

## 7. Theory of Computation

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### Two fundamental questions.

- What can a computer do?
- What can a computer do with limited resources?

### General approach.

e.g., Pentium M running Linux kernel 2.6.15

- Don't talk about specific machines or problems.
- Consider minimal abstract machines.
- Consider general classes of problems.

### Why Learn Theory?

#### In theory ...

- Deeper understanding of what is a computer and computing.
- Foundation of all modern computers.
- Pure science.
- Philosophical implications.

#### In practice ...

- Web search: theory of pattern matching.
- Sequential circuits: theory of finite state automata.
- Compilers: theory of context free grammars.
- Cryptography: theory of computational complexity.
- Data compression: theory of information.

In theory there is no difference between theory and practice. In practice there is. - Yogi Berra

## Regular Expressions and DFAs

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## Pattern Matching Applications

### Test if a string matches some pattern.

- Process natural language.
- Scan for virus signatures.
- Search for information using Google.
- Access information in digital libraries.
- Retrieve information from Lexis/Nexis.
- Search-and-replace in a word processors.
- Filter text (spam, NetNanny, ads, Carnivore, malware).
- Validate data-entry fields (dates, email, URL, credit card).
- Search for markers in human genome using PROSITE patterns.

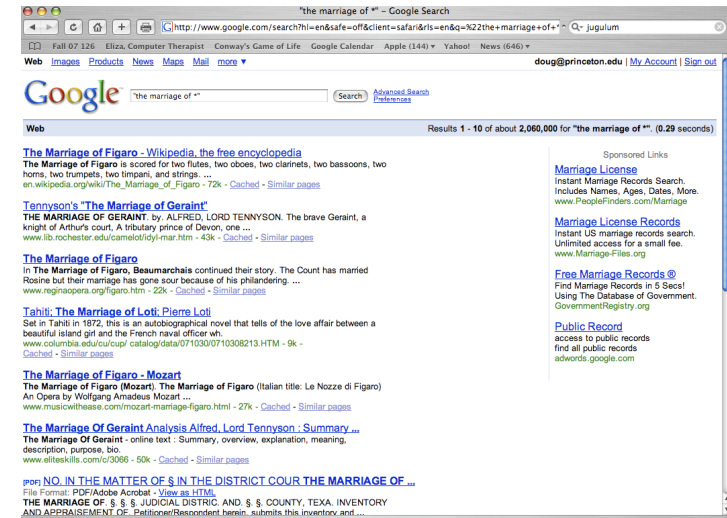
### Parse text files.

- Compile a Java program.
- Crawl and index the Web.
- Read in data stored in TOY input file format.
- Automatically create Java documentation from Javadoc comments.

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## Pattern Matching in Google

### Google. Supports \* for full word wildcard and | for union.



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## Pattern Matching in TiVo

### TiVo. WishList has very limited pattern matching.



**Using \* in WishList Searches.** To search for similar words in Keyword and Title WishList searches, use the asterisk (\*) as a special symbol that replaces the endings of words. For example, the keyword *AIRP\** would find shows containing "airport," "airplane," "airplanes," as well as the movie "Airplane!" To enter an asterisk, press the SLOW (⏪) button as you are spelling out your keyword or title.

The asterisk can be helpful when you're looking for a range of similar words, as in the example above, or if you're just not sure how something is spelled. Pop quiz: is it "irresistible" or "irresistable?" Use the keyword *IRRESIST\** and don't worry about it! Two things to note about using the asterisk:

- It can only be used at a word's end; it cannot be used to omit letters at the beginning or in the middle of a word. (For example, *AIR\*NE* or *\*PLANE* would not work.)

Reference: page 76, Hughes DirectTV TiVo manual

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## Describing a Pattern

### PROSITE. Huge database of protein families and domains.

#### Q. How to describe a protein motif?

#### Ex. [signature of the C2H2-type zinc finger domain]

- C
- Between 2 and 4 amino acids.
- C
- 3 more amino acids.
- One of the following amino acids: **LIVMFYWCX**.
- 8 more amino acids.
- H
- Between 3 and 5 more amino acids.
- H



CAASC<sup>CG</sup>PPYACGGWAGYHAGWH

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## Regular Expressions: Basic Operations

Regular expression. Notation to specify a set of strings.

