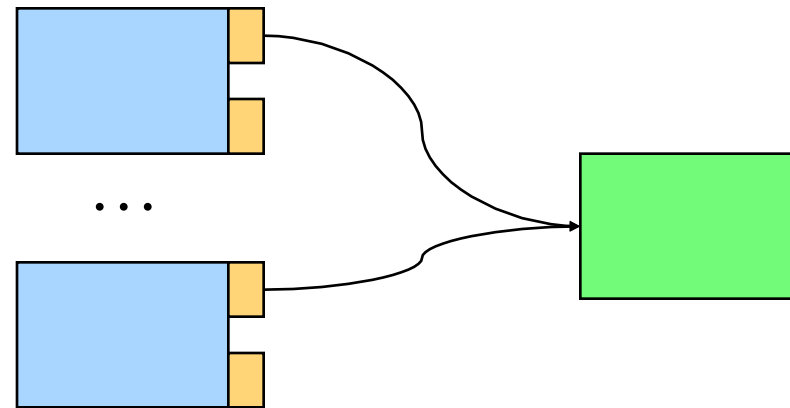
- ◆ A single buffer at the receiver

- More than one process may send messages to the receiver
- To receive from a specific sender, it requires searching through the whole buffer



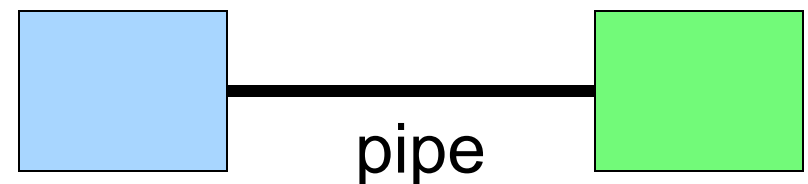
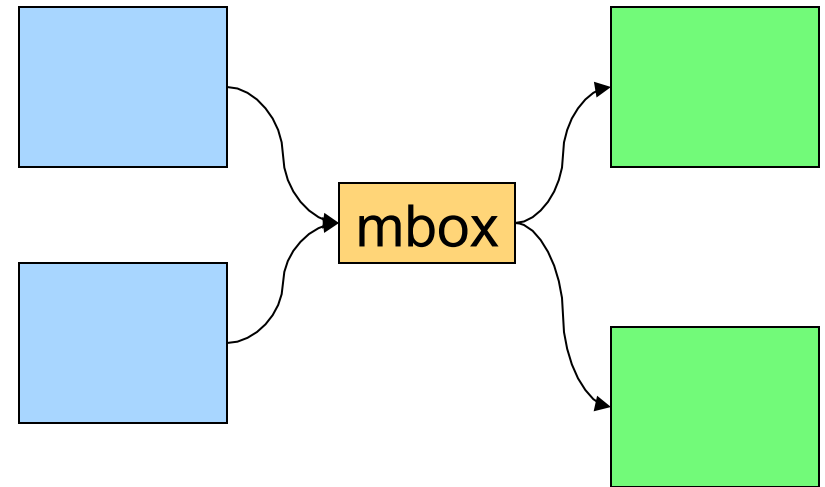
- ◆ A buffer at each sender

- A sender may send messages to multiple receivers
- To get a message, it also requires searching through the whole buffer



Indirect Communication

- ◆ Use mailbox as the abstraction
 - Allow many-to-many communication
 - Require open/close a mailbox
- ◆ Buffering
 - A buffer, its mutex and condition variables should be at the mailbox
- ◆ Message size
 - Not necessarily. One can break a large message into packets
- ◆ Mailbox vs. pipe
 - A mailbox allows many to many communication
 - A pipe implies one sender and one receiver



Synchronous vs. Asynchronous: Send

◆ Synchronous

- Block on if resource is busy
- Initiate data transfer
- Block until data is out of its source memory

```
send( dest, type, msg)
```



msg transfer resource

◆ Asynchronous

- Block if resource is busy
- Initiate data transfer and return
- Completion
 - Require applications to check status
 - Notify or signal the application

```
status = async_send( dest, type, msg )
```

```
...
```

```
if !send_complete( status )  
    wait for completion;
```

```
...
```

```
use msg data structure;
```

```
...
```



Synchronous vs. Asynchronous: Receive

◆ Synchronous

- Return data if there is a message

msg transfer resource

recv(src, type, msg)

◆ Asynchronous

- Return data if there is a message
- Return status if there is no message (probe)

```
status = async_recv( src, type, msg );  
if ( status == SUCCESS )  
    consume msg;
```

