

Myerson & Satterthwaite (4)
"Efficient Mechanisms for Bilateral Trading"

Ultimately simple situation:

1 seller, 1 buyer

value \tilde{v}_1 , value \tilde{v}_2

IPV

\tilde{v}_1 is distributed as density $f_1(\cdot) > 0$ on $[a, b]$

\tilde{v}_2

$f_2(\cdot) > 0$ on $[a_2, b_2]$

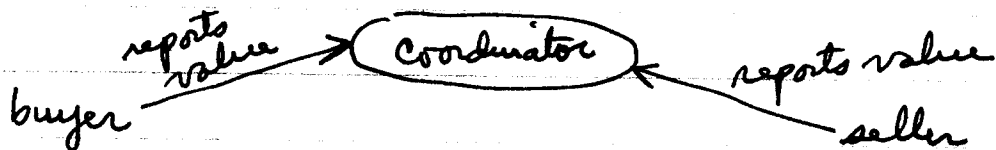
risk neutral

additively separable utility for money and object

Basic Question: Among all possible bargaining mechanisms, which have desirable economic efficiency properties?

Review of Revelation Principle:

A direct bargaining mechanism



• direct mechanism: characterization (p, x) , where

$p(v_1, v_2) = \text{prob. transfer } 1 \rightarrow 2$

$x(v_1, v_2) = \text{expected payment } 1 \rightarrow 2$

A direct mechanism is (Bayesian) incentive-compatible if honest reporting forms a BNE.

Revelation Principle

For every equilibrium in a bargaining game,
 \exists equivalent incentive compatible direct mechanism.

Proof if there were an incentive to lie when reporting then there would be an incentive to lie to oneself in the original game.

\therefore we can restrict attention to incentive-compatible direct mechanisms.

So we have (p, x) .

Some basic quantities:

$$\bar{x}_1(v_1) = \int_{a_2}^{b_2} x(v_1, t_2) f_2(t_2) dt_2 = E[\text{revenue of seller}]$$

$$\bar{p}_1(v_1) = \int_{a_2}^{b_2} p(v_1, t_2) f_2(t_2) dt_2 = \text{prob. of 1 selling to 2}$$

$$U_1(v_1) = \bar{x}_1(v_1) - v_1 \cdot \bar{p}_1(v_1) = E[\text{profit of seller, 1}]$$

& similarly,

$$U_2(v_2) = v_2 \cdot \bar{p}_2(v_2) - \bar{x}_2(v_2) = E[\text{profit of buyer, 2}]$$

In this notation, incentive-compatible means

$$U_1(v_1) \geq \bar{x}_1(\hat{v}_1) - v_1 \cdot \bar{p}_1(\hat{v}_1) \quad \forall v_1, \hat{v}_1 \in [a_1, b_1]$$

$$U_2(v_2) \geq v_2 \cdot \bar{p}_2(\hat{v}_2) - \bar{x}_2(\hat{v}_2) \quad \forall v_2, \hat{v}_2 \in [a_2, b_2]$$

Another desirable property:

A mechanism is individually rational iff

$$\begin{aligned} U_1(v_1) &\geq 0 & \forall v_1 \in [a_1, b_1], v_2 \in [a_2, b_2] \\ U_2(v_2) &\geq 0 \end{aligned}$$

And another desirable property:

A mechanism is ex post efficient iff

$$p(v_1, v_2) = \begin{cases} 1 & \text{if } v_1 < v_2 \\ 0 & \text{if } v_1 > v_2 \end{cases}$$

I.e. object is sold iff buyer values it more highly.

Main Result (Corollary 1)

If $[a_1, b_1] \cap [a_2, b_2] \neq \emptyset$ (overlap)

Then

no incentive-compatible individually rational trading mechanism can be ex post efficient!

