Function Calls

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COS 217

Reading: Chapter 4 of “Programming From the Ground Up”

(available online from the course Web site)

Goals of Today’s Lecture

- Finishing introduction to assembly language
  - EFLAGS register and conditional jumps
  - Addressing modes
- Memory layout of the UNIX process
  - Data, BSS, roData, Text
  - Stack frames, and the stack pointer ESP
- Calling functions
  - Call and ret commands
  - Placing arguments on the stack
  - Using the base pointer EBP

Detailed Example

n	%edx
count	%ecx

```
count=0;
while (n>1) {
   count++;
   if (n&1)
      n = n*3+1;
   else
      n = n/2;
}
```

```
        movl  $0, %ecx
.loop:
        cmpl  $1, %edx
        jle   .endloop
        addl  $1, %ecx
        movl  %edx, %eax
        andl  $1, %eax
        je    .else
        movl  %edx, %eax
        addl  %eax, %edx
        addl  %eax, %edx
        addl  $1, %edx
        jmp   .endif
.else:
        sarl  $1, %edx
.endif:
        jmp   .loop
.endloop:
```

Setting the EFLAGS Register

- Comparison **`cmpl`** compares two integers
  - Done by subtracting the first number from the second
    - Discarding the results, but setting the eflags register
  - Example:
    - **`cmpl $1, %edx`** (computes %edx – 1)
    - **`jle .endloop`** (looks at the sign flag and the zero flag)
- Logical operation **`andl`** compares two integers
  - Example:
    - **`andl $1, %eax`** (bit-wise AND of %eax with 1)
    - **`je .else`** (looks at the zero flag)
- Unconditional branch **`jmp`**
  - Example:
    - **`jmp .endif`** and **`jmp .loop`**

EFLAGS Register & Condition Codes

31–22	21	20	19	18	17	16	15	14	13–12	11	10	9	8	7	6	5	4	3	2	1	0
Reserved (set to 0)	ID	VIP	VIF	AC	VM	RF	0	NT	IOPL	OF	DF	IF	TF	SF	ZF	0	AF	0	PF	1	CF

- Identification flag
- Virtual interrupt pending
- Virtual interrupt flag
- Alignment check
- Virtual 8086 mode
- Resume flag
- Nested task flag
- I/O privilege level
- **Overflow flag**
- Direction flag
- Interrupt enable flag
- Trap flag
- **Sign flag**
- **Zero flag**
- Auxiliary carry flag or adjust flag
- Parity flag
- **Carry flag**

A Simple Assembly Program

```
.section .data
# pre-initialized
# variables go here

.section .bss
# zero-initialized
# variables go here

.section .rodata
# pre-initialized
# constants go here
```

```
.section .text
.globl _start
_start:
# Program starts executing
# here

# Body of the program goes
# here

# Program ends with an
# “exit()” system call
# to the operating system
movl $1, %eax
movl $0, %ebx
int $0x80
```

Main Parts of the Program

- Break program into sections (**`.section`**)
  - Data, BSS, RoData, and Text
- Starting the program
  - Making _start a global (**`.global _start`**)
    - Tells the assembler to remember the symbol **`_start`**
    - … because the linker will need it
  - Identifying the start of the program (**`_start`**)
    - Defines the value of the label **`_start`**
- Exiting the program
  - Specifying the exit() system call (**`movl $1, %eax`**)
    - Linux expects the system call number in EAX register
  - Specifying the status code (**`movl $0, %ebx`**)
    - Linux expects the status code in EBX register
  - Interrupting the operating system (**`int $0x80`**)

Function Calls

- Function
  - A piece of code with well-defined entry and exit points, and a well-defined interface
- “Call” and “Return” abstractions
  - **Call:** jump to the beginning of an arbitrary procedure
  - **Return:** jump to the instruction immediately following the “most-recently-executed” Call instruction
- The jump address in the return operation is dynamically determined

Implementing Function Calls

```
P:                # Function P

   …

   jmp R         # Call R

Rtn_point1:

   …
```

```
R:                # Function R

   …

   jmp ???       # Return
```

```
Q:                # Function Q

   …

   jmp R         # Call R

Rtn_point2:

   …
```

What should the return instruction in R jump to?

Implementing Function Calls

```
P:                # Proc P
   movl $Rtn_point1, %eax
   jmp R      # Call R
Rtn_point1:
   …
```

```
R:                # Proc R
   …
   jmp %eax   # Return
```

```
Q:                # Proc Q
   movl $Rtn_point2, %eax
   jmp R      # Call R
Rtn_point2:
   …
```

Convention: At Call time, store return address in EAX

Problem: Nested Function Calls

```
P:              # Function P
   movl $Rtn_point1, %eax
   jmp Q        # Call Q
Rtn_point1:
   …
```

```
Q:              # Function Q
   movl $Rtn_point2, %eax
   jmp R        # Call R
Rtn_point2:
   …
   jmp %eax     # Return
```

```
R:              # Function R
   …
   jmp %eax     # Return
```

- Problem if P calls Q, and Q calls R
- Return address for P to Q call is lost

Need to Use a Stack

- A return address needs to be saved for as long as the function invocation continues
- Return addresses are used in the reverse order that they are generated: Last-In-First-Out
- The number of return addresses that may need to be saved is not statically known
- Saving return addresses on a Stack is the most natural solution

Stack Frames

- Use stack for all temporary data related to each active function invocation
    - Return address
    - Input parameters
    - Local variables of function
    - Saving registers across invocations

  (the four items above together form the) **Stack Frame**

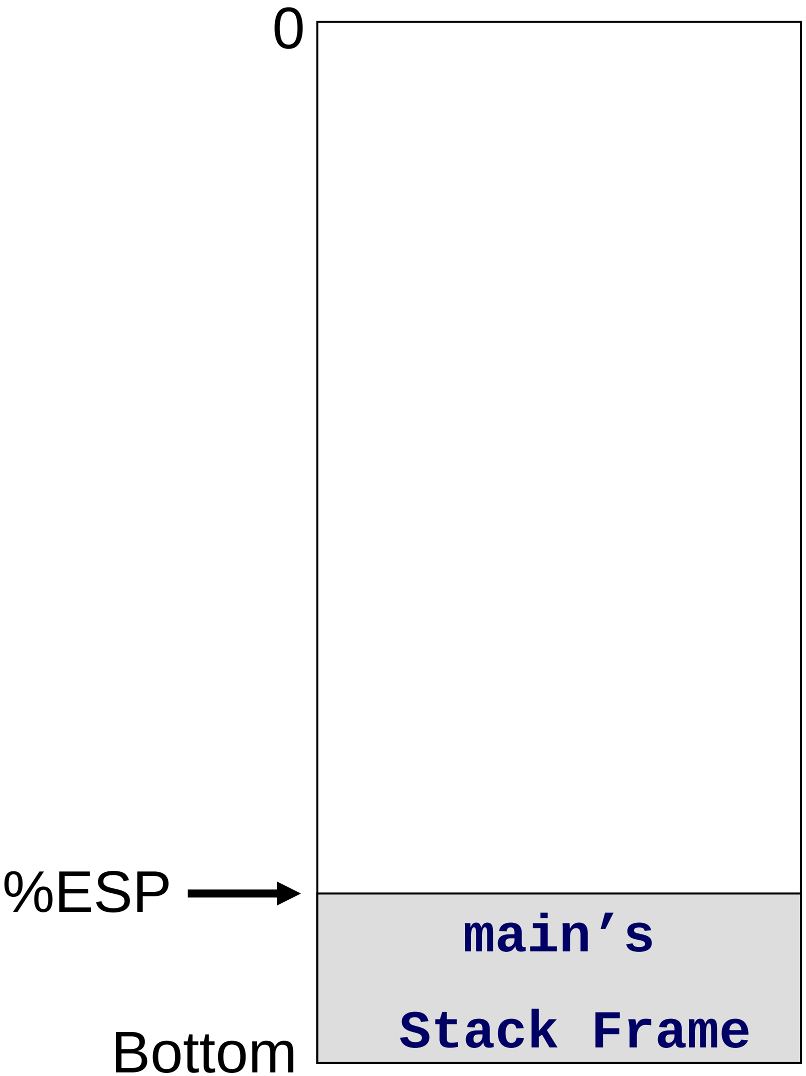
- Stack has one Stack Frame per active function invocation

High-Level Picture

- At Call time, push a new Stack Frame on top of the stack
- At Return time, pop the top-most Stack Frame

