



Local Search

**Introduction to
Artificial Intelligence**

COS302

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Administration

Questions?

Reading: Boyan thesis for more information on local search.

[http://www-](http://www-2.cs.cmu.edu/afs/cs.cmu.edu/user/jab/mosaic/thesis/)

[2.cs.cmu.edu/afs/cs.cmu.edu/user/jab/mosaic/thesis/](http://www-2.cs.cmu.edu/afs/cs.cmu.edu/user/jab/mosaic/thesis/)

chap3.ps.gz, appen2.ps.gz



Search

- Recall the search algorithms for constraint problems (CSP, SAT).**
- **Systematically check all states.**
 - **Look for satisfied state.**
 - **In finite spaces, can determine if no state satisfies constraints (complete).**



Local Search

Alternative approach:

- **Sometimes more interested in positive answer than negative.**
- **Once we find it, we're done.**
- **Guided random search approach.**
- **Can't determine unsatisfiability.**



Optimization Problems

Framework: States, neighbors (symmetric). Add notion of *score* for states: $V(s)$

Goal: Find state w/ highest (sometimes lowest) score.

Generalization of constraint framework because...



SAT as Optimization

States are complete assignments.

Score is number of satisfied clauses.

Know the maximum (if formula is truly satisfiable).



Local Search Idea

**In contrast to CSP/SAT search,
work with complete states.**

**Move from state to state
remembering very little (in
general).**

Hill Climbing

Most basic approach

- Pick a starting state, s .
- If $V(s) > \max_{s' \in N(s)} V(s')$, stop
- Else, let $s = \text{any } s' \text{ s.t.}$

$$V(s') = \max_{s'' \in N(s')} V(s'')$$

- Repeat

Like DFS. Stops at local maximum.

N-queens Optimization

**Minimize violated constraints.
Neighbors: Slide in column.**

		1				
						0
			0			
	1					
				0		
						2
0						
					2	



N-queens Search Space

How far, at most, from minimum configuration?

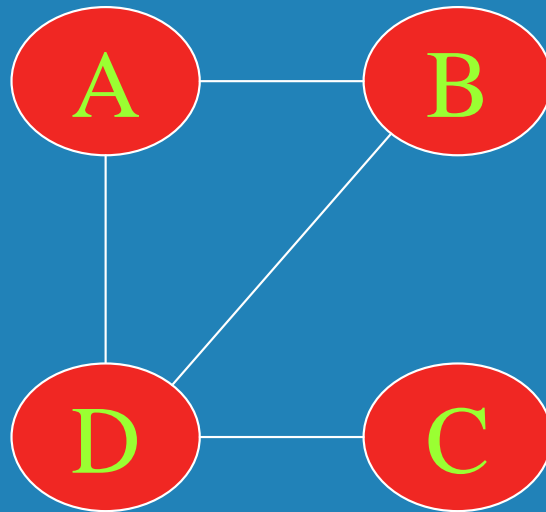
Never more than n slides from any configuration.

How many steps, at most, before local minimum reached?

Maximum score is $n(n-1)$.

Graph Partitioning

Separate nodes into 2 groups to minimize edges between them





Objective Function

Penalty for imbalance

$$V(x) = \text{cutsize}(x) + (L-R)^2$$

State space? Neighbors?

