

## Bulow & Klemperer, 94 "Rational Frenzies & Crashes"

Asset markets are volatile!

Common wisdom  $\rightarrow$  irrational behavior  
market imperfections

This paper offers a model (simple situation) in which rational behavior leads to "frenzies" and "crashes."

Key ideas use equilibrium strategy in auctions + the revenue equivalence theorem as the basis for buyer actions — in a dynamic setting.

WTP = "Willingness to Pay" changes from moment to moment & can change suddenly & drastically based on newly revealed information.

### The Model

$K$  identical units for sale

$K+L$  risk-neutral potential buyers, each wants a single unit (important)

IPVs, drawn from a distribution  $F(v)$  on  $[0, \bar{v}]$

Buyer derives surplus  $v-p$  from a purchase at price  $p$

## Dynamics

- ① Seller begins offering units at max price  $\bar{V}$  and lowers it until a purchase occurs, at price  $p$ .
- ② (NEW SALE) When a purchase occurs, every buyer gets an invitation to purchase 1 unit at price  $p$ . Either
  - (a) (FRENZY) all goods are sold at  $p$   $\rightarrow$  game ends
  - (b) (FRENZY) not all goods are sold at  $p$ , no one is left willing to buy at that price.  
 $\rightarrow$  then go to ① and continue lowering price until another NEW SALE takes place.
  - (c) (EXCESS DEMAND) More buyers want to buy at price  $p$  than there are units remaining.  
 Then if there are  $k+l$  bidders offering to buy ~~the~~ the remaining  $k$  units, we go to ① and restart the game with these  $k+l$  bidders competing for the remaining  $k$  units. All previous sales remain valid.

---

We restrict ourselves to symmetric equilibria in which bidders never bid more than their value.

---

## Solution to game    Apply auction theory!

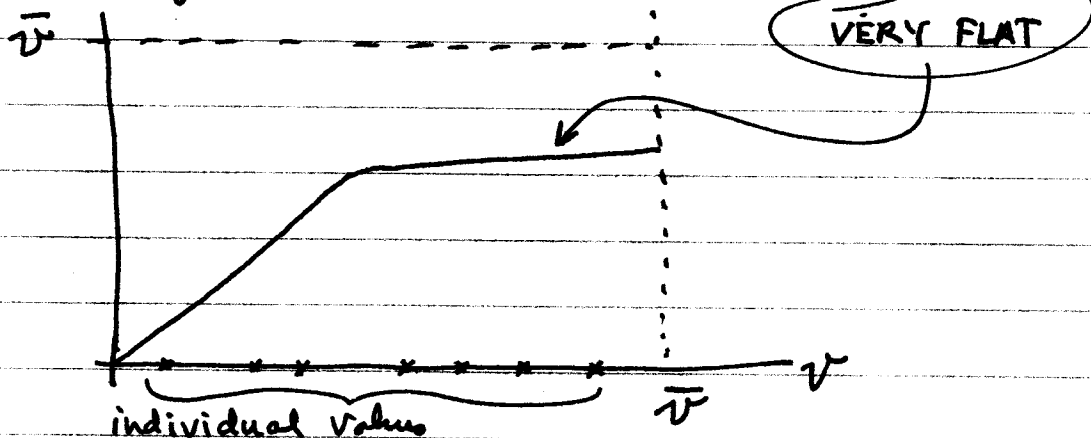
at any point,  $k = \#$  units remaining  
 $k+l = \text{total } \# \text{ of bidders remaining}$   
 ( $l=L$  unless a case (c) restart has occurred)

$\bar{v} =$  highest possible value of the remaining bidders, conditional on the bidders following equilibrium strategies.  
 $p =$  current asking price.

key fact:  $w(v) = E[\text{price a bidder would pay in a first-rejected-price auction} \mid \text{winning an object}]$

$$= E[(k+1)\text{st highest out of } (k+l) \text{ values} \mid (k+1)\text{st highest} \leq v]$$

Important feature of  $w(v)$ :



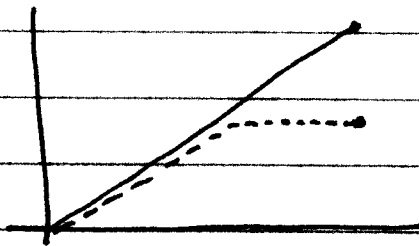
Why?  $\rightarrow$  in some range of  $v$ , buyers are fairly certain of ~~being~~ being above  $(k+1)$ st highest — doesn't change much in this range.