High-Level Picture

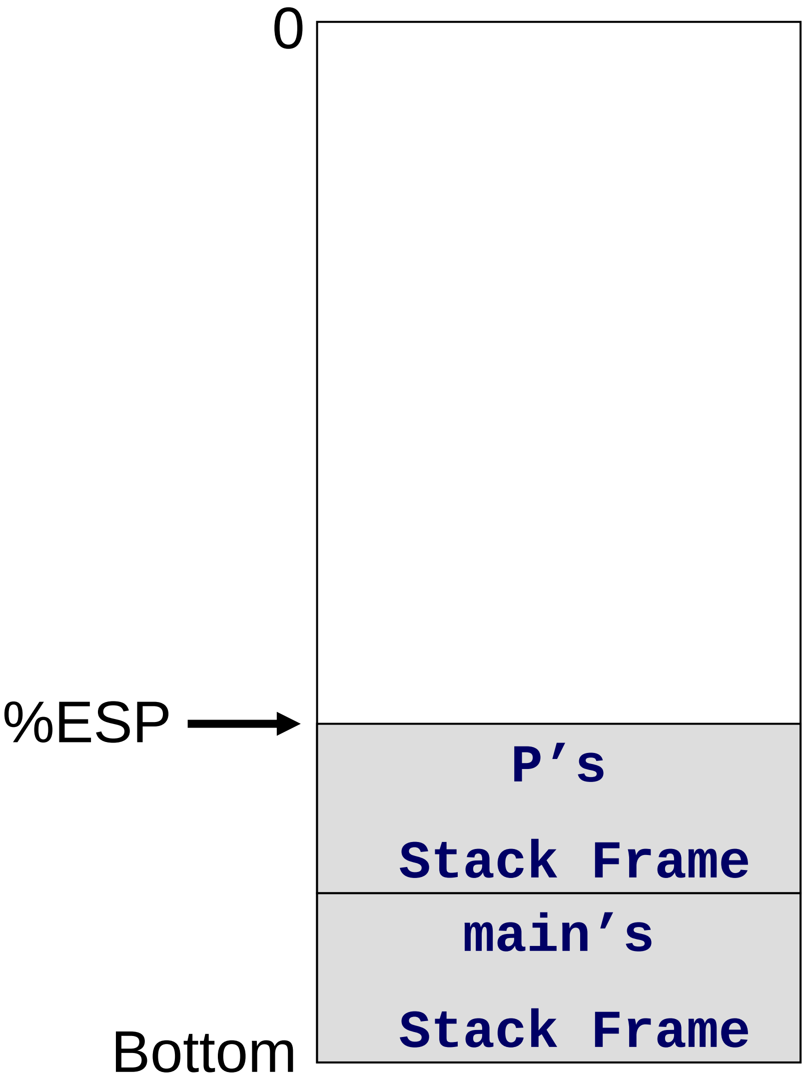
main begins executing



High-Level Picture

`main begins executing`

`main calls P`

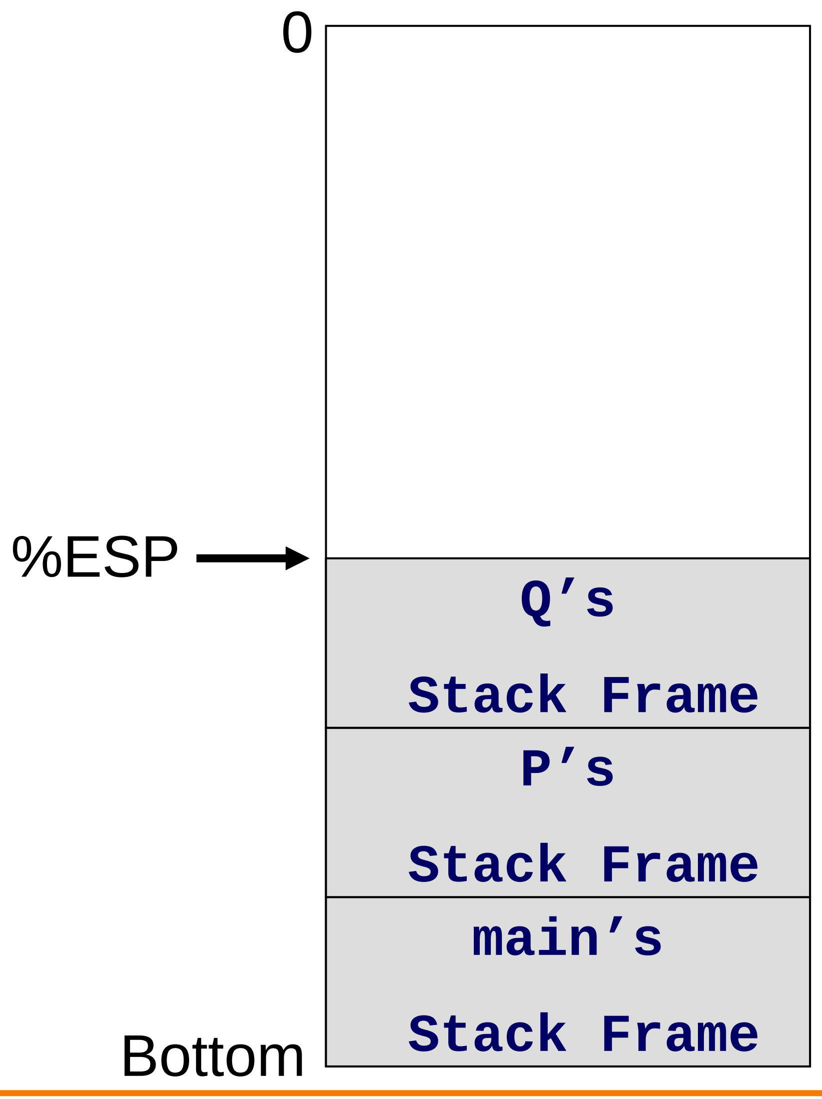


High-Level Picture

main begins executing

main calls P

P calls Q



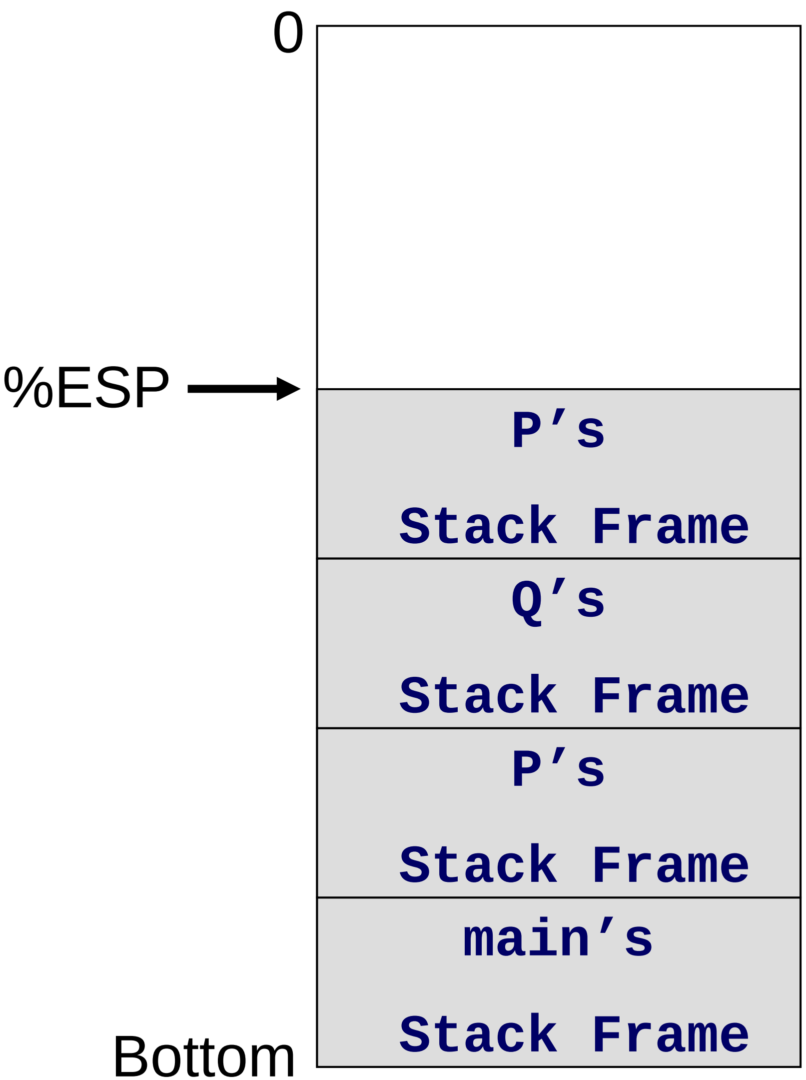
High-Level Picture

`main begins executing`

`main calls P`

`P calls Q`

`Q calls P`



High-Level Picture

```
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main calls P

P calls Q

Q calls P

P returns
```
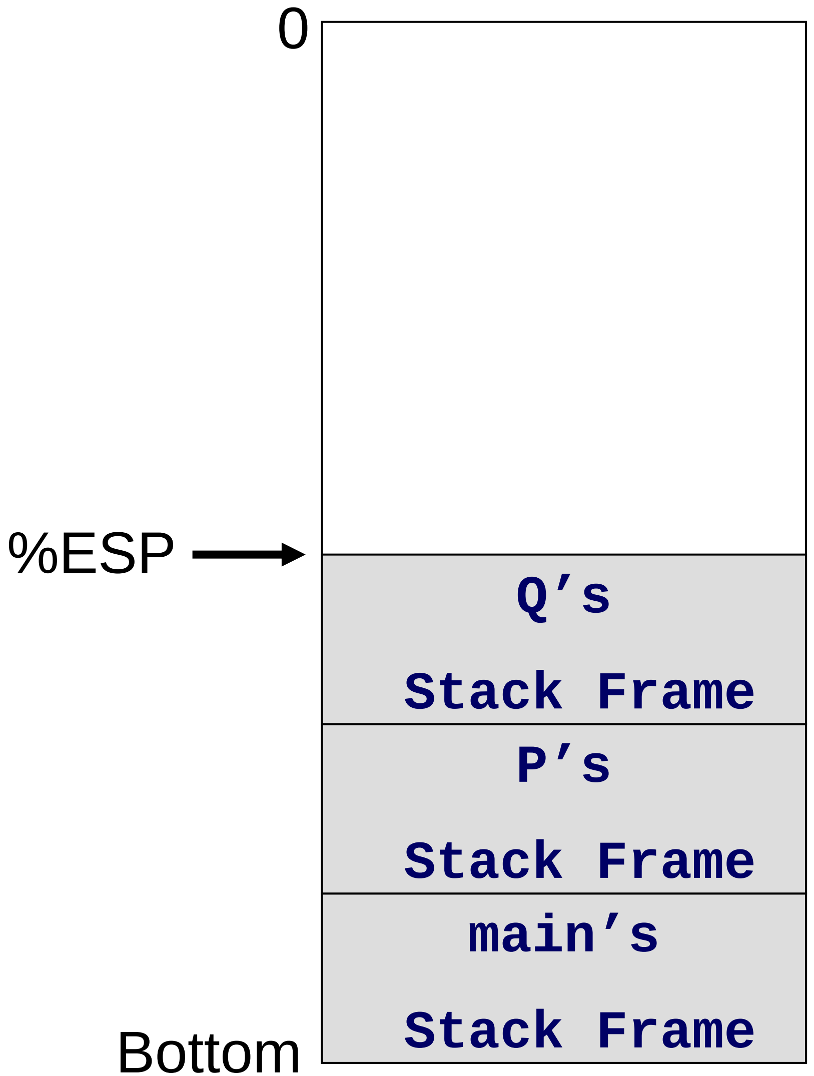