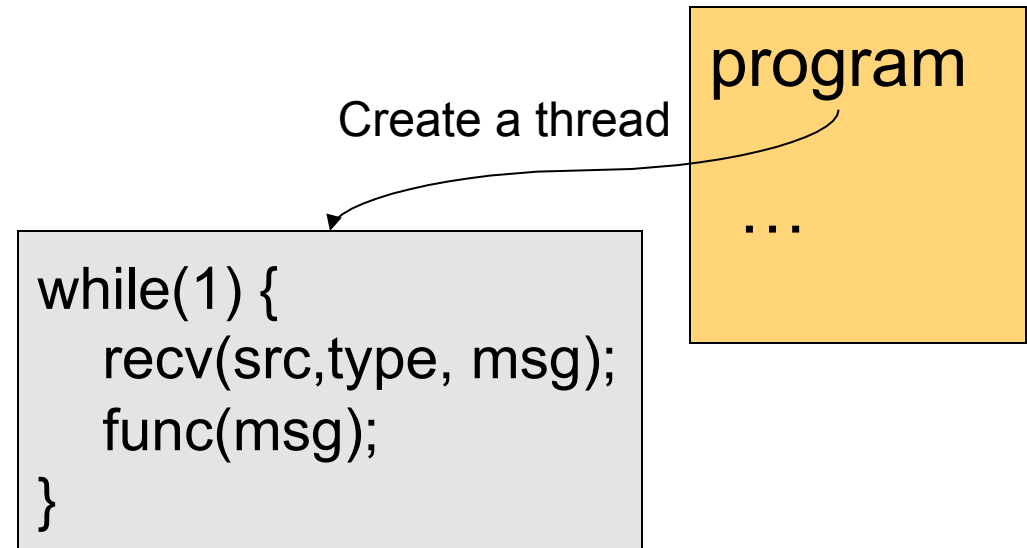
```
while ( probe(src) != HaveMSG )  
    wait for msg arrival  
recv( src, type, msg );  
consume msg;
```



Event Handler vs. Receive

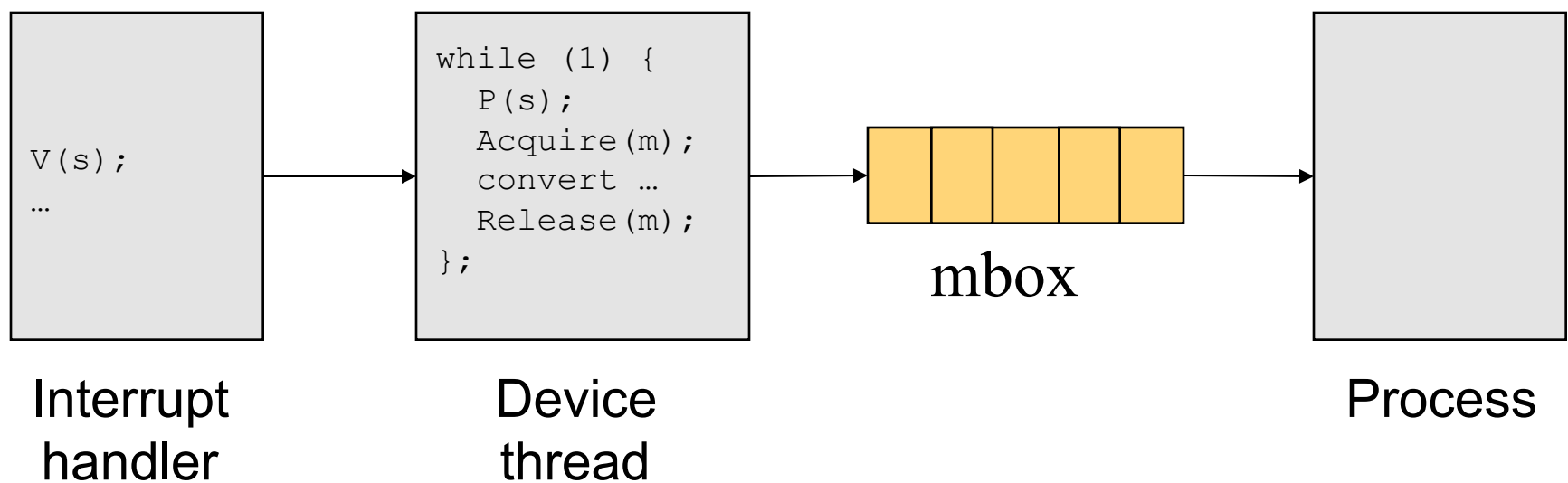
- ◆ `hrecv(src, type, msg, func)`
 - `msg` is an arg of `func`
 - Execute “`func`” on a message arrival
- ◆ Which one is more powerful?
 - `Recv` with a thread can emulate a Handler
 - Handler can be used to emulate `recv` by using Monitor
- ◆ Pros and Cons

```
void func( char * msg ) {  
    ...  
}  
  
...  
hrecv( src, type, msg, func )  
...
```