(Slight generalization of Riley & Samuelson et al.)  
Revenue Equivalence Theorem Any equilibrium in class  $\mathcal{L}$  has expected payment conditional on winning an item =  $w(v)$ .

Leads to Equilibrium Strategy: offer to buy if and only if  $w(v) \geq p$ .

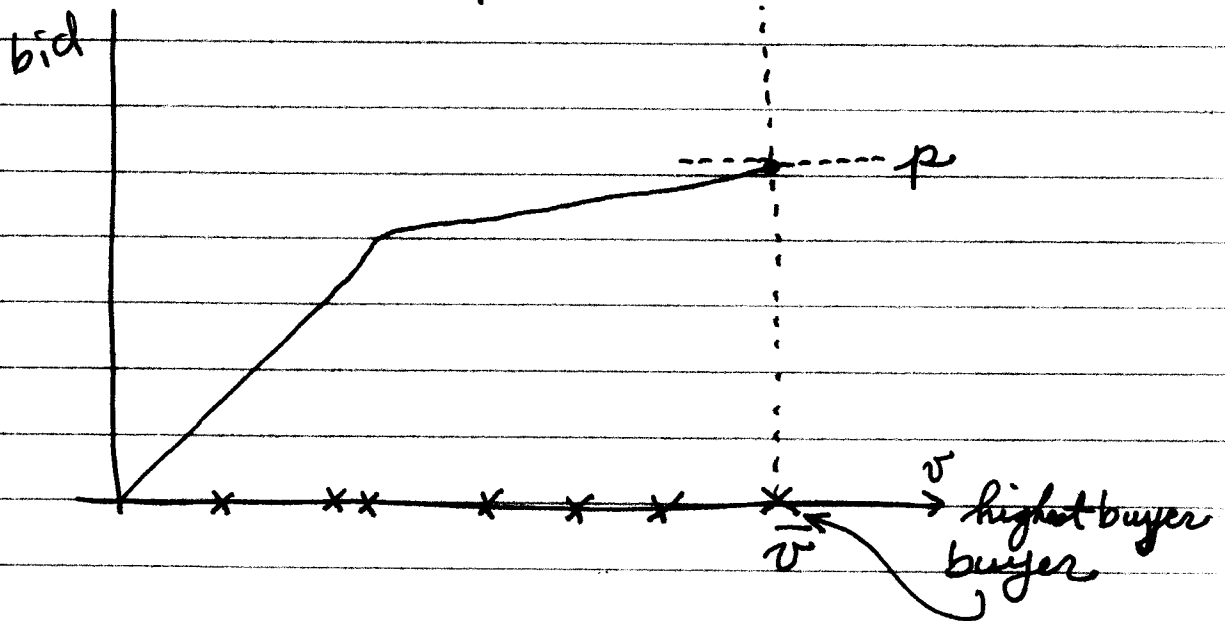
Hence to term  $w(v) = \text{WTP} = \text{willingness to Pay}$