Operation	Regular Expression	Yes	No
Concatenation	<code>aabaab</code>	aabaab	every other string
Wildcard	<code>.u.u.u.</code>	cumulus jugulum	succubus tumultuous
Union	<code>aa   baab</code>	aa baab	every other string
Closure	<code>ab*a</code>	aa abbba	ab ababa
Parentheses	<code>a(a b)aab</code>	aaaab abaab	every other string
	<code>(ab)*a</code>	a ababababa	aa abbba

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## Regular Expressions: Examples

Regular expression. Notation is surprisingly expressive.

Regular Expression	Yes	No
<code>.*spb.*</code> <i>contains the trigraph spb</i>	raspberry crispbread	subspace subspecies
<code>a*   (a*ba*ba*ba*)*</code> <i>multiple of three b's</i>	bbb aaa bbbaababbaa	b bb baabbbbaa
<code>.*0...</code> <i>fifth to last digit is 0</i>	1000234 98701234	111111111 403982772
<code>gcg(cgg agg)*ctg</code> <i>fragile X syndrome indicator</i>	gcgctg gcgcgctg gcgcgaggctg	gcgcgg cgcgcgcgctg gcgcgagctg

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## Generalized Regular Expressions

Regular expressions are a standard programmer's tool.

- Built in to Java, Perl, Unix, Python, ....
- Additional operations typically added for convenience.
- Ex: `[a-e]+` is shorthand for `(a|b|c|d|e)(a|b|c|d|e)*`.

Operation	Regular Expression	Yes	No
One or more	<code>a(bc)+de</code>	abcde abcbcde	ade bcde
Character classes	<code>[A-Za-z][a-z]*</code>	lowercase Capitalized	camelCase 4illegal
Exactly k	<code>[0-9]{5}-[0-9]{4}</code>	08540-1321 19072-5541	111111111 166-54-1111
Negations	<code>[^aeiou]{6}</code>	rhythm	decade

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## Regular Expressions in Java

Validity checking. Is `input` in the set described by the `re`?

```
public class Validate {
    public static void main(String[] args) {
        String re = args[0];
        String input = args[1];
        StdOut.println(input.matches(re));
    }
}
```

powerful string library method

```
% java Validate "C.{2,4}C...[LIVMFYWC].{8}H.{3,5}H" CAASCGGPYACGGAAGYHAGAH
true
% java Validate "[$_A-Za-z][$_A-Za-z0-9]*" ident123
true
% java Validate "[a-z]+@[([a-z]+\.)+(edu|com)]" doug@cs.princeton.edu
true
```

*C2H2 type zinc finger domain*

*legal Java identifier*

*valid email address (simplified)*

*need quotes to "escape" the shell*

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## More String Library Functions

### String searching methods.

<code>public class String</code> (Java's String library)	
<code>boolean matches(String re)</code>	does this string match the given regular expression
<code>String replaceAll(String re, String str)</code>	replace all occurrences of regular expression with the replacement string
<code>int indexOf(String r, int from)</code>	return the index of the first occurrence of the string r after the index from
<code>String[] split(String re)</code>	split the string around matches of the given regular expression

```
String s = StdIn.readAll();
s = s.replaceAll("\\s+", " ");
```

replace all sequences of whitespace characters with a single space

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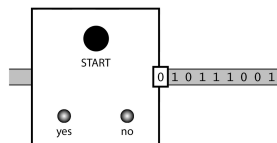
## Solving the Pattern Match Problem

### Regular expressions are a concise way to describe patterns.

- How would you implement the method `matches()` ?
- Hardware: build a deterministic finite state automaton (DFA).
- Software: simulate a DFA.