High-Level Picture

`main begins executing`

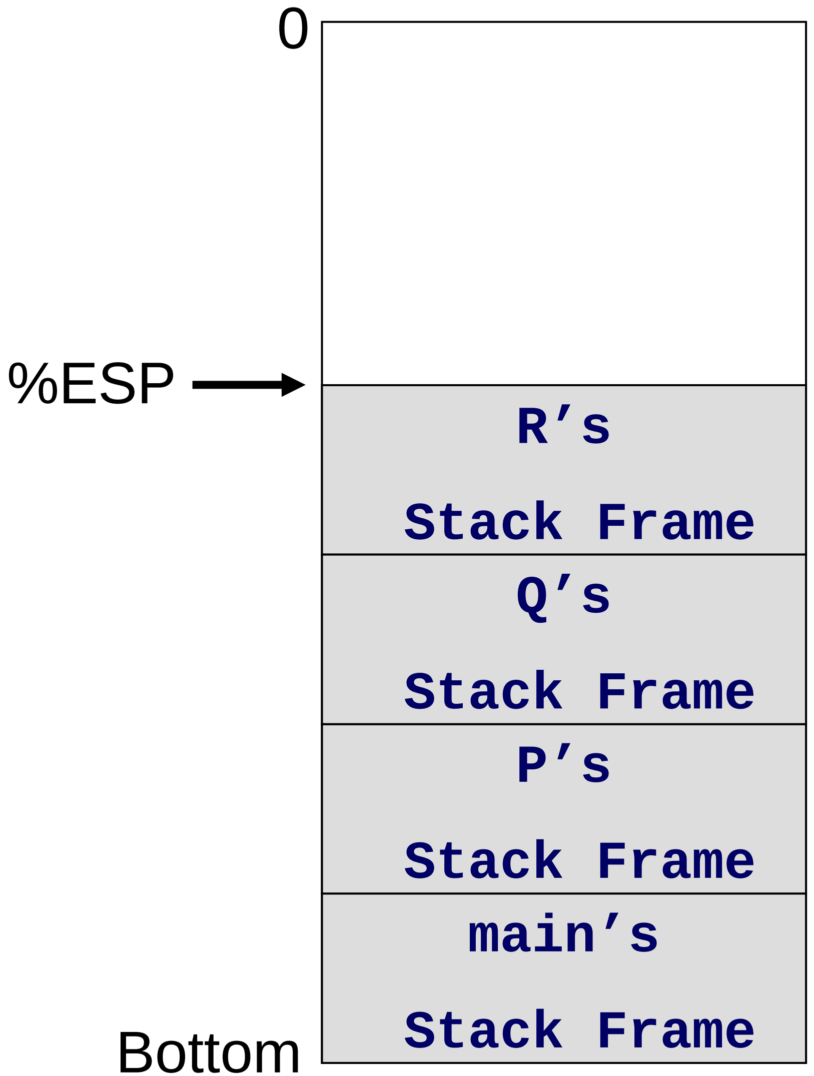
`main calls P`

`P calls Q`

`Q calls P`

`P returns`

`Q calls R`



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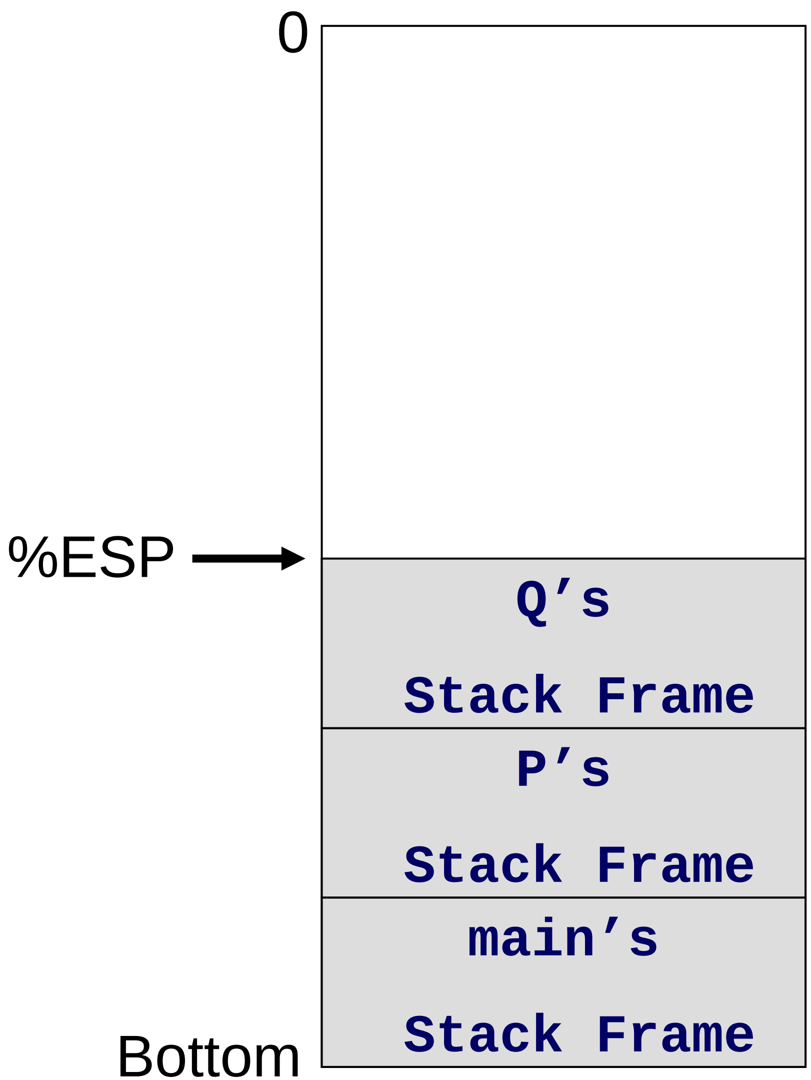

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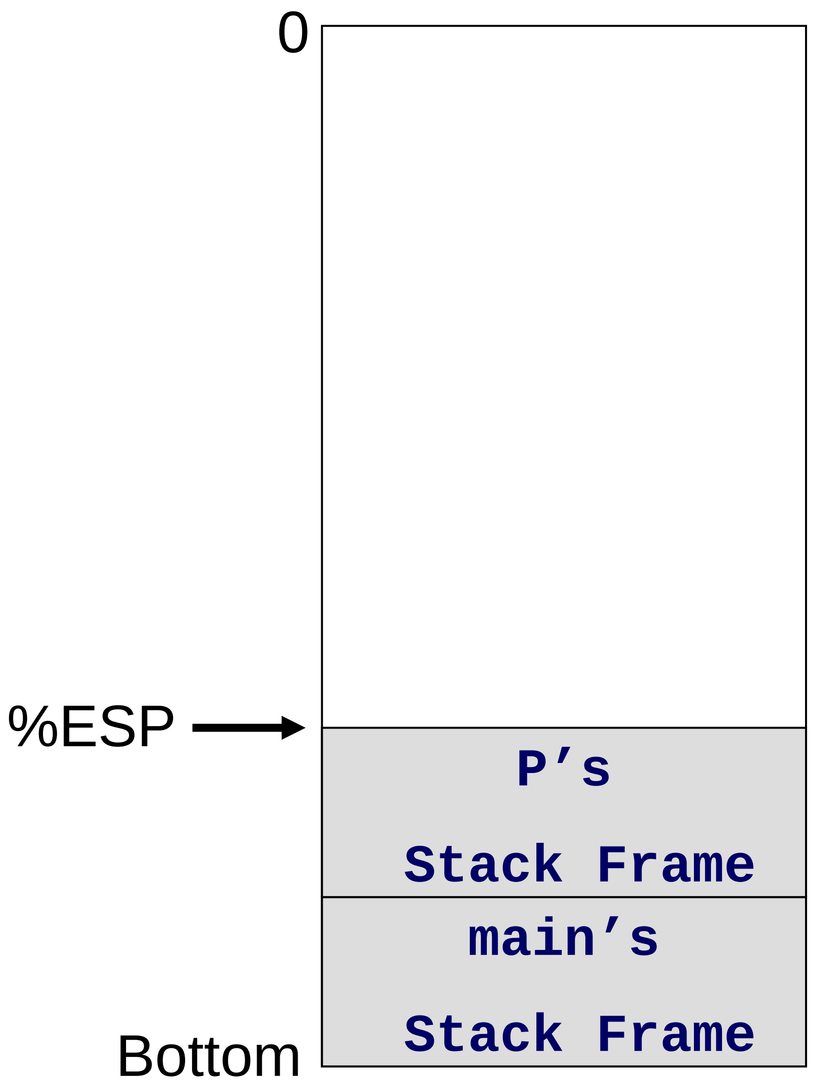

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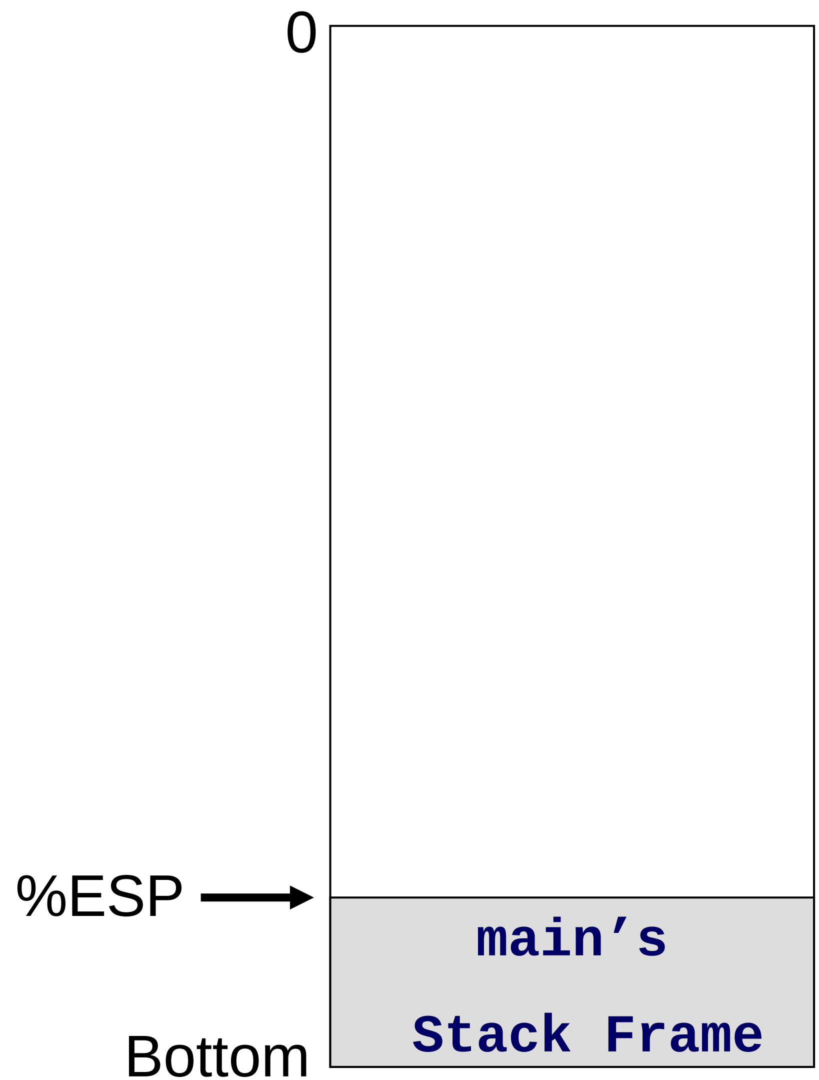