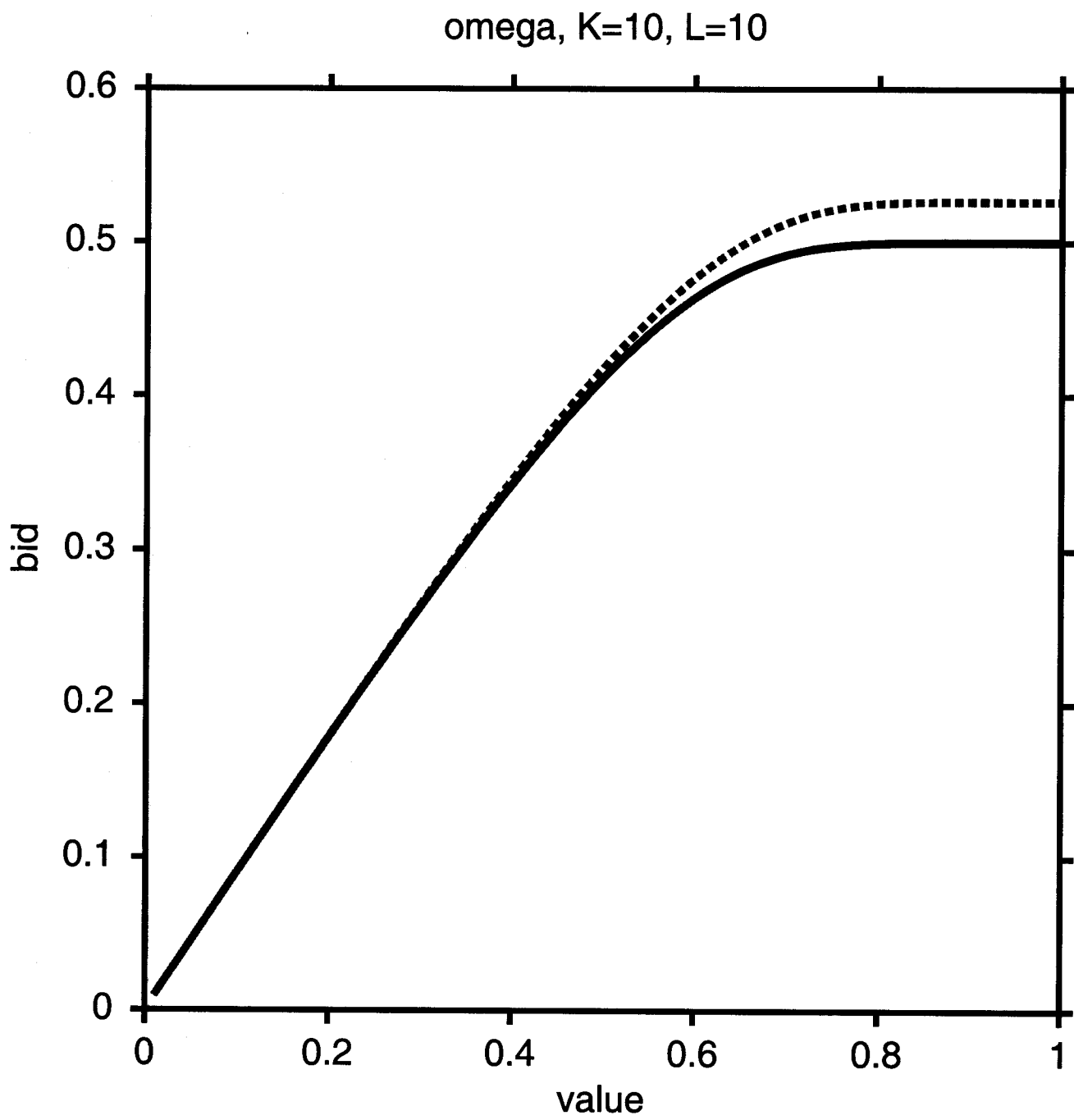
Contrast with willingness to accept take-it-or-leave-it final offers!  $\rightarrow$  just  $\hat{w}(v) = v \geq p$ .