Example: Keyboard Input

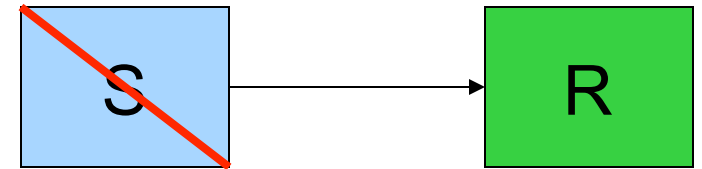
- ◆ How do you implement keyboard input?
 - Need an interrupt handler
 - Generate a mbox message from the interrupt handler
- ◆ Suppose a keyboard device thread converts input characters into an mbox message
 - How would you synchronize between the keyboard interrupt handler and device thread?
 - How can a device thread convert input into mbox messages?



Exception: Process Termination

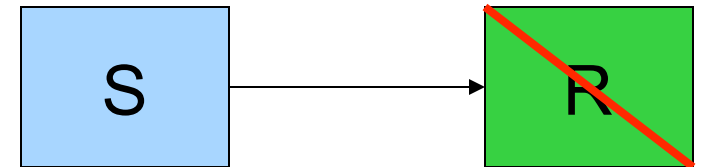
- ◆ R waits for a message from S, but S has terminated

- Problem: R may be blocked forever



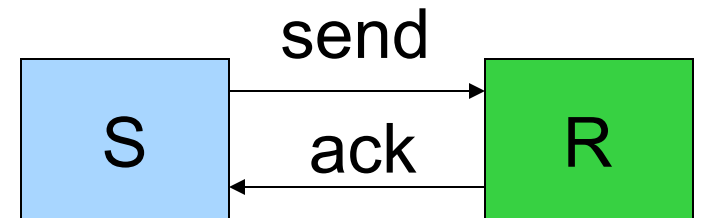
- ◆ S sends a message to R, but R has terminated

- Problem: S has no buffer and will be blocked forever



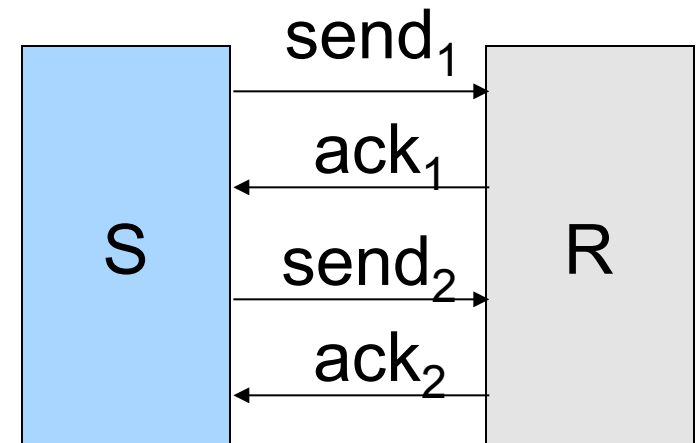
Exception: Message Loss

- ◆ Use ack and timeout to detect and retransmit a lost message
 - Require the receiver to send an ack message for each message
 - Sender blocks until an ack message is back or timeout
`status = send(dest, msg, timeout);`
 - If timeout happens and no ack, then retransmit the message
- ◆ Issues
 - Duplicates
 - Losing ack messages