State Space

0000: 16+0

0001: 4+3

0010: 4+1

0011: 0+2

0100: 4+2

0101: 0+3

0110: 0+3

0111: 4+2

1000: 4+2

1001: 0+3

1010: 0+3

1011: 4+2

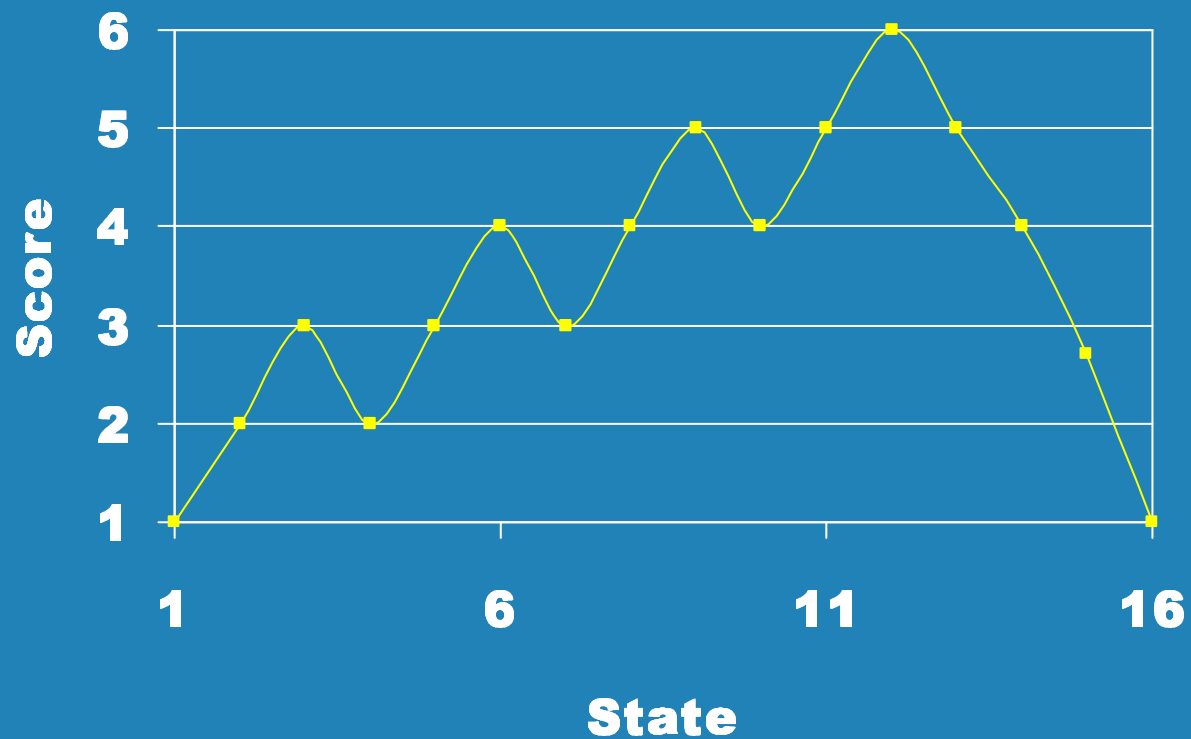
1100: 0+2

1101: 4+1

1110: 4+3

1111: 16+0

Bumpy Landscape



Simulated Annealing

Jitter out of local minima.

Let T be non-negative.

Loop:

- **Pick random neighbor s' in $N(s)$**
- **Let $D = V(s) - V(s')$ (improve)**
- **If $D > 0$, (better), $s = s'$**
- **Else with prob $e^{-D/T}$, $s = s'$**



Temperature

As T goes to zero, hill climbing

As T goes to infinity, random walk

In general, start with large T and decrease it slowly.

If decrease *infinitely* slowly, will find optimal solution.

Analogy is with metal cooling.



Aside: Analogies

Natural phenomena inspire algs

- **Metal cooling**
- **Evolution**
- **Thermodynamics**
- **Societal markets**
- **Ant colonies**
- **Immune systems**
- **Neural networks, ...**

Annealing Schedule

After each step, update T

- **Logarithmic: $T_i = \gamma / \log(i+2)$**
provable properties, slow
- **Geometric: $T_i = T_0 \gamma^{\text{int}(i/L)}$**
lots of parameters to tune
- **Adaptive**
change T to hit target accept rate

Lam Schedule

Ad hoc, but impressive (see Boyan, pg. 190, 191). Parameter is run time.

- **Start target accept rate at 100%**
- **Decrease exponentially to 44% at 15% through run**
- **Continue exponential decline to 0% after 65% through run.**



Locality Assumption

Why does local search work?

Why not jump around at random?

**We believe scores change slowly
from neighbor to neighbor.**

**Perhaps keep a bunch of states
and search around the good
ones...**

Genetic Algorithms

Search : Population genetics ::

State : individual ::

State set : population ::

Score : fitness ::

Neighbors : offspring / mutation

? : sexual reproduction, crossover



Algorithmic Options

Fitness function

Representation of individuals

Who reproduces? *

How alter individuals (mutation)?