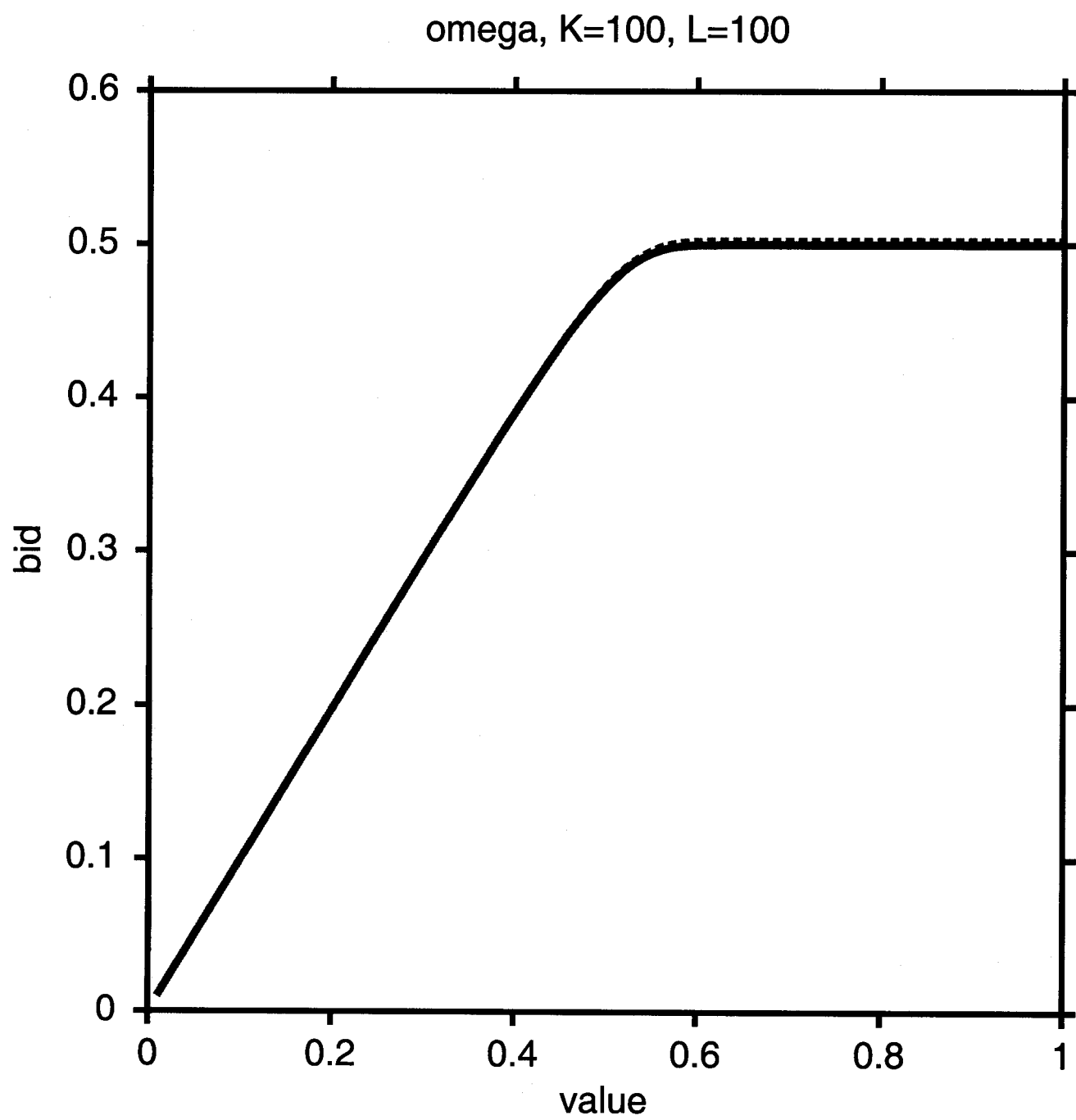


Dynamics of Frenzy: why do others jump in?