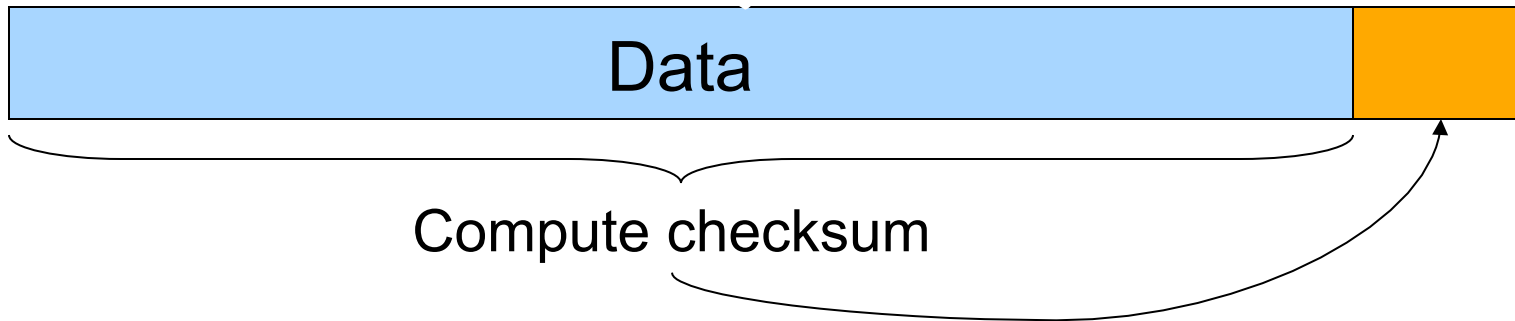


Exception: Message Loss, cont'd

- ◆ Retransmission must handle
 - Duplicate messages on receiver side
 - Out-of-sequence ack messages on sender side
- ◆ Retransmission
 - Use sequence number for each message to identify duplicates
 - Remove duplicates on receiver side
 - Sender retransmits on an out-of-sequence ack
- ◆ Reduce ack messages
 - Bundle ack messages
 - Receiver sends noack messages: can be complex
 - Piggy-back acks in send messages



Exception: Message Corruption



◆ Detection

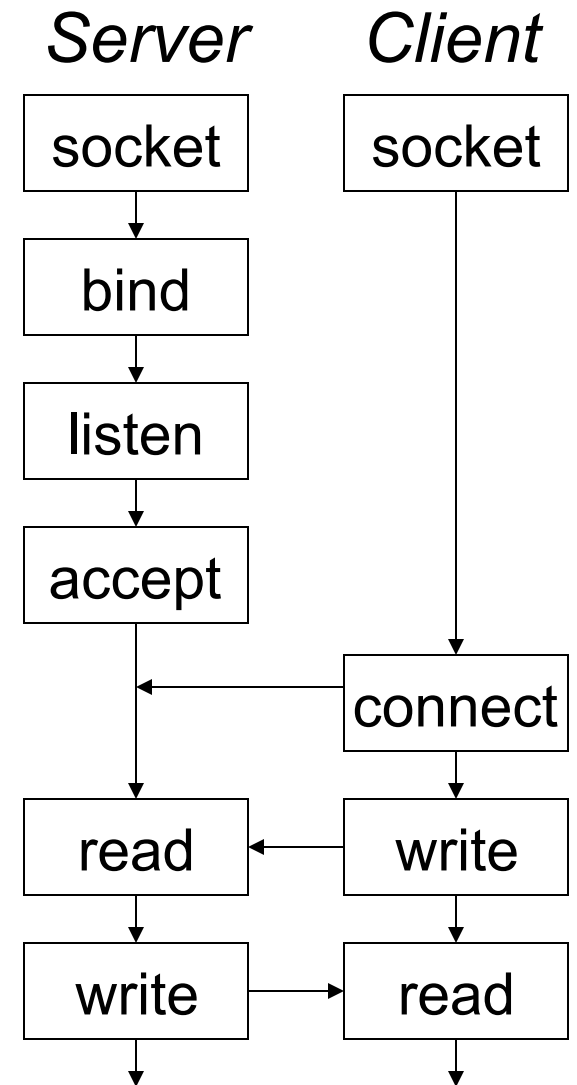
- Compute a checksum over the entire message and send the checksum (e.g. CRC code) as part of the message
- Recompute a checksum on receive and compare with the checksum in the message

◆ Correction

- Trigger retransmission
- Use correction codes to recover

Example: Sockets API

- ◆ Abstraction for TCP and UDP
 - Guest lecture by Prof. Rexford on 11/29
- ◆ Addressing
 - IP address and port number (2^{16} ports available for users)
- ◆ Create and close a socket
 - `sockid = socket(af, type, protocol);`
 - `Sockerr = close(sockid);`
- ◆ Bind a socket to a local address
 - `sockerr = bind(sockid, localaddr, addrlen);`
- ◆ Negotiate the connection
 - `listen(sockid, length);`
 - `accept(sockid, addr, length);`
- ◆ Connect a socket to destination
 - `connect(sockid, destaddr, addrlen);`



Summary

◆ Message passing

- Move data between processes
- Implicit synchronization
- API design is important

◆ Implementation issues

- Synchronous method is most common
- Asynchronous method provides overlapping but requires careful design considerations
- Indirection makes implementation flexible
- Exception needs to be carefully handled

