| cos 425: |
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| Database and Information <br> Management Systems |
| Relational model: |
| Relational algebra |
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## Modeling access

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- Have looked at modeling information as $\qquad$ data + structure
- Now: how model access to data in $\qquad$ relational model?
- Formal specification of access provides:
- Unambiguous queries $\qquad$
- Correctness of results
- Expressiveness of query languages


## Queries

- A query is a mapping from a set of relations to a $\qquad$ relation

$$
\text { Query: relations } \rightarrow \text { relation }
$$

- Can derive schema of result from schemas of $\qquad$ input relations
- Can deduce constraints on resulting relation that $\qquad$ must hold for any input relations
- Can identify properties of result relation $\qquad$
$\qquad$


## Relational query languages

- Two formal relational languages to describe mapping $\qquad$
- Relational algebra
- Procedural - lists operations to form query result $\qquad$ - Relational calculus
- Declarative - describes results of query
- Equivalent expressiveness
- Each has strong points for usefulness
- DB system query languages (e.g. SQL) take best of both
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## begin with Relational Algebra

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Basic operations of relational algebra:

1. Selection $\sigma$ :select a subset of tuples from a relation according to a condition
2. Projection $\pi$ :delete unwanted attributes (columns) from tuples of a relation
3. cross product $X$ : combine all pairs of tuples of two relations by making tuples with all attributes of both
4. Set difference - :* tuples in first relation and not in second
5. union U:* tuples in first relation or second relation
6. Renaming $\rho$ : to deal with name conflicts

* Set operations: $D_{1} \times D_{2} \ldots \times D_{k}$ of two relations must agree

Board examples $\qquad$
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## Formal definition

- A relational expression is
- A relation $R$ in the database
- A constant relation
- For any relational expressions $E_{1}$ and $E_{2}$
- $E_{1} \cup E_{2}$
- $E_{1}-E_{2}$
- $\mathrm{E}_{1} \times \mathrm{E}_{2}$
- $\sigma_{P}\left(E_{1}\right)$ for predicate $P$ on attributes of $E_{1}$
- $\pi_{S}\left(E_{1}\right)$ where $S$ is a subset of attributes of $E_{1}$
- $P\left(Q(L), E_{1}\right)$ where $Q$ is a new relation name and $L$ is a list of (old name $\rightarrow$ new name) mappings of attributes of $\mathrm{E}_{1}$
- A query in the relational algebra is a relational expression