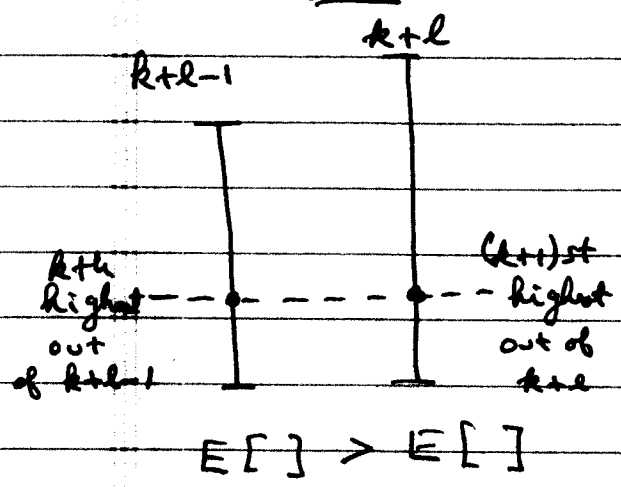
New Sale takes place at price  $p$



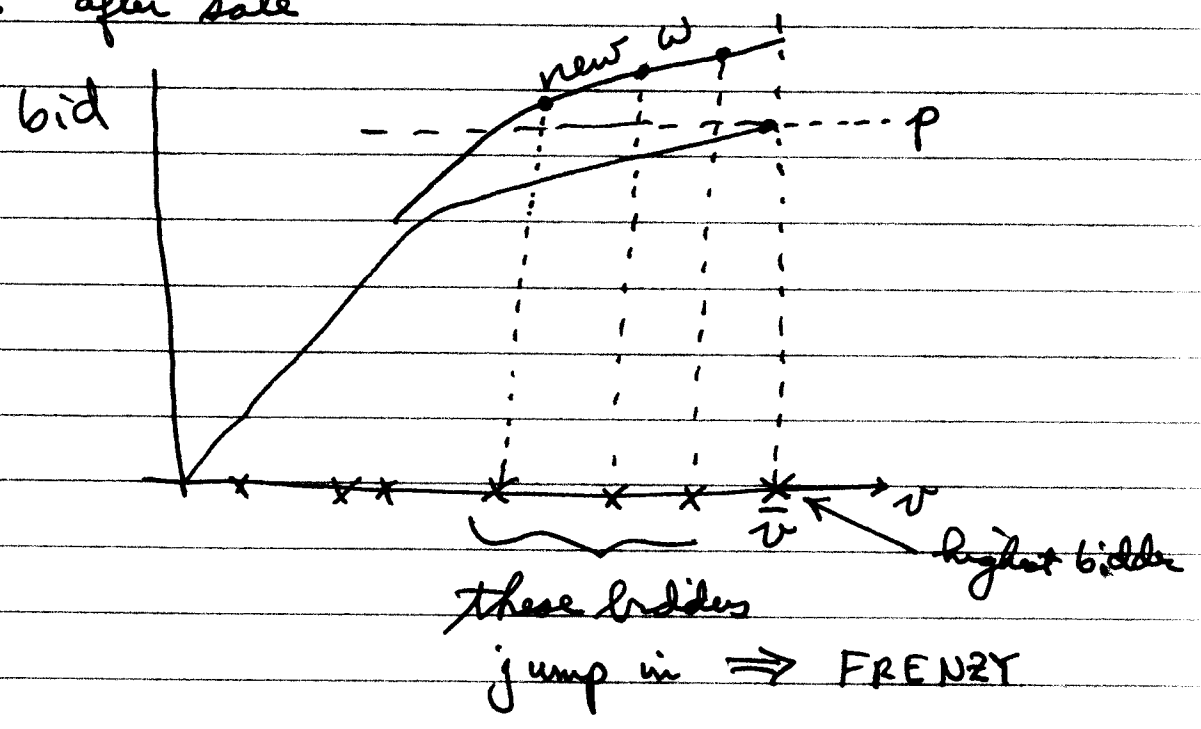




after a sale,  $k = k - 1$ .  
 this raises  $w(v) \forall v$ :



So after sale



If  $w$  is very flat, this can result in a large initial frenzy, set off by one purchase.

Corrected 4/18/01

— phone vjm & Bulow

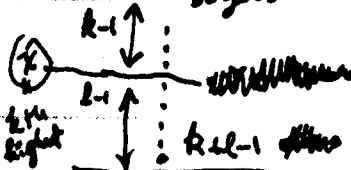
P.C

Derivation of  $w(v)$

prob. highest height is between  $x$  &  $x+dx$

$k$  items remain  
 $k+l = \text{total \# buyers}$

$= [f(x) \cdot dx] \text{ prob. [exactly } l \text{ are } > x]$



$= [f(x) \cdot dx] [\# \text{ choice for } l \text{ highest}] \cdot \text{prob. [a given set of } l \text{ are } > x \text{ \& a given set of } k-l \text{ are } < x]$

$= f(x) \cdot dx \cdot \frac{(k+l)(k+l-1) \dots (l)}{1} \cdot [1-F(x)]^{l-1} [F(x)]^{k-l-1}$

a check:  $l=1 \quad f(x) \cdot dx \cdot n \cdot [F(x)]^{n-1}$  ✓  
 $l=n \quad f(x) \cdot dx \cdot n(n-1) \dots [1-F(x)]^1 [F(x)]^{n-2}$  ✓

$n = k+l-1$

this gives  $g(x)$  if  $G(x)$  is prob. distr. of  $(k+l)$ st order statistic.

From this, the conditional expectation  $w$  is

$w(v) = E[\text{highest out of } k+l \text{ items} \mid \text{highest height} \leq v]$

$\int_0^v x \cdot g(x) dx$

$\text{prob. [highest height} \leq v]$

$\int_0^v x \cdot g(x) dx / \int_0^v g(x) dx$

$\int_0^v x \cdot f(x) \cdot [1-F(x)]^{l-1} [F(x)]^{k-l-1} dx$

$\int_0^v f(x) [1-F(x)]^{l-1} [F(x)]^{k-l-1} dx$

the highest out of other  $k+l-1$   
 $\leq v \Rightarrow v$  is among the highest

this  $k+l$  highest becomes  $(k+l)$ st highest



B&K has  $[1-F(x)]^{k-1} \dots$  ? I'll use that form

In uniform case  $F(x) = x$ ,  $f(x) = 1$ .