### DFA: simple machine that solves a pattern match problem.

- Different machine for each pattern.
- Accepts or rejects string specified on input tape.
- Focus on `true` or `false` questions for simplicity.



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```
String s = StdIn.readAll();
String[] words = s.split("\\s+");
```

create array of words in document

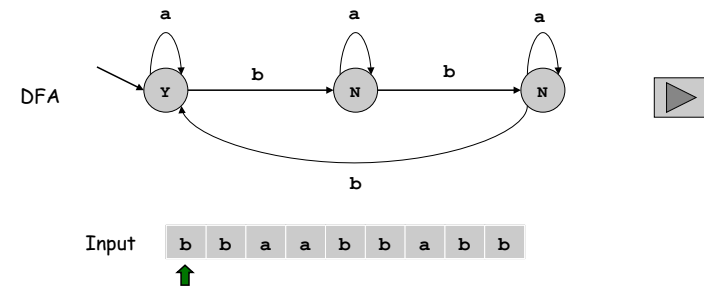
regular expression that matches any whitespace character

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## Deterministic Finite State Automaton (DFA)

### Simple machine with N states.

- Begin in start state.
- Read first input symbol.
- Move to new state, depending on current state and input symbol.
- Repeat until last input symbol read.
- Accept input string if last state is labeled Y.



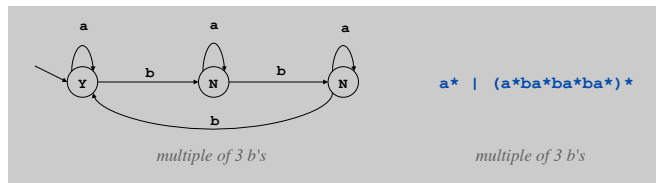
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## DFA and RE Duality

**RE.** Concise way to **describe** a set of strings.

**DFA.** Machine to **recognize** whether a given string is in a given set.

**Duality.** For any DFA, there exists a RE that describes the same set of strings; for any RE, there exists a DFA that recognizes the same set.



**Practical consequence of duality proof:** to match RE, (i) build DFA and (ii) simulate DFA on input string.

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## Implementing a Pattern Matcher

**Problem.** Given a RE, create program that tests whether given input is in set of strings described.

**Step 1.** Build the DFA.

- A compiler!
- See COS 226 or COS 320.

**Step 2.** Simulate it with given input.

```
State state = start;
while (!StdIn.isEmpty()) {
    char c = StdIn.readChar();
    state = state.next(c);
}
StdOut.println(state.accept());
```

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## Application: Harvester

Harvest information from input stream.

- Harvest patterns from DNA.

```
% java Harvester "gcg(cgg|agg)*ctg" chromosomeX.txt
gcgcggcggcggcggcggctg
gcgctg
gcgctg
gcgcggcggcggaggcggaggcggctg
```

- Harvest email addresses from web for spam campaign.

```
% java Harvester "[a-z]+@[a-z]+\.(edu|com)" http://www.princeton.edu/~cos126
rs@cs.princeton.edu
dgabai@cs.princeton.edu
doug@cs.princeton.edu
wayne@cs.princeton.edu
```

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## Application: Harvester

Harvest information from input stream.

- Use `Pattern` data type to compile regular expression to NFA.
- Use `Matcher` data type to simulate NFA.

equivalent, but more efficient representation of a DFA

```
import java.util.regex.Pattern;
import java.util.regex.Matcher;

public class Harvester {
    public static void main(String[] args) {
        String re = args[0];
        In in = new In(args[1]);
        String input = in.readAll();
        Pattern pattern = Pattern.compile(re);
        Matcher matcher = pattern.matcher(input);

        while (matcher.find()) {
            StdOut.println(matcher.group());
        }
    }
}
```

Annotations in the code block:  
 - `Pattern.compile(re)`: create NFA from RE  
 - `pattern.matcher(input)`: create NFA simulator  
 - `matcher.find()`: look for next match  
 - `matcher.group()`: the match most recently found

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Ex: parsing an NCBI genome data file.