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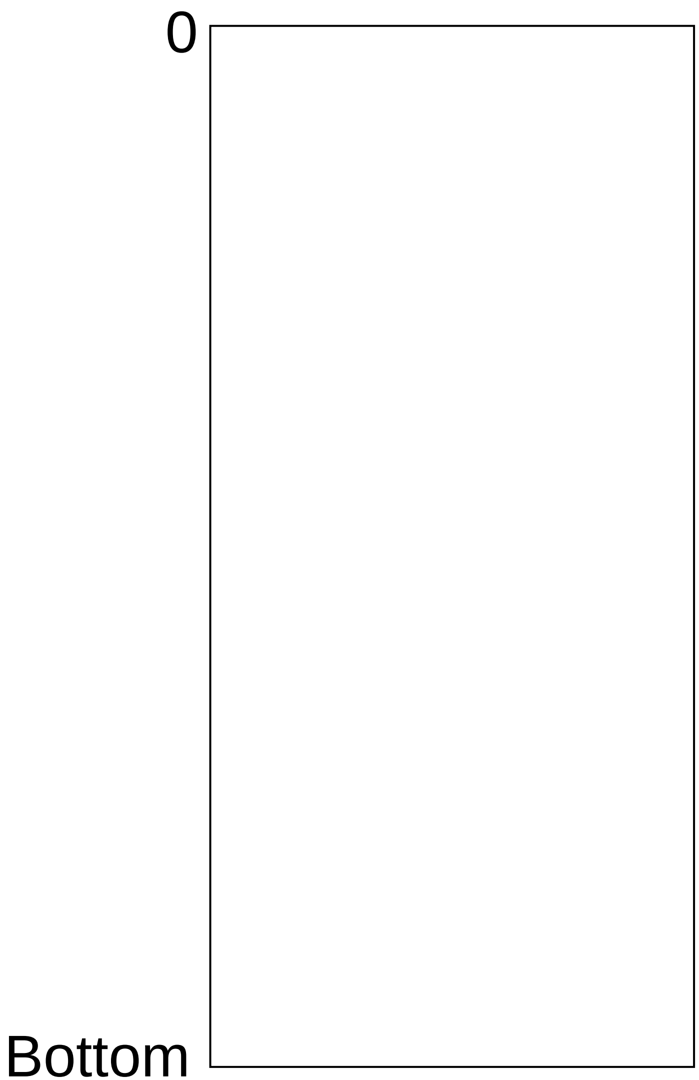
P returns

main returns

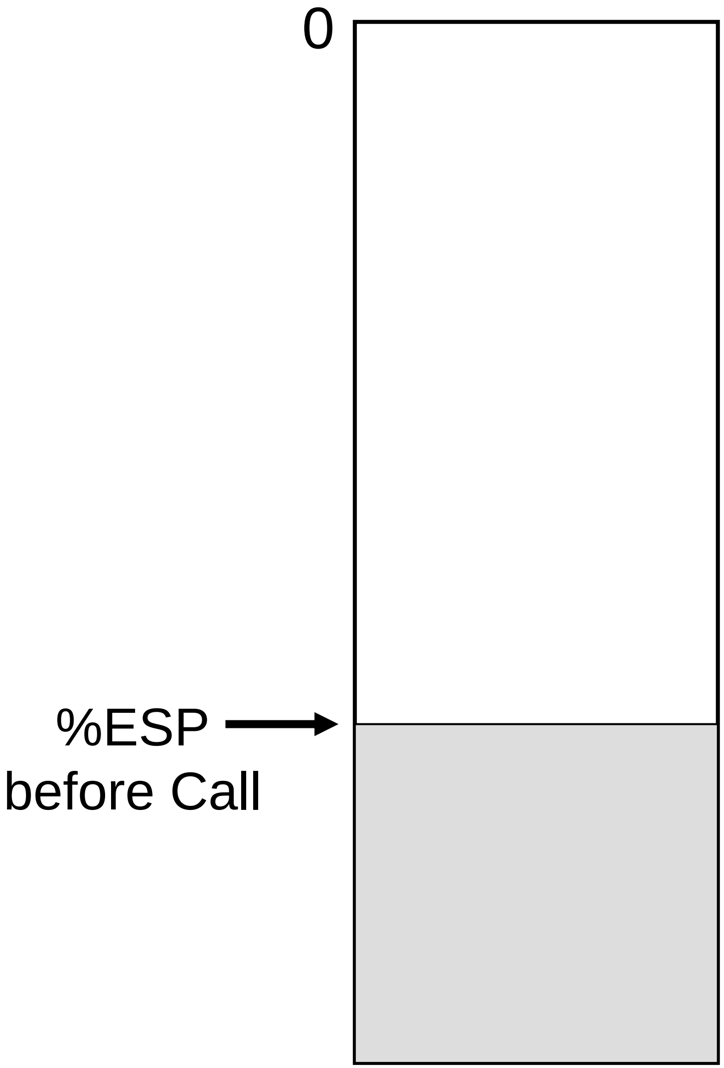


Function Call Details

- Call and Return instructions
- Argument passing between procedures
- Local variables
- Register saving conventions

Call and Return Instructions

Instruction	Function
`pushl src`	`subl $4, %esp` `movl src, (%esp)`
`popl dest`	`movl (%esp), dest` `addl $4, %esp`
`call addr`	`pushl %eip` `jmp addr`
`ret`	`pop %eip`



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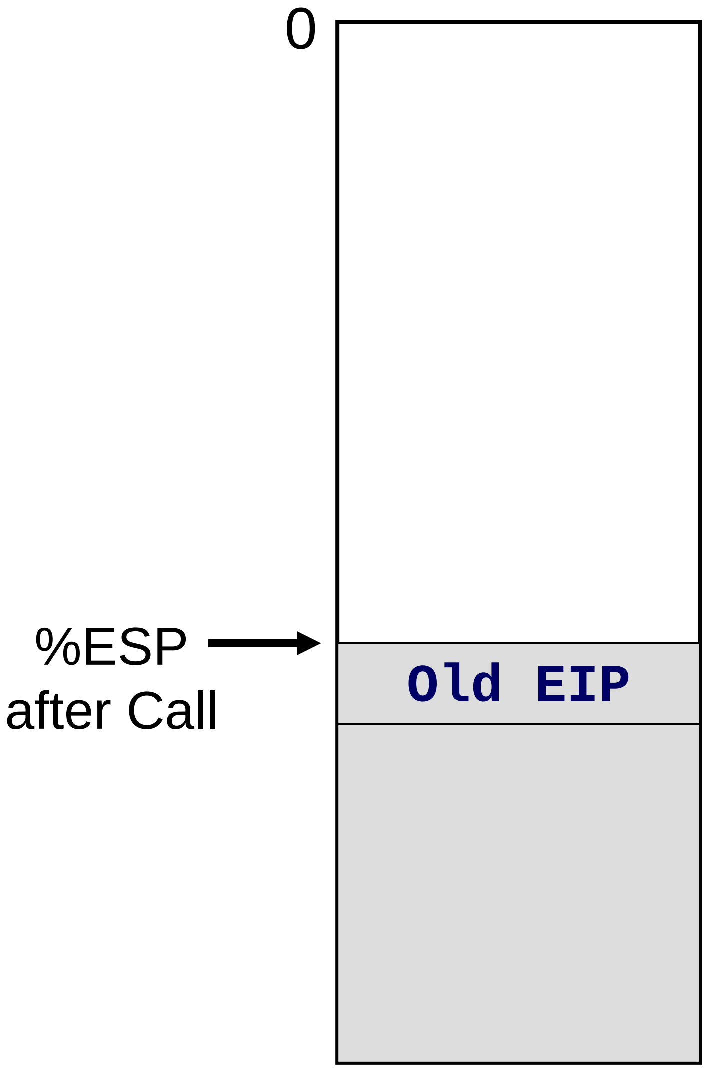

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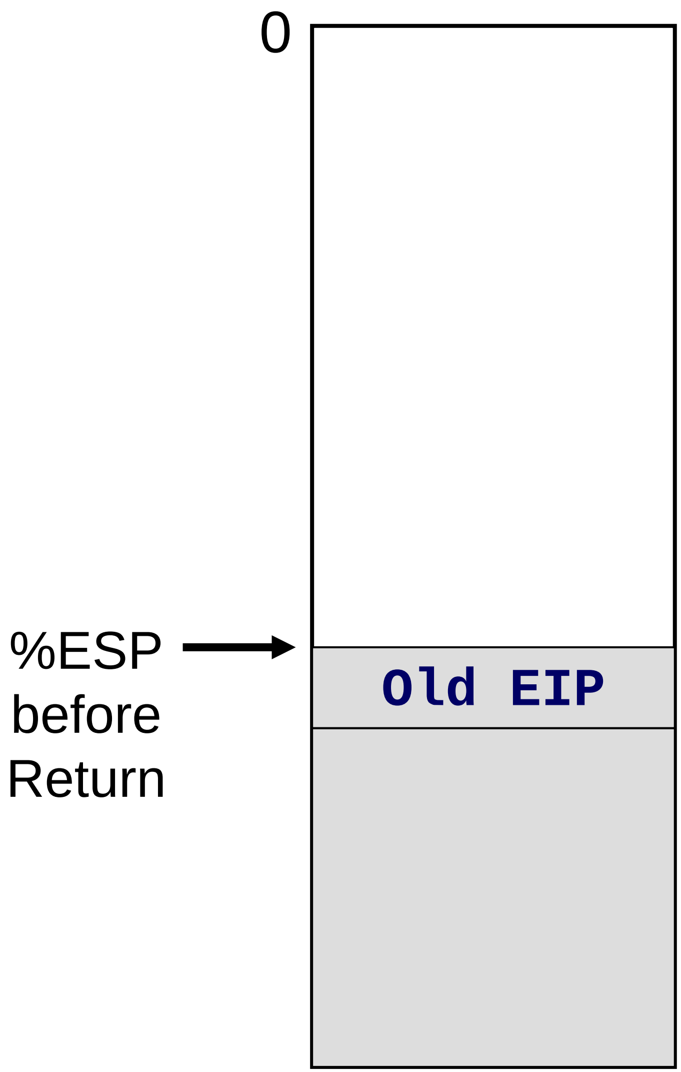
Return instruction assumes that the return address is at the top of the stack