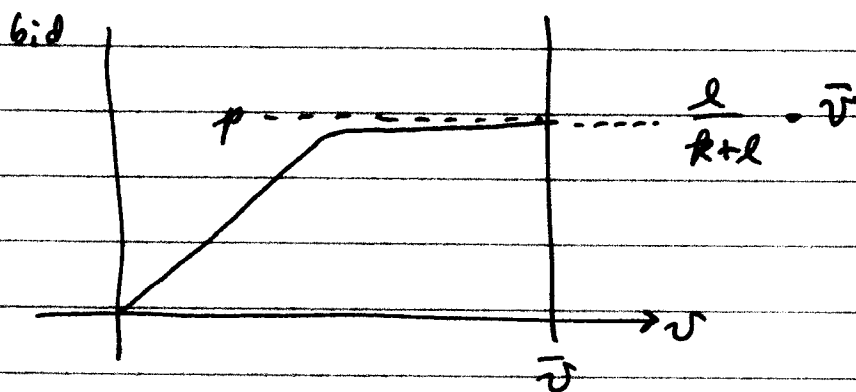
These integrals  $\int_0^{\bar{v}} [1-x]^{k-1} x^{l-1} dx$

are called incomplete  $\beta$  funcs. — go back to Laplace, talked by Perisio.

when  $v = \bar{v} = 1$ ,

$$w(\bar{v}) = w(1) = \frac{l}{k+l}$$

So this looks like



$w$  is really a func.  $w(l, k, \bar{v}, v)$ .  
 Easy to verify that  
 $w(l, k, \bar{v}, v) = \bar{v} \cdot w(l, k, 1, v/\bar{v})$

TABLES OF  
THE INCOMPLETE  
BETA-FUNCTION

*Originally prepared under the  
direction of and edited by*

KARL PEARSON

SECOND EDITION

*with a new Introduction by*

E.S. PEARSON *and* N.L. JOHNSON

CAMBRIDGE

*Published for the Biometrika Trustees*

AT THE UNIVERSITY PRESS

1968

sec per evaluation

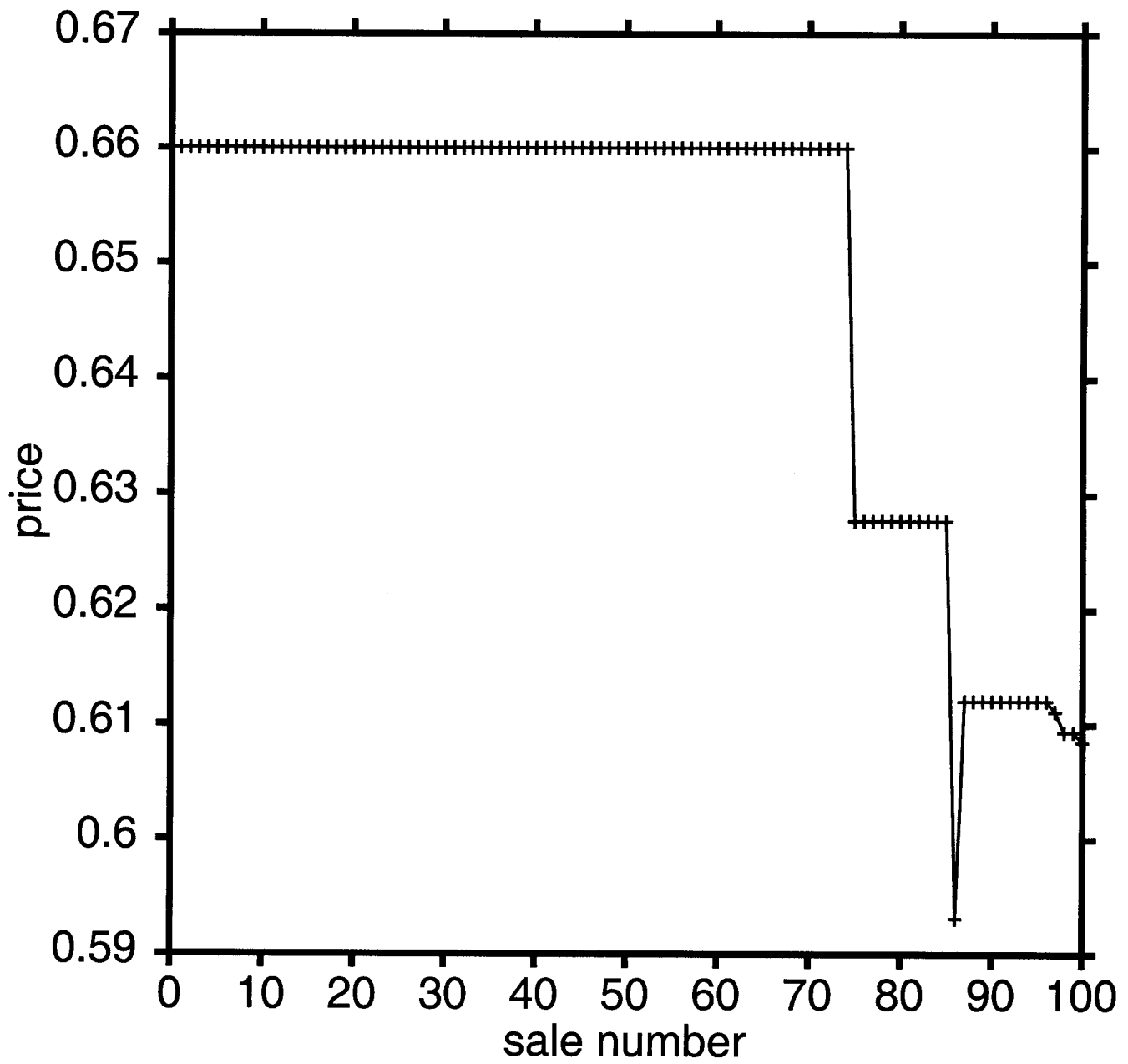
a	b	x	Simpson's Rule
---	---	---	-------------------