```
LOCUS AC146846 128142 bp DNA linear HTG 13-NOV-2003
DEFINITION Ornithorhynchus anatinus clone CLM1-393H9,
ACCESSION AC146846
VERSION AC146846.2 GI:38304214
KEYWORDS HTG; HTGS_PHASE2; HTGS_DRAFT.
SOURCE Ornithorhynchus anatinus
ORIGIN
  1 tgtatttcat ttgaccgtgc tgttttttcc oggtttttca gtaccgtgtt agggagccac
  61 gtgattctgt ttgttttatg ctgcogaata gctgctogat gaatctctgc atagacagct // a comment
 121 gccgcaggga gaaatgacca gtttgtgatg acaaaatgta ggaagctgt ttcttcataa
  ...
128101 gaaatgcca cccccaagct aatgtacagc ttcttagat tg
//
```

```
String re = "[ ]*[0-9]+([actg ])*.*";
Pattern pattern = Pattern.compile(re);
In in = new In(filename);
while (!in.isEmpty()) {
    String line = in.readLine();
    Matcher matcher = pattern.matcher(line);
    if (matcher.find()) {
        String s = matcher.group(1).replaceAll(" ", "");
        // do something with s
    }
}
```

Summary

Programmer.

- Regular expressions are a powerful pattern matching tool.
- Implement regular expressions with finite state machines.

Theoretician.

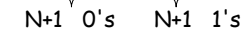
- Regular expression is a compact description of a set of strings.
- DFA is an abstract machine that solves pattern match problem for regular expressions.
- DFAs and regular expressions have limitations.

Variations

- Yes (accept) and No (reject) states sometimes drawn differently
- Terminology: Deterministic Finite State Automaton (DFA), Finite State Machine (FSM), Finite State Automaton (FSA) are the same
- DFAs can have output, specified on the arcs or in the states
  - These may not have explicit Yes and No states

No DFA can recognize the language of all bit strings with an equal number of 0's and 1's.

- Suppose an N-state DFA can recognize this language.
- Consider following input: 0000000011111111



- DFA must accept this string.
- Some state **x** is revisited during first N+1 0's since only N states

0000000011111111

x x



- Machine would accept same string without intervening 0's.

00000111111111

x

- This string doesn't have an equal number of 0's and 1's.



Fundamental Questions

Q. Are there patterns that cannot be described by any RE/DFA?

- A. Yes.
- Bit strings with equal number of 0s and 1s.
  - Decimal strings that represent prime numbers.
  - DNA strings that are Watson-Crick complemented palindromes.
  - and many, many more . . .

Q. Can we extend RE/DFA to describe richer patterns?

- A. Yes.
- Context free grammar (e.g., Java).
  - Turing machines.

# Turing Machines

**Challenge:** Design simplest machine that is "as powerful" as conventional computers.



Alan Turing (1912-1954)

## Turing Machine

**Desiderata.** Simple model of computation that is "as powerful" as conventional computers.

**Intuition.** Simulate how humans calculate.

**Ex.** Addition.

			1	2	3	4	5	6	
			+	3	1	4	1	5	9