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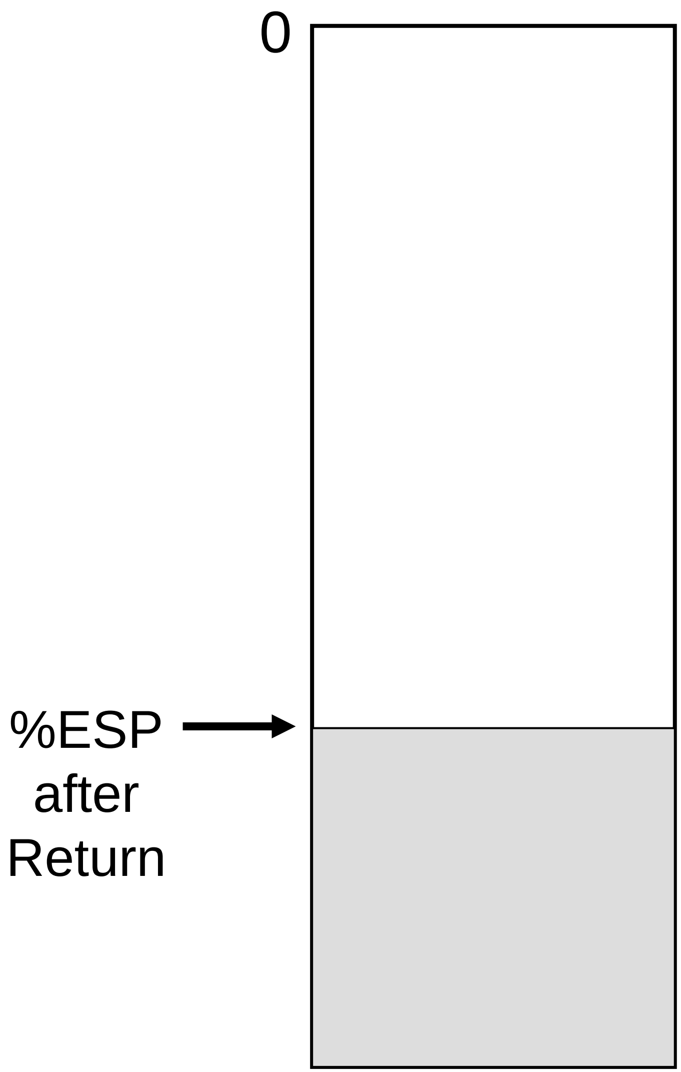
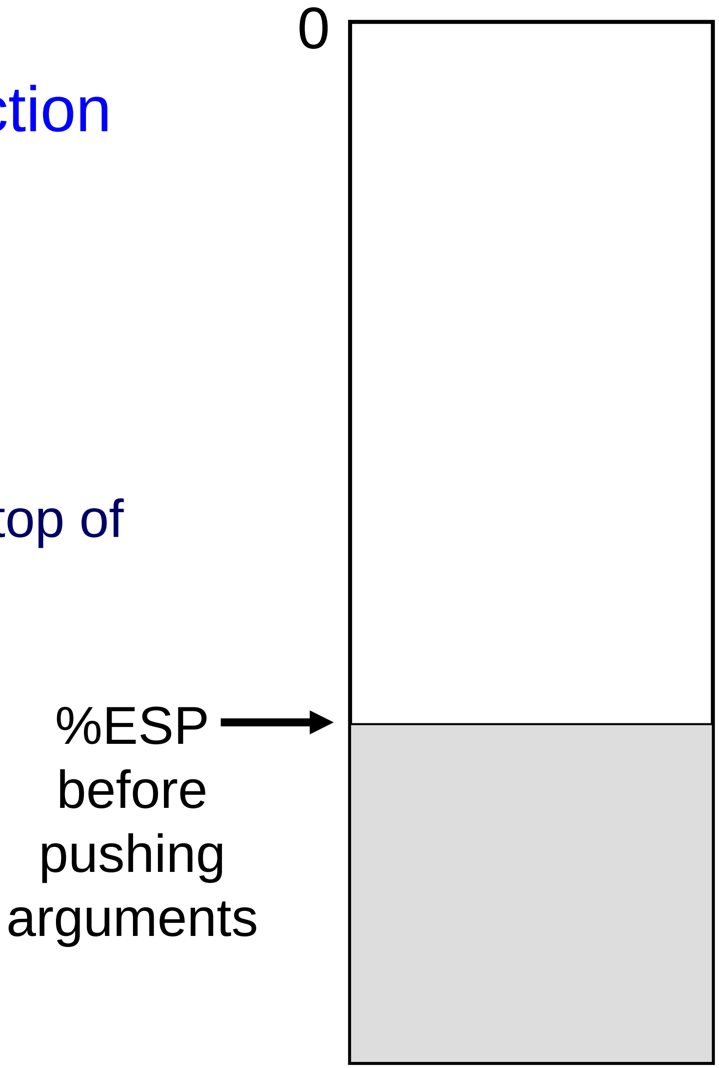


Input Parameters

- Caller pushes input parameters before executing the Call instruction
- Parameters are pushed in the reverse order
  - Push N^th argument first
  - Push 1^st argument last
  - So that the first argument is at the top of the stack at the time of the Call



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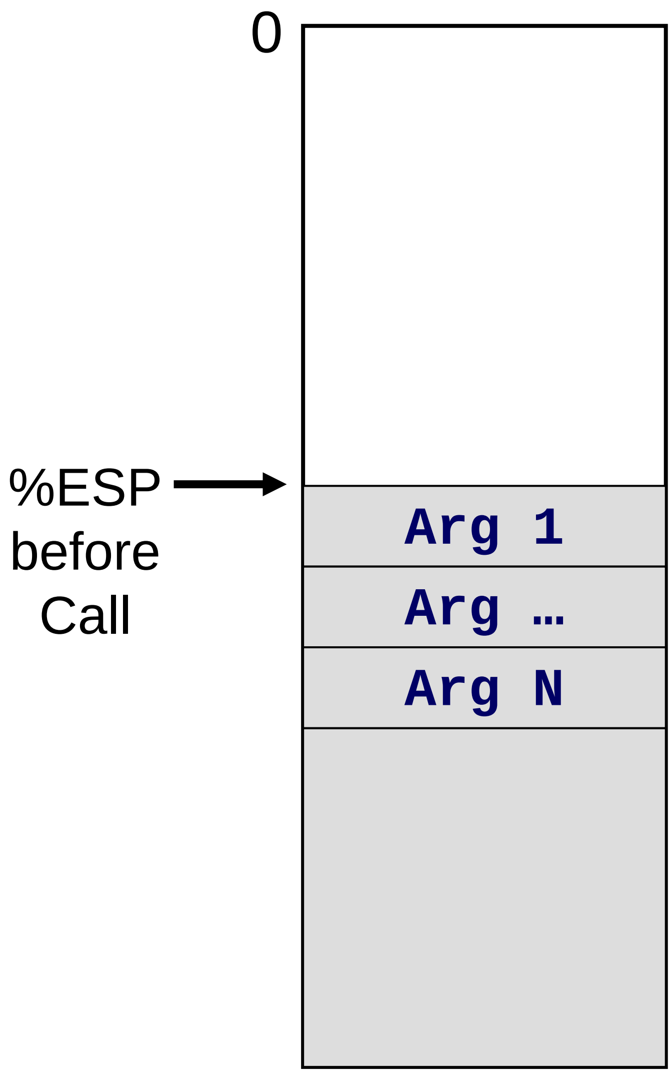

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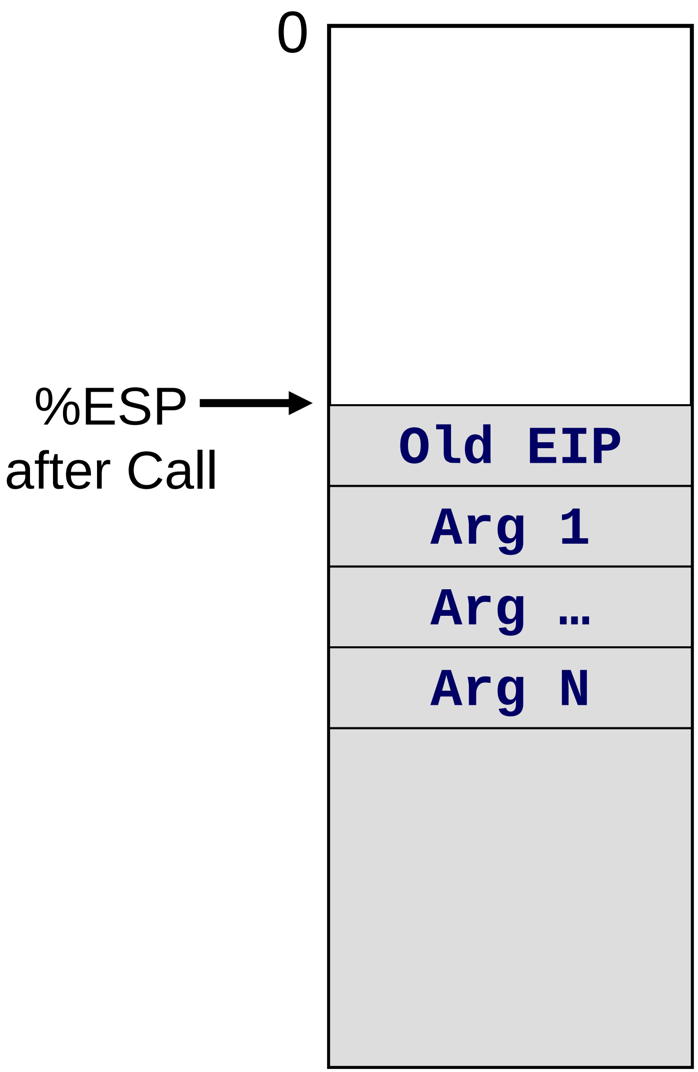
Callee can address arguments relative to ESP: Arg 1 as 4(%esp)



