COS 425: Database and Information Management Systems

Relational model: Relational algebra

## Modeling access

- Have looked at modeling information as data + structure
- Now: how model access to data in relational model?
- Formal specification of access provides:
  - Unambiguous queries
  - Correctness of results
  - Expressiveness of query languages

# Queries

• A query is a mapping from a set of relations to a relation

Query: relations  $\rightarrow$  relation

- Can derive schema of result from schemas of input relations
- Can deduce constraints on resulting relation that must hold for any input relations
- · Can identify properties of result relation

## Relational query languages

- Two formal relational languages to describe mapping
   Relational algebra
  - Procedural lists operations to form query result
    Relational calculus
    - Declarative describes results of query
- Equivalent expressiveness
- · Each has strong points for usefulness
  - DB system query languages (e.g. SQL) take best of both

## begin with Relational Algebra

#### Basic operations of relational algebra:

- 1. Selection  $\sigma$  :select a subset of tuples from a relation according to a condition
- 2. Projection  $\pi$  :delete unwanted attributes (columns) from tuples of a relation
- 3. cross product X : combine all pairs of tuples of two relations by making tuples with all attributes of both
- 4. Set difference :\* tuples in first relation and not in second
- 5. union U:\* tuples in first relation or second relation
- 6. Renaming p: to deal with name conflicts

\* Set operations: D<sub>1</sub> X D<sub>2</sub> ... X D<sub>k</sub> of two relations must agree

## **Board examples**

#### Formal definition

- A relational expression is
  A relation R in the database
  A constant relation

  - $\begin{array}{l} \mbox{ A constant relation} \\ \mbox{ For any relational expressions } E_1 \mbox{ and } E_2 \\ \cdot \mbox{ } E_1 \mbox{ } E_2 \\ \cdot \mbox{ } \sigma_{p}(E_1) \mbox{ for predicate } P \mbox{ on attributes of } E_1 \\ \cdot \mbox{ } \pi_{g}(E_1) \mbox{ where } S \mbox{ is a subset of attributes of } E_1 \\ \cdot \mbox{ } P(Q(L), \mbox{ } E_1) \mbox{ where } Q \mbox{ is a new relation name and } L \mbox{ is a list of } (\mbox{ old name} \rightarrow \mbox{ new name}) \mbox{ mappings of attributes of } E_1 \end{array}$
- A query in the relational algebra is a relational expression