**How combine individuals
(crossover)? ***

Representation

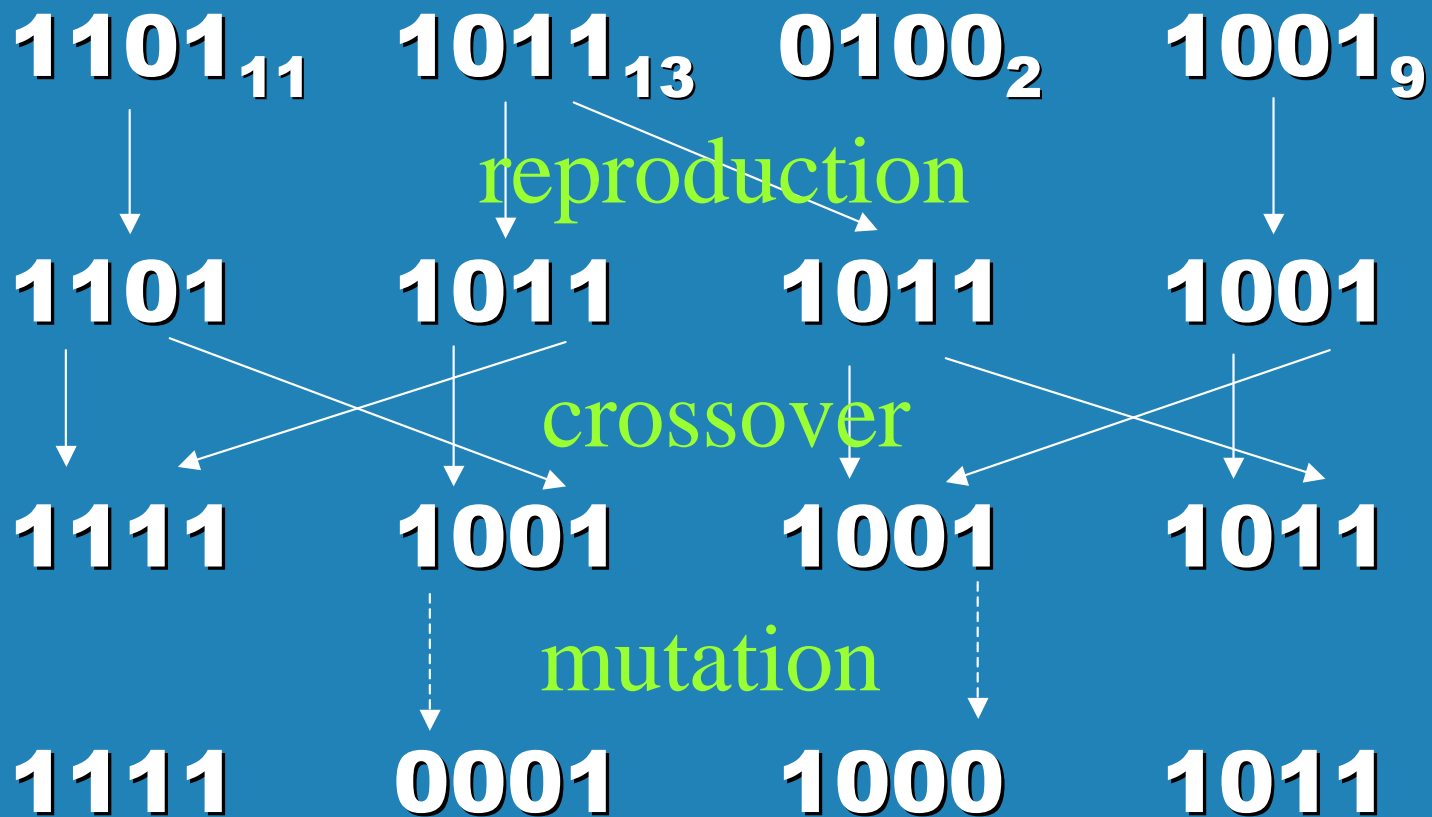
“Classic” approach: states are fixed-length bit sequences

Can be more complex objects: genetic programs, for example.

Crossover:

- **One point, two point, uniform**

GA Example



Generation to Next

G is a set of N (even) states

Let $p(s) = V(s)/\sum_{s'} V(s')$, $G' = \{ \}$

For $k = 1$ to $N/2$

choose x, y with prob. $p(x), p(y)$

randomly swap bits to get x', y'

rarely, randomly flip bits in x', y'

$G' = G' \cup \{x', y'\}$

GSAT

Score: number of unsat clauses.

- **Start with random assignment.**
- **While not satisfied...**
 - **Flip variable value to satisfy greatest number of clauses.**
 - **Repeat (until done or bored)**
- **Repeat (until done or bored)**



GSAT Analysis

**What kind of local search is this most like? Hill climbing?
Simulated annealing? GA?
Known to perform quite well for coloring problems.**

WALKSAT

On each step, with probability p , do a greedy GSAT move.

With probability $1-p$, “walk” (flip a variable in an unsatisfied clause).

Even faster for many problems (planning, random CNF).



WALKSAT Analysis

**What kind of local search is this most like? Hill climbing?
Simulated annealing? GA?**

Known to perform quite well for hard random CNF problems.