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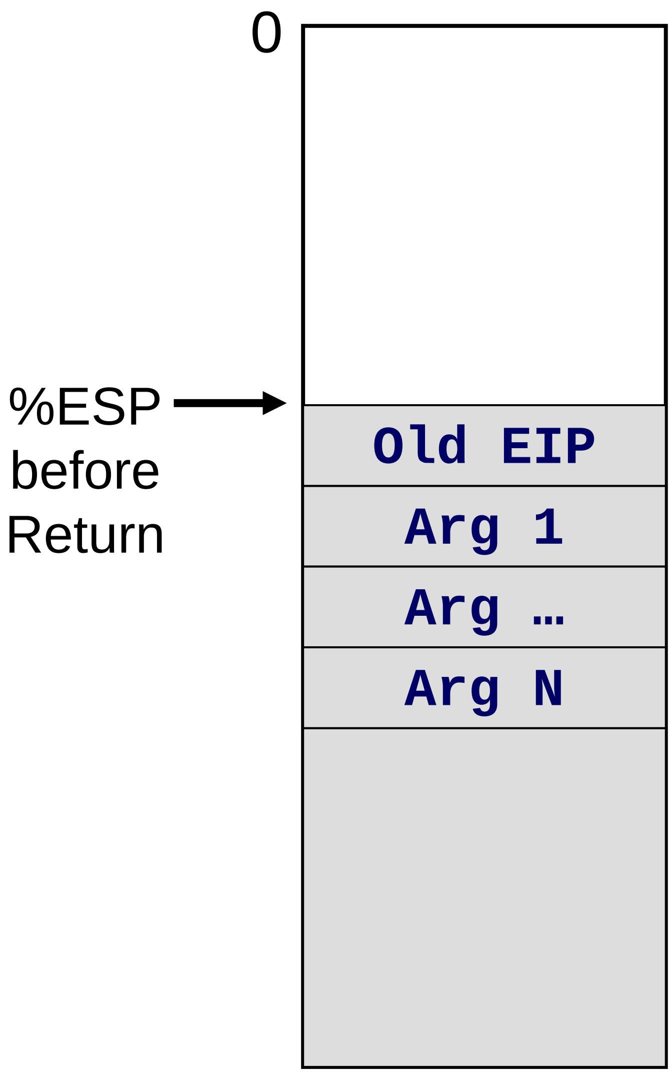

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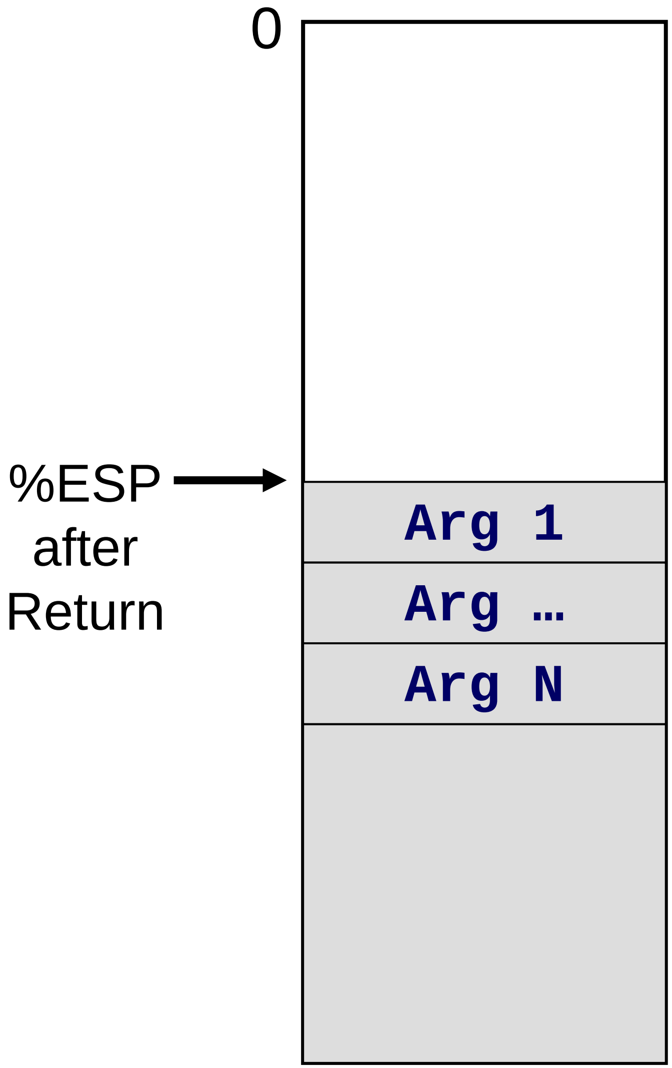
After the function call is finished, the caller pops the pushed arguments from the stack



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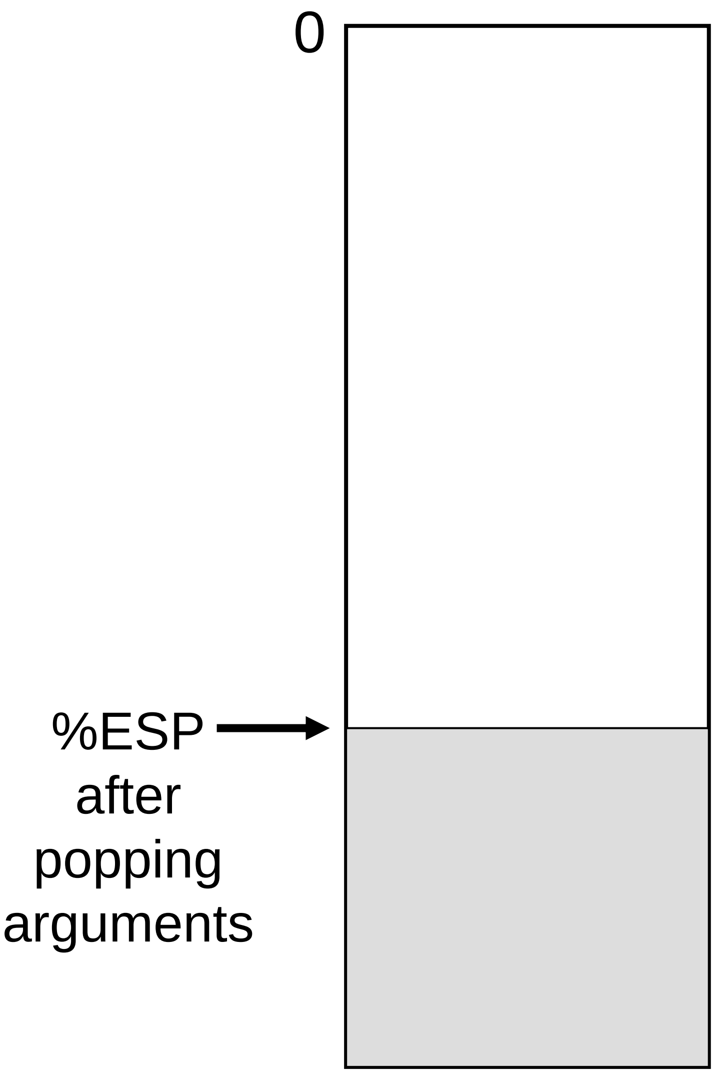
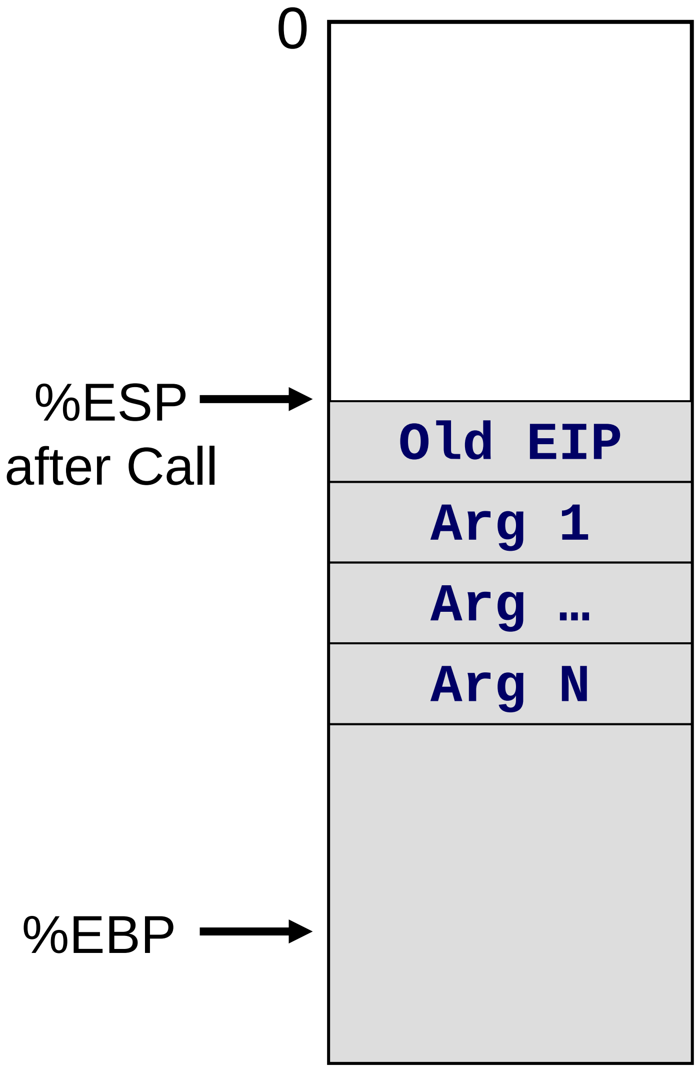


Base Pointer: EBP

- As Callee executes, ESP may change
- Use EBP as a fixed reference point to access arguments and other local variables
- Need to save old value of EBP before using EBP
- Callee begins by executing

```
pushl %ebp
movl %esp, %ebp
```