5	5	0.1	0.5e-2
10	10	0.2	0.9e-2
100	100	0.2	0.15
200	200	0.3	0.17
300	300	0.1	0.62e-1
300	300	0.2	0.44
300	300	0.3	0.4
400	100	0.3	0.9
400	100	0.1	---

sec per evaluation

a	b	x	cont. frac.
5	5	0.1	0.7e-5
10	10	0.2	0.8e-5
100	100	0.2	0.8e-5
200	200	0.3	0.9e-5
300	300	0.1	1.1e-5
300	300	0.2	0.7e-5
300	300	0.3	0.9e-5
400	100	0.3	0.7e-5
400	100	0.1	0.6e-5

K=100 L=200 SEED=100



X 96 13 3 0.6171 0.6119 0.6120  
size = 9 --> fillable positive demand

FRENZY, round 1  
f 87 13 3 0.6346 0.6124 0.6120  
...  
f 94 6 3 0.6225 0.6124 0.6120  
f 95 5 3 0.6221 0.6123 0.6120  
END FRENZY, round 1  
new vsup after round 1 0.6180  
after round 1, p is above right asymptote

X 96 4 3 0.6171 0.6115 0.6120  
size = 0

s 96 4 3 0.6171 0.6111 0.6111

vtilde after single sale= 0.6136  
X 97 3 3 0.6121 0.6103 0.6111  
size = 0

s 97 3 3 0.6121 0.6094 0.6094

vtilde after single sale= 0.6106  
- 98 2 3 0.6111 0.6097 0.6094  
X 99 2 3 0.6091 0.6084 0.6094  
size = 1 --> fillable positive demand

FRENZY, round 1  
f 98 2 3 0.6111 0.6097 0.6094  
END FRENZY, round 1  
new vsup after round 1 0.6106  
new vtilde after round 1 0.6103

X 99 1 3 0.6091 0.6085 0.6094  
size = 0

s 99 1 3 0.6091 0.6085 0.6085

K items sold, sale over  
average price per item= 0.648998

=====  
K= 100 L= 200 random seed= 100  
transaction code c: s = single sale, r = frenzy, round r, e excess demand  
c i k l vi pi p  
s 0 100 200 0.9901 0.6600 0.6600

vtilde after single sale= 0.7094  
- 1 99 200 0.9871 0.6623 0.6600  
...  
- 73 99 200 0.7179 0.6611 0.6600  
X 74 99 200 0.7093 0.6600 0.6600  
size = 73 --> fillable positive demand

FRENZY, round 1  
f 1 99 200 0.9871 0.6623 0.6600  
...  
f 72 28 200 0.7235 0.6616 0.6600  
f 73 27 200 0.7179 0.6611 0.6600  
END FRENZY, round 1  
new vsup after round 1 0.7094  
after round 1, p is above right asymptote

X 74 26 200 0.7093 0.6278 0.6600  
size = 0

s 74 26 200 0.7093 0.6277 0.6277

vtilde after single sale= 0.6486  
- 75 25 200 0.7073 0.6305 0.6277  
...  
- 84 25 200 0.6538 0.6291 0.6277  
X 85 25 200 0.6376 0.6226 0.6277  
size = 10 --> fillable positive demand