Random k-CNF

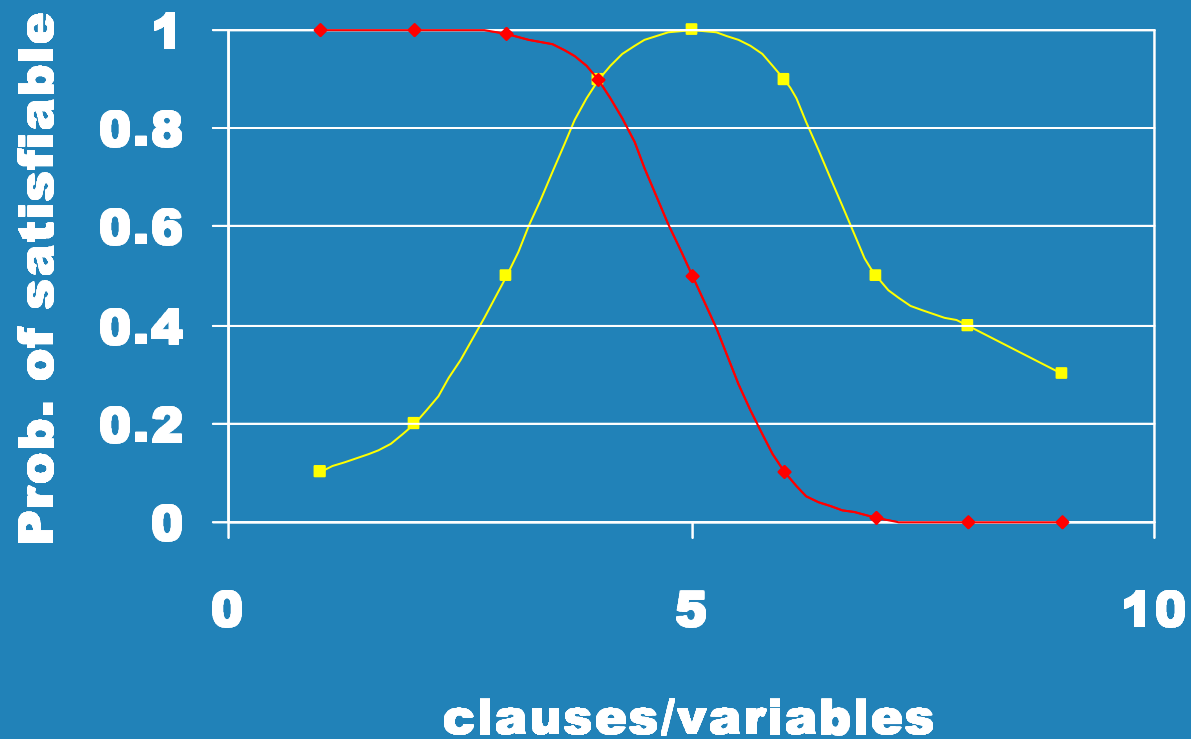
Variables: n

Clauses: m

Variables per clause: k

**Randomly generate all m clauses
by choosing k of the n variables
for each clause at random.
Randomly negate.**

Satisfiable Formulae



Explaining Results

Three sections

- $m/n < 4.2$, under constrained: nearly all satisfiable
- $m/n > 4.3$, over constrained: nearly all unsatisfiable
- m/n , under constrained: nearly all satisfiable



Phase Transition

Under-constrained problems are easy, just guess an assignment.

Over-constrained problems aren't too bad, just say "unsatisfiable" and try to verify via DPLL.

Transition region sharpens as n increases.



Local Search Discussion

Often the *second* best way to solve any problem.

Relatively easy to implement.

So, good first attempt. If good enough, stop.

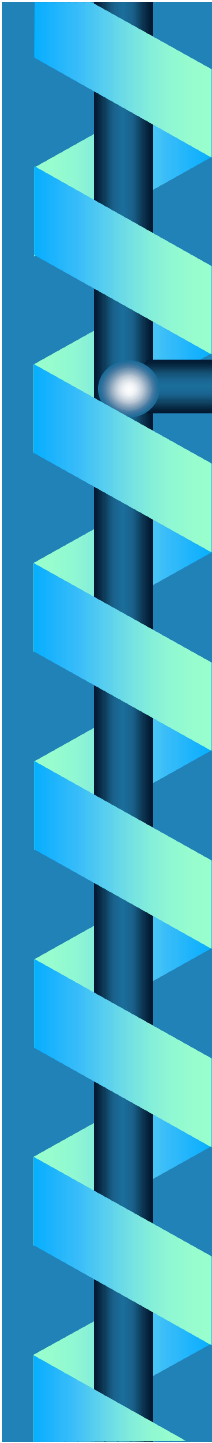


What to Learn

Definition of hill climbing, simulated annealing, genetic algorithm.

Definition of GSAT, WALKSAT.

Relative advantages and disadvantages of local search.



Programming Project 1

(due 10/22)

Solve Rush Hour problems.

Implement BFS, A* with simple blocking heuristic, A* with more advanced heuristic.

Input and output examples available on course web site

Speed and accuracy count.

Homework 4 (due 10/17)

- 1. In graph partitioning, call a *balanced* state one in which both sets contain the same number of nodes. (a) How can we choose an initial state at random that is balanced? (b) How can we define the neighbor set N so that every neighbor of a balanced state is balanced? (c) Show that standard GA crossover (uniform) between two balanced states need not be balanced. (d) Suggest an alternative crossover operator that preserves balance.**
- 2. More soon...**