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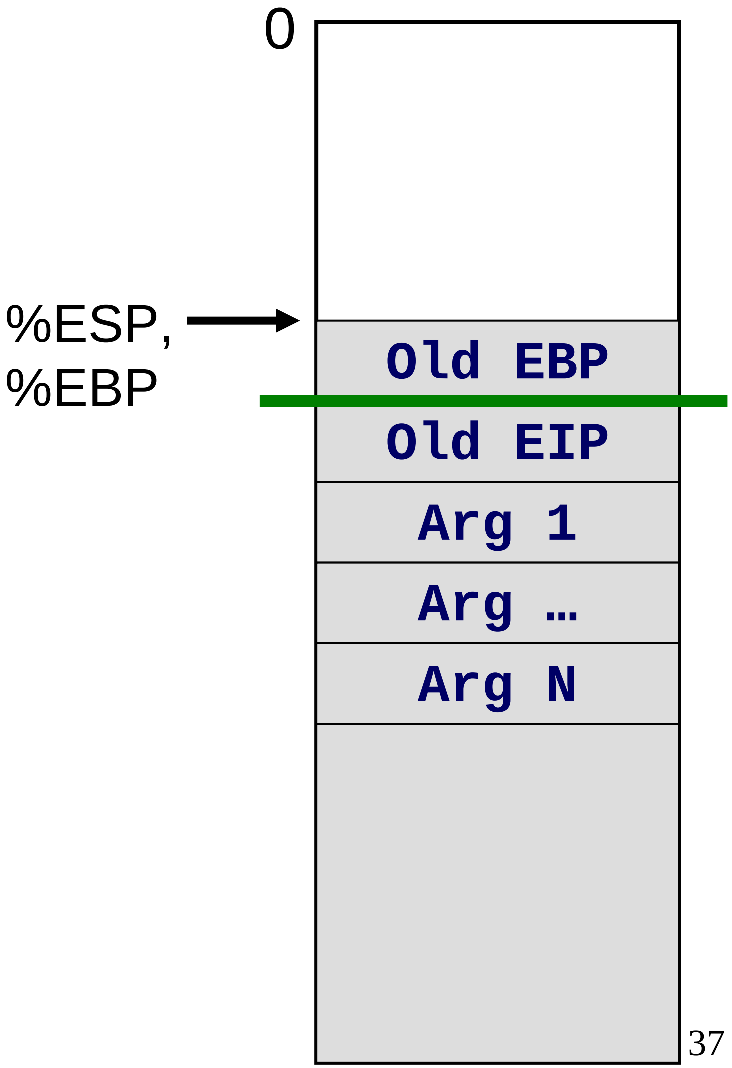
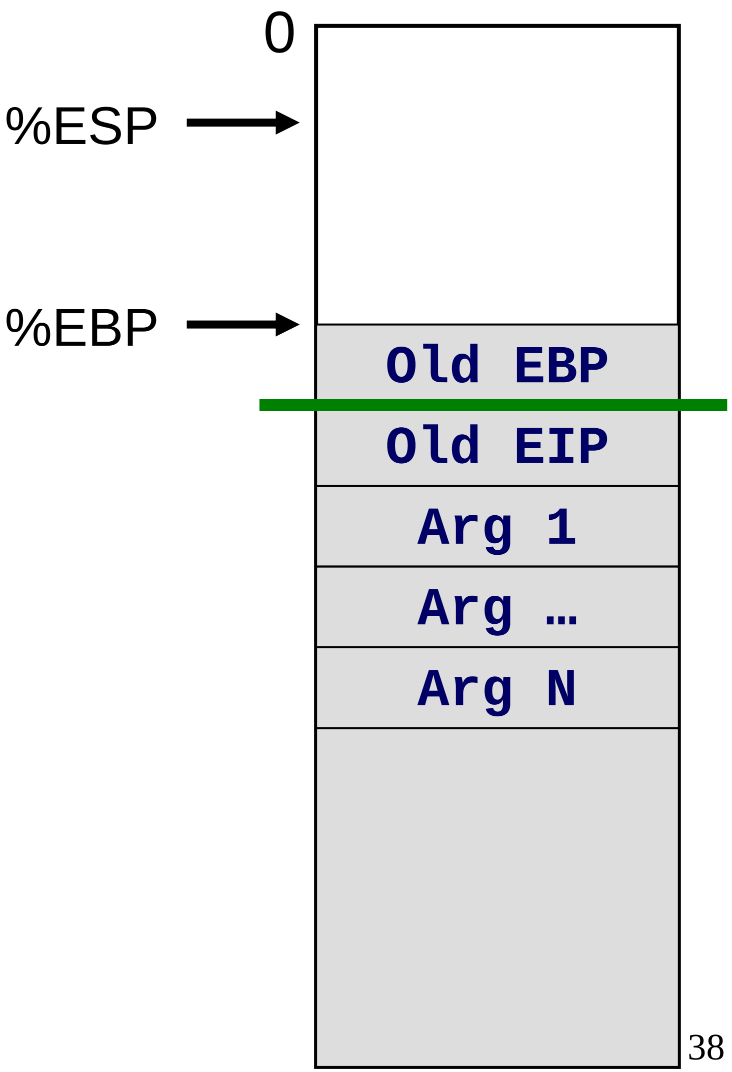
- Regardless of ESP, Callee can address Arg 1 as 8(%ebp)



Base Pointer: EBP

- Before returning, Callee must restore EBP to its old value
- Executes

```
movl %ebp, %esp
popl %ebp
ret
```



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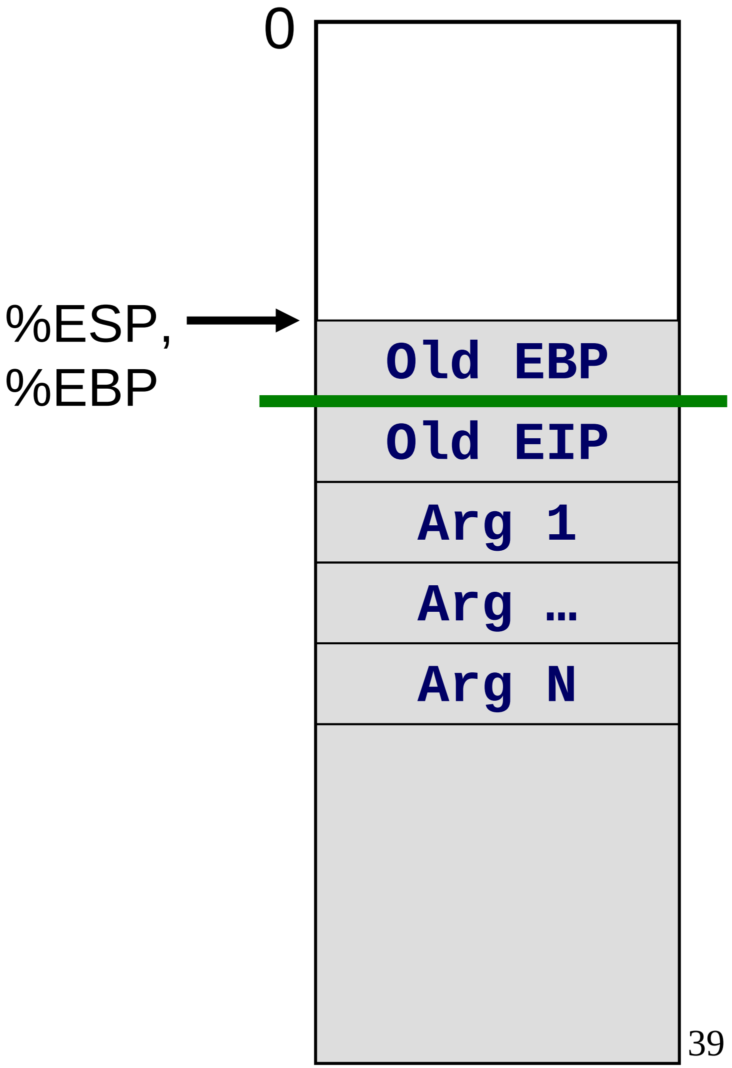

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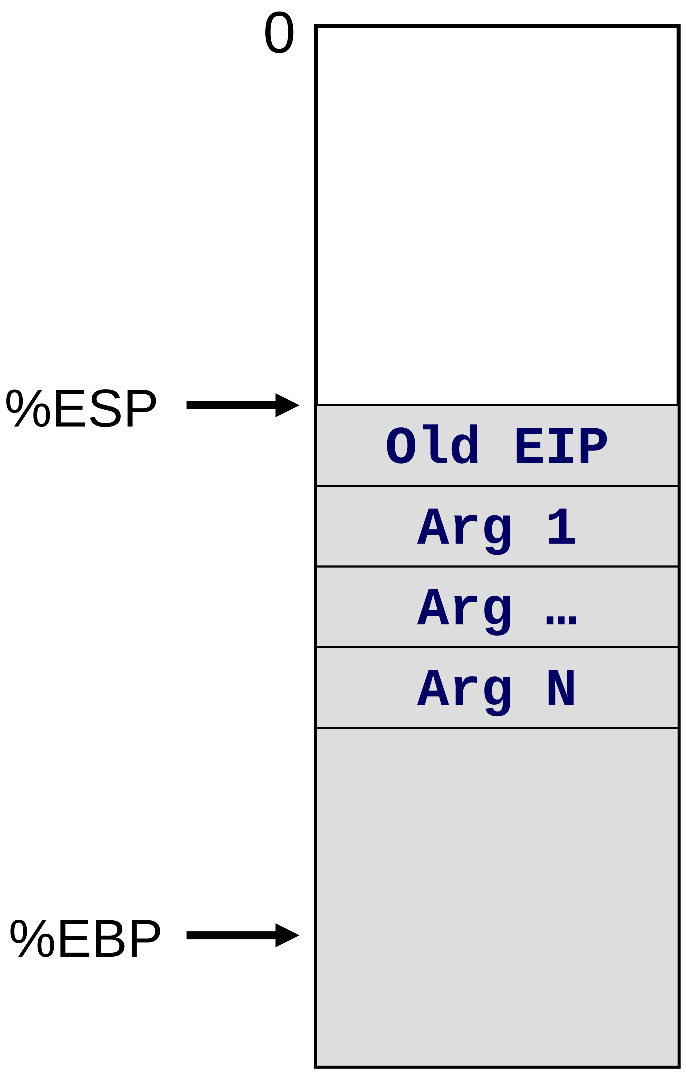