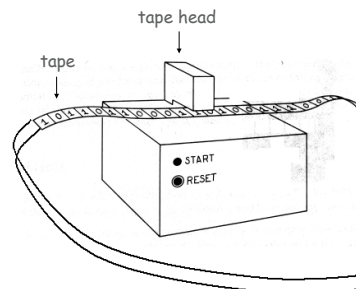
## Turing Machine: Tape

### Tape.

- Stores input, output, and intermediate results.
- One arbitrarily long strip, divided into cells.
- Finite alphabet of symbols.

### Tape head.

- Points to one cell of tape.
- Reads a symbol from active cell.
- Writes a symbol to active cell.
- Moves left or right one cell at a time.



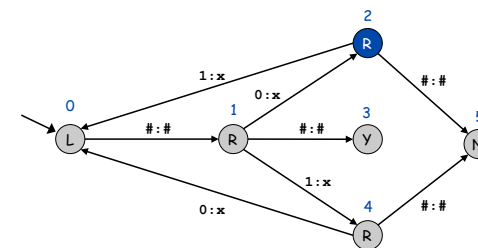
## Turing Machine: Fetch, Execute

### States.

- Finite number of possible machine configurations.
- Determines what machine does and which way tape head moves.

### State transition diagram.

- Ex. if in state 2 and input symbol is 1 then: overwrite the 1 with x, move to state 0, move tape head to left.



Before



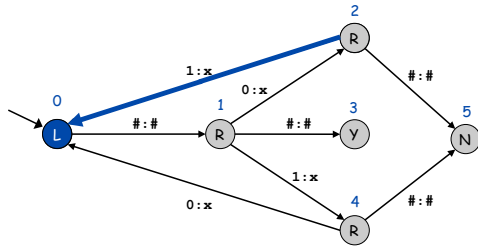
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After



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## Turing Machine: Initialization and Termination

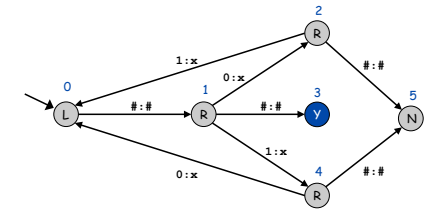
### Initialization.

- Set input on some portion of tape.
- Set tape head.
- Set initial state.



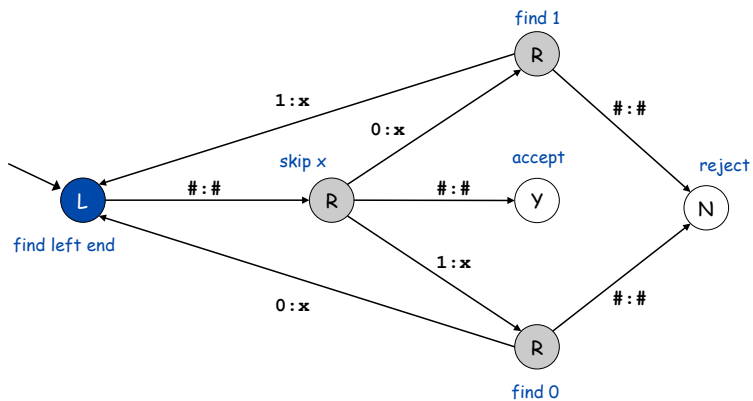
### Termination.

- Stop if enter yes, no, or halt state.
- Infinite loop possible.
  - (definitely stay tuned!)



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## Example: Equal Number of 0's and 1's



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## Turing Machine Summary

Goal: simplest machine that is "as powerful" as conventional computers.

Surprising Fact 1. Such machines are very simple: TM is enough!

Surprising Fact 2. Some problems cannot be solved by ANY computer.

↑  
next lecture

### Consequences.

- Precursor to general purpose programmable machines.
- Exposes fundamental limitations of all computers.
- Enables us to study the physics and universality of computation.
- No need to seek more powerful machines!

### Variations

- Instead of just recognizing strings, TM's can produce output: the contents of the tape
- Instead of Y and N states, TM's can have a plain Halt state

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## Alan Turing

### Alan Turing (1912-1954).

- Father of computer science.
- Computer Science's "Nobel Prize" is called the Turing Award.

FIRST HALF TERM.	MARKS	REMARKS	MASTERS
ENGLISH SUBJECTS (History, English, Theory, Geography) No. 29	23	I can judge his writing though it is the worst I have ever seen & I try to write through his handwriting, inexactitude and sloppiness, but I look in astonishment through the microscope, in a schoolroom, and I cannot suppose the possibility of his writing being the manner of the handwriting.	
LATIN No. 21	20	He ought not to be in the form of course as he is from behind, i.e. He is in a crowd behind.	

Alan's report card at 14.



Alan Turing and his elder brother.