COS 425: **Database and Information** Management Systems

Relational model continued

Understanding how to use the relational model

Summary of board example

Relational model:

books: (title, ISBN#, edition, date) PU branches: (br_name, librarian, hours)

copy: (ISBN#, copy#, condition, purchase date, br_name)

br_name not null

isbn# is a foreign key referencing books

br_name is a foreign key referencing PU branches

because ER key

ER relationship *PU book* folded into *copy* (relationship between books and weak entity copy) constraints ER relationship PU holding folded into copy

Board example: Total participation constraints?

PU branches: (br name, librarian, hours)

copy: (ISBN#, copy#, condition, purchase date, br_name)

isbn# is a foreign key referencing books

br_name is a foreign key referencing *PU branches*

"br_name not null" - constraint captures total participation of copy in PU holding Because PU holding represented within copy

total participation of PU branches in PU holding not representable in pure relational definition

Basic Paradigm Attribute R_A₁ ... Attribute R_A_m Entity 1 • Each entity becomes a relation • Relationship becomes R: { (list of attributes forming key of Entity 1 (denote L₁), list of attributes forming key of Entity 2 (denote L₂), ... list of attributes forming key of Entity k (denote L_k), Attribute R_A₁, ..., Attribute R_A_m) L₁ is a foreign key referencing Entity 1, ... L_k is a foreign key referencing Entity k }

What about constraints on relationships?

- Key constraint:
 - Simplifies key of corresponding relation
 - Allows folding of relation into key entity
- Total participation constraint:
 - In general, cannot represent in purely relational definition:
 - Domain specification
 - Keys of relations
 - Foreign keys
 - "not null"s

Constraints have in relational definition

Enforcing relational constraints

- Constraints must be satisfied at all times
- What happens when tuples in relations change?
- Action of changing a relation not part of basic relational model
- Database language implementing model enforces

Enforcement in SQL

SQL commands changing relations: INSERT, DELETE, UPDATE

- Domain constraints
 - Don't allow attribute value not in domain INSERT or UPDATE fails
- "Not null" constraints
 - Special case of domain constraints

Enforcement in SQL

- Key constraints
 - Can have other candidate keys declared as well as primary key
 - Don't allow 2nd tuple with same key value INSERT or UPDATE fails
 - Implicit "not null" for attributes in a key
 INSERT or UPDATE fails

Enforcement in SQL

• Foreign key constraints

Suppose Y denotes a set of attributes of relation B that reference the primary key of relation A.

 Don't allow tuple into B if no tuple in A with matching values for Y INSERT or UPDATE fails

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Enforcement in SQL

Foreign key constraints continued

- suppose want to remove a tuple in A
- Suppose there is a tuple in B with matching values for Y

Choices (in SQL):

Disallow deletion from A
 DELETE or UPDATE fails

Enforcement in SQL

Choices (in SQL) continued:

- 2. Ripple effect (CASCADE):
 - Remove tuple from A and all tuples from B with matching values for Y
 - DELETE or UPDATE in A causes DELETE in B
- 3. Substitute value
 - Put "null" (if Y not part of candidate key for B) or other default value for Y in B
 - DELETE or UPDATE in A causes UPDATE in B

Actions for board example?

books: (title, ISBN#, edition, date)

PU branches: (br name, librarian, hours)

copy: (ISBN#, copy#, condition, purchase date, br_name) br_name not null

isbn# is a foreign key referencing books

br_name is a foreign key referencing *PU branches*

What about constraints not expressible in ER model?

- Value based constraints?
- General functional constraints?

In relational model:

- Declaring and enforcing these depend on use of database language
- Use query semantics to check