FRENZY, round 1  
f 75 25 200 0.7073 0.6305 0.6277  
...  
f 84 16 200 0.6538 0.6291 0.6277  
END FRENZY, round 1  
new vsup after round 1 0.6486  
after round 1, p is above right asymptote

X 85 15 200 0.6376 0.6033 0.6277  
size = 0

s 85 15 200 0.6376 0.5931 0.5931

vtilde after single sale= 0.6067  
- 86 14 200 0.6372 0.5959 0.5931  
...  
- 102 14 200 0.6068 0.5931 0.5931  
X 103 14 200 0.6038 0.5917 0.5931  
size = 17 --> excess demand, new vund = 6.066653e-01

s 86 14 3 0.6372 0.6120 0.6120  
price increase after excess demand, seed= 100, 87/100 items sold  
vtilde after single sale= 0.6180  
- 87 13 3 0.6346 0.6124 0.6120  
...  
- 95 13 3 0.6221 0.6123 0.6120



10000 runs, price per item  
 k 1 avg predicted

*Tests of Revenue Equiv.*

-----

1	1	0.333347	0.333333
1	2	0.500026	0.500000
2	2	0.400157	0.400000
1	3	0.600616	0.600000
2	3	0.499878	0.500000
1	4	0.666849	0.666666
20	10	0.322530	0.322580
30	10	0.243460	0.243902
40	5	0.108449	0.108695
5	40	0.869741	0.869565
20	30	0.588372	0.588235
1	100	0.980528	0.980392
100	100	0.497485	0.497512
200	100	0.332154	0.332225

RevCheck checks total expected revenue in uniform case against what is known from revenue equivalence theorem: that  $E[\text{revenue}]$  in equilibrium = same as SP auction =  $E[(K+L)\text{st highest of } K+L \text{ bidders}]$ . When  $K=1$  object and  $L=1$ , we have  $K+L=2$  bidders and 1 object, and  $(L+1)\text{st highest out of } 3$  has exp. value  $1/3$ . In general,  $L = \text{no. of bidders} - \text{no. of objects}$ , and the exp. revenue is  $L/(K+L+1)$ .



$$\begin{aligned}
 E[\text{price}] &= \frac{(k+l)!}{k(l-1)!} \int_0^1 \left[ \int_0^v (1-x)^{k-1} x^l dx \right] dv \\
 &= \frac{(k+l)!}{k(l-1)!} \left\{ \left[ \int_0^v (1-x)^{k-1} x^l dx \right] \cdot v \Big|_0^1 - \int_0^1 v^{l+1} (1-x)^{k-1} dv \right\} \\
 &= \frac{(k+l)!}{k(l-1)!} \left[ \int_0^1 (1-x)^{k-1} x^l dx - \int_0^1 v^{l+1} (1-x)^{k-1} dv \right]
 \end{aligned}$$

$$\left[ \int_0^1 t^{z-1} (1-t)^{w-1} dt = B(z, w) = \frac{\Gamma(z)\Gamma(w)}{\Gamma(z+w)} \right]$$

$$= \frac{(z-1)!(w-1)!}{(z+w-1)!}$$

$$\begin{aligned}
 E[\text{price}] &= \frac{(k+l)!}{k(l-1)!} \left[ \frac{(k-1)! l!}{(k+l+1-1)!} - \frac{(k-1)!(l+1)!}{(k+l)!} \right] \\
 &= \frac{(k+l)!}{k(l-1)!} \left[ \frac{(k-1)! l!}{(k+l)!} - \frac{(k-1)(l+1)!}{(k+l)!} \right] \\
 &= \frac{(k+l)!}{k(l-1)! (k+l)!} \left[ l! - \frac{l+1}{k+l+1} \right] \\
 &= \frac{l}{k} \left[ 1 - \frac{l+1}{k+l+1} \right] = \frac{l}{k} \left[ \frac{k+l+1-l-1}{k+l+1} \right] = \frac{l}{k+l+1}
 \end{aligned}$$

check

$k=1, l=1$	$\frac{1}{3} \checkmark$
$k=1, l=2$	$\frac{2}{4} \checkmark$
$k=2, l=2$	$\frac{2}{5} \checkmark$
$k=1$	$\frac{l}{l+2} \checkmark$
$k=2, l=3$	$\frac{3}{6} = \frac{1}{2} \checkmark$

(OK)