Some details of proof technique:

first part (Theorem 1)

incentive-compatible

& individually rational $\Rightarrow U_1(b_1) + U(a_2) \geq 0$

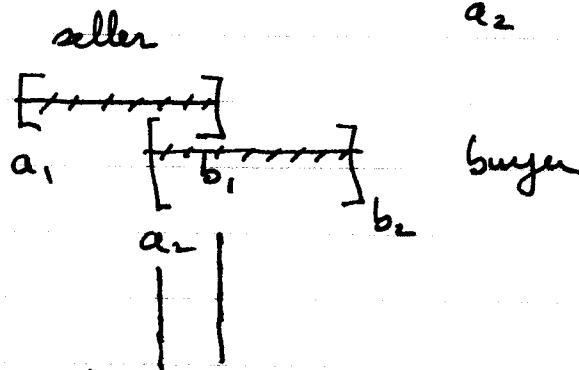
$$\begin{array}{ccc} \uparrow & & \uparrow \\ \text{min. expected} & & \text{min. expected} \\ \text{profit of seller} & & \text{profit of buyer} \end{array}$$

follows from definitions and crank turning.

second part

ex post efficient $\Rightarrow U_1(b_1) + U_2(a_2)$

$$= - \int_{a_2}^{b_1} [1 - F_2(t)] F_1(t) dt < 0$$



This quantity, $\int_{a_2}^{b_1} [1 - F_2(t)] F_1(t) dt$ is the smallest

lump-sum subsidy to create a Mechanism that is

- Incentive compatible
- individually rational
- ex post efficient

[cf the beginning of Vickrey 61!]

Example 1 Shows that it is necessary that $f_i > 0$ on respective intervals.

Counterexample Suppose

$$\text{prob}(\tilde{v}_1 = 1) = \text{prob}(\tilde{v}_1 = 4) = \text{prob}(\tilde{v}_2 = 0) = \text{prob}(\tilde{v}_2 = 3) = 1/2$$

seller f_1		↑	↑
buyer f_2		↓	↓
		1	4
		0	3

seller		buyer	
allot		2	3
1		1/4	(1/4)
4		1/4	1/4
		0	3

← profitable to both

Claim: the mechanism "sell at price 2 if both are willing, else no trade"

is incentive-compatible, individually rational, & ex post efficient.

- incentive-compatible honest reporting is a BNE (check)
- individually rational clearly $E[\text{profit}] > 0$
- efficient trade occurs ^{when &} only when $v_1 < v_2$.

Example 2 \tilde{v}_1, \tilde{v}_2 both uniform $[0, 1]$

Theorem becomes

$$U_1(b, 1) + U_2(a, 2) = \iint_{a_2, b_1}^{b_2, b_1} \left[v_2 - \frac{1 - F_2(v_2)}{f_2(v_2)} \right] - \left[v_1 + \frac{F_1(v_1)}{f_1(v_1)} \right] \cdot p(v_1, v_2) \cdot f_1(v_1) f_2(v_2) dv_1 dv_2$$

$$= \int_0^1 \int_0^1 \left\{ \left[v_2 - \frac{1 - v_2}{1} \right] - \left[v_1 + \frac{v_1}{1} \right] \right\} p(v_1, v_2) dv_1 dv_2$$

$$\text{or} \quad \int_0^1 \int_0^1 (v_2 - v_1 - \frac{1}{2}) p(v_1, v_2) dv_1 dv_2 \geq 0$$

$$\text{or} \quad \int_0^1 \int_0^1 (v_2 - v_1) p(v_1, v_2) dv_1 dv_2 - \frac{1}{2} \int_0^1 \int_0^1 p(v_1, v_2) dv_1 dv_2 \geq 0$$

$$\text{or} \quad E[v_2 - v_1 \mid \text{trade}] \geq \frac{1}{2}$$

divide \rightarrow
conditional
prob.

But in general

$$E[v_2 - v_1 \mid v_2 \geq v_1] = \frac{1}{3}$$

$E[\text{subsidy}]$ required for ex post efficiency = $\frac{1}{6}$.

check: (7) of MAS 83 $\int_{a_2}^{b_1} [1 - F_2(t)] F_1(t) dt$

$$= \int_0^1 (1-t)t dt = \frac{1}{6} \quad \checkmark$$