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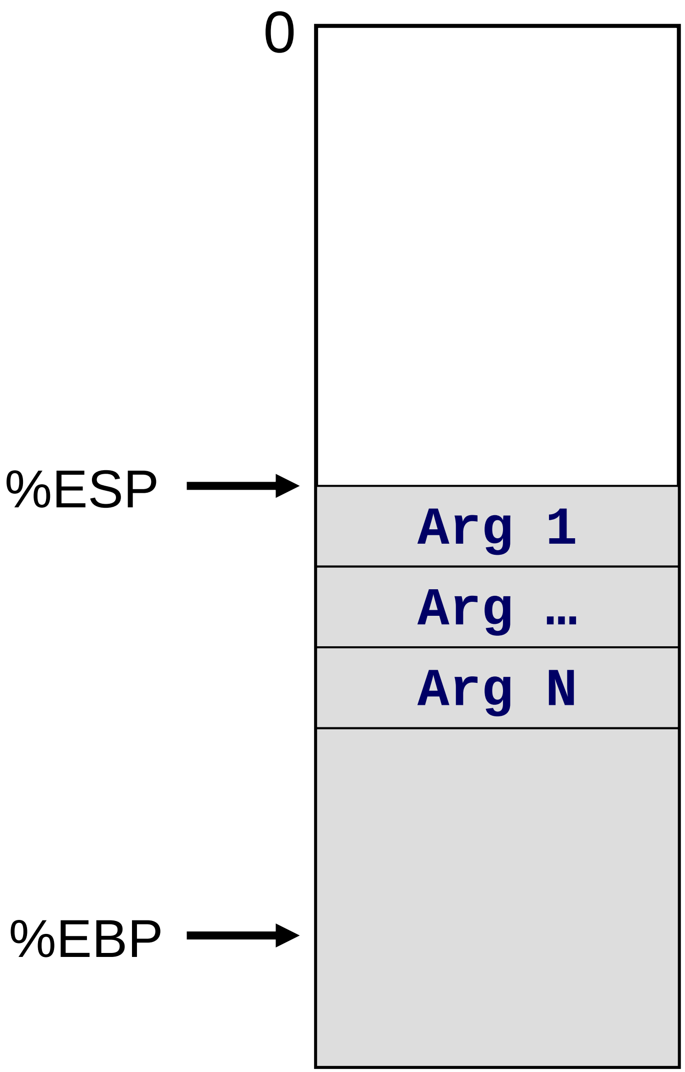
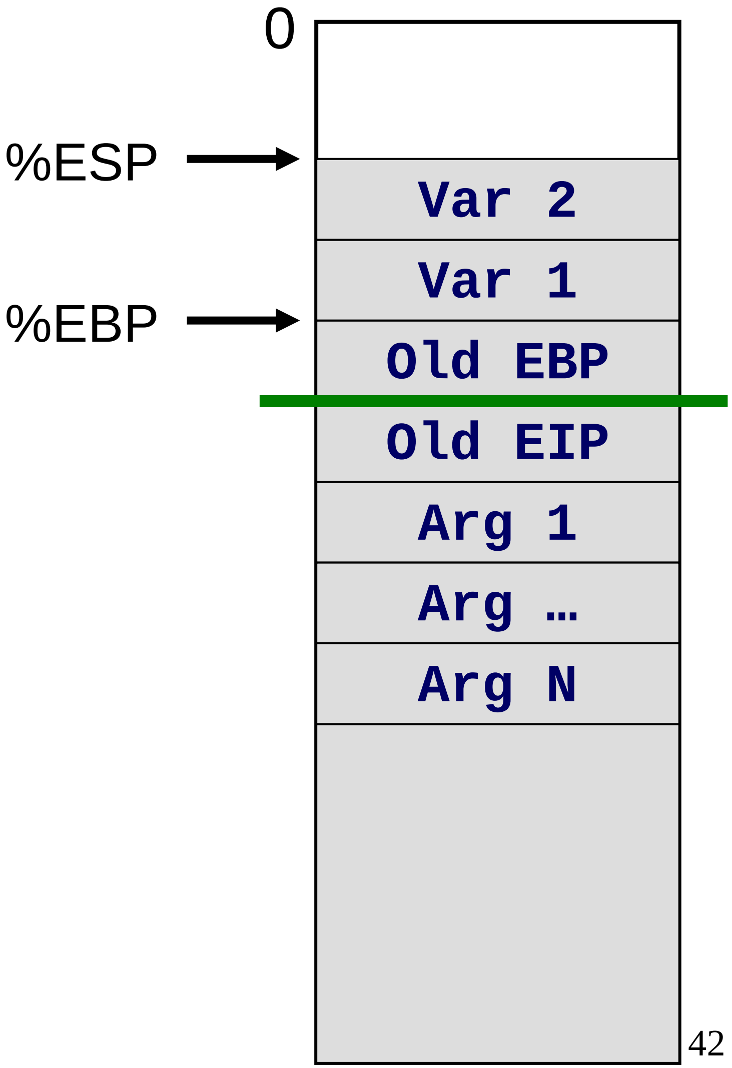


Allocation for Local Variables

- Local variables of the Callee are also allocated on the stack
- Allocation done by moving the stack pointer
- Example: allocate two integers
  - subl $4, %esp
  - subl $4, %esp
  - (or equivalently, subl $8, %esp)
- Reference local variables using the base pointer
  - -4(%ebp)
  - -8(%ebp)



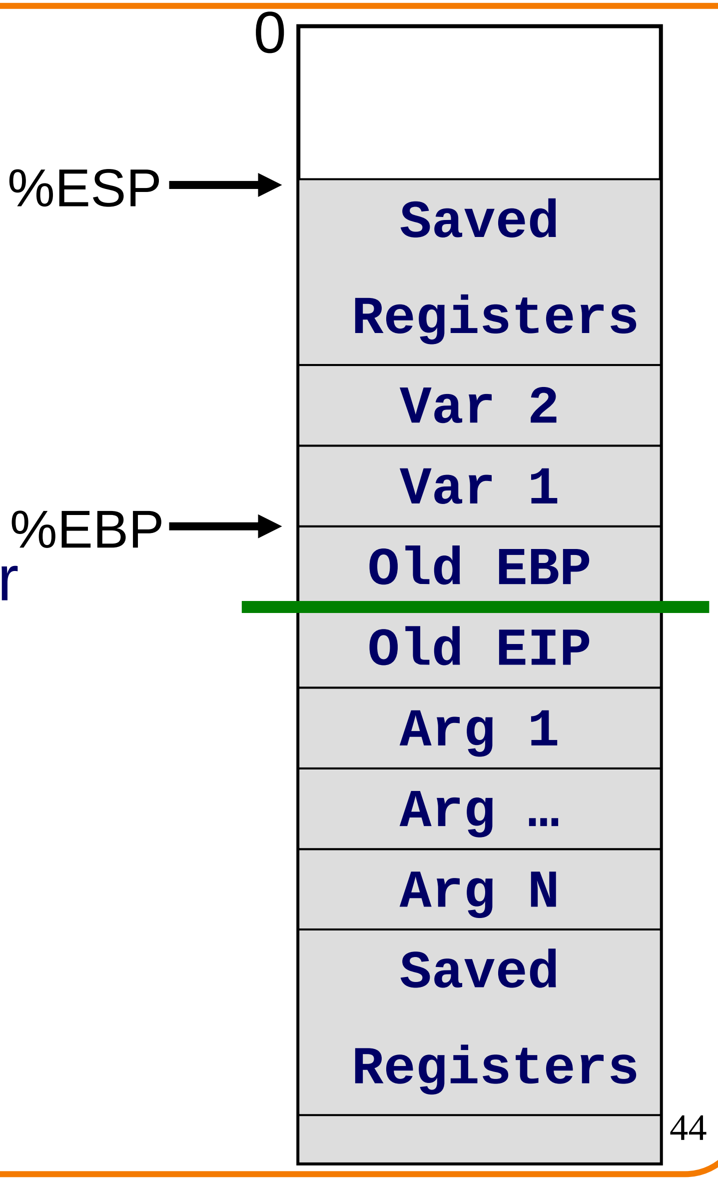
Use of Registers

- Problem: Callee may use a register that the caller is also using
  - When callee returns control to caller, old register contents may be lost
  - Someone must save old register contents and later restore

- Need a convention for who saves and restores which registers

GCC/Linux Convention

- Caller-save registers
  - **`%eax, %edx, %ecx`**
  - Save on stack prior to calling
- Callee-save registers
  - **`%ebx, %esi, %edi`**
  - Old values saved on stack prior to using
- %esp, %ebp handled as described earlier
- Return value is passed from Callee to Caller in %eax



A Simple Example

```
int add3(int a, int b, int c)
{
  int d;

  d = a + b + c;

  return d;
}
```

```
int foo(void)
{
  return add3( 3, 4, 5 );
}
```

A Simple Example

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int add3(int a, int b, int c){
  int d;
  d = a + b + c;
  return d;
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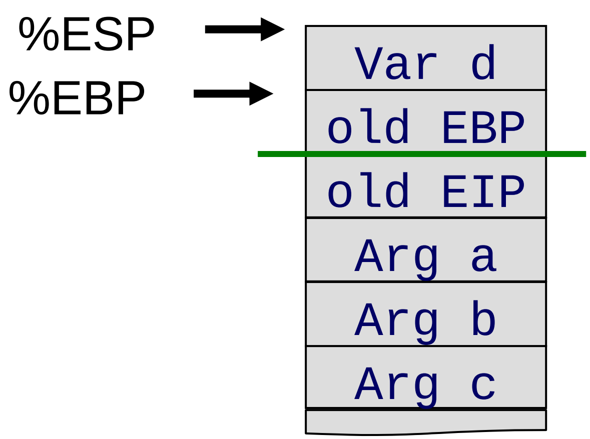
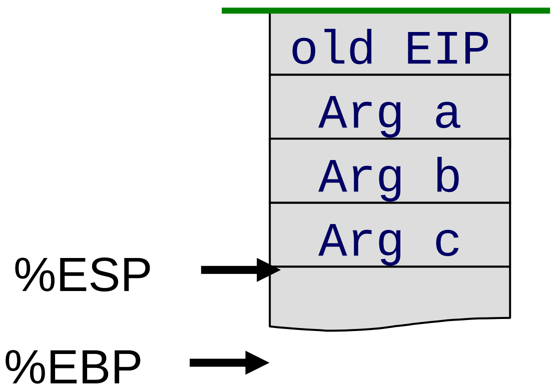



```
add3:
 # Save old ebp and set up new ebp
  pushl %ebp
  movl %esp, %ebp

 # Allocate space for d
  subl $4, $esp

# In general, one may need to push
# callee-save registers onto the stack

 # Add the three arguments
  movl 8(%ebp), %eax
  addl 12(%ebp), %eax
  addl 16(%ebp), %eax

 # Put the sum into d
  movl %eax, -4(%ebp)

 # Return value is already in eax

 # In general, one may need to pop
 # callee-save registers

# Restore old ebp, discard stack frame
  movl %ebp, %esp
  popl %ebp

# Return
  ret
```

A Simple Example

```
int foo(void) {
  return add3( 3, 4, 5 );
}
```




```
foo:
  # Save old ebp, and set-up
  # new ebp
   pushl %ebp
   movl %esp, %ebp

  # No local variables

  # No need to save callee-save
  # registers as we
  # don't use any registers

  # No need to save caller-
  # save registers either

  # Push arguments in reverse order
   pushl $5
   pushl $4
   pushl $3

   call add3

  # Return value is already in eax

  # Restore old ebp and
  # discard stack frame
   movl %ebp, %esp
   popl %ebp

  # Return
   ret
```

Conclusion

- Invoking a function
  - Call: call the function
  - Ret: return from the instruction
- Stack Frame for a function invocation includes
  - Return address,
  - Procedure arguments,
  - Local variables, and
  - Saved registers
- Base pointer EBP
  - Fixed reference point in the Stack Frame
  - Useful for referencing arguments and local variables